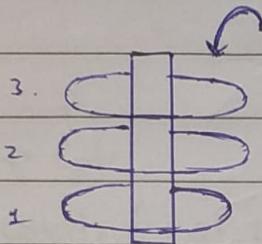


— X — X — X —

Stack.



→ Can insert at only 1 position.

→ Can access only top most disk.

→ To access any disk, I need to remove from top.

LIFO :- Last - In - ~~First~~ - Out.

FIFO :- First - In - Last - Out.

⇒ Operation on stack.

① Push :- st.push(x). → (To insert an ele in stacks)

② Pop :- st.pop(). → (To remove & print top most ele)

③ Peek :- st.peek(). → (Only to print topmost ele)

4
 3
 2
 1

st.push(1)
 st.push(2)
 st.push(3)
 st.push(4).

st.pop() → 4.
 st.peek() → 3,
 print :- 1 2 3
 st.pop() → 3.

M	T	W	T	F	S	S
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Print :- 1 2.

st.size() → 2,

→ STL for stack [Unlimited size]

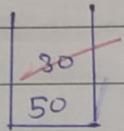
Stack <Data-type> st = new Stack<>();

e.g. Stack <Integer> st = new Stack<>();

st.push(50);

st.push(30);

sout(st.pop()); → 30



False ← st.isEmpty() → Whether stack is empty or not.

Return type is Boolean [True or False].

T.C :- O(N) → To access q^{th} element.

Comparing Arrays, Linked List & Stacks

Ques Method :-

T.C

S.C.

Array s.

O(1)

O(1)

Linked List

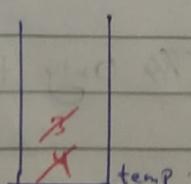
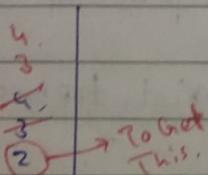
O(n).

O(1)

Stack

O(n)

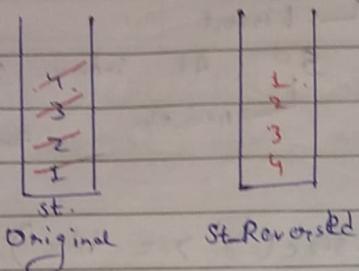
O(n).



temp.push(st.pop());
 temp.push(st.pop());
 sout(st.peek());

Reverse A Stack.

M	T	W	T	F	S	S
Page No.:	YOUVA					
Date:						



T.C :- $O(n)$

→ We pop ele from Original stack & push in reverse stack.

S.C :- $O(n)$.

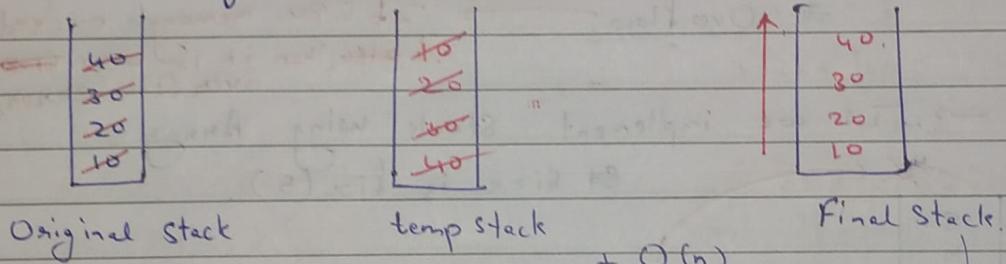
→ for auxillary stack space.

→ Code :-

```

while (!st.isEmpty()) {
    st_reversed.push(st.pop());
}
Sout(st_reversed);
    
```

Copy stack into another stack in Same Order.



$$\begin{aligned} \text{T.C} &:= O(n) + O(n) + O(n) \Rightarrow O(n) \\ \text{S.C} &:= O(n) + O(n) \Rightarrow O(n) \end{aligned}$$

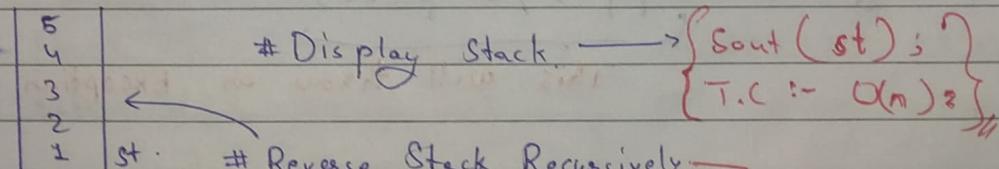
→ Code :-

```

while (!st.isEmpty()) {
    st_temp.push(st.pop());
}
    
```

```

while (!st_temp.isEmpty()) {
    st_final.push(st_temp.pop());
}
    
```



→ Code :-

```

psv displayReverse(Stack<Integer> st) {
    if (st.isEmpty()) return;
    int top = st.pop();
    Sout(top + " ");
    displayReverse(st);
    Sout(top + " ");
    st.push(top);
}
    
```

→ T.C :- $O(n)$;

→ S.C :- $O(n)$,

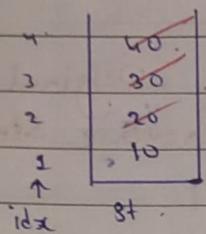
Recursive stack space.

// To print in Reverse Order {5, 4, 3, 2, 1}

// To print as it is {1, 2, 3, 4, 5}.

~~# Push Ele At Bottom / Any Index.~~

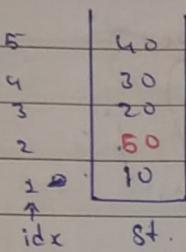
M	T	W	T	F	S	S
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st_temp .

\hookrightarrow List & insert

50 At index 2
new_ele. \downarrow
~~idx.~~



$$T.C = O(n)$$

$$O(n)$$

S.C = O(1). $out(st);$

\Rightarrow Code:-

while($st.size() \geq idx$) {

$st_temp.push(st.pop());$

}

$st.push(new_ele);$

 while($!st_temp.isEmpty()$) {

$st.push(st_temp.pop());$

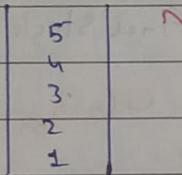
 }

~~# Overflow~~ \rightarrow If CPU memory exceeds

~~then it is also shown~~ overflow error,

\rightarrow If we implement Stack using Array,

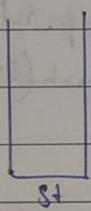
~~size is fixed(5).~~



$st.push(6);$

\rightarrow Overflow Condition //

Underflow



\rightarrow if Stack is empty & then.

$st.pop()$ / $st.peek()$

$\underbrace{\hspace{10em}}$ This will throw an Exception.

Array Implementation.

=> Code :-

```
public static class Stack {
    private int[] arr = new int[5];
    int n = arr.length;
    private int idx = 0;
```

// Push Method.

```
void push(int x) {
    if (isFull()) {
        cout("Stack is Full");
        return;
    }
    int top = arr[idx];
    arr[idx] = x;
    idx++;
}
```

// Peek Method.

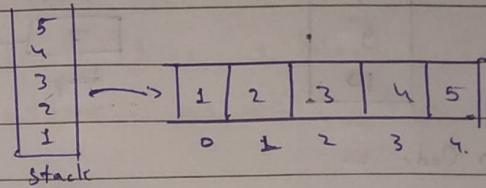
```
int peek() {
    if (idx == 0) {
        cout("The Stack is Empty");
        return -1;
    }
    return arr[idx - 1];
}
```

// Size Method.

```
int size() {
    return idx;
}
```

// display Any Idx ele.

```
int displayEle(int idx) {
    if (x < 0 || x > n) {
        cout("Error in Index..");
        return -1;
    }
    return arr[idx - 1];
}
```



// Pop Method.

```
int pop() {
    if (isEmpty()) {
        cout("Stack is Empty");
        return -1;
    }
    int top = arr[idx - 1];
    arr[idx - 1] = 0;
    idx--;
    return top;
}
```

// Display Method.

```
void display() {
    for (int i = 0; i < idx; i++) {
        cout(arr[i] + " ");
    }
    cout();
}
```

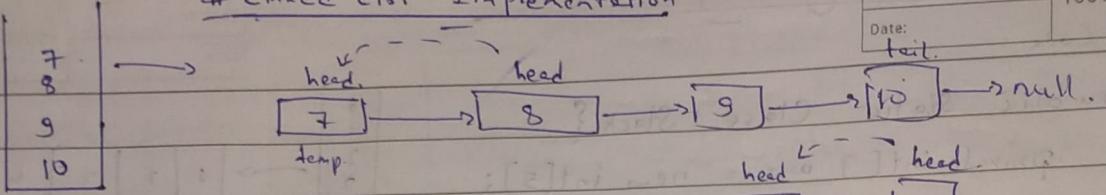
// isEmpty Method.

```
boolean isEmpty() {
    if (idx == 0) return true;
    else return false;
}
```

// isfull Method.

```
boolean isFull() {
    if (n == idx) return true;
    return false;
}
```

Linked List Implementation



→ Code :-

```
public static class Node {
    int val;
    Node next;
    Node (int val) {
        this.val = val;
    }
}
```

```
public static class Stack {
```

```
    Node head;
```

```
    Node tail;
```

```
    int size = 0;
```

// Push Method.

```
void push(int x) {
```

```
    Node temp = new Node(x);
```

```
    if (size == 0) {
```

```
        head = tail = temp;
```

```
        size++;
```

```
}
```

```
else {
```

```
    temp.next = head;
```

```
    head = temp;
```

```
    size++;
```

```
}
```

```
}
```

// display Method

```
void display() {
    displayRec(head);
}
```

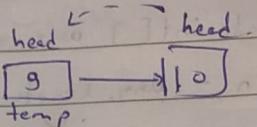
```
void displayRec(Node h) {
```

```
    if (h == null) return;
```

```
    displayRec(h.next);
```

```
    sout(h.val + " ");
```

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Date:	tail.					



- (1) Create temp Node.
- (2) Point temp to head.
- (3) make temp as head.
- (4) Size++;

// Pop Method

```
int pop() {
    if (head == null) {
        sout("Stack is empty...");
```

```
    return -1;
```

```
    int x = head.val;
```

```
    head = head.next;
```

```
    size--;
```

```
    return x;
```

// Peek Method

```
int peek() {
```

```
    if (head == null) {
```

```
        sout("Stack is empty...");
```

```
    return -1;
```

```
    return head.val;
```

// isEmpty() Method

```
boolean isEmpty() {
```

```
    if (size == 0) return true;
```

```
    return false;
```

// Size Method

```
int size() {
```

```
    return size;
```

Linked-List V/s Array Implementation.

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→ Advantages of Array Reprn.

(1) Size :- for every ele space taken by is 1 block. In LL, we have to store address as well which takes ~~more~~ more memory.

(2) Display :- S.C :- $O(1)$.

→ DisAdvantages :- (1) Size is fixed (Overflow).

→ Advantages of LL Reprn:- (1) Size is Unlimited.

→ Dis Advantages :-

(1) ~~size of~~ memory taken to store data is more.

(2) Display :- S.C = $O(n)$ [Recursive Space]

Infix Expression.

$$9 - 5 + 3 * 4 / 6.$$

$$9 - 5 + \frac{12}{6}$$

$$\underbrace{9 - 5}_{4} + \underbrace{2}_{2}$$

DODMAS

D, M > A, S

left to Right

$$\begin{aligned} & 3 * 4 / 2 \\ & 12 / 2 \\ & \Rightarrow 6 \end{aligned}$$

⇒ Evaluate Infix Using Stacks:-

→ Rules :- we need 2 stacks.

- Number stack. (Single digit numbers).
- Operator stack

(1) If character is a number, push it in numbers stack.

(2) If character is operator

(a) If operator stack is empty, then push it

(b) If op stack is not empty, we see operator on top of op stack. If its precedence is \geq current operator, we first complete operation for operator on top of stack & then push the current operator.

$$9 - 5 + 3 * 4 / 6$$

Number	operator
9	-
5	+
3	*
4	/
6	

$V_2 = \text{numst.pop}()$

M	T	W	T	F	S	S
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$V_1 = \text{numst.pop}()$

$OP = \text{opst.pop}()$

$\text{resout} = V_1 \text{ OP } V_2$

$V_2 = 5 \quad OP = -$

$V_1 = 3$

$\text{resout} = 9 - 5 = 4 //$

$V_2 = 4 \quad OP = *$

$V_1 = 3 \quad \text{res} = 4 * 3 = 12 //$

Till opstack is not empty {

$V_2 = \text{numst.pop}()$

$V_1 = \text{numst.pop}()$

$OP = \text{opst.pop}()$

$x = V_1 \text{ OP } V_2$

$\text{numst.push}(x);$

$V_2 = 6 \quad OP = /$

$V_1 = 12$

$\text{res} = 12 / 6 = 2 //$

$V_2 = 2 \quad OP = +$

$V_1 = 4$

$\text{res} = 2 + 4 = 6 //$

} return numst.peek();

~~Stack~~

Infix with Brackets.

$St = "9 - (5 + 3) * 4 / 6".$

\Rightarrow

$\Rightarrow 9 - 8 * 4 / 6.$

\Rightarrow

$\Rightarrow 9 - 32 / 6.$

\Rightarrow

$\Rightarrow 9 - 5 \rightarrow 4 //$

$4 \rightarrow Ans.$

Extra.

if (op.size == 0 || ch == '(')

|| ~~ch == op.peek() == '('~~

op.push(ch);

if (ch == ')') {

while (op.peek() != ')') {

work.

num	OP
9	-
5	
3	
8	*
3	
5	/
8	

$V_2 = 3 \quad OP = +$

$V_1 = 5 \quad x = 3 + 5 = 8 //$

$V_2 = 4 \quad OP = *$

$V_1 = 8 \quad x = 4 * 8 = 32 //$

$V_2 = 5 \quad OP = -$

$V_1 = 9$

$x = 9 - 5 = 4 //$

$V_2 = 6 \quad OP = /$

$V_1 = 32$

$x = 32 / 6 = 5 //$

=> Code :- psym() {

String str = "g-(5+3)*4/6";

Stack<Integer> val = new Stack<>();

Stack<Character> op = new Stack<>();

for (int i = 0; i < n; i++) {

 char ch = str.charAt(i);

 int ascii = (int) ch;

 if (ascii >= 48 && ascii <= 57) {

 val.push(ascii - 48);

 else if (op.size() == 0 || ch == '(' || op.peek() == '(')

 op.push(ch);

 else if (ch == ')') {

 while (op.peek() != '(') {

 int v2 = val.pop();

 int v1 = val.pop();

 if (op.peek() == '-') val.push(v2 - v1);

 if (op.peek() == '+') val.push(v1 + v2);

 if (op.peek() == '*') val.push(v1 * v2);

 if (op.peek() == '/') val.push(v1 / v2);

 op.pop();

}

 op.pop();

 → while (val.size() > 1) {

 // Work

}

 else {

 if (ch == '+' || ch == '-') {

 // Work

 op.push(ch);

 return val.peek();

 else if (ch == '*' || ch == '/') {

 if (op.peek() == '*' || op.peek() == '/') {

 // Work. No Need to add '+' & '-' Lines

 op.push(ch);

 else op.push(ch);

→ Continue .

Prefix Expression

Infix : $a * b \rightarrow 1 * 2 * 3$

Prefix : $* ab \rightarrow * 1 2 3$

Postfix : $ab * \rightarrow 1 2 3 *$

$$g = 5 + 3 * 4 / 6.$$

$$g = 5 + \{ 3 * 4 \} / 6$$

$$V_2 = 4 \quad op = *$$

$$V_2 = 6 \quad op = /$$

$$V_1 = 3 \quad * 34$$

$$V_1 = * 34 \quad / * 349 \quad g - 5 + / * 346$$

$$V_2 = 5 \quad op = -$$

$$V_2 = / * 346 \quad op = + \quad - 95 \oplus / * 346.$$

$$V_1 = 8 \quad - 95$$

$$V_1 = - 95 \quad + - 95 / * 346$$

$$\text{Op} + V_1 + V_2$$

$$+ - 95 / * 346 \rightarrow \text{Ans.}$$

Postfix Expression

Infix : $V_1 \text{ Op } V_2$

$$9 - (5 + 3) * 4 / 6.$$

Prefix : $\text{Op } V_1 \text{ } V_2$

$$9 - 53 + * 4 / 6.$$

Postfix : $V_1 \text{ } V_2 \text{ Op.}$

$$9 - \underline{\underline{53 + 4}} * 6 /$$

$$V_2 = 3 \quad op = +.$$

$$9 - 53 + 4 * 6 /$$

$$V_1 = 5 \quad \Rightarrow 53 +$$

$$953 + 4 * 6 / -$$

$$V_2 = 4 \quad op = *$$

$$V_2 = 6$$

$$\rightarrow \text{Ans.}$$

$$V_1 = 53 + \Rightarrow 53 + 4 *$$

$$V_1 = 53 + 4 * \quad op = /$$

$$\Rightarrow 53 + 4 * 6 /$$

$$V_2 = 53 + 4 * 6 /$$

$$V_1 = 9 \quad op = -$$

$$\Rightarrow 953 + 4 * 6 / - \rightarrow \text{Ans.}$$

\Rightarrow Code :- Psym()

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for (int i = 0; i < n; i++) {

 char ch = str.charAt(i);

 int ascii = (int) ch;

 if (ascii >= 48 && ascii <= 57) {

 String s = "" + ch;

 val.push(s);

}

 else if (op.size() == 0 || ch == '(' || op.peek() == '(') {

 op.push(ch);

}

 else if (ch == ')') {

 while (op.peek() != '(') {

 String v2 = val.pop();

 String v1 = val.pop();

 char o = op.pop();

 String t = o + v1 + v2; }

 val.push(t); }

}

 else {

 if (ch == '+' || ch == '-') {

 Work;

 op.push(ch);

 else if (ch == '*' || ch == '/') {

 if (op.peek() == '*' || op.peek() == '/') {

 Work;

 op.push(ch);

 } else op.push(ch); }

}

 while (val.size() > 1) {

 Work;

 return val.peek(); }

Evaluation of Postfix Expressions.

M	T	W	T	F	S	S
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str = " 9 5 3 + 4 * 6 / - ";

4	Ans.	$v_1, 0 \vee v_2$	$\{ i \text{ from } 0 \text{ to } n-1 \}$
5		$v_2 = 3$	
6		$v_1 = 5 \Rightarrow 5+3 \Rightarrow 8$	
32			
4			
8		$v_2 = 4$	$v_2 = 6$
3		$v_1 = 8$	$v_1 = 32 \Rightarrow 32/6 = 5$
5			
9		$0 = *$ $\Rightarrow 8 * 4 \Rightarrow 32$	$0 = /$

Stack:

$$v_2 \quad v_2 = 5 \quad 0 = - \\ v_1 = 8 \quad \Rightarrow 8 - 5 = 3$$

Evaluation of Prefix Expressions

str = "- 9 /* + 5 3 4 6";

9	5	32	8	4	6	0 v_1, v_2	$\{ i \text{ from } n-1 \text{ to } 0 \}$

\Rightarrow Code :- psym() {

String str = "9 5 3 + 4 * 6 / -";

Stack<Integer> val = new Stack<>();

for (int i = 0 ; i < n ; i++) {

char ch = str.charAt(i);

int ascii = (int) ch;

if (ascii >= 48 & & ascii <= 57) : val.push(ascii - 48);

else {

int v2 = val.pop();

int v1 = val.pop();

if (ch == '+') val.push(v1 + v2);

if (ch == '-') val.push(v1 - v2);

if (ch == '*') val.push(v1 * v2);

if (ch == '/') val.push(v1 / v2);

}

sout(val.peek());

for Prefix Evaluation

Changes :-

\rightarrow for (int i = n-1 to i >= 0)

\rightarrow first v1 then v2

\rightarrow first v2 then v1

Conversion

(1) Prefix to Postfix Conversion :-

$$\text{Stg} = " - 9 /* + 5346 " \quad \longleftrightarrow \quad V_1 V_2 \text{ Op}$$

953+4*6/-

~~9~~
~~53+4*6~~

~~53+4*~~

~~53+~~

~~5~~

~~3~~

~~4~~

~~6~~

$$V_1 = 5 \quad \text{Op} = +$$

$$V_2 = 3 \Rightarrow 53 +$$

$$V_1 = 53 + 4 * \quad \text{Op} = /$$

$$V_2 = 6 \Rightarrow 53 + 4 * 6 /$$

$$V_1 = 53 + \quad \text{Op} = *$$

$$V_2 = 4 \Rightarrow 53 + 4 *$$

$$V_1 = 9 \quad \text{Op} = -$$

$$V_2 = 53 + 4 * 6 /$$

$$\Rightarrow 953 + 4 * 6 / -$$

String Stack

Ans.

(2) Prefix to Infix Conversion :-

$$\text{Stg} = " - 9 /* + 5346 " \quad \longleftrightarrow \quad V_1 \text{ Op } V_2$$

(9 - ((5+3)*4)/6))

~~((5+3)*4)/6)~~

~~((5+3)*4)~~

~~(5+3)~~

~~5~~

~~3~~

~~4~~

~~6~~

$$V_1 = 5 \quad \text{Op} = +$$

$$V_2 = 3 \Rightarrow (5 + 3) \quad V_1, \text{Op}, V_2$$

$$V_1 = ((5+3)*4)$$

$$V_2 = 6 \quad \text{Op} = /$$

$$\Rightarrow ((5+3)*4) / 6$$

$$V_1 = (5+3) \quad \text{Op} = *$$

$$V_2 = 4 \Rightarrow ((5+3)*4)$$

$$V_1 = 9 \quad \text{Op} = -$$

$$V_2 = (((5+3)*4) / 6) \Rightarrow (9 - ((5+3)*4) / 6))$$

\Rightarrow Code :- # Prefix to Postfix

```
= for (int i = n-1; i >= 0; i--) { for Prefix to
    |   char ch = Stg.charAt(i); Infix.
    |   int ascii = (int) ch;
    |   if (ascii >= 48 & & ascii <= 57) { Just change
    |       |   val.push(ch + ")");
    |   }
    |   else {
    |       |   Stg String V1 = val.pop();
    |       |   Stg String V2 = val.pop();
    |       |   Stg String t = V1 + V2 + ch;
    |       |   val.push(t);
    |   }
    |   Sout(val.peek());
}
```

$$t = "C" + V1 + ch + V2 + "$$

for Prefix
to Infix.

$$V1 + V2 + ch;$$

$$val.push(t);$$

Sout(val.peek());

① Postfix To Infix Conversion.

Str = "9 5 3 + 4 * 6 / -" # V₁, op V₂.

(9 - ((5+3)*4)/6)

((5+3)*4)/6

((5+3)*4)
4
(5+3)
3
5
9

V₂ = 3 op = +
V₁ = 5 $\Rightarrow (5+3)$

V₂ = 4 op = *.
V₁ = (5+3) $\Rightarrow ((5+3)*4)$

V₂ = 6 op = /.

V₁ = ((5+3)*4) $\Rightarrow (((5+3)*4)/6)$.

V₂ = (((5+3)*4)/6).

V₁ = 9 op = -

} $\Rightarrow (9 - (((5+3)*4)/6))$.

② Postfix to Prefix.

Str = "9 5 3 + 4 * 6 / -" # Op V₁, V₂.

-9/*+5346

+5346

6

*+534

+53

3

5

9

V₂ = 3 op = +.

V₁ = 5 $\Rightarrow +53$

V₂ = 6 op = /.

V₁ = *+534

V₂ = 9 op = -

$\Rightarrow / * + 5346$

V₂ = 4 op = *.

V₁ = +53

V₂ = / * + 5346.

V₁ = 9. op = -

$\Rightarrow -9/* + 5346$

$\Rightarrow * + 5346.$

$\Rightarrow -9/* + 5346$

\Rightarrow Code 11 // postfix To Infix.

```
for(int i=0 ; i < n ; i++) {
```

```
    char ch = str.charAt(i);
```

```
    int ascii = (int)ch;
```

```
    if(ascii >= 48 && ascii <= 57) {
```

```
        val.push(ch + ")");
    }
```

```
    else {
```

```
        String v2 = val.pop();
```

```
        String v1 = val.pop();
```

```
        char op = ch;
```

```
        String t = "(" + v1 + op + v2 + ")";
```

```
        val.push(t);
```

```
}
```

```
sout(val.peek());
```

for postfix to Prefix

change only this.

t = op + v1 + v2;

for
Postfix
to
Prefix.