

Improved Learning-Augmented Algorithms for the Multi-Option Ski Rental Problem via Best-Possible Competitive Analysis

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Learning-Augmented Algorithms

- Take **ML prediction** as added input
- No assumption on prediction
- Recent success in **online optimization**

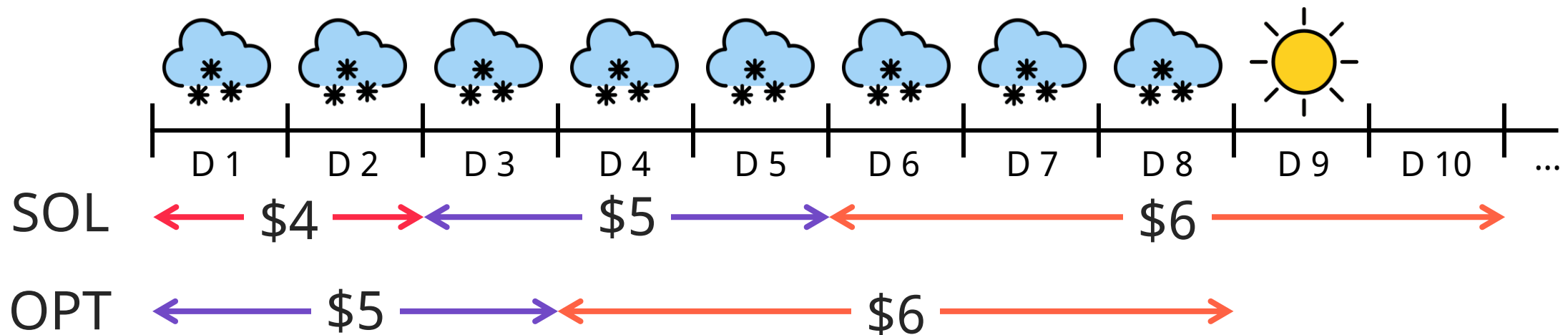
Performance Measure

- An online algorithm is *c-competitive* if $SOL \leq c \cdot OPT$.
- A learning-augmented algorithm is
 - *χ -consistent* if $SOL \leq \chi \cdot OPT$ when prediction is accurate.
 - *ρ -robust* if $SOL \leq \rho \cdot OPT$ unconditionally.

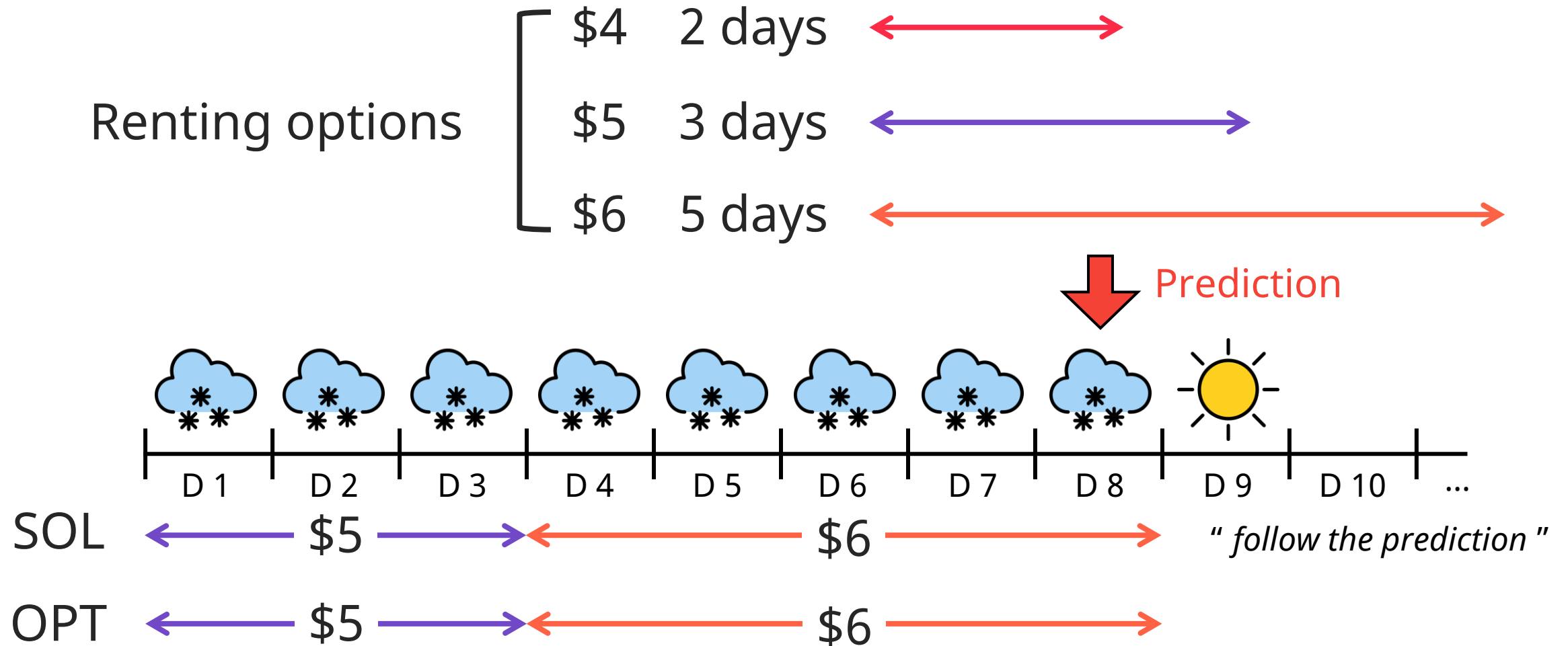
Multi-Option Ski Rental



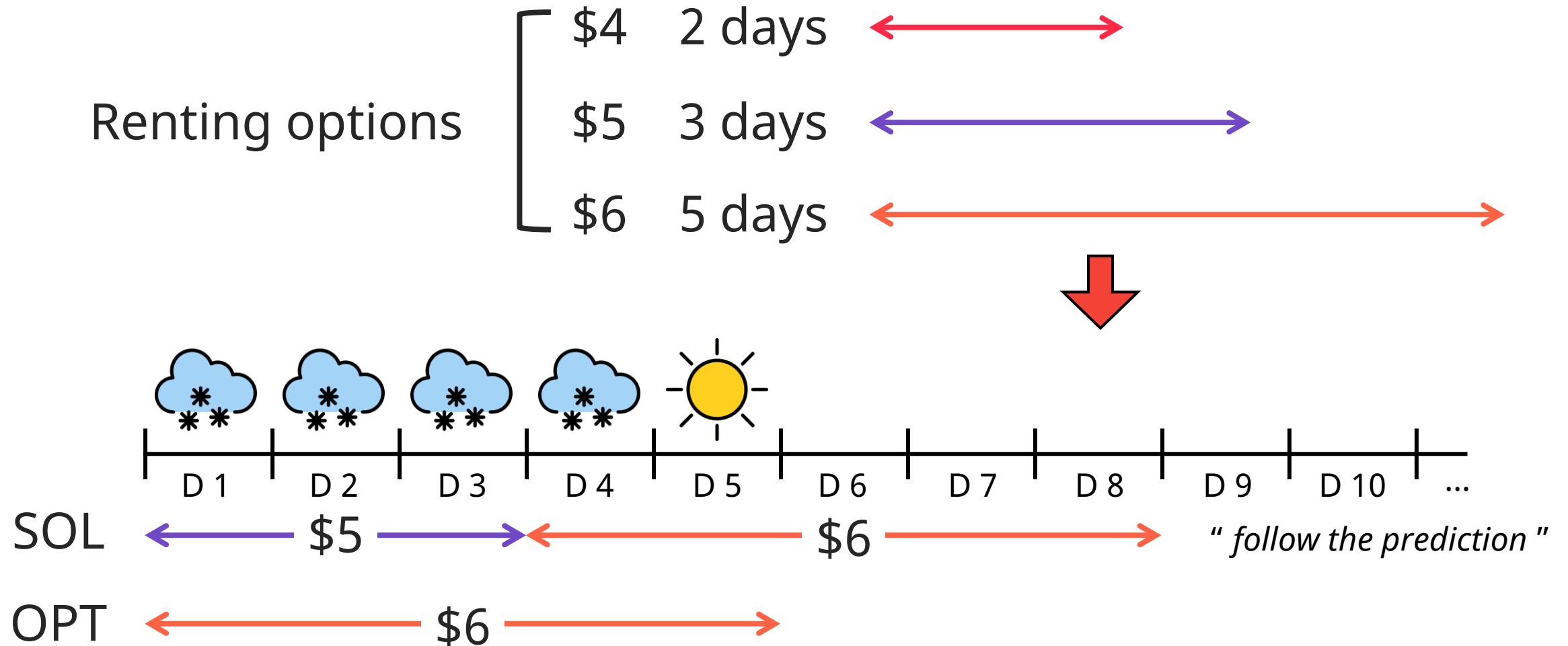
Multi-Option Ski Rental



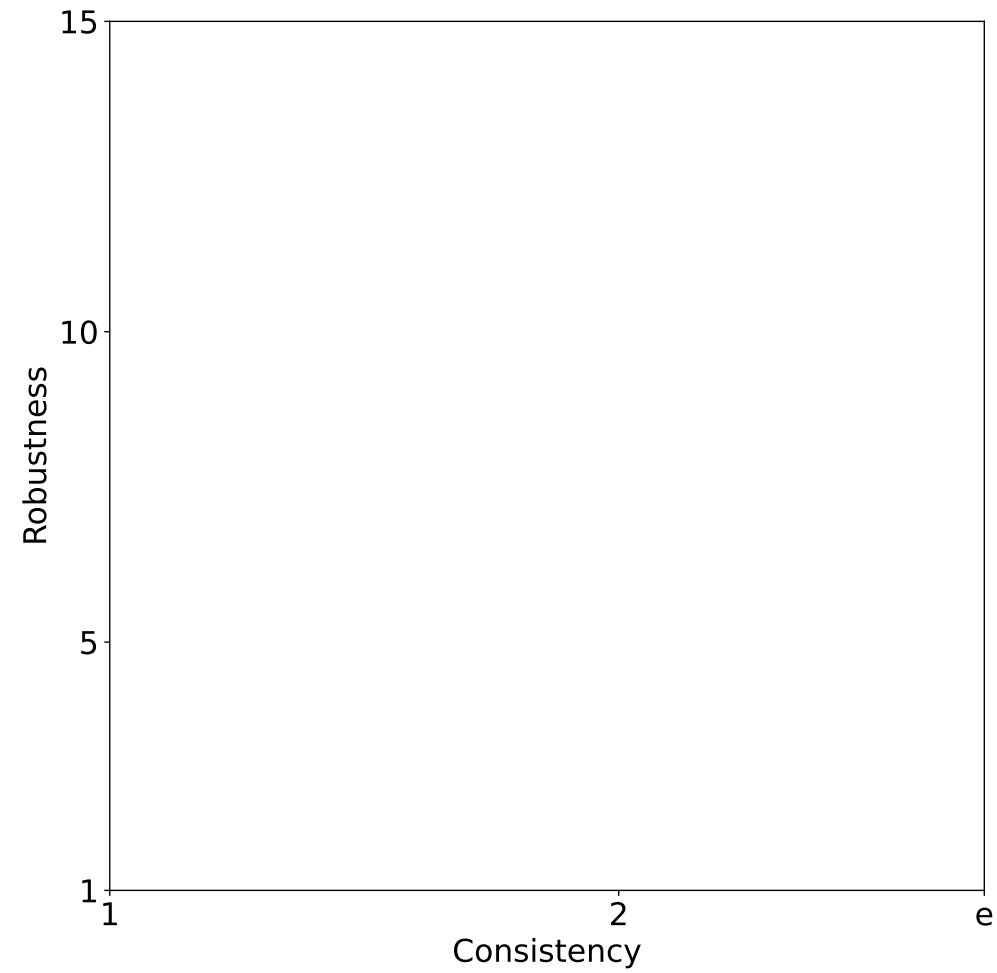
Learning-Augmented Multi-Option Ski Rental



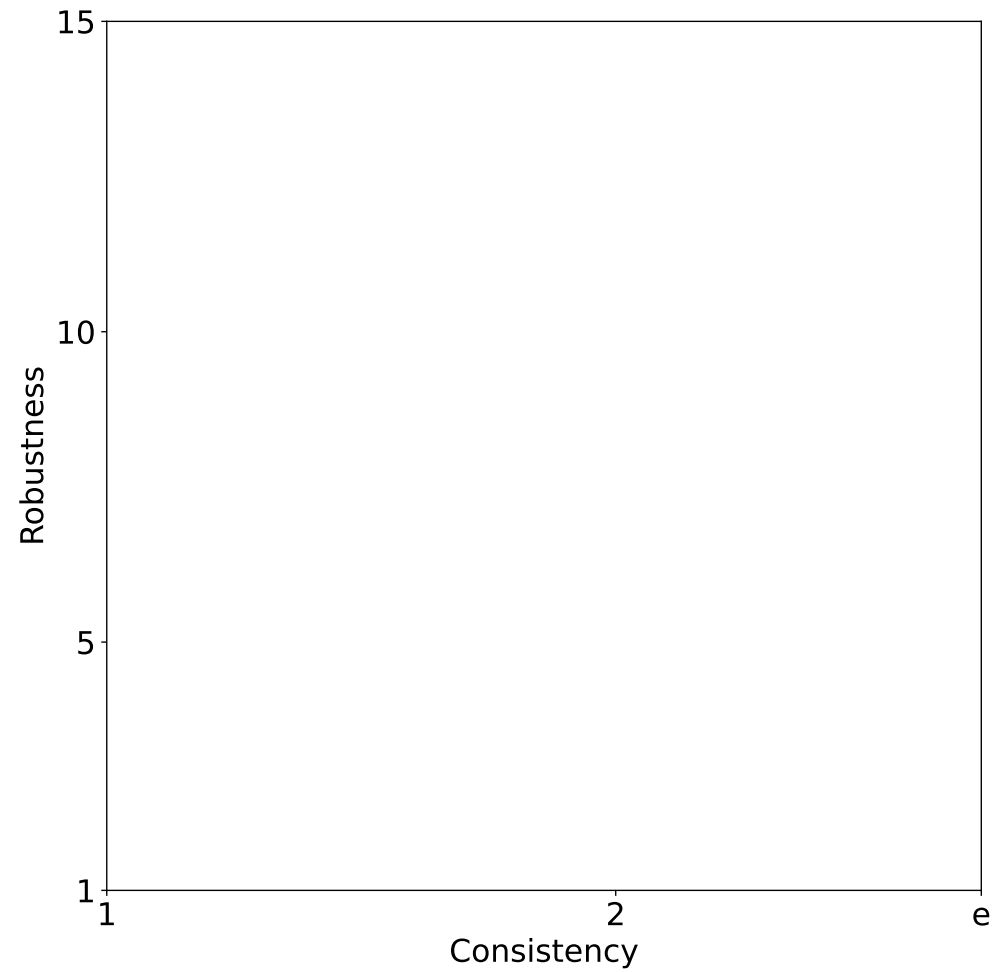
Learning-Augmented Multi-Option Ski Rental



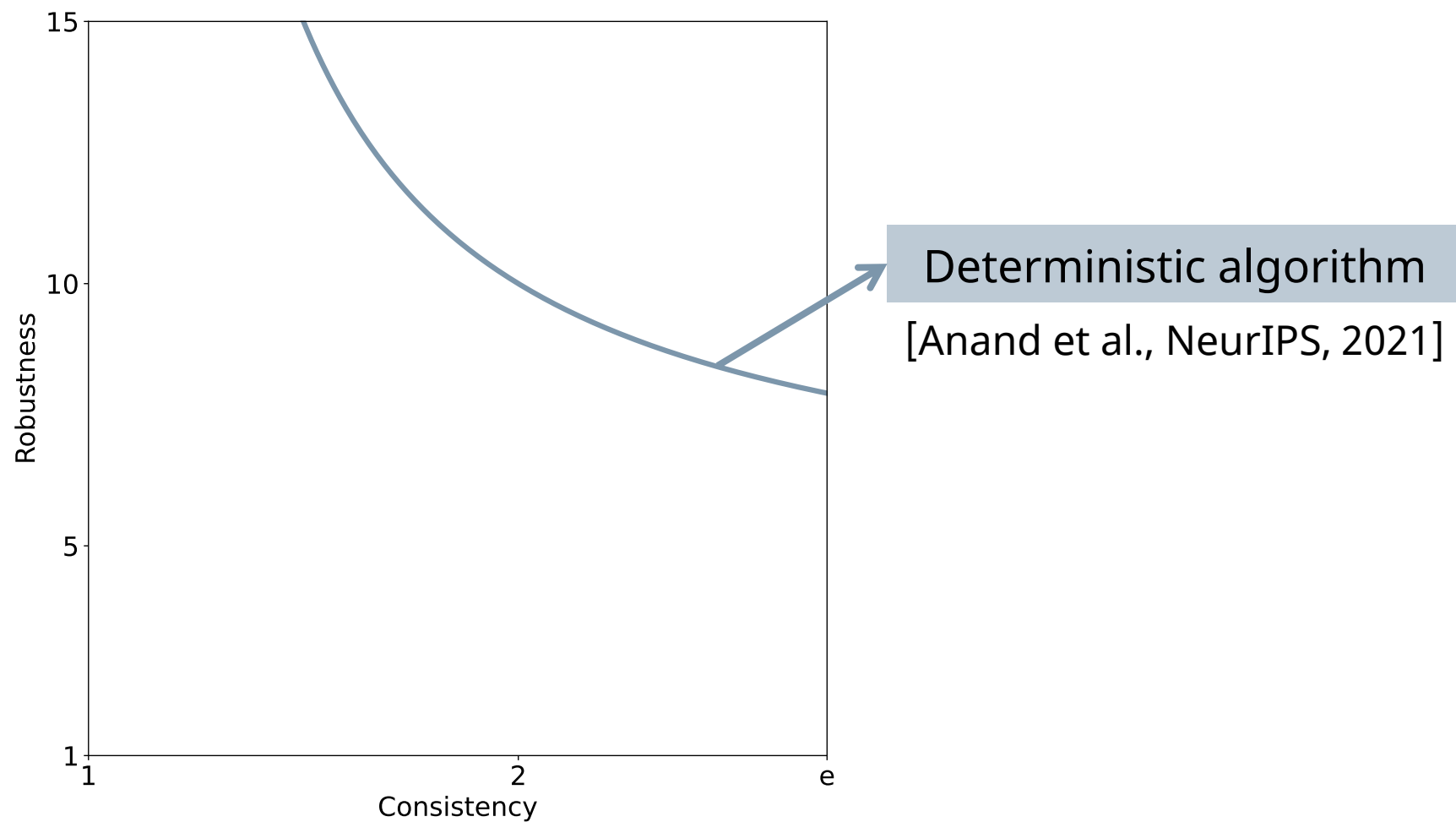
Previous Work



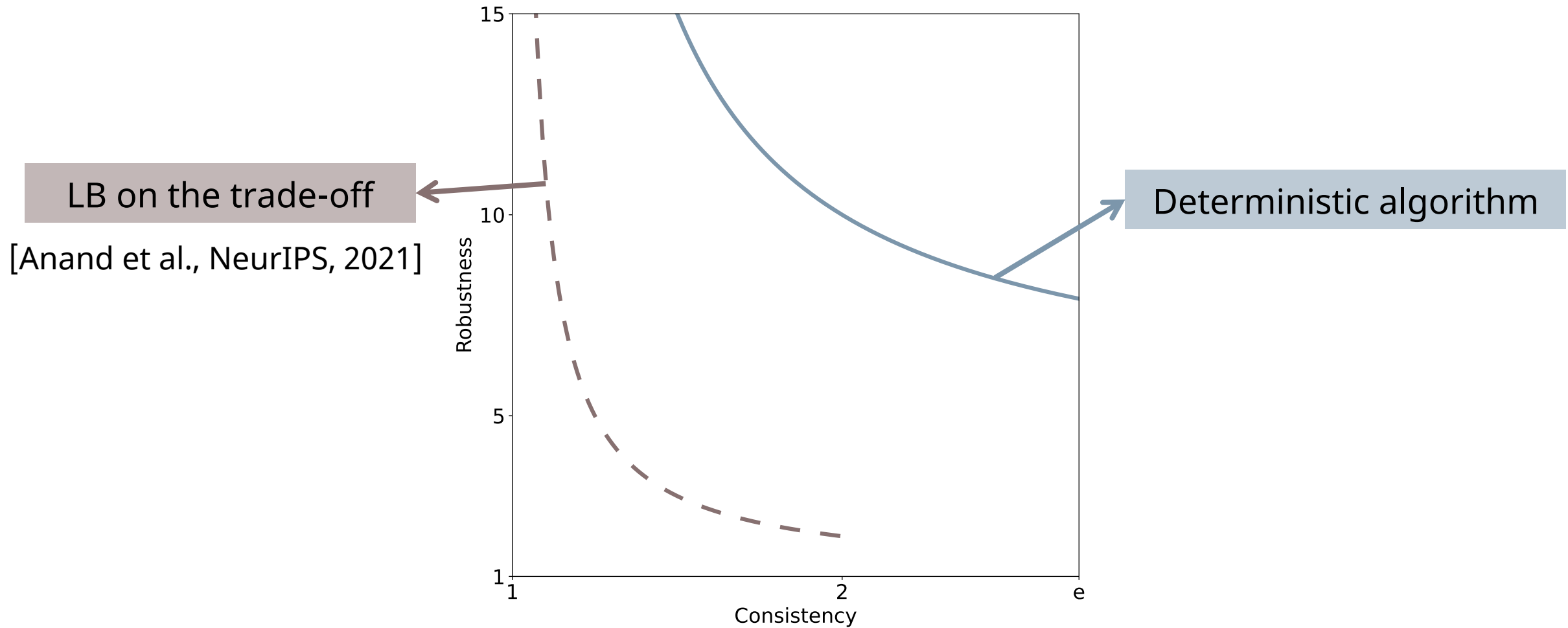
Previous Work (Deterministic)



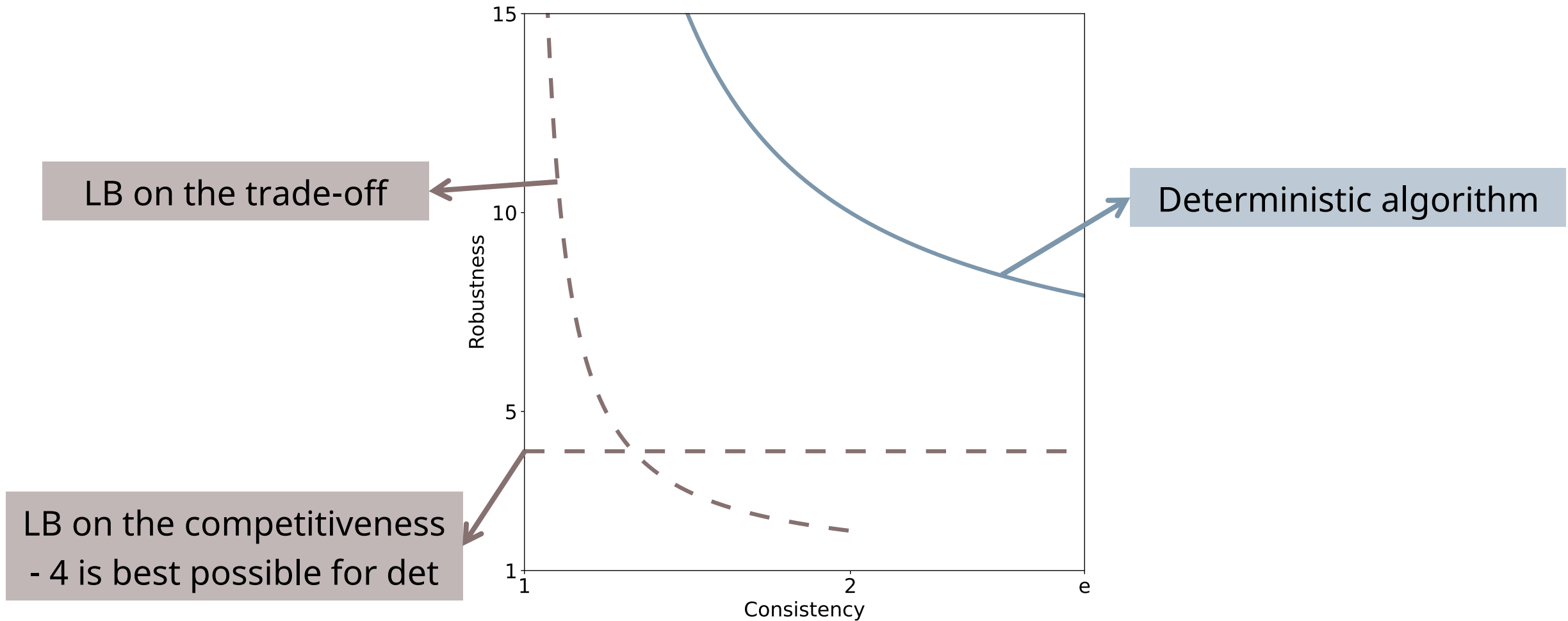
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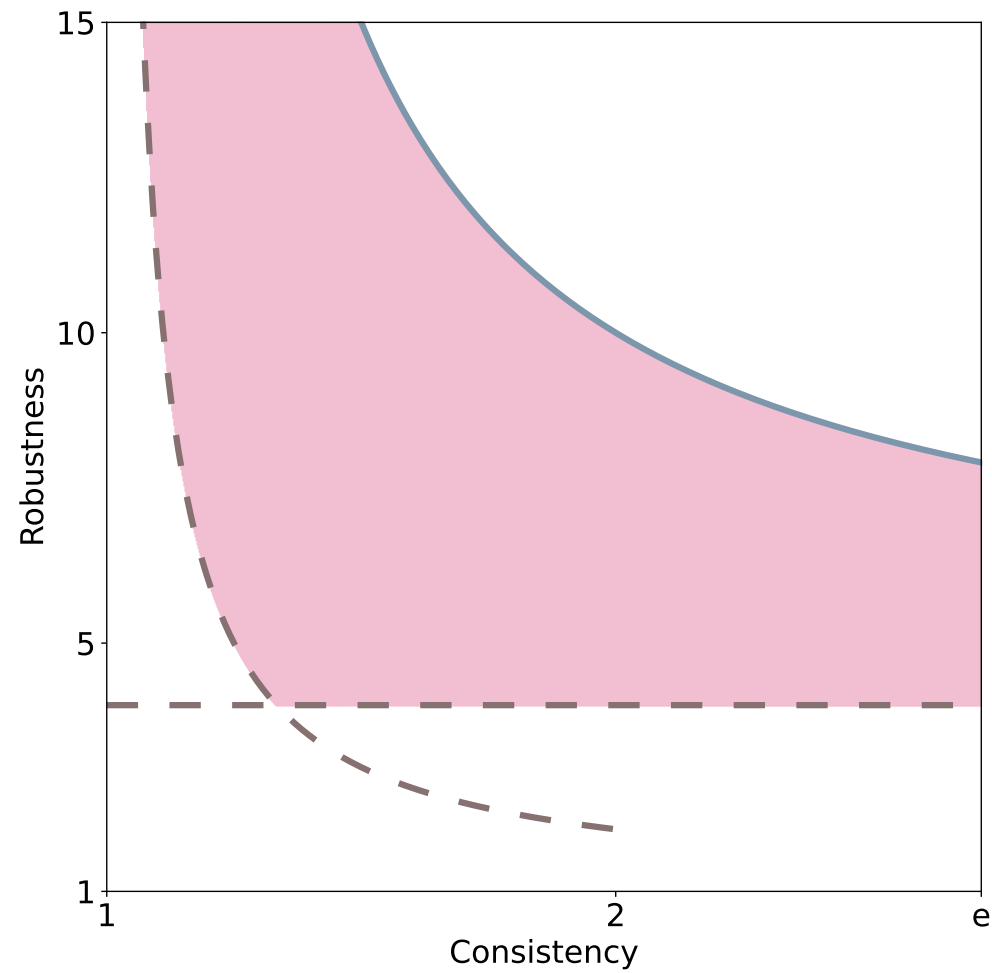


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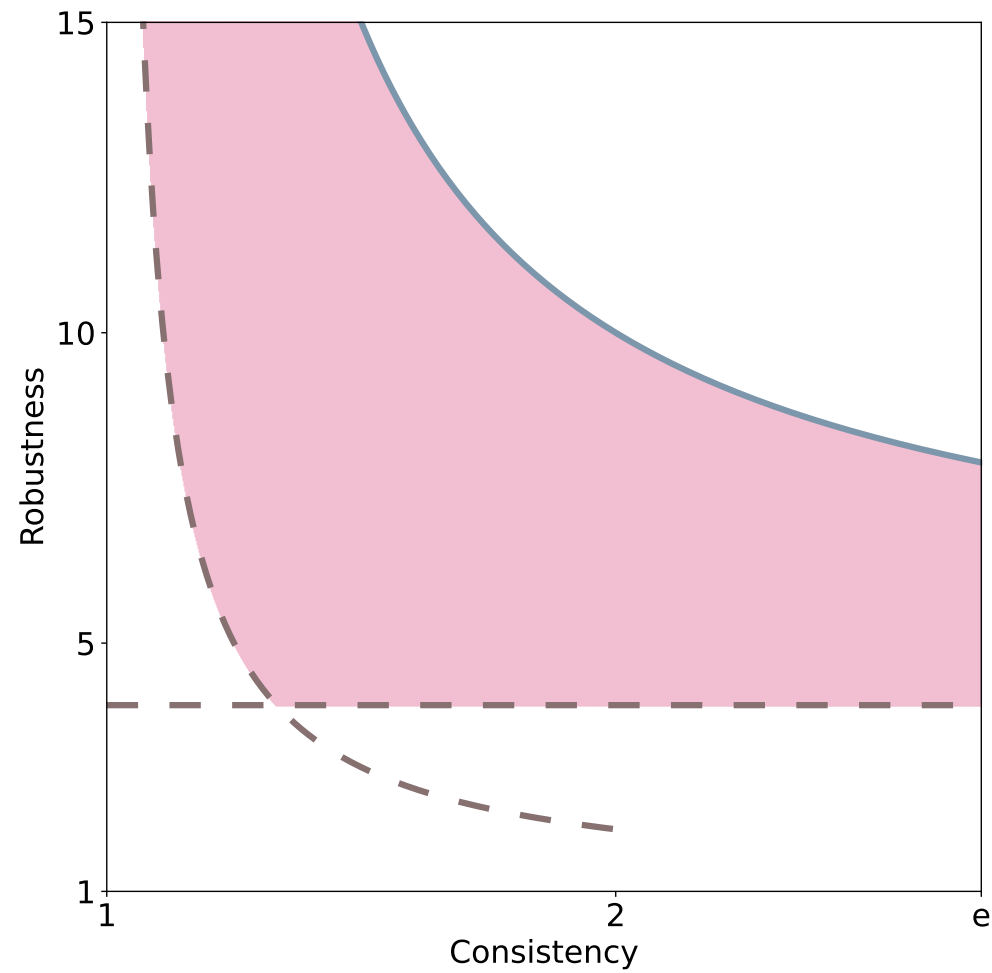


[Zhang et al., Inf. Process Lett, 2011]

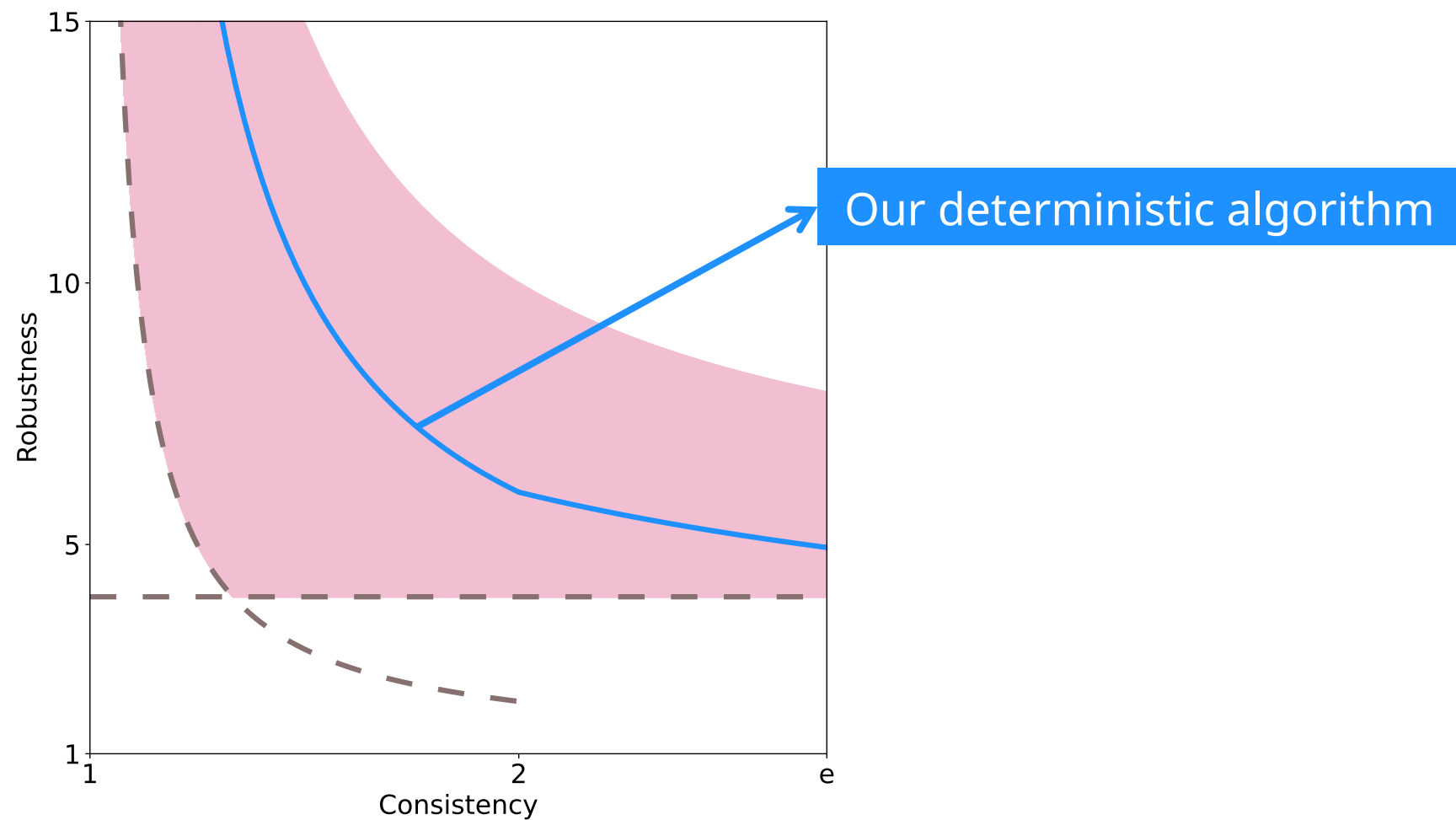
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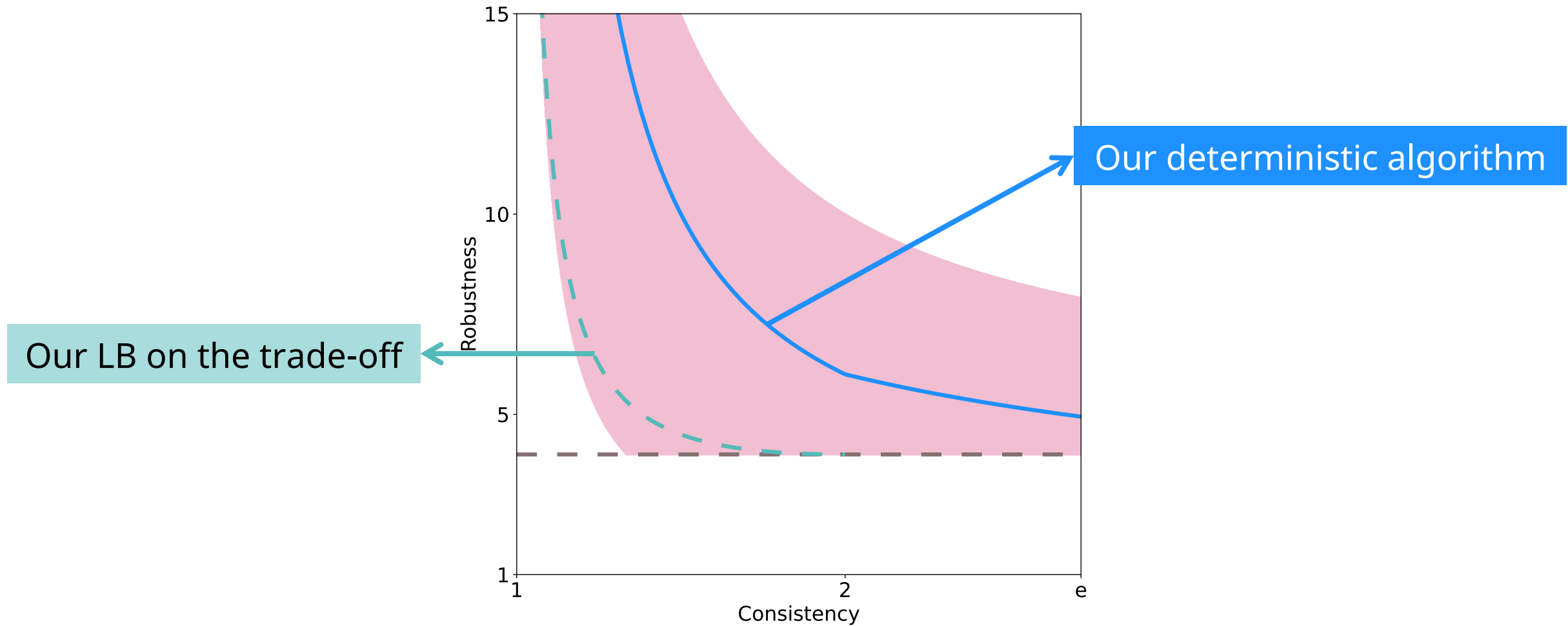
Our Results (Deterministic)



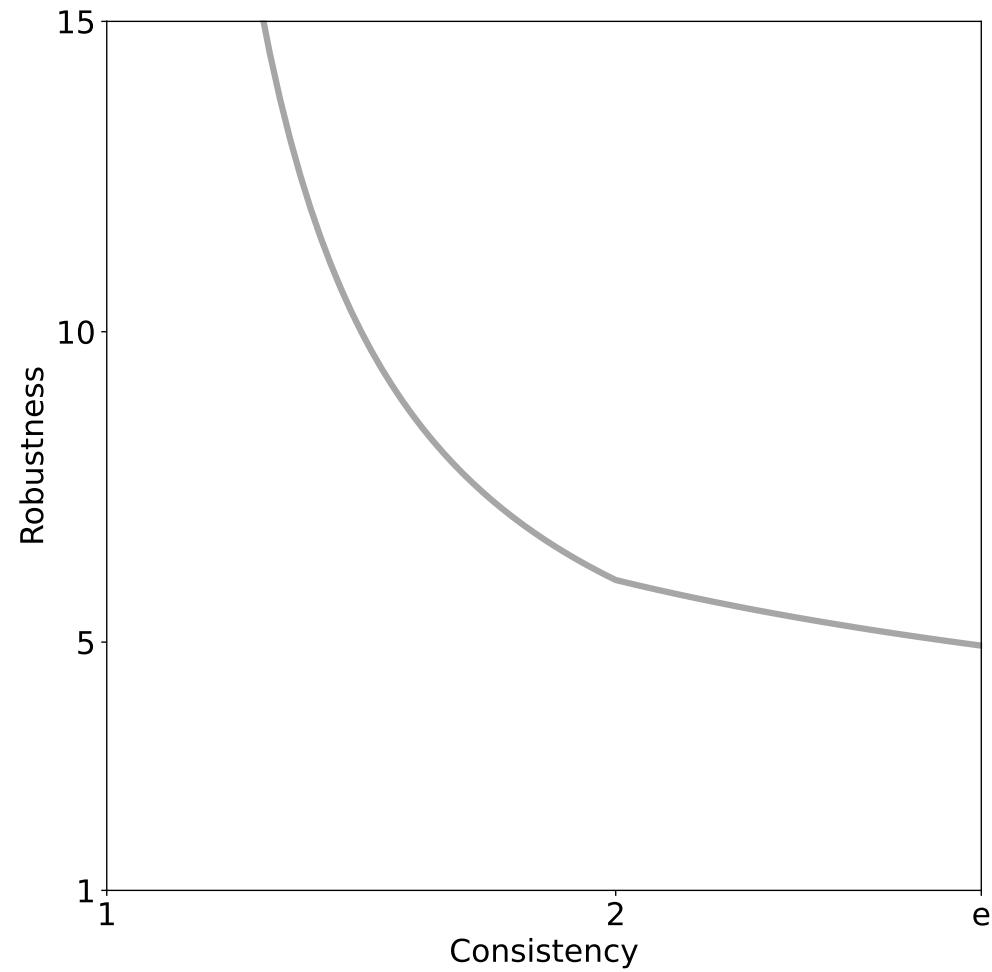
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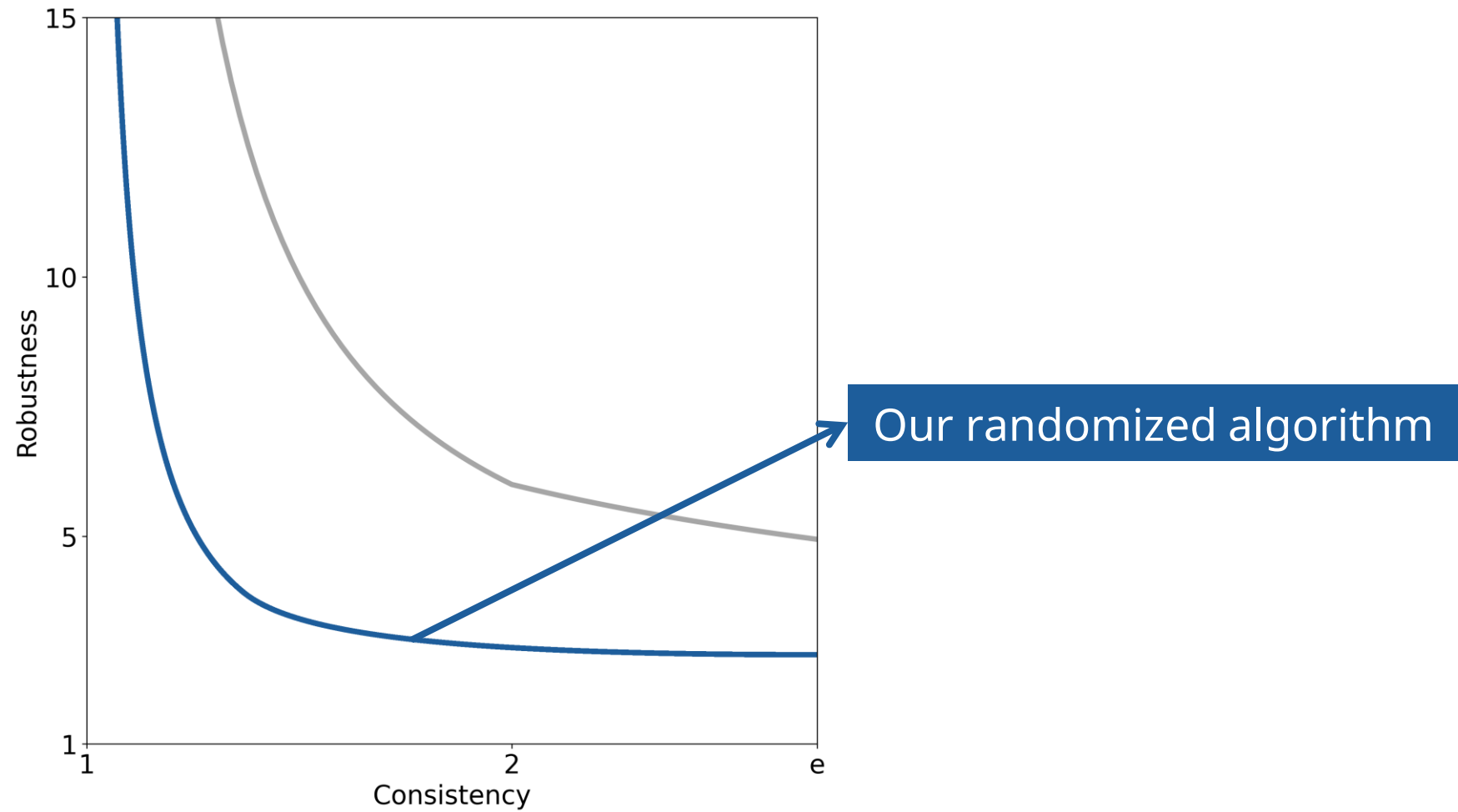
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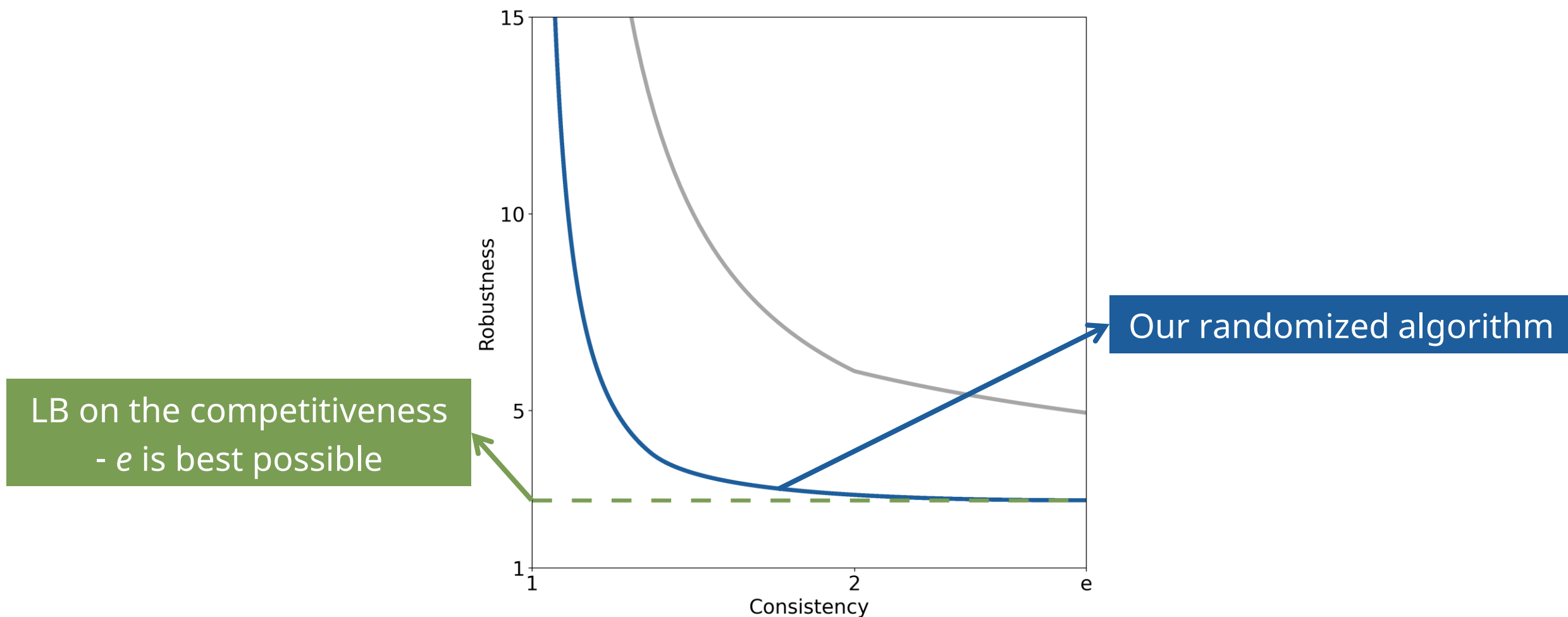
Our Results (Randomized)



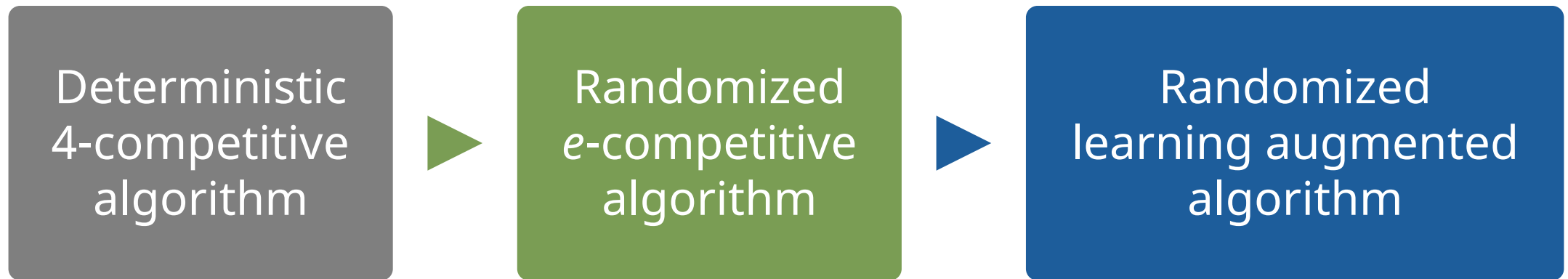
Our Results (Randomized)



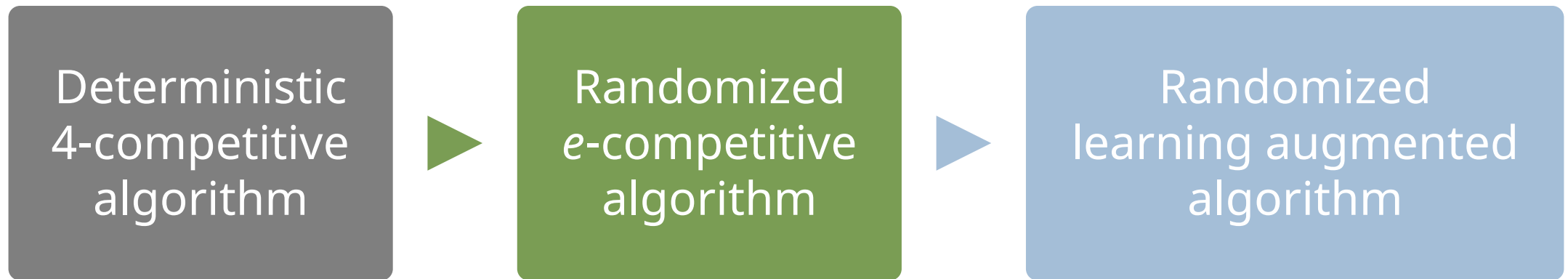
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Roadmap of Algorithms



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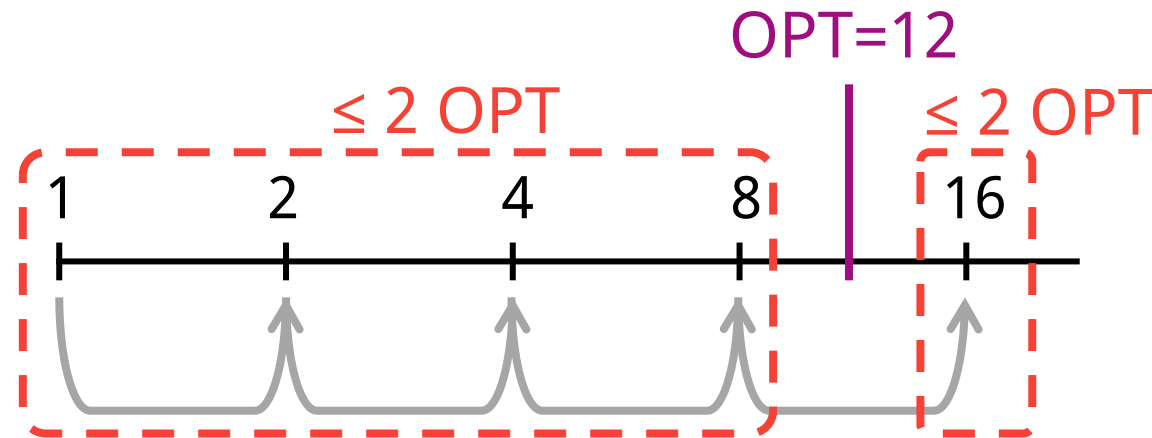


Deterministic Competitive Algorithm

for each round $i = 0, 1, \dots$:
 set *budget* to 2^i
 choose “best” options within the *budget*

Deterministic Competitive Algorithm

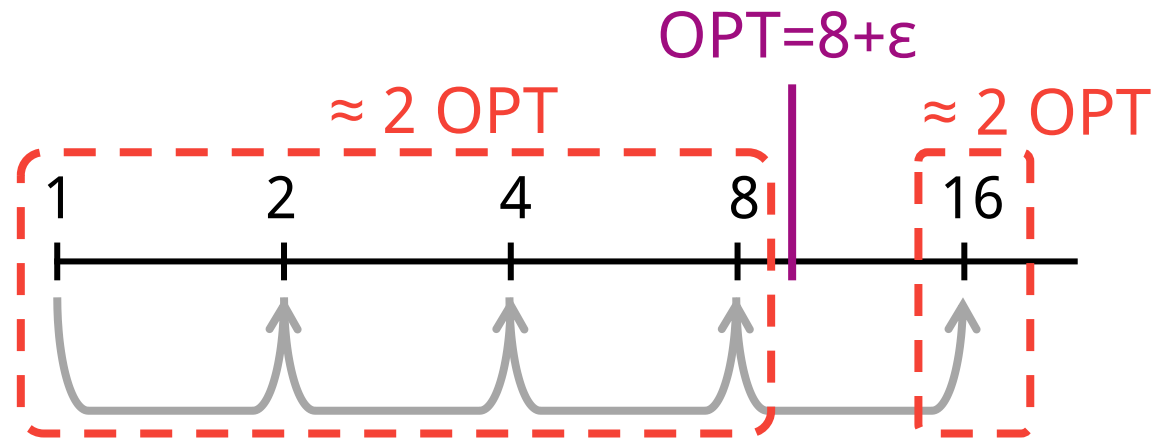
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Theorem. (Anand et al.) This algorithm is 4-competitive

Worst Case of Deterministic Algorithm

for each round $i = 0, 1, \dots$:
 set *budget* to 2^i
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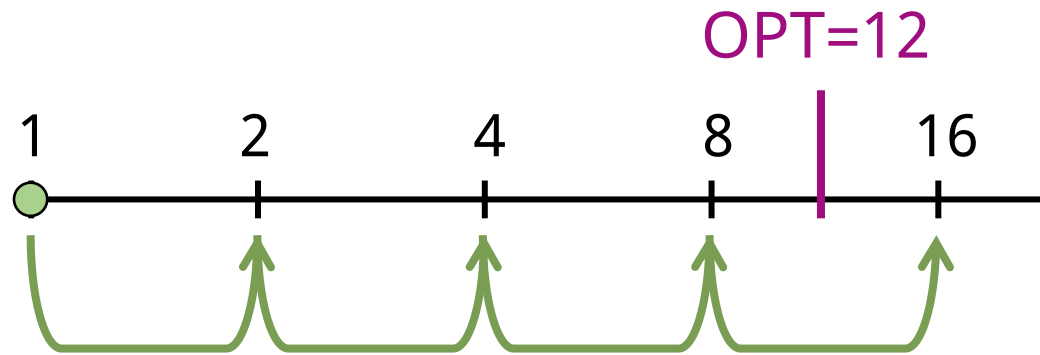
Randomized Initial Budget

sample initial *budget* $\alpha \in [1, 2)$

for each round $i = 0, 1, \dots$:

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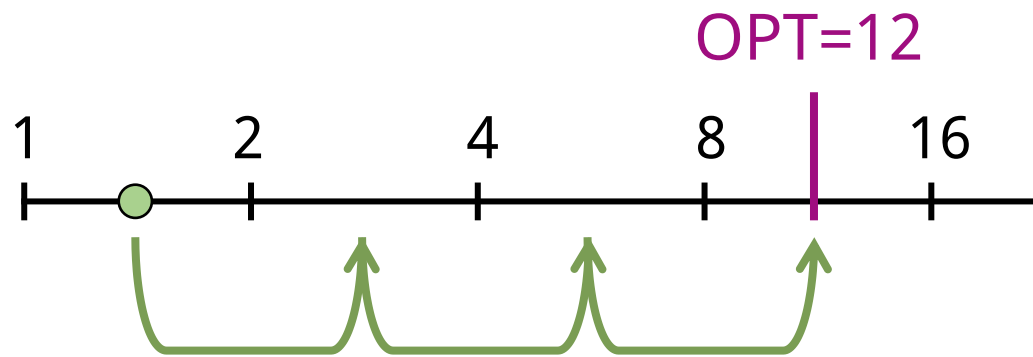
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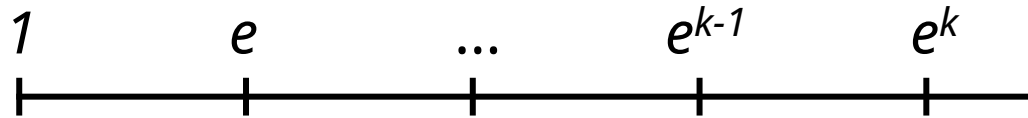
Randomized Competitive Algorithm

sample initial *budget* $\alpha \in [1, e)$ from pdf $f(\alpha) = 1/\alpha$

for each round $i = 0, 1, \dots$:

 set *budget* to $\alpha \cdot e^i$

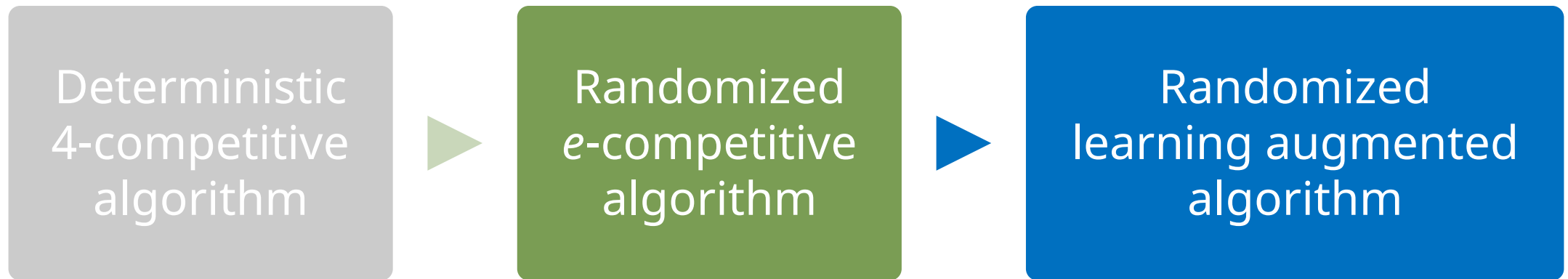
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Theorem. (SLLA, 2023) This algorithm is e -competitive

Theorem. (SLLA, 2023) This algorithm is best possible

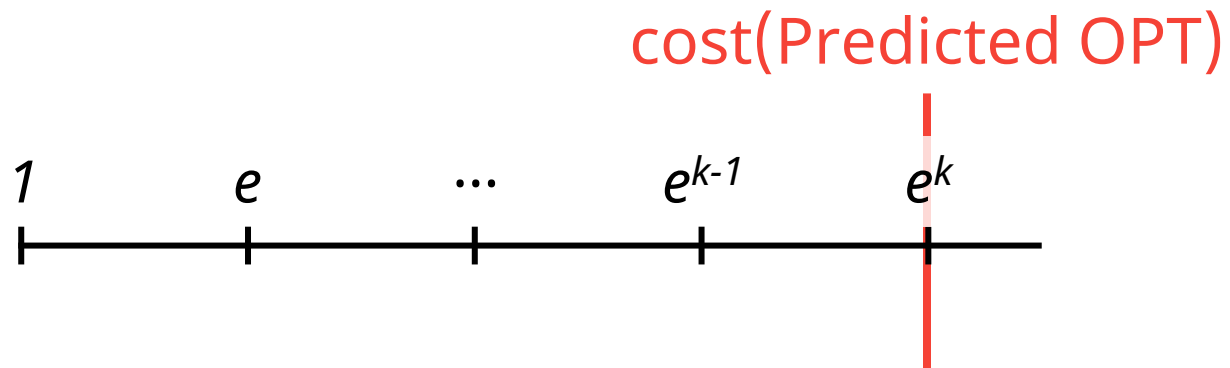
Roadmap of Algorithms



Randomized Learning Augmented Algorithm

Given: prediction on the last day

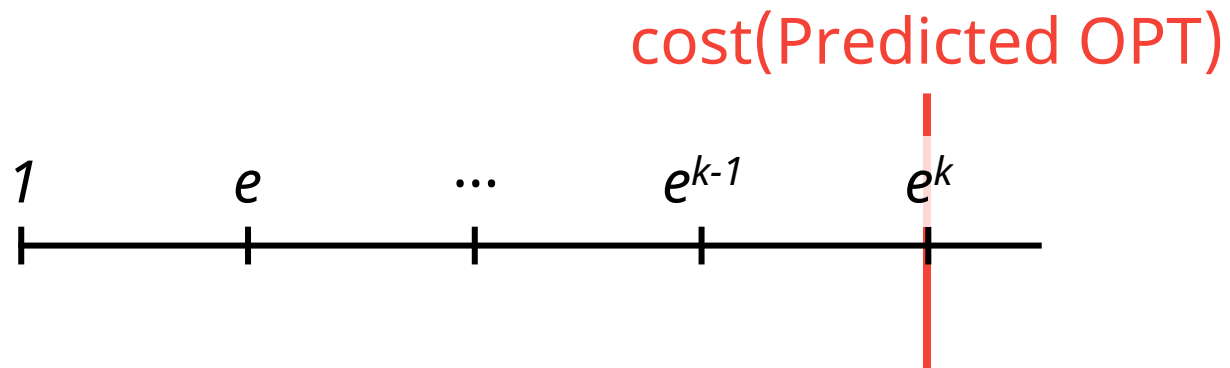
assume $\text{cost}(\text{"predicted opt soln"}) = e^k$ for some int k



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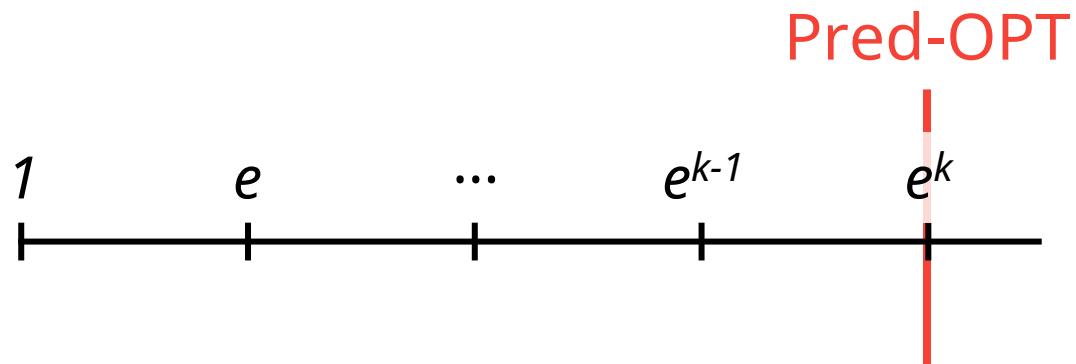
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Randomized Learning Augmented Algorithm

Given: prediction on the last day

Pred-OPT = e^k for some int k , WLOG



Randomized Learning Augmented Algorithm

Given: prediction on the last day, "level of confidence" $\lambda \in [0,1]$

Pred-OPT = e^k for some int k , WLOG

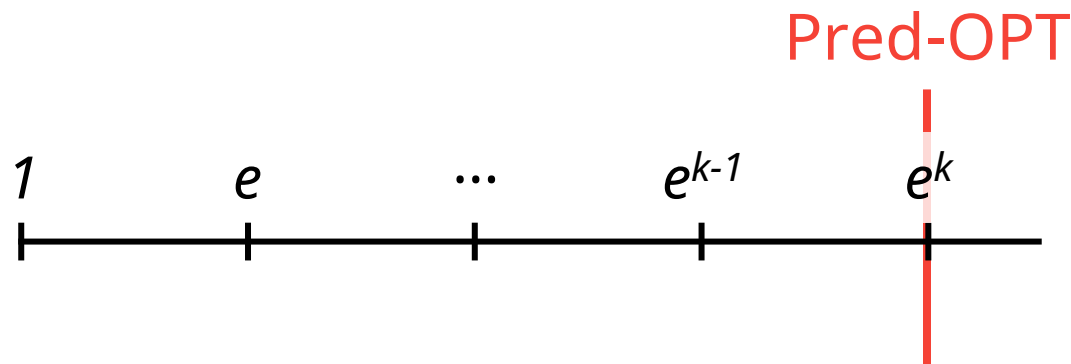
sample initial *budget* $\alpha \in [1, e)$ from pdf $f(\alpha) = 1/\alpha$

for each round $i = 0, 1, \dots$:

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if the *budget* $\in [\lambda \cdot \text{Pred-OPT}, \text{Pred-OPT})$, "*follow the prediction*" (at most once)

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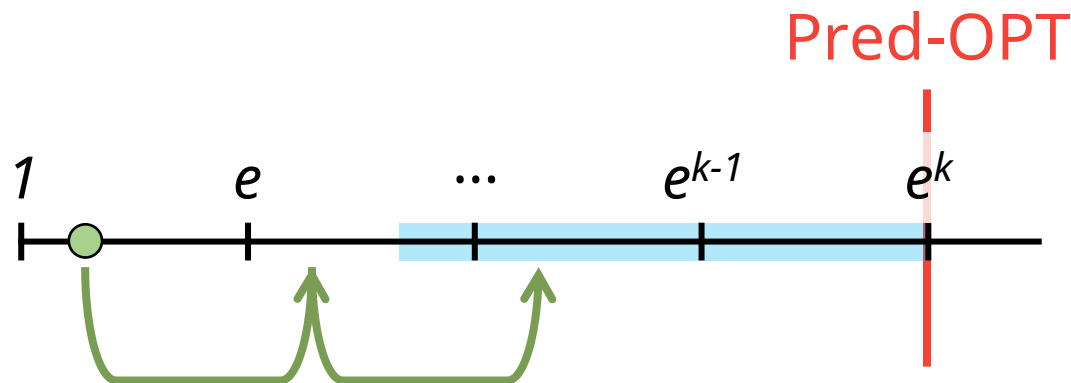
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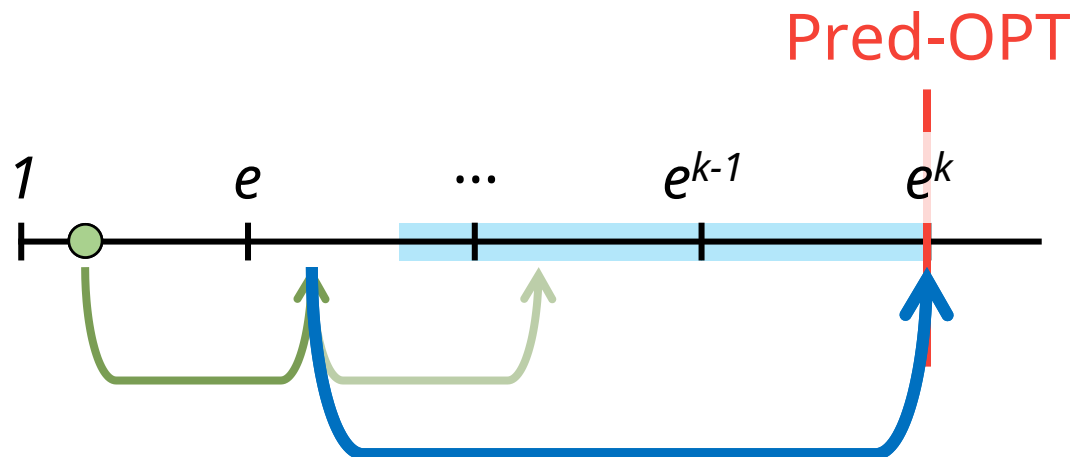
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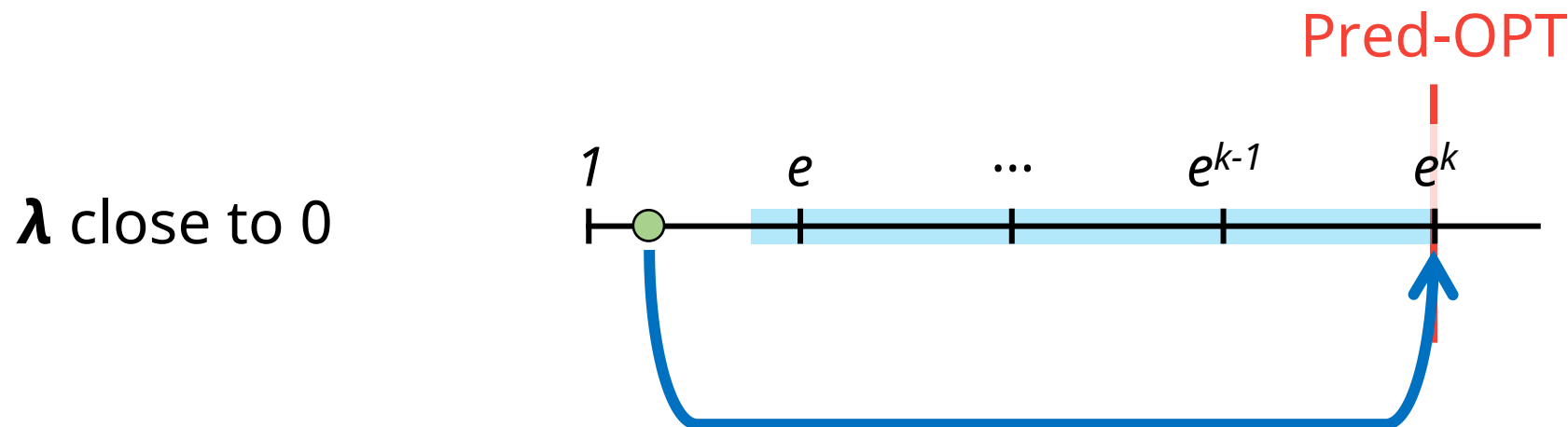
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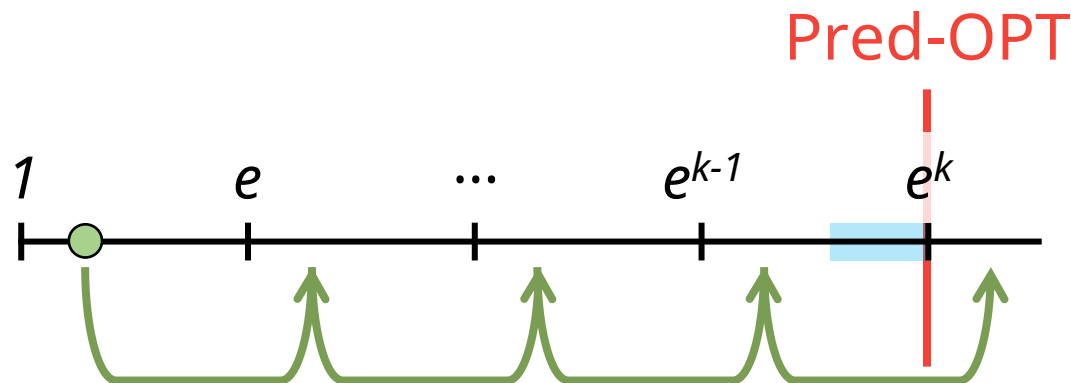
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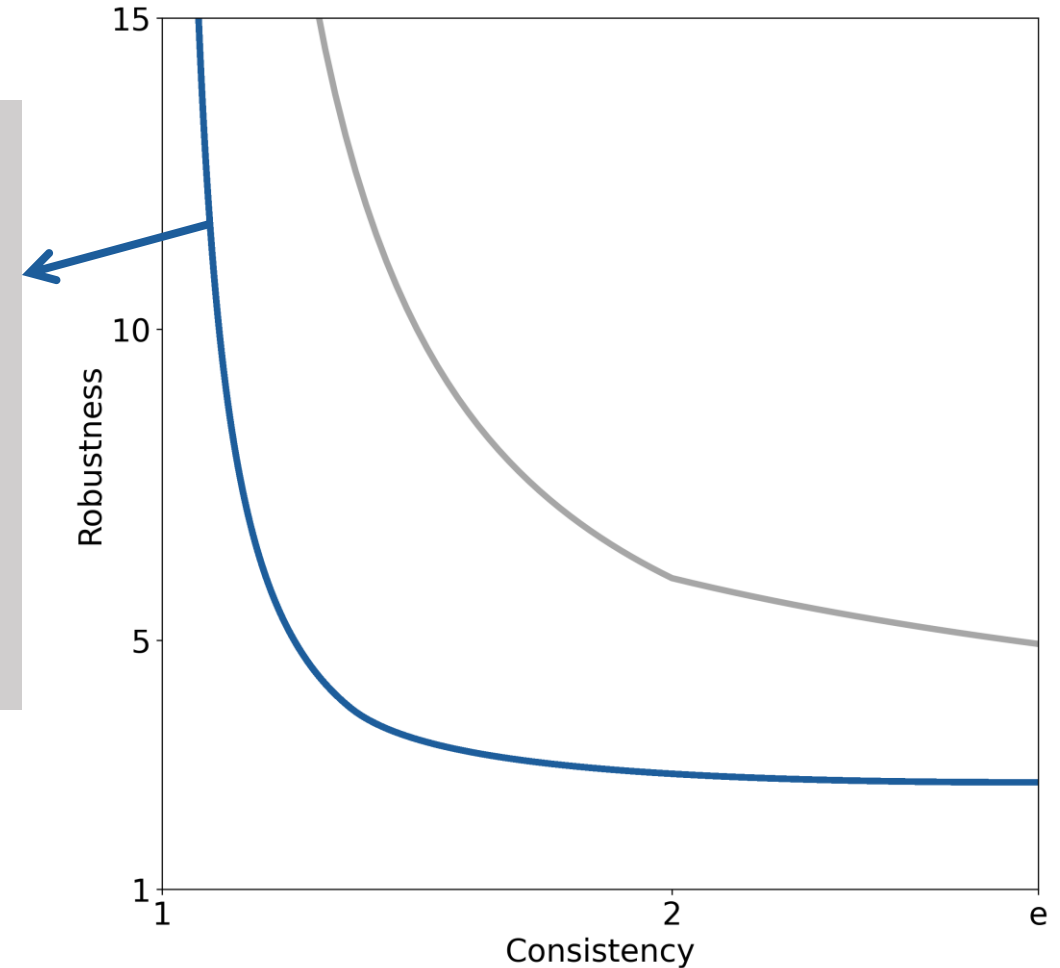
λ close to 1



Randomized Learning-Augmented Algorithm

Theorem. (SLLA, 2023) For $\lambda \in [0,1]$, this algorithm is $\chi(\lambda)$ -consistent (e^λ/λ) -robust, where

$$\chi(\lambda) := \begin{cases} 1 + \lambda, & \lambda < 1/e, \\ (e + 1)\lambda - \ln \lambda - 1, & \lambda \geq 1/e. \end{cases}$$



Lower Bounds

- Button problem
 - is (almost) as hard as multi-option ski rental
- Competitive ratio of any algorithm for the button problem \geq

minimize

γ

subject to

$$\sum_{j=1}^m x_j = 1,$$

$$\sum_{j=t+1}^m y_{t,j} = x_t + \sum_{j=1}^{t-1} y_{j,t}, \quad \forall t = 1, \dots, m-1,$$

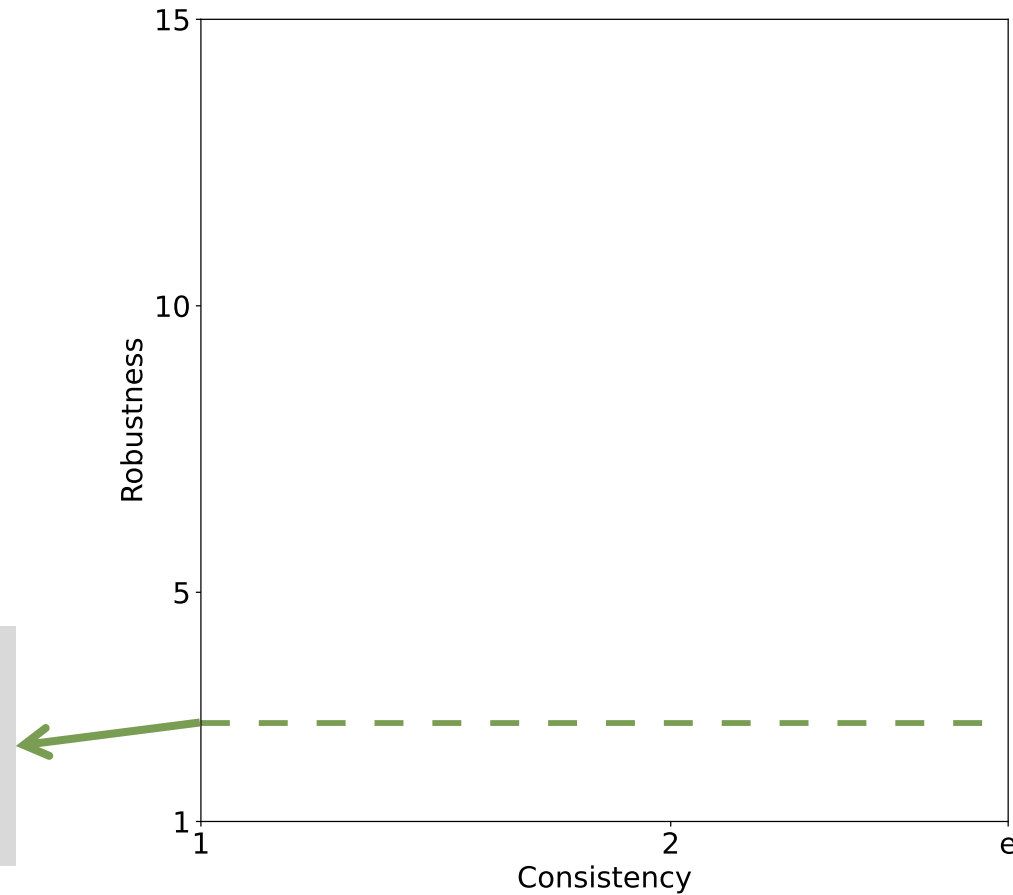
$$\sum_{j=1}^m b_j \cdot \left(x_j + \sum_{t=1}^{\min(J,j)-1} y_{t,j} \right) \leq \gamma \cdot b_J, \quad \forall J = 1, \dots, m,$$

$$x, y \geq 0.$$

- Construct a family of instances and consider the dual solutions

Lower Bounds for Button Problem

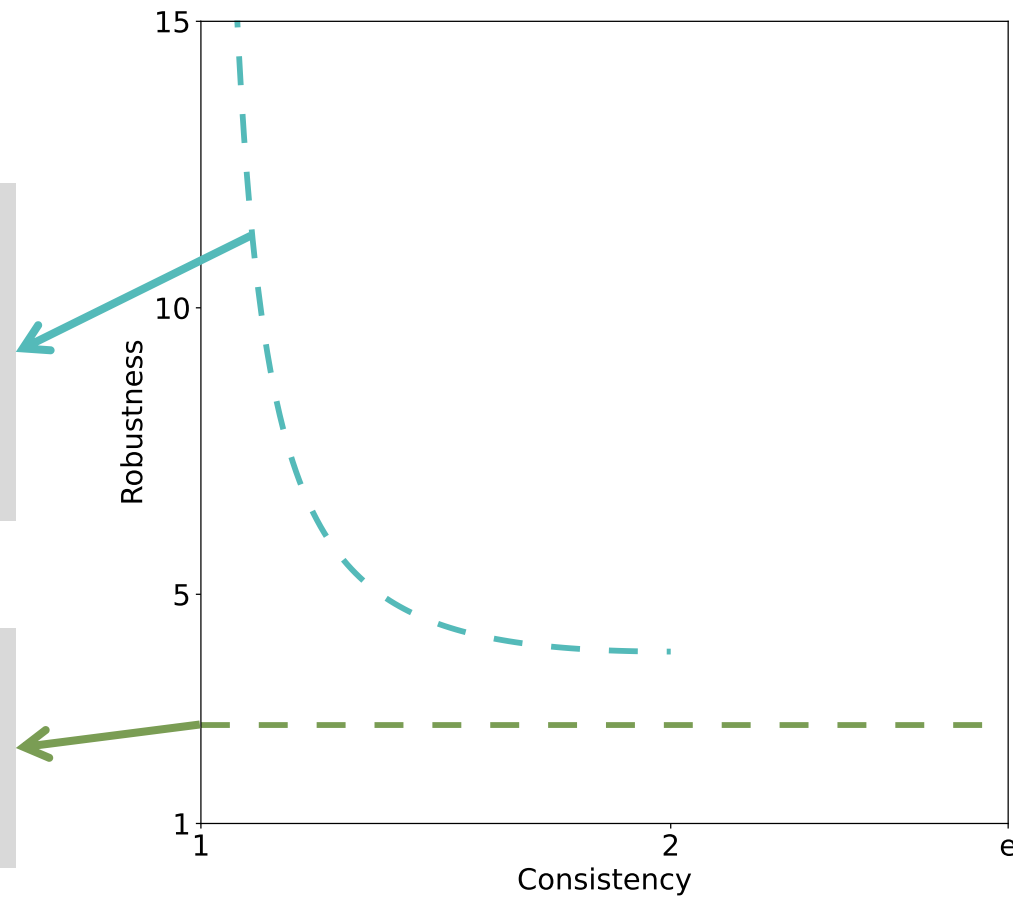
Theorem. (SLLA, 2023)
LB on the competitiveness



Lower Bounds for Button Problem

Theorem. (SLLA, 2023)
LB on the trade-off
for the deterministic algorithm

Theorem. (SLLA, 2023)
LB on the competitiveness



Thank you for listening!