Final Grammar: Left Factoring and Eliminating Left Recursion

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A \rightarrow \underline{let} B \underline{in} I \underline{end}
B \rightarrow C B \mid \epsilon
C \rightarrow D \mid F
D \rightarrow var ID D
D \rightarrow := K \mid : E
E → int | string | void
F \rightarrow \underline{\text{function}} \text{ ID (} F
F \rightarrow G): E = I \underline{end} \mid): E = I \underline{end}
G \rightarrow H G
G^{*} \rightarrow H G^{*} \mid \epsilon
H \rightarrow ID : E
I \rightarrow J I
I^{\cdot} \rightarrow ; J I^{\cdot} \mid \epsilon
J \rightarrow ID := J' \mid ID (J'' \mid return J''' \mid printint (K) \mid printstring (K)
K )
J^{*} \rightarrow K \mid getint()
J, \rightarrow ) | N )
J^{*} \rightarrow K \mid \epsilon
K \rightarrow L K
K^{\cdot} \rightarrow + L K^{\cdot} | - L K^{\cdot} | \epsilon
L \rightarrow M L
L^{*} \rightarrow * M L^{*} | / M L^{*} | \epsilon
M → ( K ) | NUMBER | STRING_LITERAL | ID M`
M^{\cdot} \rightarrow () | (N) | \epsilon
N \rightarrow K N
N^{*} \rightarrow K N^{*} \mid \epsilon
```

Original Language

```
program = A \rightarrow \underline{let} B \underline{in} I \underline{end}
decs = B \rightarrow C B | \epsilon
dec = C \rightarrow D \mid F
var dec = D \rightarrow var ID := K \mid var ID : E
type = E \rightarrow \underline{int} \mid \underline{string} \mid \underline{void}
function_dec = F \rightarrow \underline{\text{function}} \ \text{ID} \ (G) : E = I \underline{\text{end}} \ | \underline{\text{function}} \ \text{ID} \ (\ ) : E =
I end
parameters = G \rightarrow G, H \mid H
parameter = H \rightarrow ID : E
statements = I \rightarrow I; J \mid J
statement = J \rightarrow ID := K \mid \underline{printint}(K) \mid \underline{printstring}(K) \mid ID :=
getint() | ID () | ID (N) | return K | return
expr = K \rightarrow K + L \mid K - L \mid L
term = L \rightarrow L * M \mid L / M \mid M
factor = M → ( K ) | NUMBER | STRING_LITERAL | ID | ID ( ) | ID ( N )
expr_list = N \rightarrow N, K \mid K
```