Project Report

Group Member

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Checkpoint 1

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Project Goal

Implementing Immix garbage collection on Dart VM, and comparing performances with original generation garbage collection that Dart VM is currently using.

Plan

(Done on 09.20.18) Study high-level of knowledge of Immix.

- Reading paper.
- Creating slides.

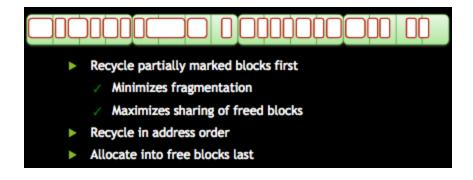
(Done on 09.27.18) Building VMs

- Jikes RVM
- Dart VM
 - Observatory (monitoring heap)

(On Going) Understanding Source Code

- Jikes RVM
 - 1. How Immix is structured with blocks and lines.
- Dart VM
 - 1. How heap is structured.
 - 2. How allocation works.
 - 3. Study write and read barrier(StoreIntoObject- how to bypass the barrier)

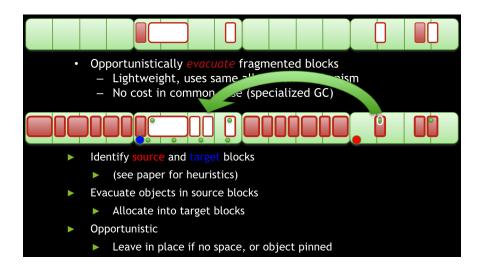
Overview of Immix



Allocation

- Heap is divided into blocks, and the blocks are divided into lines.
 - Blocks
 - Object cannot be span over blocks.
 - Should 4 times larger than max object size.

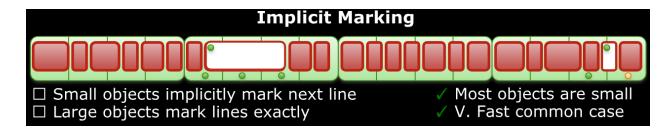
- Lines
 - Object can span lines.
 - All the lines should be marked that object spans.
- Because blocks are divided into multiple lines, by marking lines and checking the markings, it can prevent fragmentations.
- Objects get allocated on recyclable blocks first, and then fill in free blocks.



Opportunistic Defragmentation

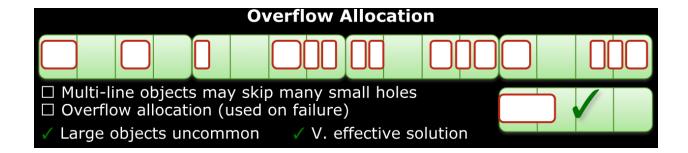
- After sweeping out, there needs defragmentation to make more free blocks and to make object allocation in contiguous order.
- Identification of source and target blocks.
 - It checks histograms(Mark and available).
 - Source: Select a block with the greatest # of holes.
 - Target: Find out a block that source block can fit in.
- After identification, it moves the objects in source block to target block.

Other Optimizations



- Implicit Marking

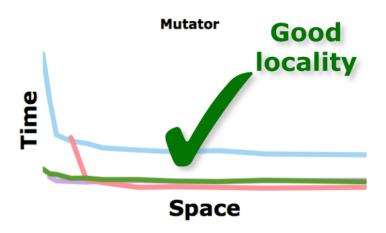
- For small objects span lines, it uses special marking scheme called implicit marking.
- Small object size is less than a line. So if an object spans a line, there will be large free space left in next line still.
- To use the space, it marks next line with implicit mark.
- When allocating small objects later, it checks implicit marked lines first whether the object can fit in.



- Overflow Allocation

- For a case that all recyclable blocks do not have a space to store current object, it gets a new free block and gets stored there.
- This works fine because large objects are uncommon.

Performances



- Good Locality

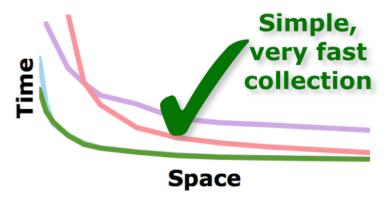
- Immix GC uses bump allocation that results good locality.



- Space Efficiency

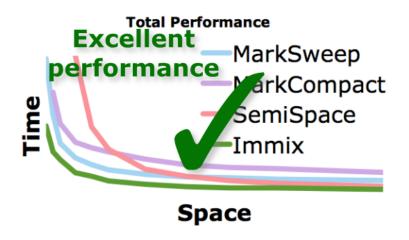
- Immix defragments objects frequently, and that makes space efficient.
- Also, optimization skills (such implicit marking and

Garbage Collection



- Garbage Collections Speed

- Opportunistic defragmentation happens with marking. (Total 2 Passes)
- So, performance is same as MarkSweep even with defragmentation.

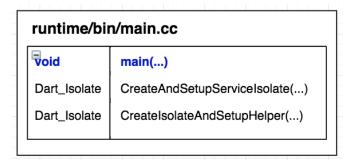


- Overall

It outperforms on all 3 performance tests.
 (Goot locality, Space Efficiency, and Garbage Collection time)

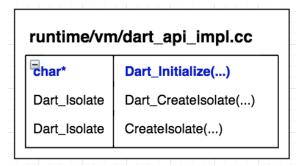
Overview of Dart VM

Classes and Heap-related Method Calls



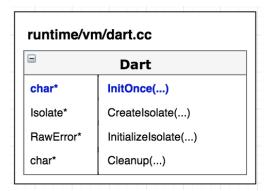
runtime/bin/main.cc

- Methods
 - void main(...) → Dart_Initialize(...)
 - static Dart_Isolate CreateAndSetupServiceIsolate(...) → Dart_CreateIsolate(...)
 - static Dart_Isolate CreateIsolateAndSetupHelper(...) → Dart_CreateIsolate(...)



runtime/vm/dart_api_impl.cc

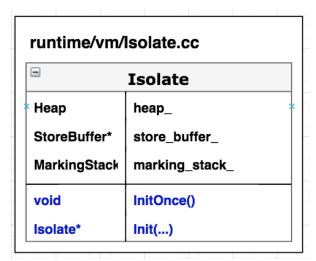
- Methods
 - DART EXPORT char* Dart Initialize(...) → Dart:InitOnce()
 - DART_EXPORT Dart_Isolate Dart_CreateIsolate(...) → CreateIsolate(...)
 - static Dart_Isolate CreateIsolate(...) → Dart::CreateIsolate(...)



runtime/vm/dart.cc

- Dart
 - Methods
 - char* Dart::InitOnce(...) → Isolate::InitOnce(), Isolate::Init(...),

 Object::InitNull(vm_isolate_), Object::InitOnce(vm_isolate_), SemiSpace::Init(...)
 - Isolate* Dart::CreateIsolate(...) → Isolate::Init(...)
 - RawError* Dart::InitializeIsolate(...) → Object::Inint(..)
 - char* Dart::Cleanup() → SemiSpace::Cleanup()



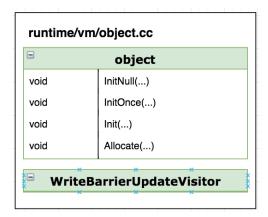
runtime/vm/isolate.cc

- Isolate
 - Variables: Heap heap_, StoreBuffer* store_buffer_, MarkingStack* marking_stack_
 - Methods
 - void Isolate::InitOnce()
 - Isolate* Isolate::Init(...) ------ Pause Here -----> Heap::Init(...)

runtime/vm/heap.cc		
=	heap	
PageSpace*	old_space_	
Scavenger*	new_space_	
void	Init()	
uword	Allocate()	
uword	AllocateOld()	
uword	AllocateNew()	
void	CollectOldSpaceGarbage()	
void	CollectNewSpaceGarbage()	
void	EvacuateNewSpace()	
void	AllocateExternal()	
void	FreeExternal()	
void	WriteProtect()	

runtime/vm/heap/heap.cc

- Heap
 - Space: kNew, kOld, KCode
 - GCType: kScavenge, kMarkSweep, kMarkCompact
 - Variables: Scavenger* new_space_, PageSpace* old_space_
 - Methods
 - void Heap::Init(...) → Heap::Heap(...)
 - Heap::Heap(...) → new_space_(...), old_space_(...), barrier_done_(...)
 - uword Allocate(...) → AllocateOld(...), AllocateNew(...)
 - uword Heap:AllocateOld(...) → old_space_.TryAllocate(...)
 - uword Heap:AllocateNew(...) → AllocateOld(...) or new_space_.TryAllocateInTLAB(...)
 - void Heap::CollectNewSpaceGarbage(...) → new_space_.Scavenge()
 - void Heap::EvacuateNewSpace(...) → new_space_.Evacuate()
 - void Heap::WriteProtect → old_space_.WriteProtect(...),
 new_space_.WriteProtect(...)
 - void Heap::AllocateExternal → old_space_.AllocateExternal(...),
 new_space_.AllocateExternal(...)
 - void Heap::FreeExternal(...) → old_space_.FreeExternal(...),
 new_space_.FreeExternal(...)



runtime/vm/object.cc

- Object
 - Methods
 - void Object::InitNull(...) → Heap::Allocate()
 - void Object::InitOnce(...) → Heap::Allocate()
 - void Object::Init(...)
 - void Object::Allocate(...) → Heap::Allocate()

- WriteBarrierUpdateVisitor

- Methods
 - WriteBarrierUpdateVIsitor visitor(...)

)	PageSpace		HeapPage
HeapPage*	pages_	HeapPage*	next_
uword	bump_top_	uword	object_end_
uword	bump_end_	uword	used_in_bytes_
HeapPage*	AllocatePage()	HeapPage*	Allocate()
HeapPage*	AllocateLargePage()	void	Deallocate()
uword	TryAllocate()	void	FreeForwardingPage()
uword	TryAllocateInternal()	void	WriteProtect()
void	CollectGarbage()		
void	CollectGarbageAtSafepoint()	□ Pag	geSpaceController
void	BlockingSweep()		
void	ConcurrentSweep()		
void	Compact()		
void	FreePage()		
void	FreePages()		
void	AllocateExternal()		
void	FreeExternal()		
void	WriteProtect()		

runtime/vm/heap/pages.h & pages.cc

- **HeapPage**: Contains old generation objects.
 - Size: 256KB
 - Variables: HeapPage* next_, uword object_end_, uword used_in_bytes_
 - Methods
 - HeapPage* HeapPage::Allocate(...)
 - void HeapPage::Deallocate()
 - void HeapPage::WriteProtect(...): Behaves depending on whether it is read_only or not
 - void HeapPage::FreeForwardingPage(...)
- PageSpace
 - Variables: HeapPage* pages_, uword bump_top_, uword bump_end_
 - Methods
 - HeapPage* PageSpace::AllocatePage(...) → HeapPage::Allocate(...)
 - HeapPage* PageSpace::AllocateLargePage(...) → HeapPage::Allocate(...)
 - uword TryAllocate(...) → PageSpace::TryAllocateInternal(...)
 - uword PageSpace::TryAllocateInternal(...)
 - void PageSpace::AllocateExternal(...)
 - void PageSpace::FreeExternal(...)
 - void PageSpace::CollectGarbage(...): Collect the garbage in the page space using mark-sweep or mark-compact.
 - → void PageSpace::CollectGarbageAtSafepoint(...)
 - void PageSpace::CollectGarbageAtSafepoint(...)
 - → GCSweeper::SweepPage(...),

GCSweeper::SweepLargePage(...),GCMarker::StartConcurrentMark(...),GCMarker::MarkObjects(...)

- void PageSpace::BlockingSweep(...): Start concurrent sweeper task. → GCSweeper::SweepPage(...)
- void PageSpace::ConcurrentSweep(...): Start concurrent sweeper task. →
 GCSweeper::SweepConcurrent(...)
- void PageSpace::Compact(...) \rightarrow GCCompactor::Compact(...)
- void PageSpace::WriteProtect(...) → HeapPage::WriteProtect(...)
- void PageSpace::FreePage(...), void PageSpace::FreePages(...): remove the page from the list of data pages. → HeapPage::Deallocate()
- PageSpaceController: Controls heap size.

runtime/vm/scavenger.cc

□ SemiSpace		
VirtualMemory*	reserved_	
MemoryRegion	region_	
void	Init()	
void	Cleanup()	
void	Delete()	
void	WriteProtect()	

□ ScavengerVisitor		
void	ScavengePointer()	
	•	

-	Scavenger
Heap*	heap_
SemiSpace*	to_
uword	top_
uword	end_
uword	AllocateGC()
uword	TryAllocateInTLAB()
void	Scavenge()
void	Evacuate()
void	AllocateExternal()
void	FreeExternal()
void	FlushTLS()
void	WriteProtect()

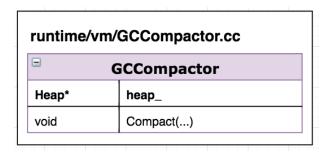
runtime/vm/heap/scavenger.h & scavenger.cc

- Scavenger
 - Variables: uword top_, uword end_, SemiSpace* to_, Heap* heap_
 - Methods
 - uword AllocateGC(...)
 - uword TryAllocateInTLAB(...)
 - void Scavenger::Scavenge(): Collect the garbage in this scavenger.
 - → Scavenger::FlushTLS(): Prepare for a scavenge,
 - ScavengerVisitor.visitor(...): Setup the visitor and run the scavenge.
 - void Scavenger::Evacuate(): Promote all live objects. → Scavenger::Scavenge()
 - void Scavenger::WriteProtect(...) → SemiSpace::WriteProtect(...)
 - void Scavenger::AllocateExternal(...)
 - void ScavengerFreeExternal(...)
 - void Scavenger::FlushTLS(): If mutator thread is scheduled, set to top on Scavenger. Used on all visiting or finding object methods in Scavenger class.
- ScavengerVisitor
 - Methods
 - void ScavengePointer(...) → scavenger_.AllocateGC(...)
- ScavengeStats: Statistics for a particular scavenge.
- SemiSpace: Wrapper around virtual memory that adds caching and handles the empty case.
 - Variables: VirtualMemory* reserved_, MemoryRegion region_
 - Methods
 - void SemiSpace::Init(): Create Mutex
 - void SemiSpace::Cleanup(): Delete cache
 - void SemiSpace::Delete()
 - void SemiSpace::WriteProtect(...)

runtime/vm/sweeper.cc	
	GCSweeper
bool	SweepPage()
intptr_t	SweepLargePage()
void	SweepConcurrent()

runtime/vm/heap/sweeper.h & sweeper.cc

- GCSweeper
 - Methods
 - bool SweepPage(...)
 - intptr_t SweepLargePage(...)
 - static void SweepConcurrent(...)



runtime/vm/heap/compactor.h & compactor.cc

- GCCompactor
 - Variables: Heap* heap_
 - Methods
 - void GCCompactor::Compact(...) → HeapPage::Deallocate(), HeapPage::FreeForwardingPage(...)

runtime/vm/marker.cc		
□ GCMarker		
void	StartConcurrentMark()	
void	MarkObjects()	

runtime/vm/heap/marker.h & marker.cc

- GCMarker
 - Methods
 - void GCMarker::StartConcurrentMark(...)
 - void GCMarker::MarkObjects(...)

<u>Understanding about each method</u>

Heap

- Garbage Collection
 - 1. CollectOldSpaceGarbage \rightarrow PageSpace.CollectGarbage
 - $2. \quad Collect New Space Garbage \rightarrow Scavenger. Scavenge \\$
 - 3. CollectGarbage
 - If GCType == Scavenge: CollectNewSpaceGarbage
 - else if GCType == kMarkSweep or KMarkCompact:
 CollectOldSpaceGarbage

Garbage Collections

Heap::CollectOldSpaceGarbage(...)

- HeapPage::CollectGarbage(...)

${\bf Heap::Collect NewSpace Garbage (...)}$

Scavenger::Scavenge()

Heap::CollectAllGarbage(...): When dart_api_impl.cc or isolate calls NotifyLowMemory(), or when no space when allocating old space.

- Heap::EvacuateNewSpace
- Heap::CollectOldSpaceGarbage

Scavenger

- Scavenge()
 - Prologue(): Creates new SemiSpace & Swap with to_, now from_ is previous SemiSpace.
 - ProcessToSpace(ScavengerVisitor):
 - ScavengePointer(): Mark objects.
 - Move Marked objects to new SemiSpace
 - Epilogue(): Delete from(SemiSpace)

Object → Raw_object

- IsOldObject()
- OldAndNotMarkedBit()

Pages

- PageSpace
 - TryAllocateInternal(): Calls AllocatePage()

Isolate

- Independent worker similar to thread, but not sharing memory.
- Sends message each other

Overview of Jikes Immix.

IMMIX Plan:

• Immix Collector Class

• Immix Mutator Class

This class implements *per-mutator thread* behavior and state for the *Immix* plan, which implements a full-heap immix collector. Specifically, this class defines *Immix* mutator-time allocation and per-mutator thread collection semantics (flushing and restoring per-mutator allocator state).

getAllocatorFromSpace()

- The allocator instance associated with this plan instance which is allocating into space, or null if no appropriate allocator can be established

Alloc()

- This class handles the default allocator from the mark sweep space, and delegates everything else to the superclass.

Policy:

• Block Class

This class defines operations over block-granularity meta-data

sweepOneBlock()

- Returns number of lines marked.
- Line Class

mark()

To mark a line, takes params - address and markvalue

• Collector Local Class

This class implements unsynchronized (local) elements of an immix collector. Marking is done using both a bit in each header's object word, and a mark byte. Sweeping is performed lazily.

resetLineMarksAndDefragStateTable()

- Called on prepare phase of Immix Collector class

sweepAllBlocks()

- Called on release in Immix Collector class. Finish up after a collection. Helps sweeping all the blocks in parallel.

• Immix Space Class

Each instance of this class corresponds to one immix space. Each of the instance methods of this class may be called by any thread (i.e. synchronization must be explicit in any instance or class method). This contrasts with the SquishLocal, where instances correspond to *plan* instances and therefore to kernel threads. Thus unlike this class, synchronization is not necessary in the instance methods of SquishLocal.

decideWhetherToDefrag()

Determine the collection kind.

markLines()

Mark the line/s associated with a given object. This is distinct from the above tracing code because line marks are stored separately from the object headers (thus both must be set), and also because we found empirically that it was more efficient to perform the line mark of the object during the scan phase (which occurs after the trace phase), presumably because the latency of the associated memory operations was better hidden in the context of that code.

getAvailableLines()

Establish the number of recyclable lines lines available for allocation during defragmentation, populating the spillAvailHistogram, which buckets available lines according to the number of holes on the block on which the available lines reside.

getUsableLinesInRegion()

Return the number of lines usable for allocation during defragmentation in the address range specified by start and end. Populate a histogram to indicate where the usable lines reside as a function of block hole count.

• Object Header Class

This class has the object lifespan details. Also, methods to test and mark object header.

testAndMark()

Non-atomically test and set the mark bit of an object.

```
Parameters:
object - The object whose mark bit is to be written
markState - The value to which the mark bits will be set
Returns:
the old mark state
```

testMarkState()

Return true if the mark count for an object has the given value.

Parameters:

object - The object whose mark bit is to be tested value - The value against which the mark bit will be tested *Returns*:

true if the mark bit for the object has the given value.

writeMarkState()

Write the allocState into the mark state fields of an object non-atomically. This is appropriate for collection time initialization.

```
public void prepare(boolean majorGC) {
    if (majorGC) {
        markState = ObjectHeader.deltaMarkState(markState, true);
        lineMarkState++;
        if (VM.VERIFY_ASSERTIONS) VM.assertions._assert(lineMarkState <=
        MAX_LINE_MARK_STATE);
    }
    chunkMap.reset();
    defrag.prepare(chunkMap, this);
    inCollection = true;
}</pre>
```

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Checkpoint 2

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Summary

- Immix in Jikes RVM

- We spent couple of weeks to find out how to debug runtime behavior. We finally got the solution by sending emails to Jikes RVM developer group. (Issue tracker contains the details of the challenge.)
- By using the logging function Jikes RVM developer group suggested, we found out all static OFF set value needed for Dart Immix.
- We analyzed methods calls and relationship between methods in each class.
- We are planning to draw dependency diagram just as the diagram we made for Dart VM. Methods calls in Jikes RVM are much more complex than in Dart VM.
- The detailed information will be added at the end of this report.

Immix in Dart VM

 We spent another couple of weeks to find out how to add new files into build dependencies. Luke helped us to find out where to add file names by looking at each file. (Issue tracker contains the details of the challenge.)

- It took some time to find and to understand which structure to build Immix heap.
- Now, we have implemented blocks and lines.
- Next step is allocation.

Dart Immix Implementation

Modified files

- Isolate.cc (line #: 1090-1129) and isolate.h
 - Only for a main isolate, it creates ImmixHeap.
- heap_sources.gni
 - Include newly added source files in build dependencies
- <u>verifier.cc</u> and <u>verifier.h</u>
 - For include dependencies

New files

- <u>line.h</u>
 - It contains a size for one line and line mark values.
- block.h and block.cc
 - It contains # of lines in block, block size, maximum object size, and Block state values.
 - Methods: getLineTable, getBlockState
- immix heap.h and immix heap.cc
 - It initializes ImmixHeap and its blocks (also line tables and block states).
 - It also contains a debugging method printing states of blocks and lines.

I tested with 10 blocks and 11 lines (first line is not printed since it is used for metadata).

dot means empty. Currently, allocations is not implemented.

 Current setting for block and lines are 10 each for debugging. However, we will use 1000 blocks and 128 lines per block.

Details of Jikes RVM Method Call Dependencies

Line.java

- Constants
 - static final int **LOG_BYTES_IN_LINE_STATUS** = 0;
 - static final int BYTES_IN_LINE_STATUS = 1 << LOG_BYTES_IN_LINE_STATUS = 1;
 - static final int LINE_MARK_TABLE_BYTES = LINES_IN_CHUNK << LOG_BYTES_IN_LINE_STATUS = 16384;
 - static final int LOG_LINE_MARK_BYTES_PER_BLOCK = LOG_LINES_IN_BLOCK + LOG_BYTES_IN_LINE_STATUS = 7;
 - static final int LINE_MARK_BYTES_PER_BLOCK = (1 << LOG_LINE_MARK_BYTES_PER_BLOCK) = 128;
- Methods
 - public static Address align(Address ptr)
 - public static **boolean isAligned**(Address **address**)
 - static int getChunkIndex(Address line)
 - static void mark(Address address, final byte markValue) →
 Line.getMarkAddress(...), Block.isUsed(...)

- static void markMultiLine(Address start, ObjectReference object, final byte markValue) → Line.mark(...), Line.align(...)
- public static **Address getChunkMarkTable**(Address **chunk**) → Line.getMarkAddress(...)
- public static Address getBlockMarkTable(Address block) →
 Line.getMarkAddress(...)
- public static int getNextUnavailable(Address baseLineAvailAddress, int line, final byte unavailableState)
- public static int getNextAvailable(Address baseLineAvailAddress, int line, final byte unavailableState)
- private static Address getMetaAddress(Address address, final int tableOffset) →
 Line.getChunkIndex(...)
- private static Address getMarkAddress(Address address) →
 Line.getMetaAddress(...)

Block.java

- Constants
 - private static final short UNALLOCATED_BLOCK_STATE = 0;
 - private static final short UNMARKED_BLOCK_STATE = (short)
 (MAX_BLOCK_MARK_STATE + 1) = 129;
 - private static final short REUSED_BLOCK_STATE = (short)
 (MAX_BLOCK_MARK_STATE + 2) = 130;
 - private static final short BLOCK_IS_NOT_DEFRAG_SOURCE = 0;
 - private static final short BLOCK_IS_DEFRAG_SOURCE = 1;
 /* block states */
 - static final int LOG_BYTES_IN_BLOCK_STATE_ENTRY = LOG_BYTES_IN_SHORT = 1; // use a short for now
 - static final int BYTES_IN_BLOCK_STATE_ENTRY = 1 <<
 LOG_BYTES_IN_BLOCK_STATE_ENTRY = 2;</pre>
 - static final int **BLOCK_STATE_TABLE_BYTES** = BLOCKS_IN_CHUNK << LOG_BYTES_IN_BLOCK_STATE_ENTRY = 256; /* per-block defrag state */
 - static final int LOG_BYTES_IN_BLOCK_DEFRAG_STATE_ENTRY = LOG_BYTES_IN_SHORT = 1;
 - static final int BYTES_IN_BLOCK_DEFRAG_STATE_ENTRY = 1 << LOG_BYTES_IN_BLOCK_DEFRAG_STATE_ENTRY = 2;</p>
 - static final int BLOCK_DEFRAG_STATE_TABLE_BYTES = BLOCKS_IN_CHUNK <<
 LOG_BYTES_IN_BLOCK_DEFRAG_STATE_ENTRY = 256;</pre>
- Methods
 - static **Address align**(final Address **ptr**)
 - public static **boolean isAligned**(final Address **address**)
 - private static int getChunkIndex(final Address block)
 - public static boolean isUnused(final Address address) →
 Block.getBlockMarkState(...)
 - static boolean isUnusedState(Address cursor)
 - static short getMarkState(Address cursor)
 - static void setState(Address cursor, short value)
 - public static short getBlockMarkState(Address address) →
 Block.getBlockMarkStateAddress(...)
 - static void setBlockAsInUse(Address address) → Block.inUsed(...),
 Block.setBlockState(...)
 - public static void setBlockAsReused(Address address) → Block.inUsed(...),
 Block.setBlockState(...)
 - static void setBlockAsUnallocated(Address address) → Block.inUsed(...),
 Block.getBlockMarkStateAddress(...)
 - private static void setBlockState(Address address, short value) →
 Block.getBlockMarkStateAddress(...)
 - static **Address getBlockMarkStateAddress**(Address **address**) → Block.getChunkIndex(...)
 - static **short sweepOneBlock**(Address **block**, int[] **markHistogram**, final byte **markState**, final boolean **resetMarkState**) → Line.getBlockMarkTable(...), Block.inUsed(...), Block.getDefragStateAddress(...)
 - public static **boolean isDefragSource**(Address **address**) →
 Block.getDefragStateAddress(...)
 - static **void clearConservativeSpillCount**(Address **address**) → Block.getDefragStateAddress(...)

- static short getConservativeSpillCount(Address address) →
 Block.getDefragStateAddress(...)
- static Address getDefragStateAddress (Address address) → Block.getChunkIndex(...)
- static void resetLineMarksAndDefragStateTable(short threshold, Address markStateBase, Address defragStateBase, Address lineMarkBase, int block)

Chunk.java

- Constants
 - private static final int LOG_BYTES_IN_HIGHWATER_ENTRY = LOG_BYTES_IN_ADDRESS;
 - private static final int HIGHWATER_BYTES = 1 << LOG_BYTES_IN_HIGHWATER_ENTRY;
 - private static final int LOG_BYTES_IN_MAP_ENTRY = LOG_BYTES_IN_INT;
 - private static final int MAP_BYTES = 1 << LOG_BYTES_IN_MAP_ENTRY;
 /* byte offsets for each type of metadata */
 - static final int LINE_MARK_TABLE_OFFSET = 0;
 - static final int BLOCK_STATE_TABLE_OFFSET = LINE_MARK_TABLE_OFFSET + Line.LINE_MARK_TABLE_BYTES = 16384;
 - static final int BLOCK_DEFRAG_STATE_TABLE_OFFSET =
 BLOCK_STATE_TABLE_OFFSET + Block.BLOCK_STATE_TABLE_BYTES = 16640;
 - static final int HIGHWATER_OFFSET = BLOCK_DEFRAG_STATE_TABLE_OFFSET + Block.BLOCK_DEFRAG_STATE_TABLE_BYTES = 16896;
 - static final int MAP_OFFSET = HIGHWATER_OFFSET + HIGHWATER_BYTES = 16900;
 - static final int METADATA_BYTES_PER_CHUNK = MAP_OFFSET + MAP_BYTES = 16904;
 - /* FIXME we round the metadata up to block sizes just to ensure the underlying allocator gives us aligned requests */
 - private static final int BLOCK_MASK = (1 << LOG_BYTES_IN_BLOCK) 1;</p>
 - static final int ROUNDED_METADATA_BYTES_PER_CHUNK =
 (METADATA BYTES PER CHUNK + BLOCK MASK) & ~BLOCK MASK = 32768;
 - static final int ROUNDED_METADATA_PAGES_PER_CHUNK =
 ROUNDED_METADATA_BYTES_PER_CHUNK >> LOG_BYTES_IN_PAGE = 8;
 - public static final int FIRST_USABLE_BLOCK_INDEX =
 ROUNDED_METADATA_BYTES_PER_CHUNK >> LOG_BYTES_IN_BLOCK = 1;
- Methods
 - public static **Address align**(Address **ptr**)
 - static boolean isAligned(Address ptr)
 - static int getByteOffset(Address ptr)
 - static int getRequiredMetaDataPages()
 - static void sweep(Address chunk, Address end, ImmixSpace space, int[]
 markHistogram, final byte markValue, final boolean resetMarks)

 → Block.isUnused(...), Block.isUnusedState(...), Block.setState(...),
 Block.getBlockMarkState(...), Block.getBlockMarkStateAddress(...),

 $Block.sweepOneBlock(...),\ Block.isDefragSource(...),\ Chunk.getFirstUsableBlock(...), \\ ImmixSpace.inImmixDefragCollection()$

- static void clearMetaData(Address chunk) → Chunk.checkMetaDataCleared(...)
- private static void checkMetaDataCleared(Address chunk, Address value) →
 Block.isUnused(...), Chunk.getHighWater(...)
- static void updateHighWater(Address value) → Chunk.setHighWater(...),
 Chunk.getHighWater(...)
- private static void setHighWater(Address chunk, Address value)
- public static Address getHighWater(Address chunk)
- static void setMap(Address chunk, int value)
- static int getMap(Address chunk)
- static void resetLineMarksAndDefragStateTable(Address chunk, short threshold) → Line.getChunkMarkTable(...), Block.getBlockMarkStateAddress(...), Block.getDefragStateAddress(...), Block.resetLineMarksAndDefragStateTable(...)
- static Address getFirstUsableBlock(Address chunk)

ChunkList.java

- Constants
 - private static final int LOG_PAGES_IN_CHUNK_MAP_BLOCK = 0;
 - private static final int LOG_ENTRIES_IN_CHUNK_MAP_BLOCK = LOG_BYTES_IN_PAGE + LOG_PAGES_IN_CHUNK_MAP_BLOCK -LOG_BYTES_IN_ADDRESS;

- private static final int ENTRIES_IN_CHUNK_MAP_BLOCK = 1 <<
 LOG_ENTRIES_IN_CHUNK_MAP_BLOCK;
- private static final int CHUNK_MAP_BLOCKS = 1 << 4 = 16;
- private static final int MAX_ENTRIES_IN_CHUNK_MAP =
 ENTRIES_IN_CHUNK_MAP_BLOCK * CHUNK_MAP_BLOCKS;
- private final AddressArray chunkMap = AddressArray.create(CHUNK_MAP_BLOCKS);
- private int chunkMapLimit = -1;
- private int chunkMapCursor = -1;
- Methods
 - void reset()
 - public **Address getHeadChunk()** → ChunkList.getMapAddress(...)
 - public **Address getTailChunk()** → ChunkList.getMapAddress(...)
 - void addNewChunkToMap(Address chunk) → ChunkList.getChunkIndex(...),
 Chunk.setMap(...), ChunkList.getChunkMap(...), ChunkList.checkMap(),
 ChunkList.consolidateMap()
 - void removeChunkFromMap(Address chunk) → Chunk.setMap(...),
 Chunk.getMap(...), ChunkList.getMapAddress(...), ChunkList.checkMap()
 - private int getChunkIndex(int entry)
 - private int getChunkMap(int entry)
 - private Address getMapAddress(int entry) → ChunkList.getChunkIndex(...),
 ChunkList.getChunkMap(...)
 - public Address nextChunk(Address chunk) → ChunkList.getMapAddress(...)
 - private Address nextChunk(final Address chunk, final Address limit) →
 Chunk.getMap(...)
 - public Address nextChunk(final Address chunk, final int start, final int stride) →
 Chunk.getMap(...)
 - private Address nextChunk(int entry, final int start, final int stride)
 - public Address firstChunk(int ordinal, int stride) → ChunkList.getMapAddress(...)
 → ChunkList.nextChunk(...), ChunkList.checkMap()
 - $private \ void \ checkMap() \rightarrow Chunk.getMap(...)$, ChunkList.getMapAddress(...)
 - public void consolidateMap() → Chunk.setMap(...), ChunkList.getMapAddress(...),
 ChunkList.checkMap()

ImmixConstants.java

- Constants
 - public static final boolean **BUILD_FOR_STICKYIMMIX** =

Plan.NEEDS_LOG_BIT_IN_HEADER;

/* start temporary experimental constants --- should not be allowed to lurk longer than necessary */

- public static final int TMP_MIN_SPILL_THRESHOLD = 2;
- public static final boolean PREFER_COPY_ON_NURSERY_GC = true;
 /* end temporary experimental constants */
- static final byte MAX_LINE_MARK_STATE = 127;
- static final byte RESET_LINE_MARK_STATE = 1;
- public static final boolean MARK_LINE_AT_SCAN_TIME = true; // else do it at mark time
- public static final boolean SANITY_CHECK_LINE_MARKS = false &&
 VM.VERIFY_ASSERTIONS;
- public static final float **DEFAULT_LINE_REUSE_RATIO** = (float) 0.99;
- public static final float **DEFAULT_DEFRAG_LINE_REUSE_RATIO** = (float) 0.99;
- public static final float **DEFAULT_SIMPLE_SPILL_THRESHOLD** = (float) 0.25;
- public static final int **DEFAULT_DEFRAG_HEADROOM** = $\frac{0}{7}$; // number of pages.
- public static final float **DEFAULT_DEFRAG_HEADROOM_FRACTION** = (float) 0.020;
- public static final int **DEFAULT_DEFRAG_FREE_HEADROOM** = 0; // number of pages. This should only deviate from zero for analytical purposes. Otherwise the defragmenter is cheating!
- public static final float **DEFAULT_DEFRAG_FREE_HEADROOM_FRACTION** = (float)
 0.0;

/* sizes etc */

- static final int LOG_BYTES_IN_BLOCK = (LOG_BYTES_IN_PAGE > 15 ?
 LOG_BYTES_IN_PAGE : 15) = 15;
- public static final int **BYTES_IN_BLOCK** = 1 << LOG_BYTES_IN_BLOCK = 32768;
- static final int LOG_PAGES_IN_BLOCK = LOG_BYTES_IN_BLOCK -LOG_BYTES_IN_PAGE = 3;
- static final int **PAGES_IN_BLOCK** = 1 << LOG_PAGES_IN_BLOCK = 8;
- static final int **LOG_BLOCKS_IN_CHUNK** = LOG_BYTES_IN_CHUNK LOG_BYTES_IN_BLOCK = 7;

```
- static final int BLOCKS_IN_CHUNK = 1 << LOG_BLOCKS_IN_CHUNK = 128;
```

- public static final int LOG_BYTES_IN_LINE = 8;
- static final int LOG_LINES_IN_BLOCK = LOG_BYTES_IN_BLOCK LOG_BYTES_IN_LINE = 7;
- public static final short LINES_IN_BLOCK = (short) (1 << LOG_LINES_IN_BLOCK) =
 128;
- static final int LOG_LINES_IN_CHUNK = LOG_BYTES_IN_CHUNK LOG_BYTES_IN_LINE = 14;
- static final int **LINES_IN_CHUNK** = 1 << LOG_LINES_IN_CHUNK = 16384;
- public static final int BYTES_IN_LINE = 1 << LOG_BYTES_IN_LINE = 256;
- public static final int MAX_IMMIX_OBJECT_BYTES = BYTES_IN_BLOCK >> 1 = 16384;
- private static final int LOG_BLOCKS_IN_RECYCLE_ALLOC_CHUNK = 4; // 3 + 15
 -> 19 (512KB);
- private static final int LOG_BYTES_IN_RECYCLE_ALLOC_CHUNK = LOG_BLOCKS_IN_RECYCLE_ALLOC_CHUNK + LOG_BYTES_IN_BLOCK;
- static final int BYTES_IN_RECYCLE_ALLOC_CHUNK = 1 <<
 LOG_BYTES_IN_RECYCLE_ALLOC_CHUNK = 524288;
- public static final short MAX_BLOCK_MARK_STATE = LINES_IN_BLOCK = 128;
- static final short MAX_CONSV_SPILL_COUNT = (short) (LINES_IN_BLOCK / 2) =
 64;
- public static final short **SPILL_HISTOGRAM_BUCKETS** = (short) (MAX_CONSV_SPILL_COUNT + 1) = <mark>65</mark>;
- public static final short MARK_HISTOGRAM_BUCKETS = (short) (LINES_IN_BLOCK + 1) = 129;
- public static final Word RECYCLE_ALLOC_CHUNK_MASK =
 Word.fromIntZeroExtend(BYTES_IN_RECYCLE_ALLOC_CHUNK 1) = 0x0007ffff =
 524287;
- protected static final Word CHUNK_MASK =
 Word.fromIntZeroExtend(BYTES_IN_CHUNK 1) = 0x003fffff = 4194303;
- public static final Word **BLOCK_MASK** = Word.fromIntZeroExtend(BYTES_IN_BLOCK 1) = 0x00007fff = 32767;
- protected static final Word LINE_MASK = Word.fromIntZeroExtend(BYTES_IN_LINE 1) = 0x000000ff = 255;

CollectorLocal.java

- Methods
 - public **CollectorLocal**(ImmixSpace **space**)
 - public void prepare(boolean majorGC) →
 CollectorLocal.resetLineMarksAndDefragStateTable(...),
 ImmixSpace.inImmixDefragCollection()
 - private void resetLineMarksAndDefragStateTable(int ordinal, final short threshold) → Chunk.resetLineMarksAndDefragStateTable(...),
 ChunkList.nextChunk(...), ChunkList.firstChunk(...),
 ImmixSpace.inImmixDefragCollection()
 - public **void release**(boolean **majorGC**) → CollectorLocal.sweepAllBlocks(...)
 - private void sweepAllBlocks(boolean majorGC) → Chunk.sweep(...),
 Chunk.getHighWater(...), ChunkList.nextChunk(...), ChunkList.firstChunk(...),
 Defrag.getAndZeroSpillMarkHistogram(...)

Defrag.java

- Methods
 - Defrag(FreeListPageResource pr)
 - void prepareHistograms()
 - boolean inDefrag()
 - void prepare(ChunkList chunkMap, ImmixSpace space) →
 ChunkList.consolidateMap(), Defrag.establishDefragSpillThreshold(...)
 - void globalRelease()
 - int getDefragHeadroomPages() → Defrag.prepare(...)
 - void decideWhetherToDefrag(boolean emergencyCollection, boolean collectWholeHeap, int collectionAttempt, boolean userTriggered, boolean exhaustedReusableSpace)
 - boolean determined (boolean in Defrag)
 - void getBlock()
 - private void establishDefragSpillThreshold(ChunkList chunkMap, ImmixSpace
 space) → ImmixSpace.getAvailableLine(...)
 - boolean spaceExhausted()

int[] getAndZeroSpillMarkHistogram(int ordinal)

ImmixSpace.java extends Space

- Methods
 - public ImmixSpace(String name, VMRequest vmRequest)
 - public ImmixSpace(String name, boolean zeroed, VMRequest vmRequest) → Chunk.getRequiredMetaDataPages(...)
 - public **void initializeDefrag()** → Defrag.prepareHistograms()
 - public void prepare(boolean majorGC) → ChunkList.reset(), Defrag.prepare(...),
 ObjectHeader.deltaMarkState(...)
 - public boolean release(boolean majorGC) → ChunkList.reset(),
 ChunkList.getHeadChunk(), Defrag.globalRelease(),
 ImmixSpace.isRecycleAllocChunkAligned(...), ObjectHeader.pinObject(...)
 - public void decideWhetherToDefrag(boolean emergencyCollection, boolean collectWholeHeap, int collectionAttempt, boolean userTriggeredCollection)
 → Defrag.decideWhetherToDefrag(...)
 - public int defragHeadroomPages() → Defrag.getDefragHeadroomPages()
 - public **boolean inImmixCollection()**
 - public **boolean inImmixDefragCollection()** → Defrag.inDefrag()
 - public int getPagesAllocated()
 - public static **short getReusuableMarkStateThreshold**(boolean **forDefrag**)
 - public Address getSpace(boolean hot, boolean copy, int lineUseCount)
 → Block.setBlockAsInUse(...), Block.isDefragSource(...), Chunk.updateHighWater(...),
 Defrag.getBlock()
 - public void growSpace(Address start, Extent bytes, boolean newChunk) →
 Chunk.clearMetaData(...), ChunkList.addNewChunkToMap(...)
 - public Address acquireReusableBlocks() → Chunk.getHighWater(...),
 ChunkList.nextChunk(...), ImmixSpace.isRecycleAllocChunkAligned(...)
 - public void release(Address block) → Block.setBlockAsUnallocated(...),
 Defrag.inDefrag()
 - public **int releaseDiscontiguousChunks**(Address **chunk**) → ChunkList.removeChunkFromMap(...)
 - public void postAlloc(ObjectReference object, int bytes) →
 ObjectHeader.isNewObject(...), ObjectHeader.markAsStraddling(...)
 - public void postCopy(ObjectReference object, int bytes, boolean majorGC) →
 ObjectHeader.writeMarkState(...)
 - public ObjectReference traceObject(TransitiveClosure trace, ObjectReference object, int allocator) → Defrag.determined(...), Defrag.spaceExhausted(), ImmixSpace.traceObjectWithoutMoving(...), ImmixSpace.traceObjectWithOpportunisticCopy(...), ImmixSpace.isDefragSource(...), ObjectHeader.traceObject(...)
 - public ObjectReference fastTraceObject(TransitiveClosure trace, ObjectReference object) → Defrag.determined(...), ImmixSpace.traceObjectWithoutMoving(...)
 - public ObjectReference nurseryTraceObject(TransitiveClosure trace,
 ObjectReference object, int allocator)
 - $\rightarrow {\sf Defrag.inDefrag(), ImmixSpace.traceObjectWithOpportunisticCopy(...),} \\ ObjectHeader.isMatureObject(...)$
 - public ObjectReference traceObject(TransitiveClosure trace, ObjectReference object)
 - private void traceObjectWithoutMoving(TransitiveClosure trace, ObjectReference object)
 - → Defrag.inDefrag(), Defrag.spaceExhausted(), ImmixSpace.isDefragSource(...), ObjectHeader.testAndMark(...)
 - private ObjectReference traceObjectWithOpportunisticCopy(TransitiveClosure trace, ObjectReference object, int allocator, boolean nurseryCollection)
 - → Defrag.inDefrag(), Defrag.determined(...), Defrag.spaceExhausted(), ImmixSpace.isDefragSource(...), ObjectHeader.setMarkStateUnlogAndUnlock(...), ObjectHeader.testMarkState(...), ObjectHeader.isMatureObject(...), ObjectHeader.traceObject(...), ObjectHeader.returnToPriorStateAndEnsureUnlogged(...)
 - public void markLines(ObjectReference object) → Line.mark(...), Line.markMultiLine(...), ObjectHeader.isStraddlingObject(...)
 - public int getNextUnavailableLine(Address baseLineAvailAddress, int line) →
 Line.getNextUnavailable(...)
 - public int getNextAvailableLine(Address baseLineAvailAddress, int line) →
 Line.getNextAvailable(...)
 - int getAvailableLines(int[] spillAvailHistogram) → ChunkList.getHeadChunk(),
 ImmixSpace.getUsableLinesInRegion(...)

```
private int getUsableLinesInRegion(Address start, Address end, int[]
              spillAvailHistogram)
               → Block.getBlockMarkStateAddress(...), Chunk.getHighWater(...),
              Block.getConservativeSpillCount(...), Chunk.getByteOffset(...),
              ChunkList.nextChunk(...)
              public boolean isLive(ObjectReference object) → Defrag.inDefrag(),
              ImmixSpace.isDefragSource(...), ObjectHeader.testMarkState(...)
              public boolean copyNurseryIsLive(ObjectReference object) →
              ObjectHeader.testMarkState(...)
              public boolean fastIsLive(ObjectReference object) → Defrag.inDefrag(),
              ObjectHeader.testMarkState(...)
              public boolean willNotMoveThisGC(ObjectReference object) → Defrag.inDefrag(),
              ObjectHeader.traceObject(...)
              public boolean willNotMoveThisNurseryGC(ObjectReference object) →
              ObjectHeader.isMatureObject(...)
              private boolean isDefragSource(ObjectReference object) →
              Block.isDefragSource(...)
              public boolean willNotMoveThisGC(Address address) → Defrag.inDefrag(),
              Defrag.spaceExhausted(), ImmixSpace.isDefragSource(...)
              public boolean isDefragSource(Address address)
              private void lock()
                      Handle depending on whether GC or mutator
              private void unlock()
                     Handle depending on whether GC or mutator
              public static boolean isRecycleAllocChunkAligned(Address ptr)
Space.java
       Constants
              /* Class variables */
              private static boolean DEBUG = false;
              private static final boolean FORCE_SLOW_MAP_LOOKUP = false;
              private static final int PAGES = 0;
              private static final int MB = 1;
              private static final int PAGES_MB = 2;
              private static final int MB_PAGES = 3;
              private static int spaceCount = 0;
              private static Space[] spaces = new Space[MAX_SPACES];
              private static Address heapCursor = HEAP_START;
              private static Address heapLimit = HEAP_END;
       Methods
              protected Space(String name, boolean movable, boolean immortal, boolean
              zeroed, VMRequest vmRequest)
              public static Address getDiscontigStart()
              public static Address getDiscontigEnd()
              public final String getName()
              public final Address getStart()
              public final Extent getExtent()
              public final int getDescriptor()
              public final int getIndex()
              public final boolean isImmortal()
              public boolean isMovable()
              public final int reservedPages()
              public final int committedPages()
              public final int availablePhysicalPages()
              public static long cumulativeCommittedPages()
              public static boolean isImmortal(ObjectReference object)
              public static boolean isMovable(ObjectReference object)
              public static boolean isMappedObject(ObjectReference object)
              public static boolean isMappedAddress(Address address)
              public static boolean isInSpace(int descriptor, ObjectReference object)
              public static boolean isInSpace(int descriptor, Address address)
              public static Space getSpaceForObject(ObjectReference object)
              public static Space getSpaceForAddress(Address addr)
               public void setZeroingApproach(boolean useNT, boolean concurrent)
              public void skipConcurrentZeroing()
              public void triggerConcurrentZeroing()
```

public final Address acquire(int pages)

```
- public Address growDiscontiguousSpace(int chunks)
```

- public static int requiredChunks(int pages)
- public void growSpace(Address start, Extent bytes, boolean newChunk)
- public int releaseDiscontiguousChunks(Address chunk)
- public Address getHeadDiscontiguousRegion()
- public void releaseAllChunks()
- public abstract void release(Address start);
- private static int getPagesReserved()
- public static void printUsageMB()
- public static void printUsagePages()
- public static void printVMMap()
- public static void visitSpaces(SpaceVisitor v)
- public static void eagerlyMmapMMTkSpaces()
- public static void eagerlyMmapMMTkContiguousSpaces()
- public static void eagerlyMmapMMTkDiscontiguousSpaces()
- private static void printUsage(int mode)
- private static void printPages(int pages, int mode)
- public abstract ObjectReference traceObject(TransitiveClosure trace, ObjectReference object);
- public boolean isReachable(ObjectReference object)
- public abstract boolean isLive(ObjectReference object);
- public static Extent getFracAvailable(float frac)
- public static int getSpaceCount()
- public static Space[] getSpaces()

MutatorLocal.java extends ImmixAllocator

- Methods
 - public MutatorLocal(ImmixSpace space, boolean hot)
 - public void prepare()
 - public void release()

ObjectHeader.java

- Constants
 - /* number of header bits we may use */
 - static final int **AVAILABLE_LOCAL_BITS** = 8 HeaderByte.USED_GLOBAL_BITS = 8; /* header requirements */
 - public static final int LOCAL_GC_BITS_REQUIRED = AVAILABLE_LOCAL_BITS = 8;
 - public static final int GLOBAL_GC_BITS_REQUIRED = 0;
 - public static final int GC_HEADER_WORDS_REQUIRED = 0;
 /* local status bits */
 - static final byte NEW_OBJECT_MARK = 0; // using zero means no need for explicit initialization on allocation
 - public static final int **PINNED_BIT_NUMBER** = ForwardingWord.FORWARDING_BITS
 = 2;
 - public static final byte PINNED_BIT = 1 << PINNED_BIT_NUMBER = 4;
 - private static final int $STRADDLE_BIT_NUMBER = PINNED_BIT_NUMBER + 1;$
 - public static final byte STRADDLE_BIT = 1 << STRADDLE_BIT_NUMBER = 8;
 /* mark bits */</pre>
 - private static final int MARK_BASE = STRADDLE_BIT_NUMBER + 1;
 - static final int MAX_MARKCOUNT_BITS = AVAILABLE_LOCAL_BITS MARK_BASE =
 4;
 - private static final byte MARK_INCREMENT = 1 << MARK_BASE;</p>

 - private static final byte MARK_AND_FORWARDING_MASK = (byte) (MARK_MASK | ForwardingWord.FORWARDING_MASK);
 - public static final byte **MARK_BASE_VALUE** = MARK_INCREMENT = 16;
- Methods
 - static **byte testAndMark**(ObjectReference **object**, byte **markState**)
 - static void setMarkStateUnlogAndUnlock(ObjectReference object, byte gcByte, byte markState)
 - static **boolean testMarkState**(ObjectReference **object**, byte **value**)
 - static boolean testMarkState(byte gcByte, byte value)
 - static boolean isNewObject(ObjectReference object)
 - static boolean isMatureObject(ObjectReference object)
 - static void markAsStraddling(ObjectReference object)
 - static boolean isStraddlingObject(ObjectReference object)

- public static **void pinObject**(ObjectReference **object**)
- static boolean isPinnedObject(ObjectReference object)
- static void writeMarkState (ObjectReference object, byte markState, boolean straddle)
- static void returnToPriorStateAndEnsureUnlogged(ObjectReference object, byte status)
- static byte deltaMarkState(byte state, boolean increment)

ImmixAllocator.java extends Allocator

- Methods
 - public ImmixAllocator(ImmixSpace space, boolean hot, boolean copy)
 - public void reset()
 - public final Address alloc(int bytes, int align, int offset) →
 ImmixAllocator.overflowAlloc(...), ImmixAllocator.allocSlowHot(...)
 - public final Address overflowAlloc(int bytes, int align, int offset)
 - public final boolean getLastAllocLineSt raddle()
 - protected final **Address allocSlowOnce**(int **bytes**, int **align**, int **offset**) → ImmixSpace.getSpace(...), ImmixAllocator.alloc(...), ImmixSpace.getSpace()
 - private Address allocSlowHot(int bytes, int align, int offset) →
 ImmixAllocator.alloc(...), ImmixAllocator.acquireRecyclableLines(...)
 - private boolean acquireRecyclableLines(int bytes, int align, int offset)
 → Block.isUnused(...), Block.isDefragSource(...),
 ImmixSpace.getNextUnavailableLine(...), ImmixSpace.getNextAvailableLine(...),
 ImmixAllocator.acquireRecyclableBlock()
 - private boolean acquireRecyclableBlock() →
 ImmixAllocator.acquireRecyclableBlockAddressOrder()
 - private boolean acquireRecyclableBlockAddressOrder() →
 Block.getBlockMarkState(...), Block.setBlockAsReused(...), Block.isDefragSource(...),
 ImmixSpace.getReusableMarkStateThreshold(...),
 ImmixSpace.acquireReusableBlocks(), ImmixSpace.isRecycleAllocChunkAligned(...)
 - private void zeroBlock(Address block)
 - public final Space getSpace()
 - public final void show()

Allocator.java

- Methods
 - public static **int determineCollectionAttempts**()
 - protected abstract Space getSpace();
 - public static Address alignAllocation(Address region, int alignment, int offset, int knownAlignment, boolean fillAlignmentGap)
 - public static void fillAlignmentGap(Address start, Address end)
 - public static **Address alignAllocation**(Address **region**, int **alignment**, int **offset**)
 - public static Address alignAllocationNoFill(Address region, int alignment, int offset)
 - public static int getMaximumAlignedSize(int size, int alignment)
 - public static int getMaximumAlignedSize(int size, int alignment, int knownAlignment)
 - protected abstract Address allocSlowOnce(int bytes, int alignment, int offset);
 - public final Address allocSlow(int bytes, int alignment, int offset)
 - public final Address allocSlowInline(int bytes, int alignment, int offset)

Plan/immix/Immix.java

- Methods
 - public void collectionPhase(short phaseId) → ImmixSpace.prepare(...),
 ImmixSpace.release(...), ImmixSpace.decideWhetherToDefrag(...),
 ImmixSpace.release(...)
 - public boolean lastCollectionWasExhaustive()
 - public int getPagesUsed()
 - $\quad \text{public } \textbf{int getCollectionReserve()} \rightarrow \text{ImmixSpace.defragHeadroomPages()} \\$
 - public boolean willNeverMove(ObjectReference object)
 - protected void registerSpecializedMethods()
 - public **void preCollectorSpawn()** → ImmixSpace.initializeDefrag()

plan/immix/ImmixCollector.java

- Methods

- public ImmixCollector()
- public Address allocCopy(ObjectReference original, int bytes, int align, int offset, int allocator) → ImmixSpace.inImmixDefragCollection()
- public void postCopy(ObjectReference object, ObjectReference typeRef, int bytes, int allocator) → ImmixSpace.postCopy(...)
- public void collectionPhase(short phaseId, boolean primary) →
 CollectorLocal.release()
- private static Immix global()
- public final **TraceLocal getCurrentTrace**()

plan/immix/ImmixMutator.java

- Methods
 - public **ImmixMutator**()
 - public Address alloc(int bytes, int align, int offset, int allocator, int site)
 - public void postAlloc(ObjectReference ref, ObjectReference typeRef, int bytes, int allocator) → ImmixSpace.postAlloc(...)
 - public Allocator getAllocatorFromSpace(Space space)
 - public **void collectionPhase**(short **phaseId**, boolean **primary**)

plan/immix/ImmixTraceLocal.java

- Methods
 - public **ImmixTraceLocal**(Trace **trace**, ObjectReferenceDeque **modBuffer**)
 - public **boolean isLive**(ObjectReference **object**) \rightarrow ImmixSpace.fastIsLive(...)
 - public ObjectReference traceObject(ObjectReference object)
 - public **boolean willNotMoveInCurrentCollection**(ObjectReference **object**)
 - protected **void scanObject**(ObjectReference **object**) → ImmixSpace.markLines(...)
 - protected void processRememberedSets()

plan/immix/ImmixDefragTraceLocal.java

- Methods
 - public ImmixDefragTraceLocal(Trace trace, ObjectReferenceDeque modBuffer)
 - public boolean isLive(ObjectReference object) →
 ImmixSpace.inImmixDefragCollection(), ImmixSpace.isLive(...)
 - public ObjectReference traceObject(ObjectReference object) →
 ImmixSpace.inImmixDefragCollection(), ImmixSpace.traceObject(...)
 - public **boolean willNotMoveInCurrentCollection**(ObjectReference **object**) → ImmixSpace.inImmixDefragCollection(), ImmixSpace.willNotMoveThisGC(...)
 - protected void scanObject(ObjectReference object) →
 ImmixSpace.inImmixDefragCollection(), ImmixSpace.markLines(...)
 - protected **void processRememberedSets()**

Ch	ec	kŗ	00	int	: 3

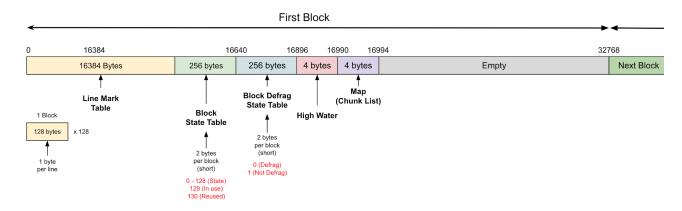
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Summary

- Jikes RVM
 - Created an overall structure of Jikes RVM and Chunk.
- <u>Dart</u>
 - Implemented an array of block addresses.
 - Implemented allocation and tested for couple of different object sizes.

Jikes RVM

Chunk (4MB) = 128 Blocks, Block (32KB) = 128 Lines, Line (256B)

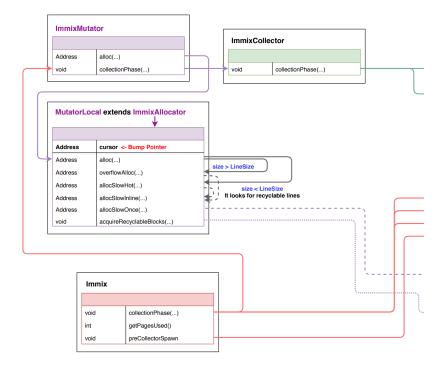


- There was another container called chunk.
- A chunk contains 128 blocks and tables for blocks.
- Chunk is that contains all metadata for lines and blocks.
- On top, there is a chunk map that contains addresses of all chunks created.

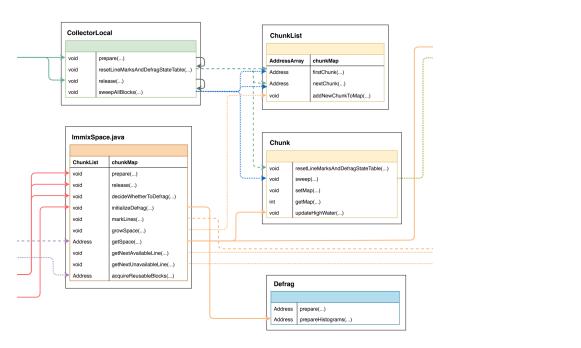
Structure

- <u>Immix</u>
 - Immix creates a ImmixSpace that handles all chunks, blocks, and lines.
 - Immix also initiate collectionPhase for garbage collection.
- <u>ImmixMutator</u> and <u>ImmixAllocator</u>
 - These classes handle actual allocations. Depending on object size, whether to call overflowAlloc (bigger than a line) or overSlowHot (smaller than a line).
 - It also initiates ImmixCollector to garbage collect.
- ImmixSpace
 - It has all the functions handling chunks, blocks, and lines.
 (e.g., marking lines, growing spaces, getting available lines, and etc.)
- Chunk, Block, and Line
 - These are all pointer addresses, and they all have each own helper functions and offsets.

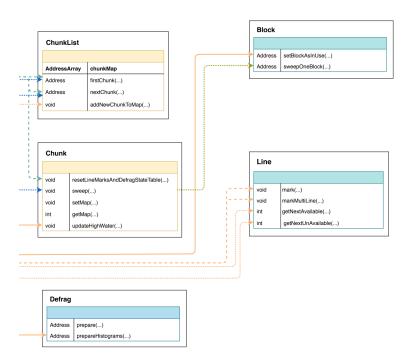
There are actual structure images below (Cut into 3 pieces due to limited space).



.....



.....

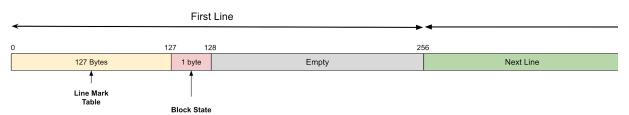


Dart Immix

blockAddresses_*

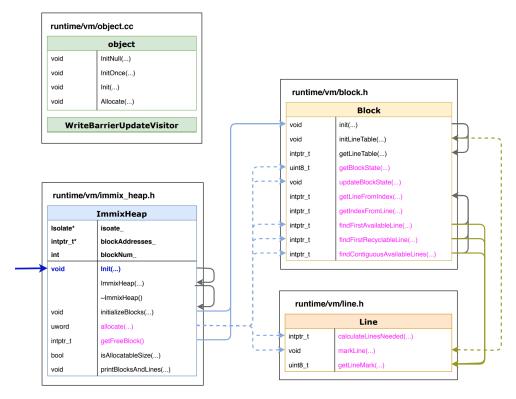
→ Contains all addresses of blocks

Block (32KB) = 128 Lines, Line (256B)



- Different with Immix in Jikes RVM, we keep metadata (e.g., line mark tables and block state) in each first line of a block.
- In ImmixHeap, we keep addresses of all the blocks in blockAddresses_*. This will be updated when getting a new block.

Allocation



- All methods colored with pink are implemented after checkpoint 2, and those methods are related with allocate function in ImmixHeap.
- All methods in Block and Line are helper functions dealing with pointer addresses.
- getFreeBlock() method in ImmixHeap gets another free heap when there are no more spaces to put certain object. An address of a new block is added to blockAddresses_, and also blockNum_ increases by one.
- As shown above, object.cc does not call allocate method of our ImmixHeap yet. This is a future work.

Test Results (To make it simple, I created 3 blocks and 11 lines for each block)

Allocating 10 small objects.

```
/* Test for 10 smalll allocations */
printf(" --> %ld\n", immixHeap->allocate(60));
printf(" --> %ld\n", immixHeap->allocate(237));
printf(" --> %ld\n", immixHeap->allocate(123));
printf(" --> %ld\n", immixHeap->allocate(43));
printf(" --> %ld\n", immixHeap->allocate(62));
printf(" --> %ld\n", immixHeap->allocate(190));
printf(" --> %ld\n", immixHeap->allocate(23));
printf(" --> %ld\n", immixHeap->allocate(203));
printf(" --> %ld\n", immixHeap->allocate(87));
printf(" --> %ld\n", immixHeap->allocate(193));
result->immixHeap()->printBlocksAndLines();
```

- It allocates 10 small object (size is less than 256 bytes = line)
- Each mark means different status (#: Unavailable, *: Recyclable, .: Not used)

- We can see that 10 objects were allocated in first block, and not the block is unavailable.
- Currently even object is smaller than a line, it gets regular marking.
- Allocating a big objects

```
/* Test for big allocations */
printf(" --> %ld\n", immixHeap->allocate(970));
printf(" --> %ld\n", immixHeap->allocate(800));
printf(" --> %ld\n", immixHeap->allocate(2500));
printf(" --> %ld\n", immixHeap->allocate(1000));
printf(" --> %ld\n", immixHeap->allocate(200));
result->immixHeap()->printBlocksAndLines();
```

```
- Allocate an object (970 bytes): 4 lines
--> 94222561690112
- Allocate an object (800 bytes): 4 lines
--> 94222561691136
- Cannot allocate an object (2500 bytes)
--> 0
- Allocate an object (1000 bytes): 4 lines
--> 94222561692928
- Allocate an object (200 bytes): 1 Line
--> 94222561692160
----- Printing Blocks & Lines -----
[#] # # # # # # # #
[*] # # # # # # # .
[*] # # # # # # # .
```

- The allocation continues from the blocks in previous test.
- Third object is larger than half size of a block (max object size possible), and it gives an 0 because it cannot be allocated.
- It starts allocate other 4 objects from a second block. After allocating 2 objects, 1000 bytes object is allocated on next line because there are no enough contiguous lines available
- Last object can be allocated in second block and gets allocated there.
- Requesting a new block

```
/* Test for requesting new block */
printf(" --> %ld\n", immixHeap->allocate(500));
printf(" --> %ld\n", immixHeap->allocate(500));
printf(" --> %ld\n", immixHeap->allocate(1000));
result->immixHeap()->printBlocksAndLines();
exit(0);
```

```
- Allocate an object (500 bytes): 2 lines
--> 94222561693952
- Allocate an object (500 bytes): 2 lines
--> 94222561694464
- Allocate an object (1000 bytes): 4 lines
--> ImmixHeap::getFreeBlock()
--> 94222561695744

---- Printing Blocks & Lines ----
[#] # # # # # # # #
[*] # # # # # # # # .
[*] # # # # # # # .
[*] # # # # # # # .
```

- After allocating first two objects in 3rd block, it requests a new free block to allocate 1000 bytes object.

Future Work

- Object.cc in Dart
 - Intercept original allocate calls and bypass to our *ImmixHeap* allocate method.
 - It will not be simple because there are lots of methods that interact with original heap. For all of those methods, we need to handle with our own ImmixHeap.
 - Goal is printing "Hello World" only with allocation methods, it will crash at some point.
- Implicit marking for single line
 - We already created helper functions for implicit marking, but it currently marks only normally.

-	When implementing, <i>cursor</i> and <i>limit</i> should be also implemented to track exact address where to start allocating.
	g.