

# Problem Set 7

---

Justin Ely

615.202.81.FA15 Data Structures

20 October, 2015

---

1a)

- The 0th level (root) node has exactly 0 ancestors
- the first level will have exactly 1 ancestor, the root node
- the second level will have exactly 2 ancestors (1 node at level 1 and the root)
- ...

Following this trend, a node at level  $n$  will have exactly  $n$  ancestors.

1b)

- A tree with 1 leaf has 1 node
- A tree with 2 leaves has 3 total nodes
- A tree with 3 leaves has 5 nodes
- A tree with 4 leaves has 7 nodes
- ...

Following this trend, a regular binary tree with  $n$  leaves has  $2n-1$  total nodes.

2)

3)

For a general  $m$ -ary tree, the total number of pointers set aside will be equal to  $n \times m$ . However, each node  $n$  (apart from the root) will remove 1 null pointer from this count as the pointer will connect to its parent.

$$N_{null} = (n * m) - (n - 1) \quad (1)$$

$$= (n * m) - n + 1 \quad (2)$$

$$= n * (m - 1) + 1 \quad (3)$$

4)

```
class Tree
    data = Array()

    def maketree(value)
        newNode = Node()

        newNode.value = value
        newNode.left = Null
        newNode.right = Null

        return newNode

    def traverse(Node)
        If not Node == null:
            traverse(Node.left)
            print Node.value
            traverse(Node.right)

    def setleft(value, parent)
        if parent.left != null:
            raise Error

        child = maketree(value)
        child.right = parent

        data[IndexOf(parent) * 2] = child

    def setright(value, parent)
        child = maketree(value)
        child.right = parent.right

        data[IndexOf(parent)*2 +1] = child
```

**5a)**

```
def fib(n)
    if n = 1:
        Here.left = 0
    else:
        Here.left = fib(n-1)
        Here.right = fib(n-2)
```

**5b)**

No, the 1 node will only have an n-1 child, and not an n-2 child.

**5c)**

An order N fibonacci tree will have  $\text{fib}(n-1) + \text{fib}(n-2)$  leaves.

**5d)**

A fibonacci tree of order n has depth n.