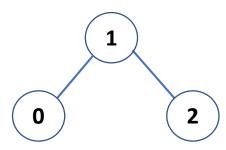
# Balanced BSTs

<b>Abstract Data Type</b>	Operations	Data Structures
Dictionary		

<b>Abstract Data Type</b>	Operations	Data Structures
Dictionary	Search, Insert, Delete	

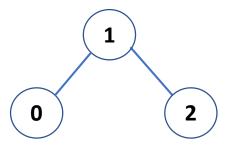
Operations	Data Structures
Search, Insert, Delete	BSTs

Suppose we start with a "balanced" BST



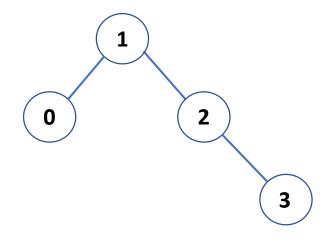
Suppose we start with a "balanced" BST

Insert keys 3, 4, 5, ..., n



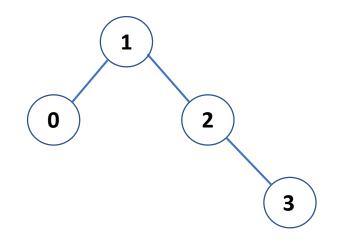
Suppose we start with a "balanced" BST

Insert keys 3, 4, 5, ..., n



Suppose we start with a "balanced" BST

Insert keys 3, 4, 5, ..., n



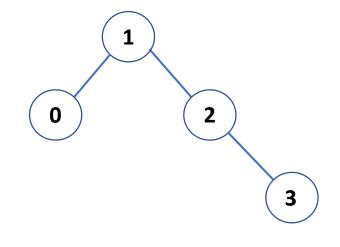
height =  $\Theta(n)$ 

n

Suppose we start with a "balanced" BST

Insert keys 3, 4, 5, ..., n

**Search** takes  $\Theta(n)$ !



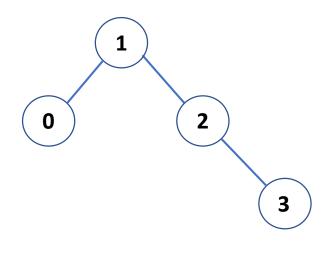
height =  $\Theta(n)$ 

Suppose we start with a "balanced" BST

Insert keys 3, 4, 5, ..., n

**Search** takes  $\Theta(n)$ !

Intuitively, we want BST trees with height  $\Theta(\log n)$ 



height =  $\Theta(n)$ 

n

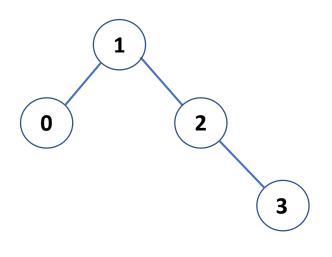
Suppose we start with a "balanced" BST

Insert keys 3, 4, 5, ..., n

**Search** takes  $\Theta(n)$ !

Intuitively, we want BST trees with height  $\Theta(\log n)$ 

We achieve this using balanced BSTs



height =  $\Theta(n)$ 

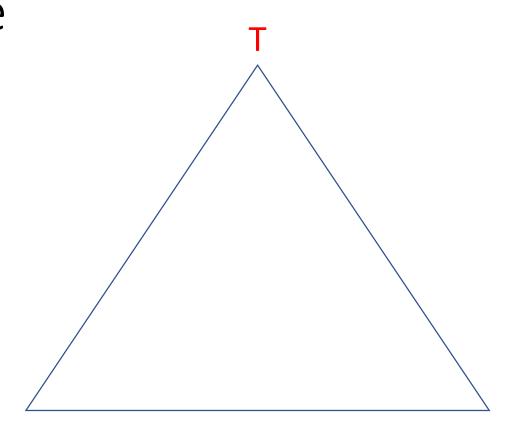
n

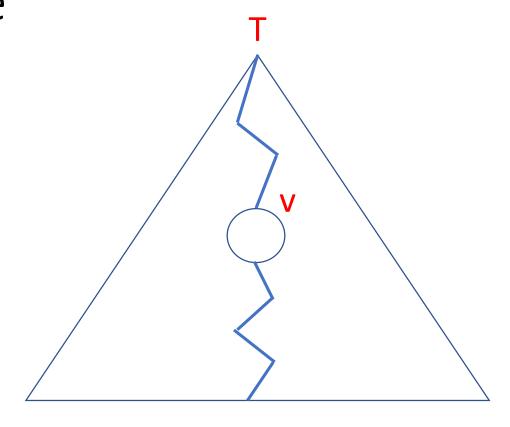
Abstract Data Type	Operations	Data Structures
Dictionary	Search, Insert, Delete	BSTs Balanced BSTs

<b>Abstract Data Type</b>	Operations	Data Structures
Dictionary	Search, Insert, Delete	BSTs Balanced BSTs: - 2-3 trees - Red-Black Trees - B-Trees - AVL Trees

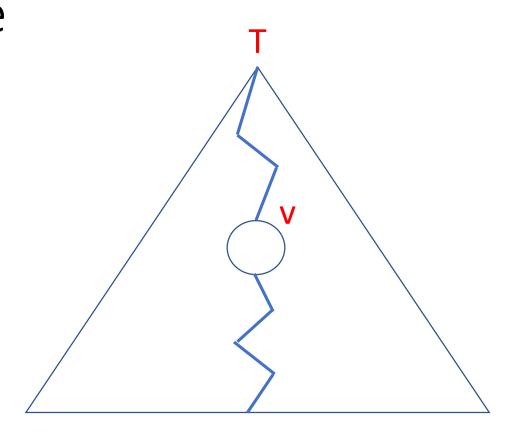
<b>Abstract Data Type</b>	Operations	Data Structures
Dictionary	Search, Insert, Delete	BSTs Balanced BSTs: - 2-3 trees - Red-Black Trees - B-Trees - AVL Trees

## **AVL Trees**

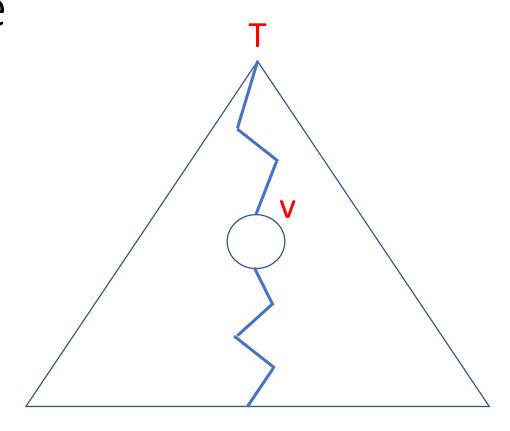




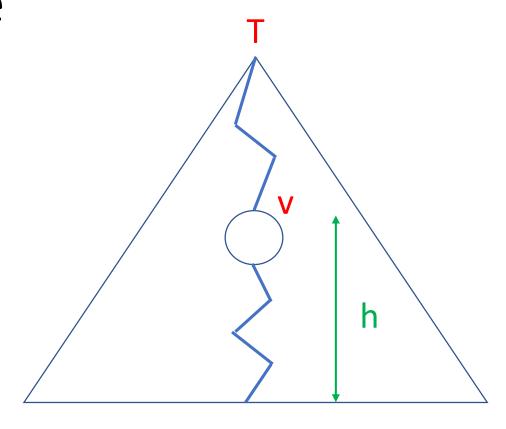
• height(v):



 height(v): Number of edges in the longest path from v to a leaf

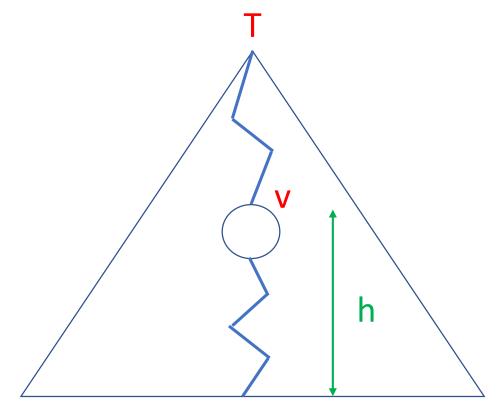


 height(v): Number of edges in the longest path from v to a leaf



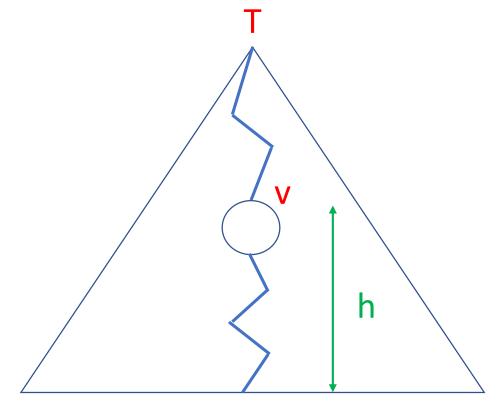
 height(v): Number of edges in the longest path from v to a leaf

height(T): height of the root node of T



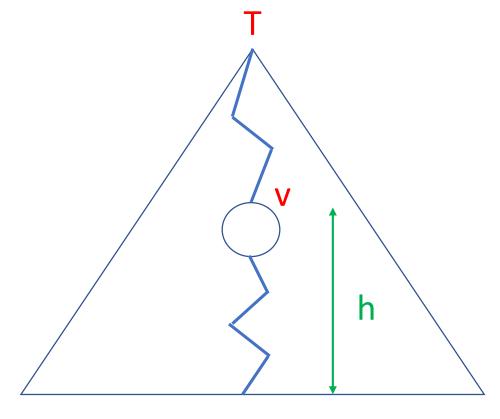
 height(v): Number of edges in the longest path from v to a leaf

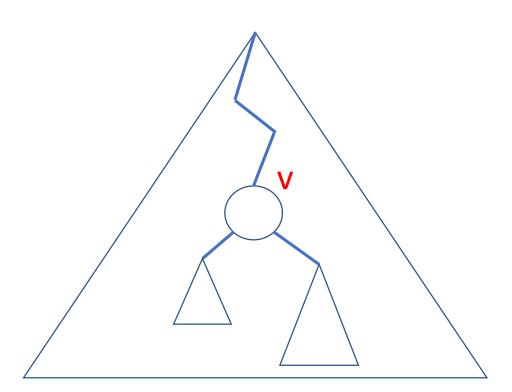
height(T): height of the root node of T



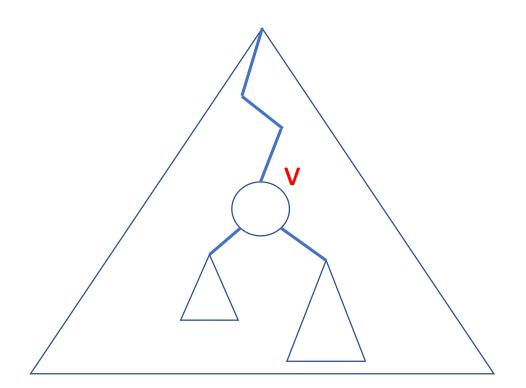
 height(v): Number of edges in the longest path from v to a leaf

height(T): height of the root node of T

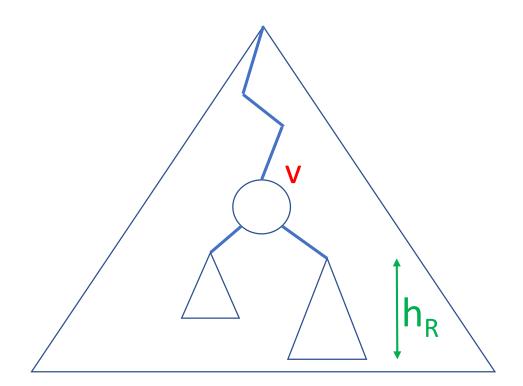




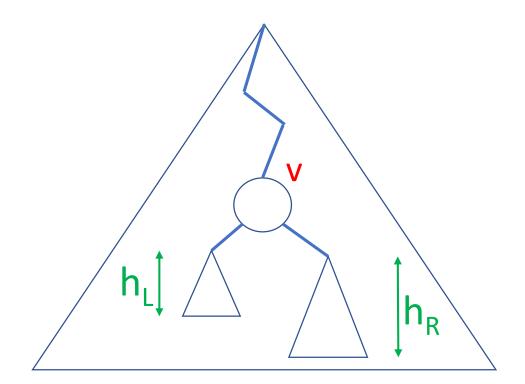
$$BF(v) =$$



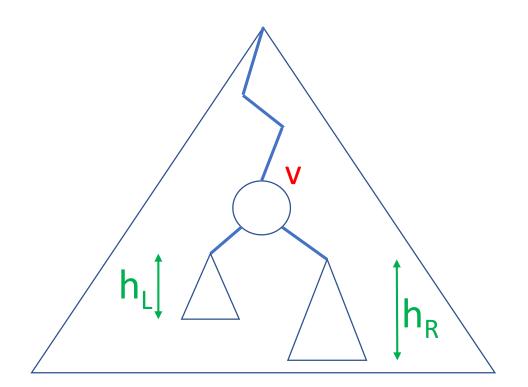
BF(v) = height(right subtree of v)



BF(v) = height(right subtree of v) - height(left subtree of v)

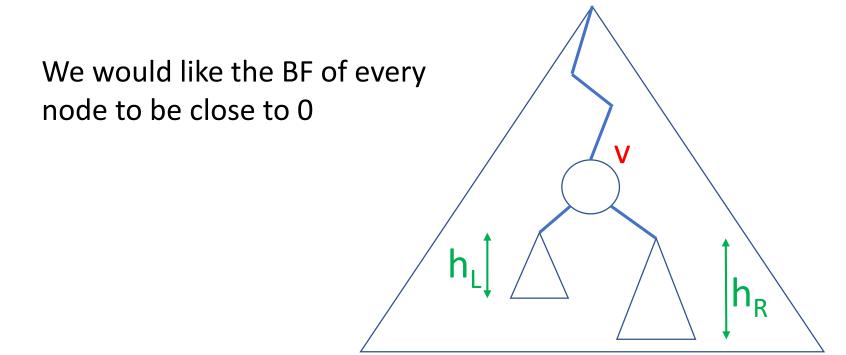


BF(v) = height(right subtree of v) - height(left subtree of v) $BF(v) = h_R - h_I$ 



BF(v) = height(right subtree of v) - height(left subtree of v)

$$BF(v) = h_R - h_L$$



$$-1 \leq BF(v) \leq +1$$

$$-1 \leq BF(v) \leq +1$$

$$BF(\mathbf{v}) = \begin{cases} +1 \\ 0 \\ -1 \end{cases}$$

$$-1 \leq BF(v) \leq +1$$

$$BF(v) = \begin{cases} +1 & (v \text{ is "right-heavy"}) \\ 0 & \\ -1 & \end{cases}$$

$$-1 \leq BF(v) \leq +1$$

```
BF(v) = \begin{cases} +1 & (v \text{ is "right-heavy"}) \\ 0 & (v \text{ is "balanced"}) \end{cases}
```

# Adelson-Velski-Landis Trees

An AVL tree T is a BST where for every node  $v \in T$ :

$$-1 \leq BF(v) \leq +1$$

```
H1 (v is "right-heavy")

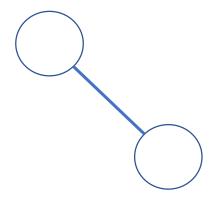
BF(v) = 0 (v is "balanced")

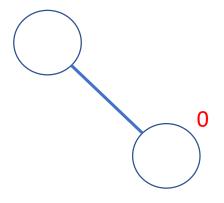
-1 (v is "left-heavy")
```

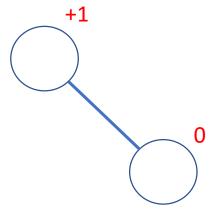


# Is this an AVL tree? Yes!

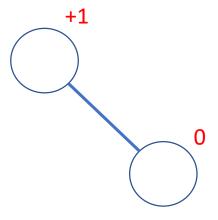


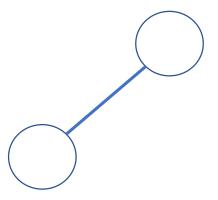


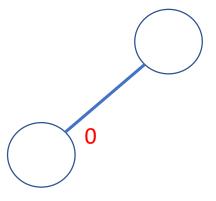


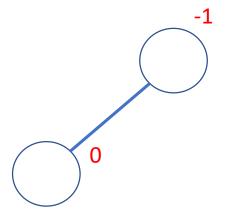


### Is this an AVL tree? Yes!

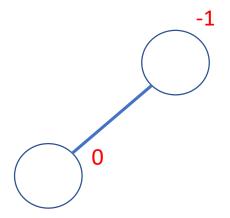


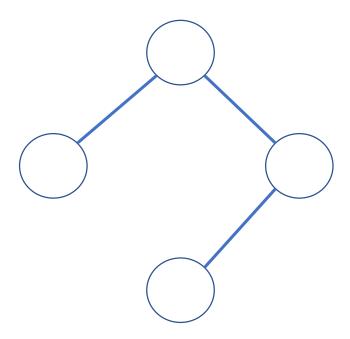


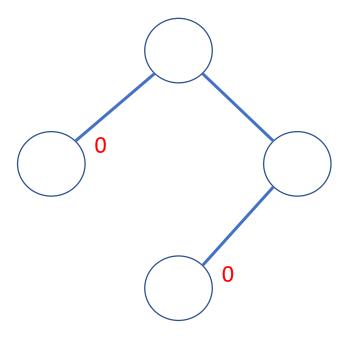


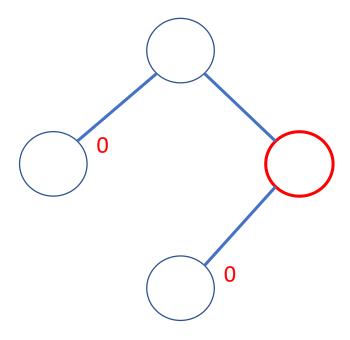


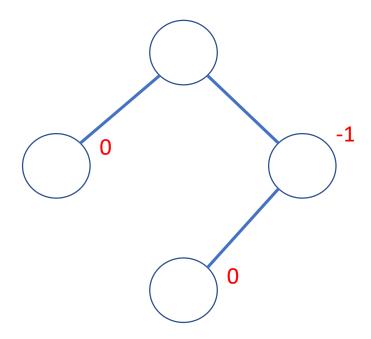
### Is this an AVL tree? Yes!

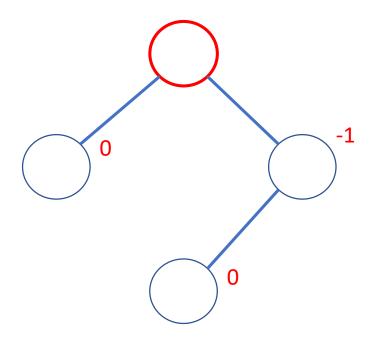


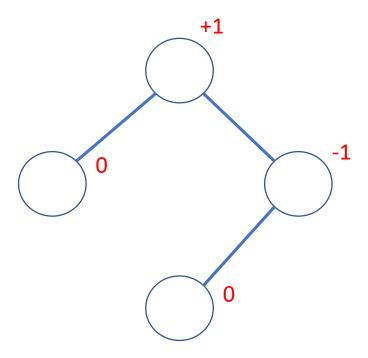




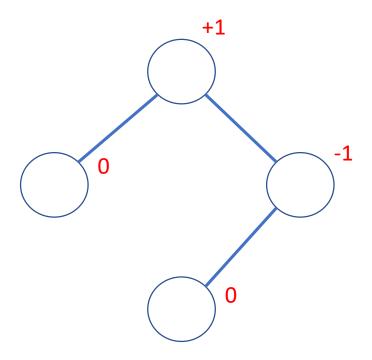


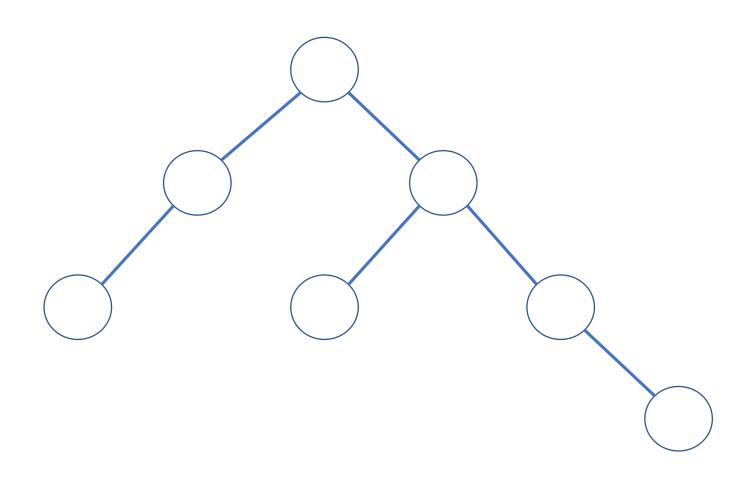


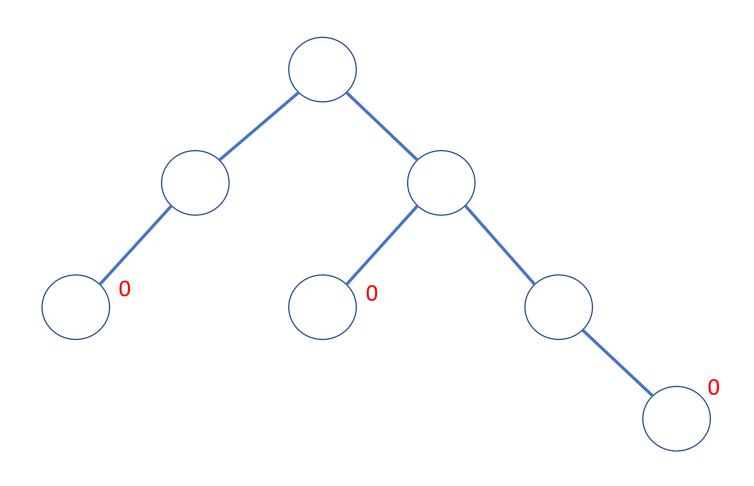


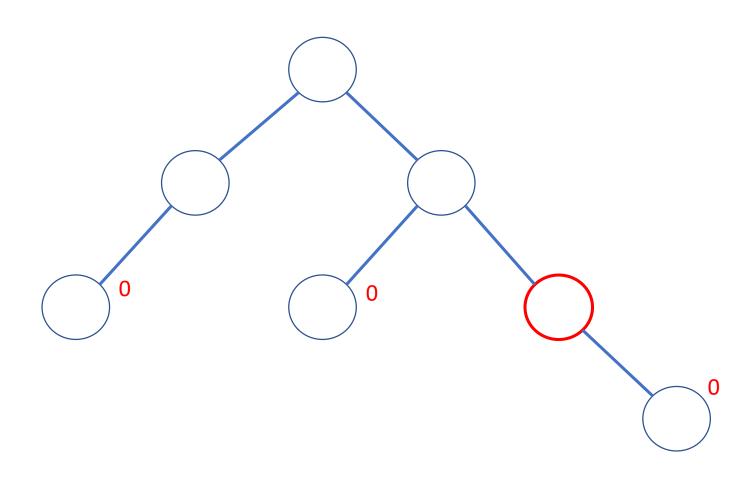


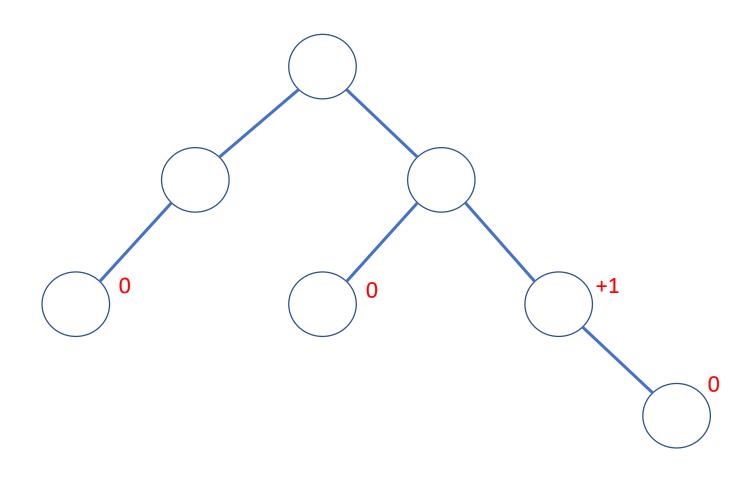
# Is this an AVL tree? Yes!

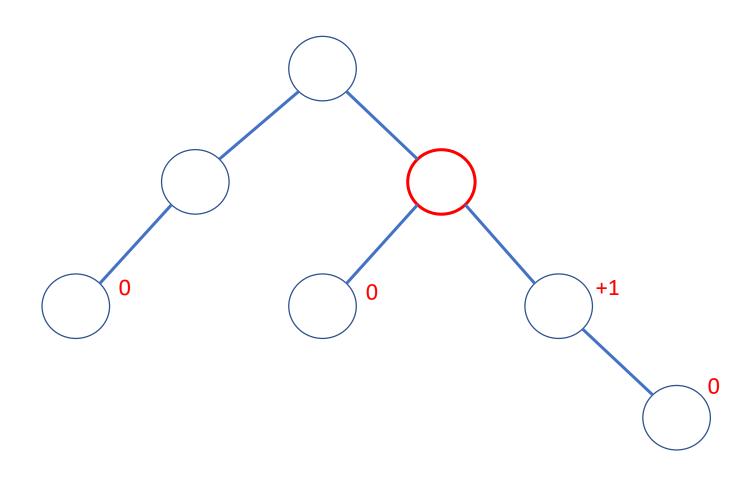


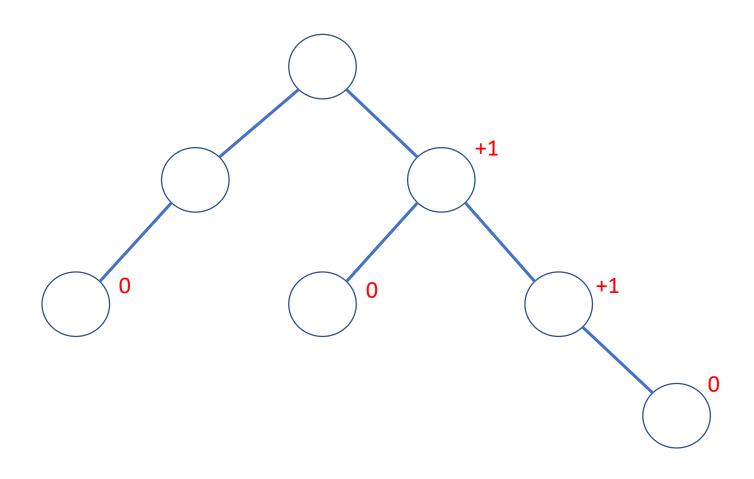


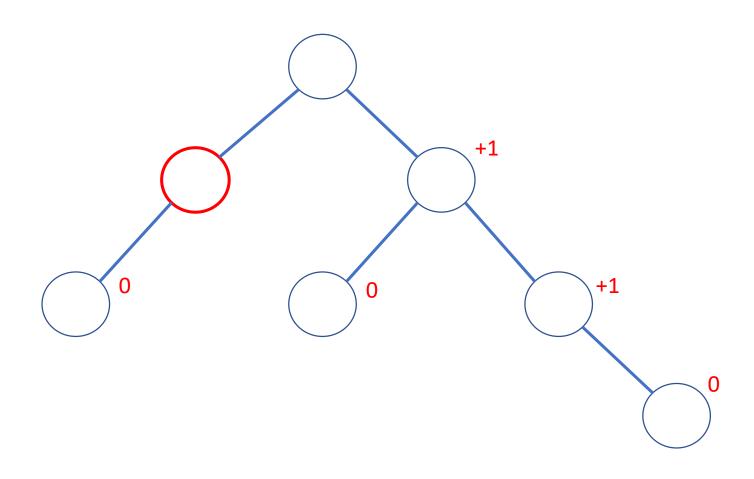


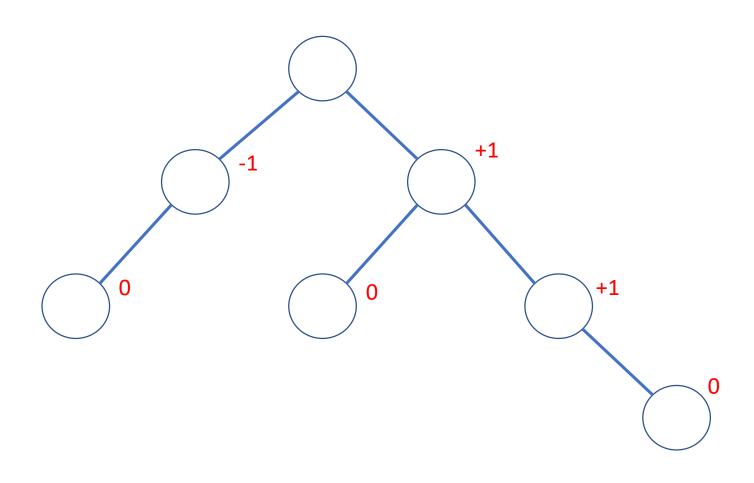


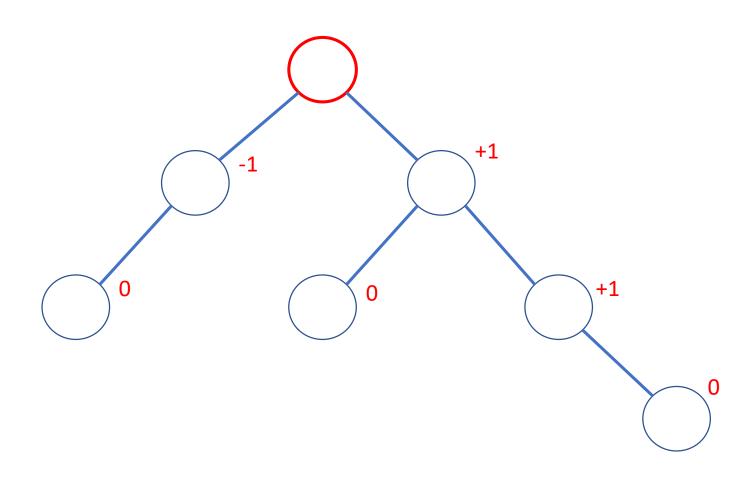


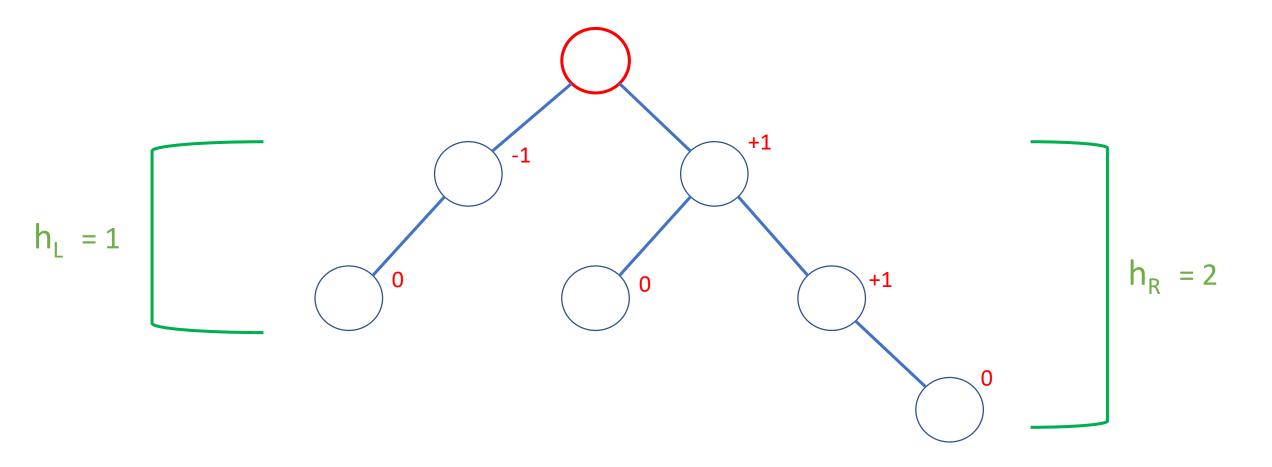


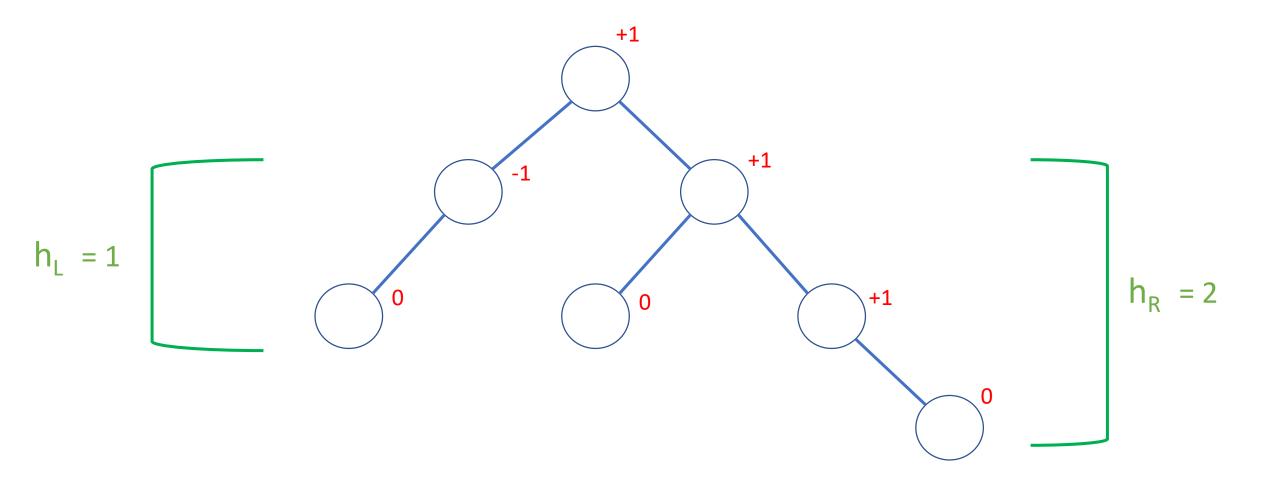




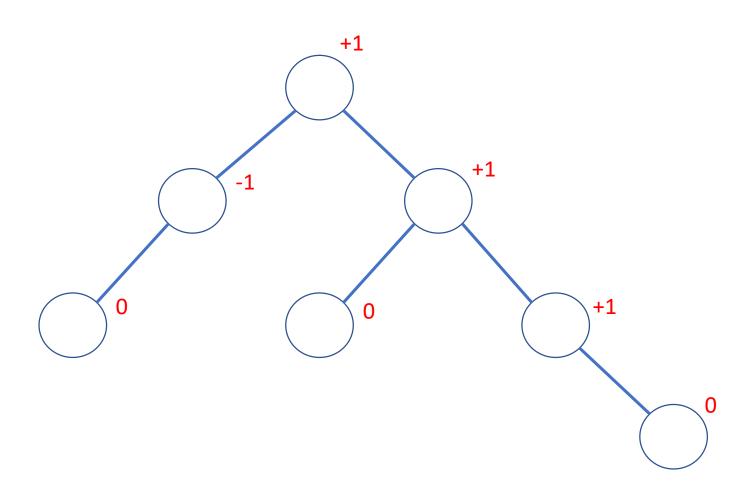


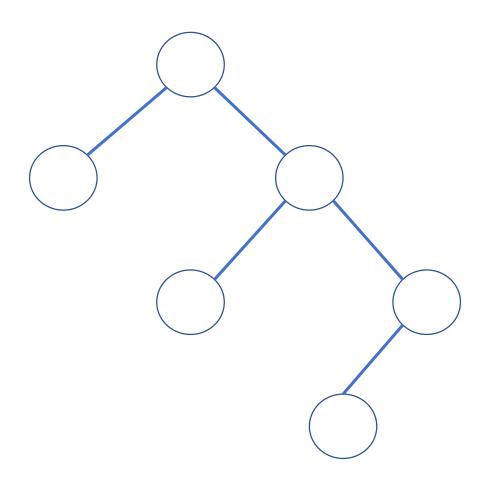


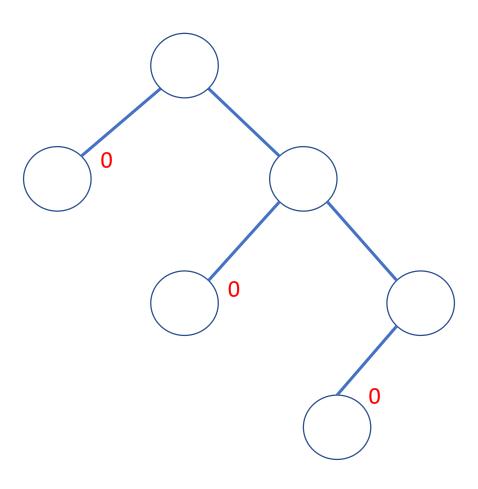


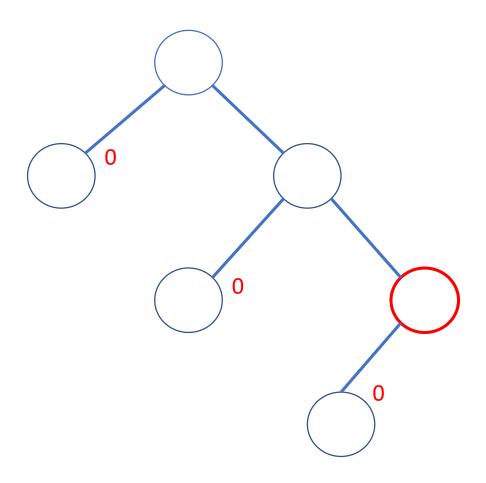


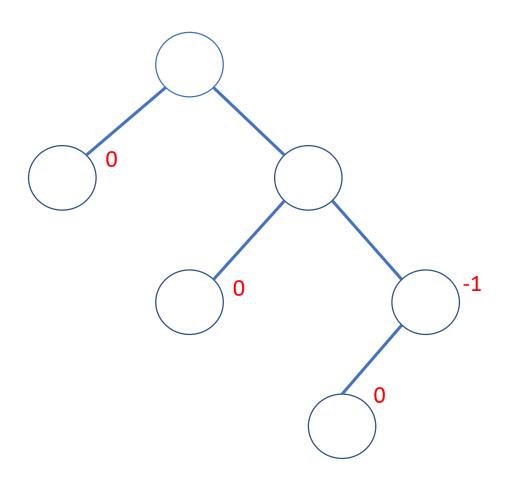
### Is this an AVL tree? Yes!

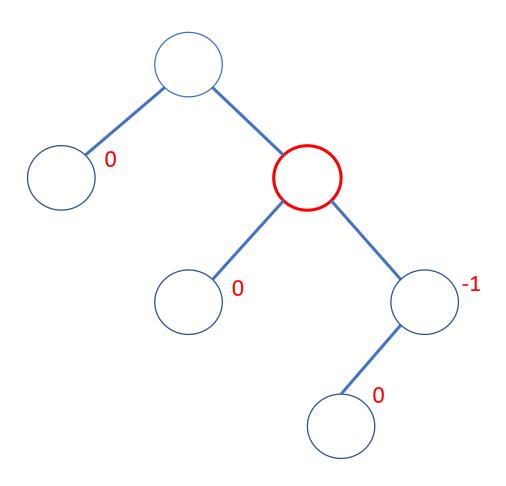


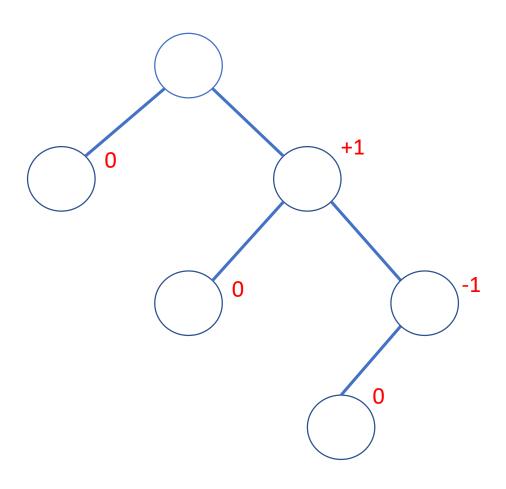




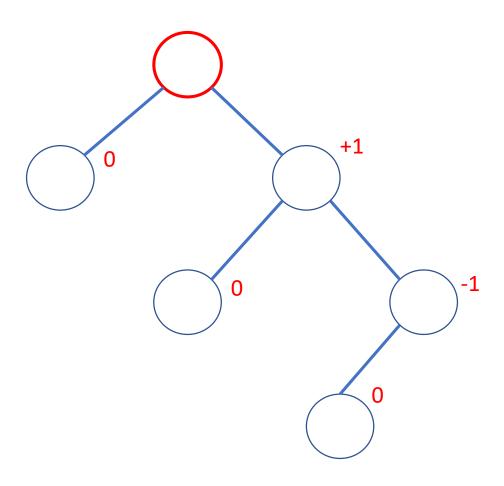




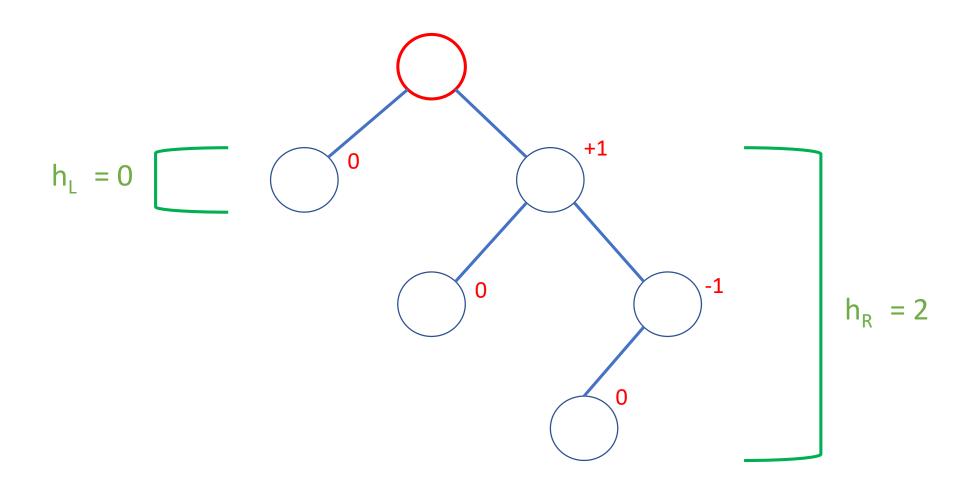




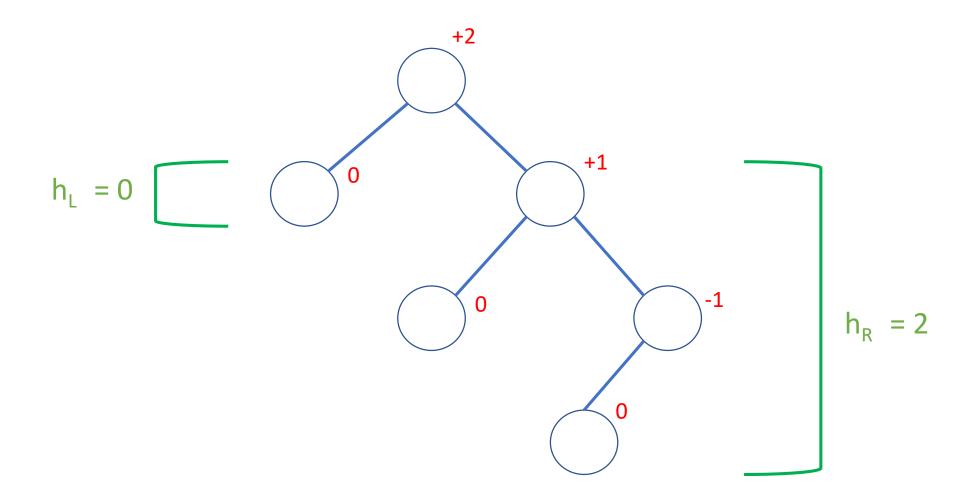
#### Is this an AVL tree?



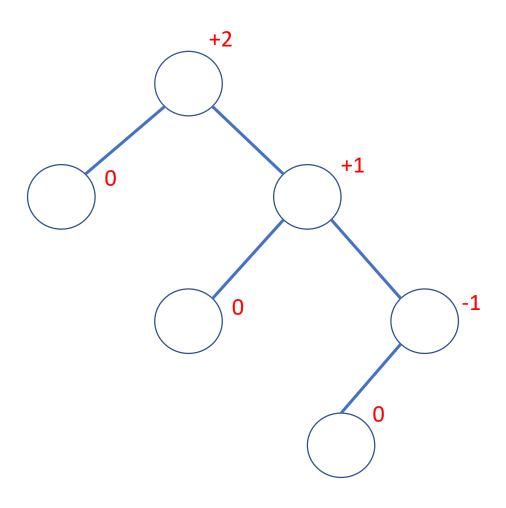
#### Is this an AVL tree?



#### Is this an AVL tree?



#### Is this an AVL tree? No!



• AVL trees of n nodes has height  $\Theta(\log n)$  [height  $\leq 1.44 \log_2(n+2)$ ]

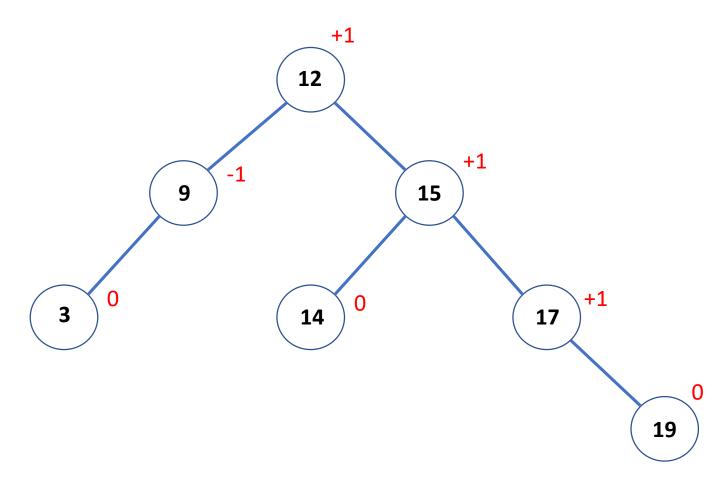
- AVL trees of n nodes has height  $\Theta(\log n)$  [height  $\leq 1.44 \log_2(n+2)$ ]
- Can do inserts, deletes while maintaining the tree balance in  $\Theta(\log n)$  time

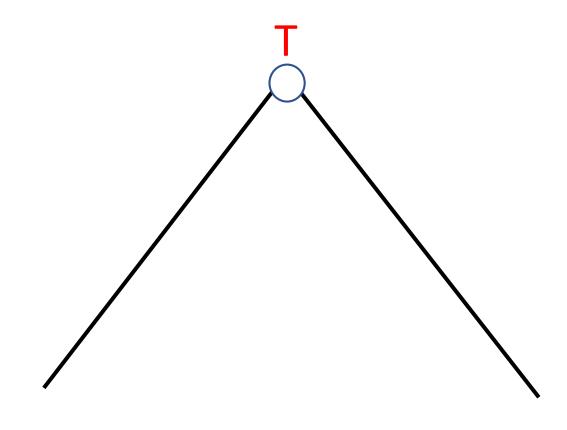
- AVL trees of n nodes has height  $\Theta(\log n)$  [height  $\leq 1.44 \log_2(n+2)$ ]
- Can do inserts, deletes while maintaining the tree balance in  $\Theta(\log n)$  time

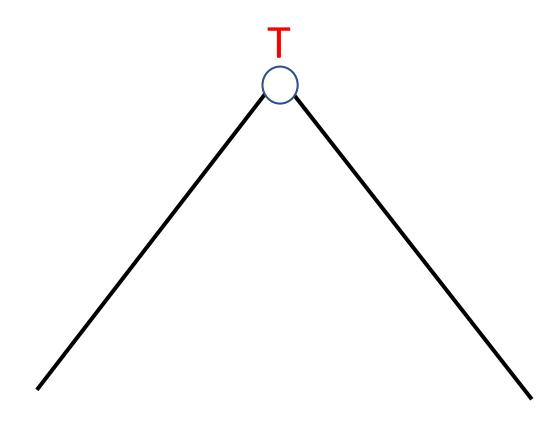
Elegant, relatively clean and simple, works well in practice

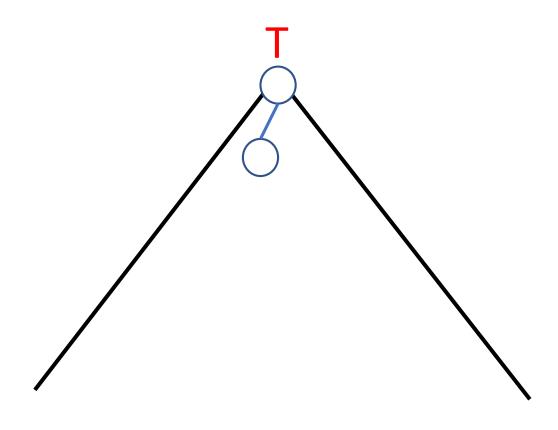
#### AVL tree operations

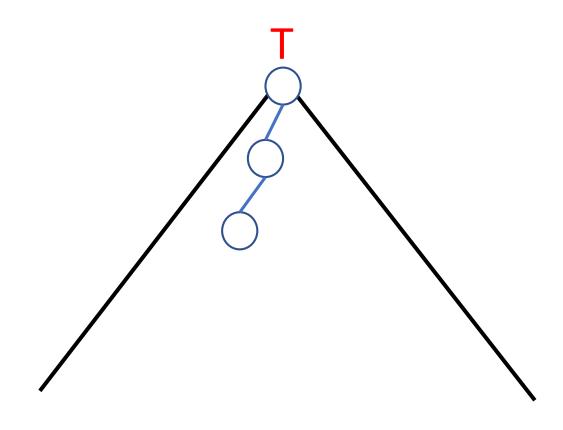
- Search(T, x)
- Insert(T, x)
- Delete(T, x)

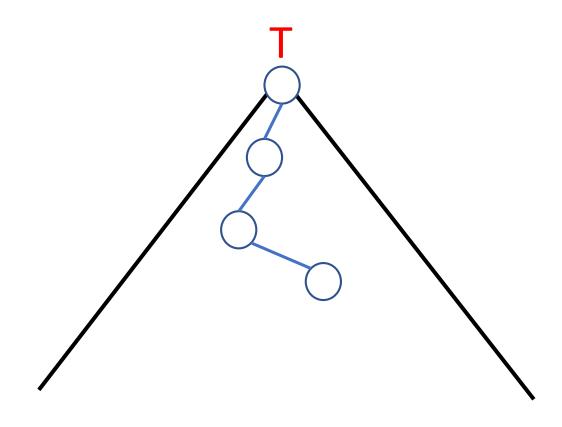


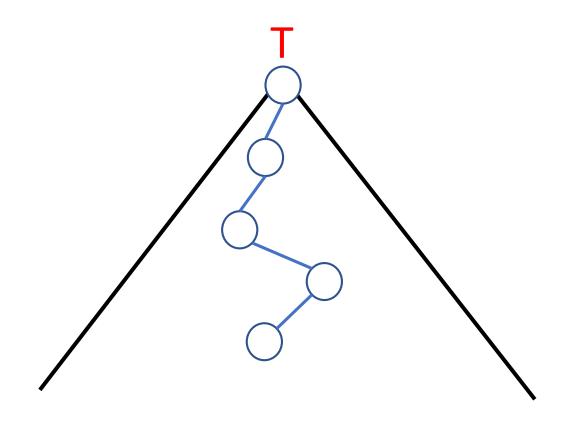


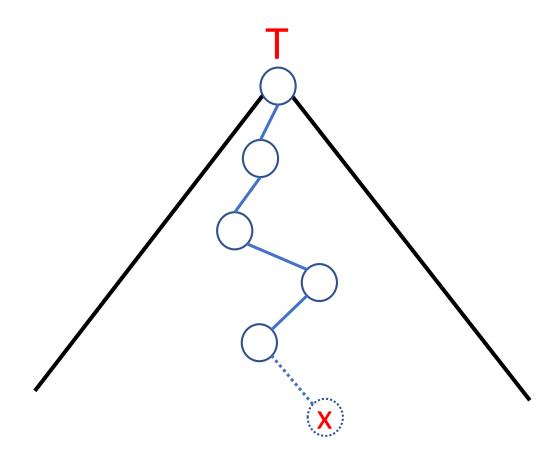




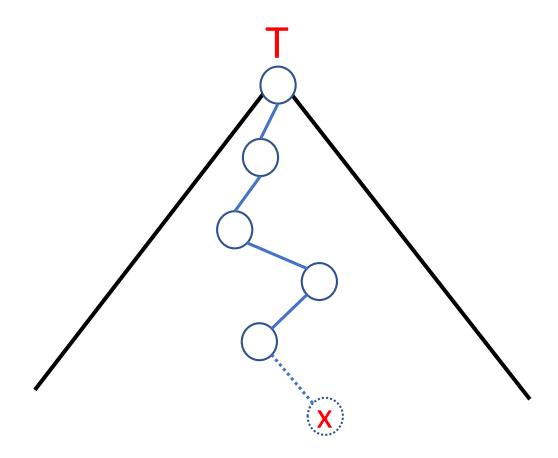




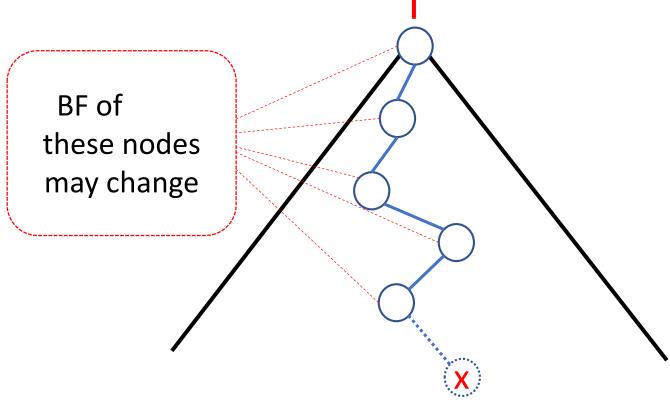




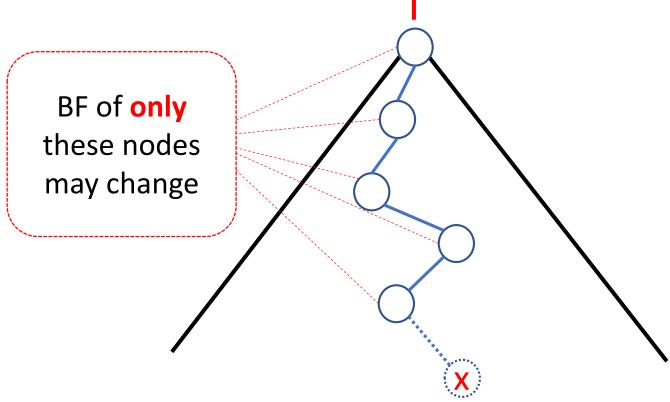
- Insert x into T as in any BST :
  - x is now a leaf



- Insert x into T as in any BST :
  - x is now a leaf

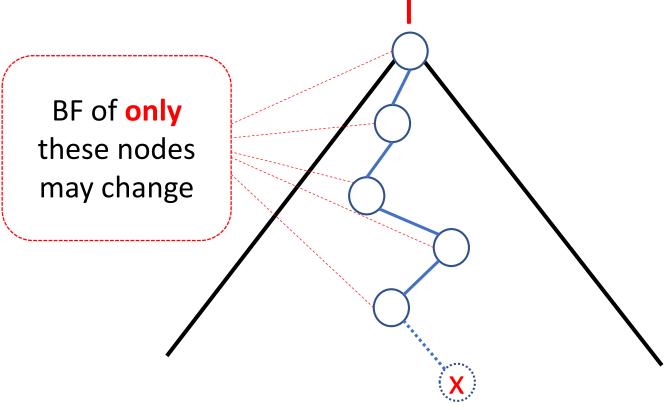


- Insert x into T as in any BST :
  - x is now a leaf



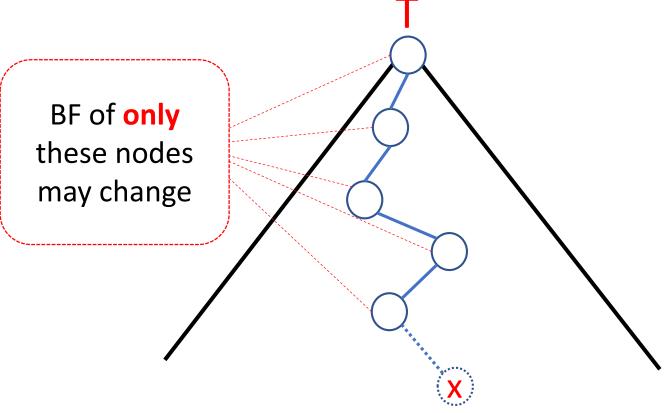
- Insert x into T as in any BST :
  - x is now a leaf

• Go up from x to the root



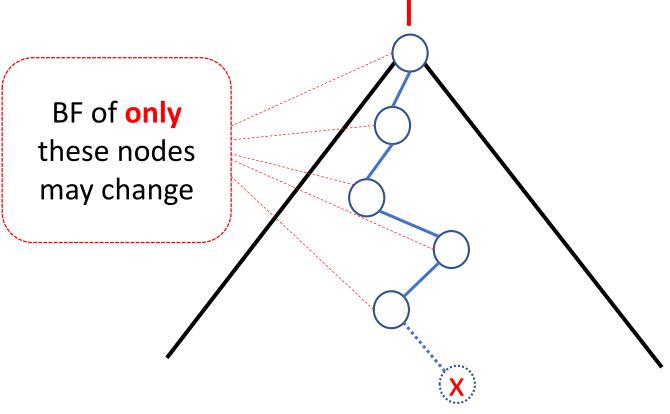
- Insert x into T as in any BST :
  - x is now a leaf

 Go up from x to the root and for each node :



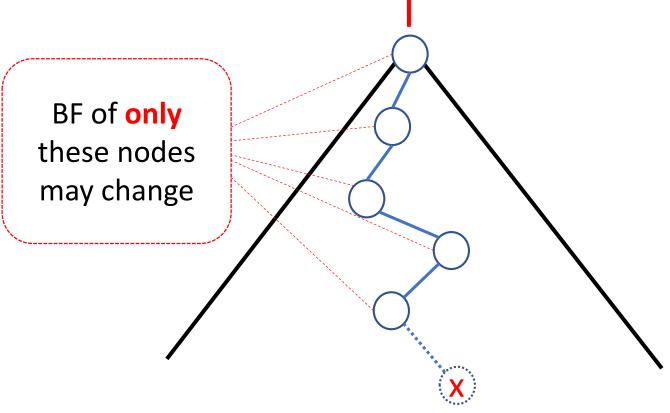
- Insert x into T as in any BST :
  - x is now a leaf

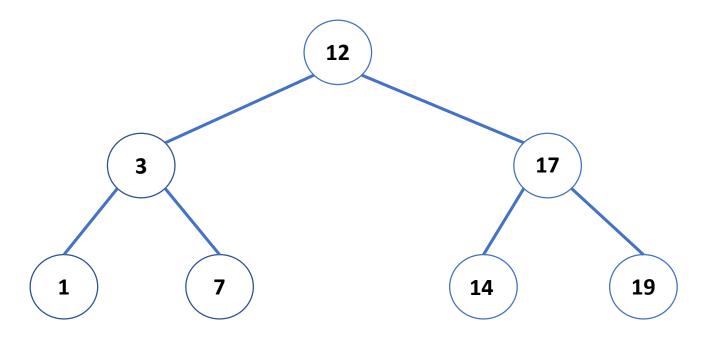
- Go up from x to the root and for each node :
  - Adjust the BF

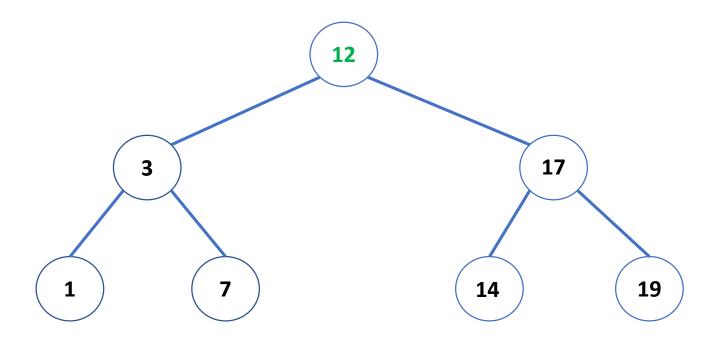


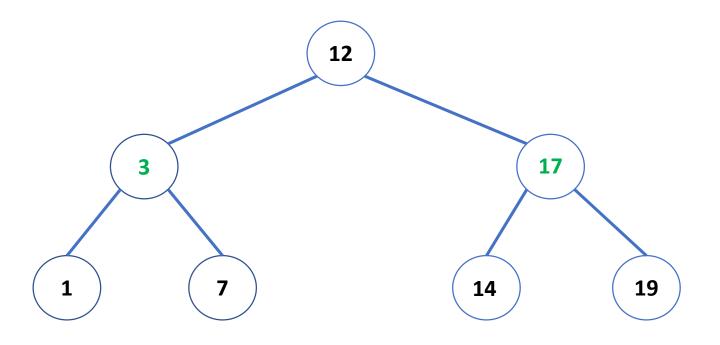
- Insert x into T as in any BST :
  - x is now a leaf

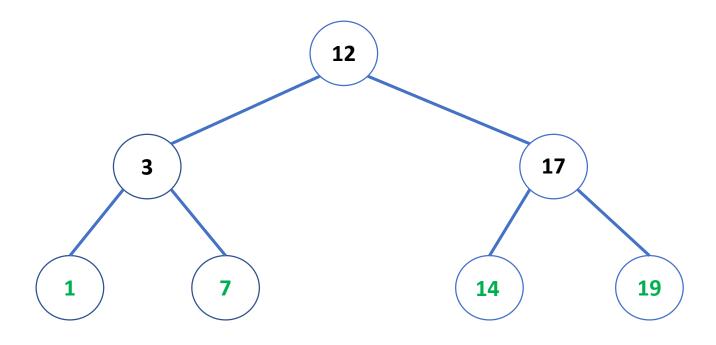
- Go up from x to the root and for each node :
  - Adjust the BF
  - "Rebalance" if necessary i.e. if BF > 1 or BF < -1

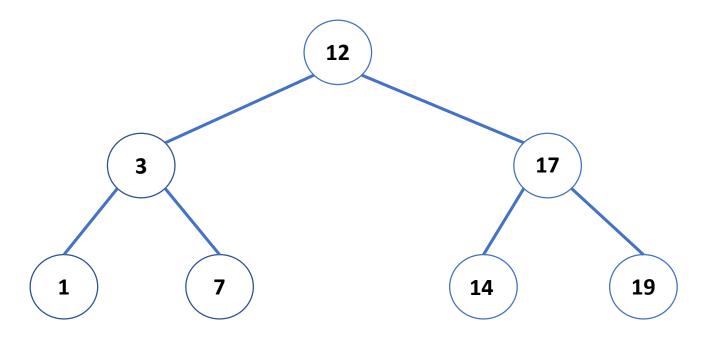




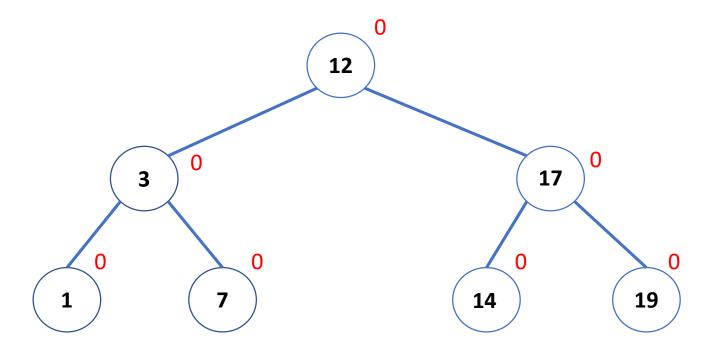


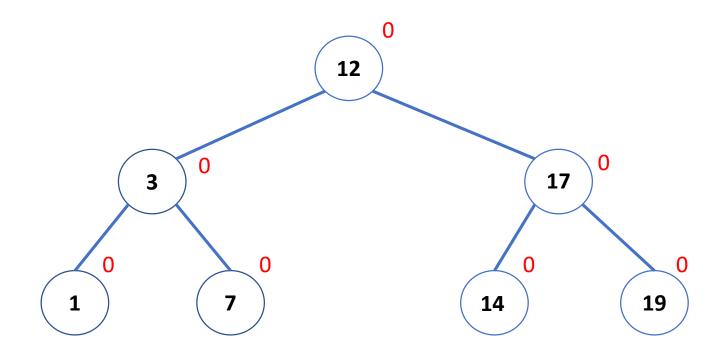


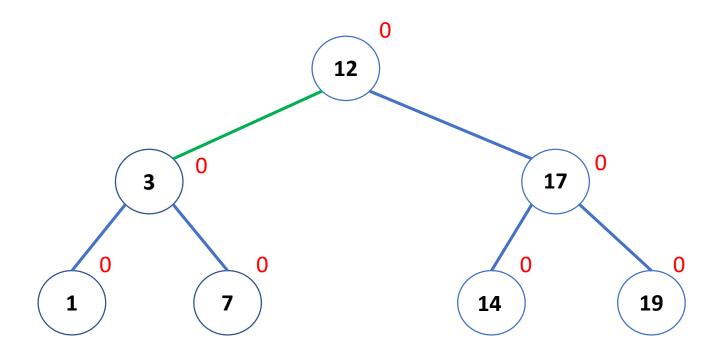


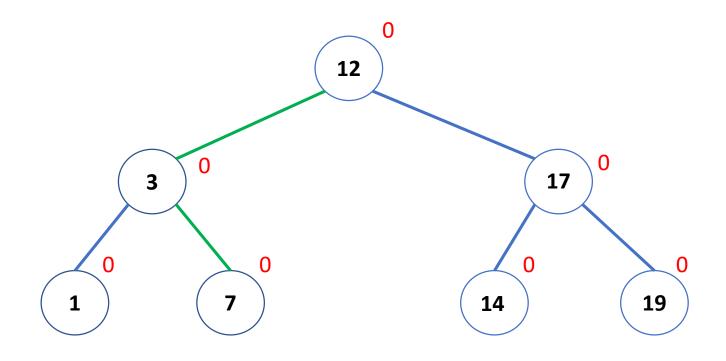


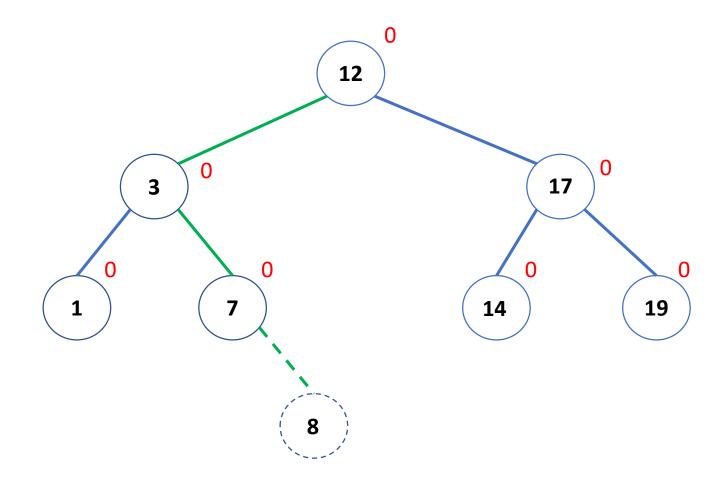
Example: AVL of {1, 3, 7, 12, 14, 17, 19}

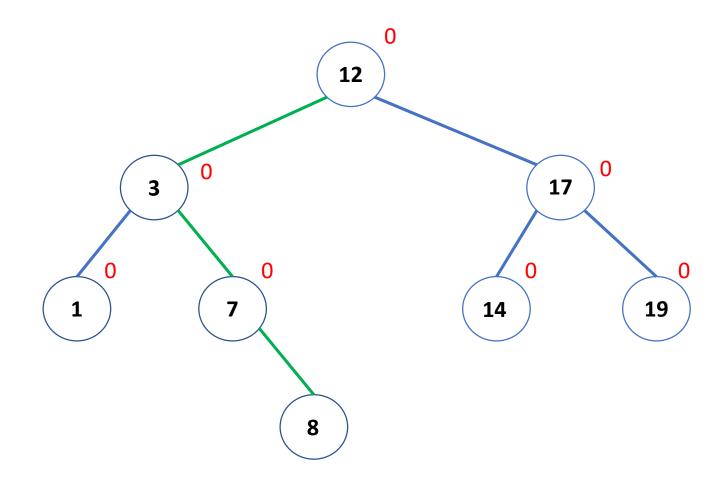


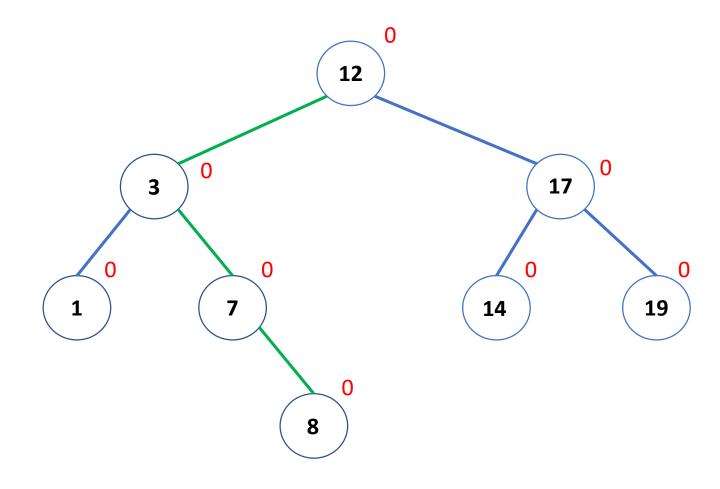


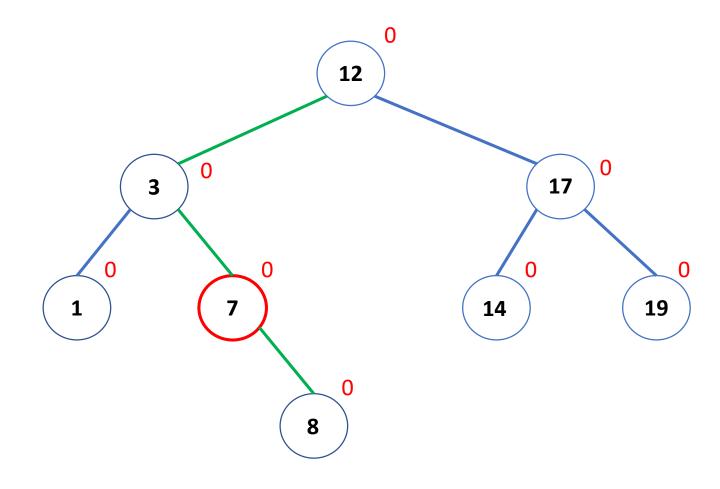


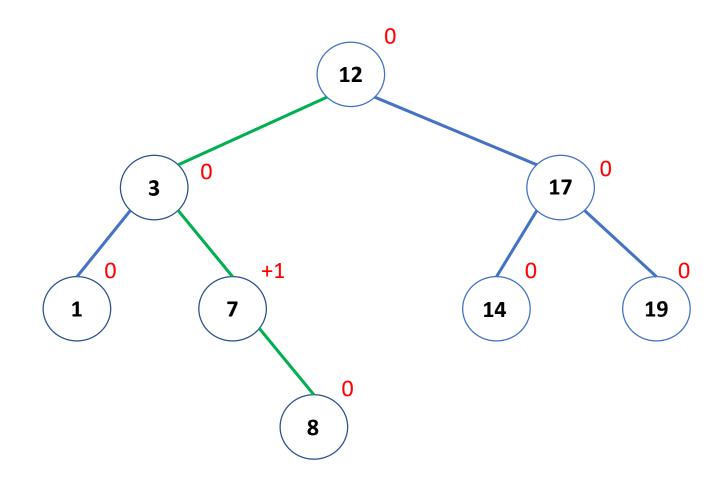


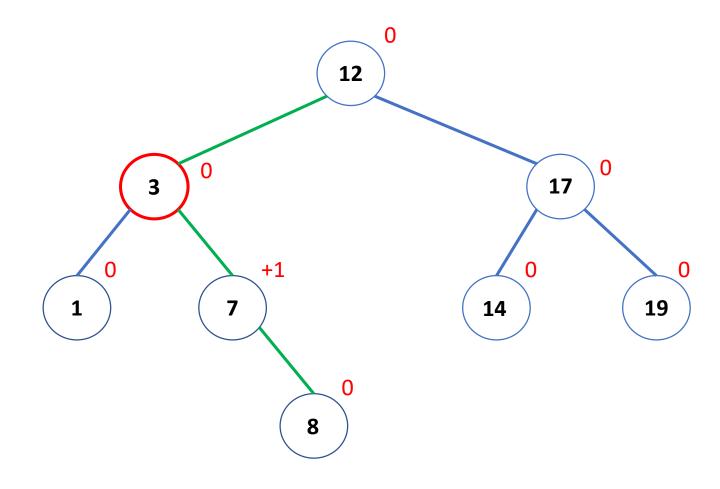


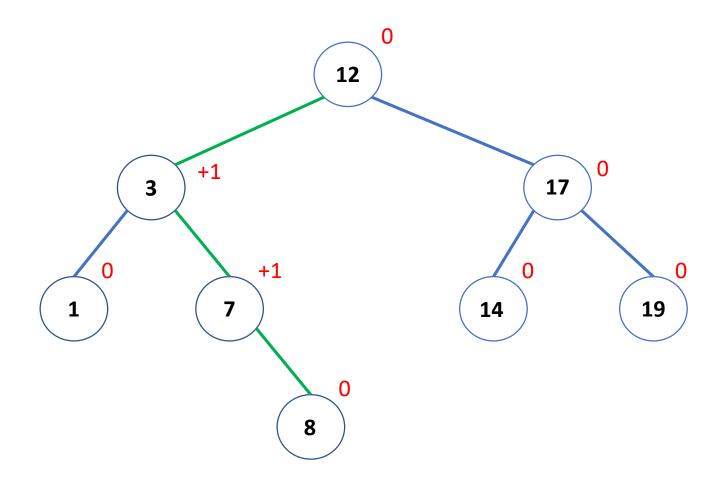


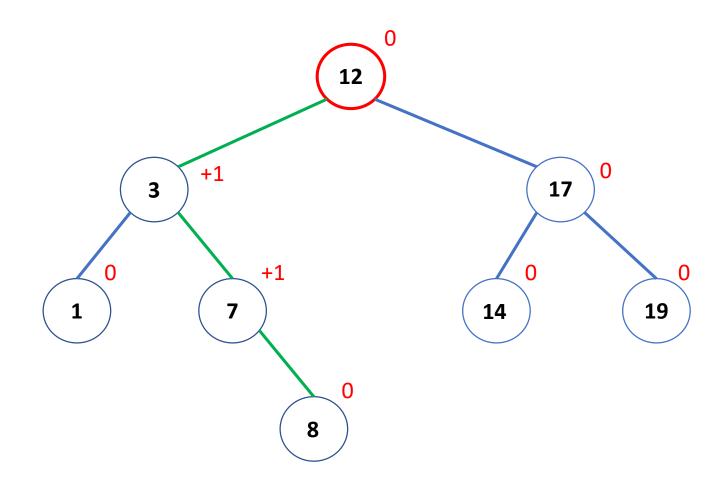


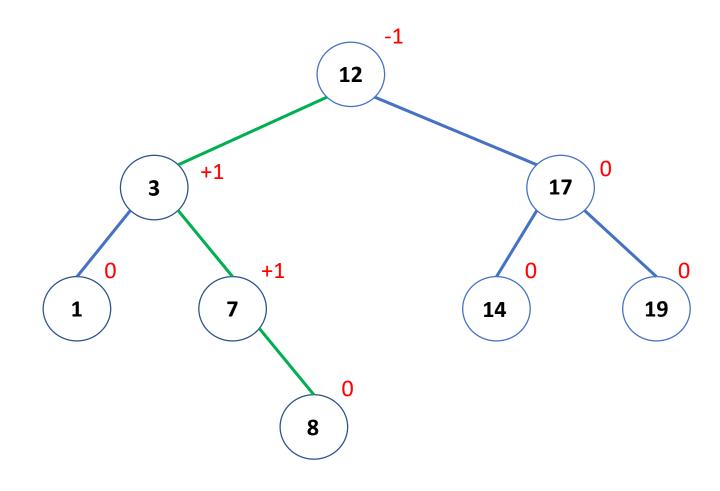


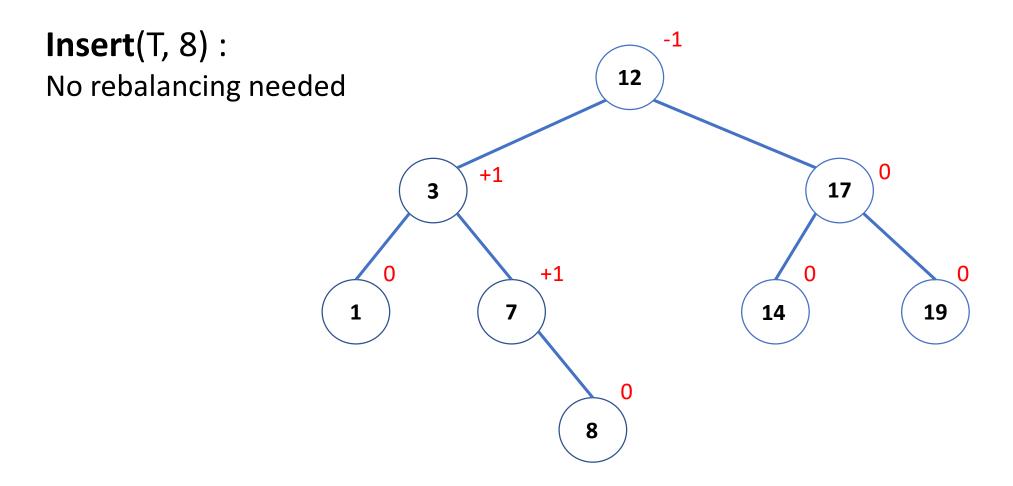


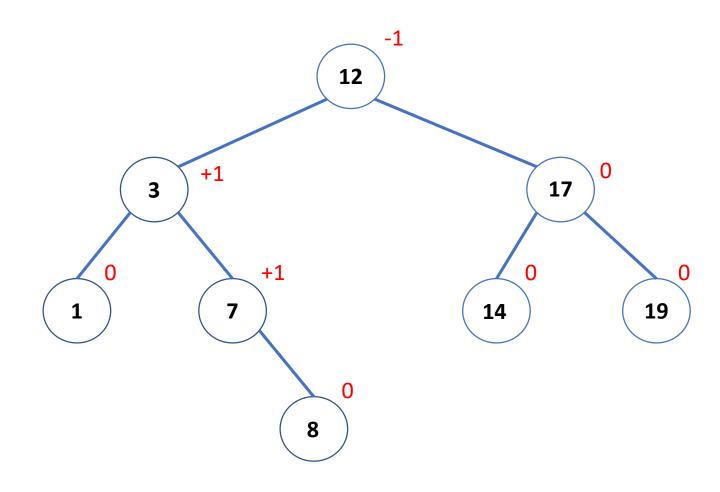


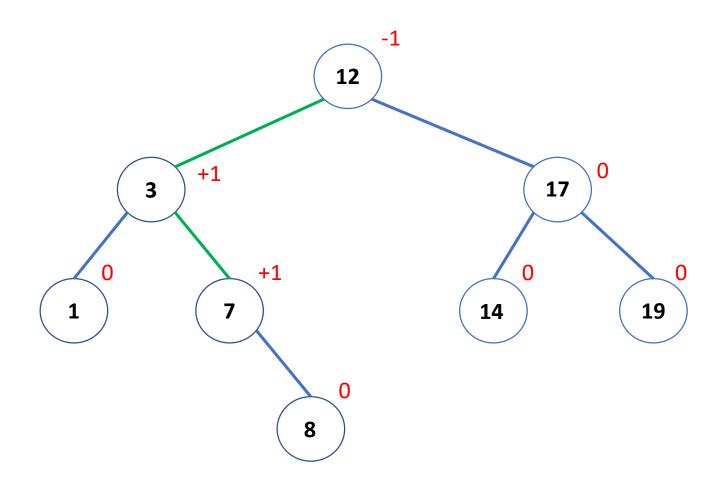


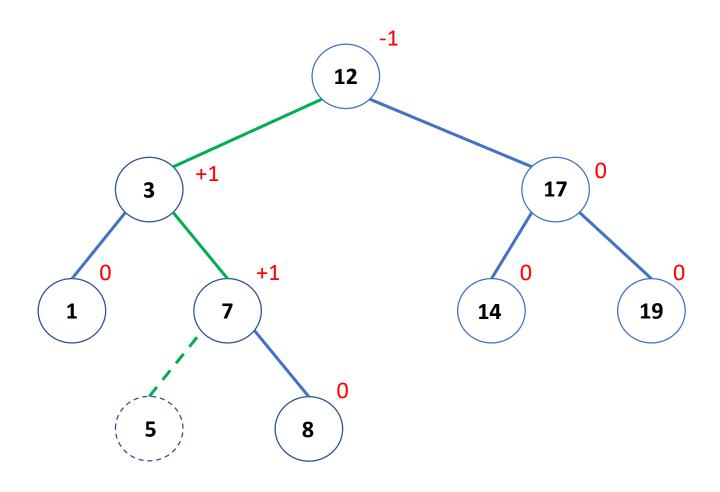


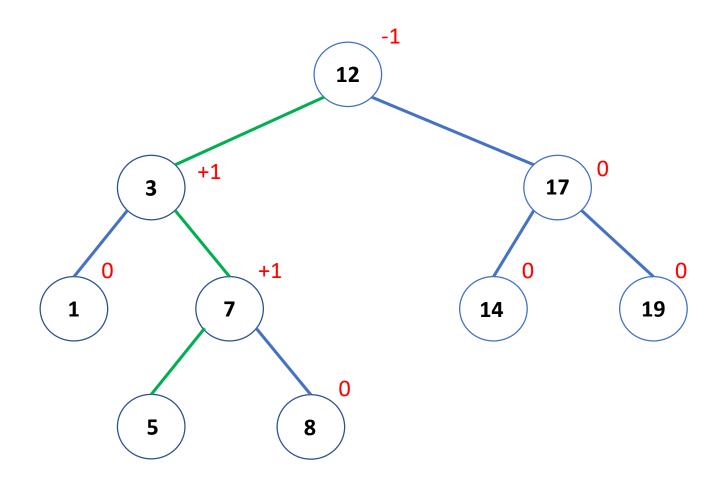


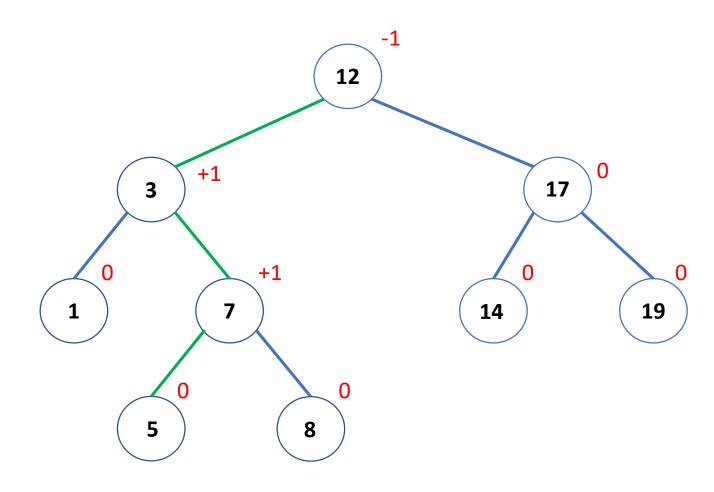


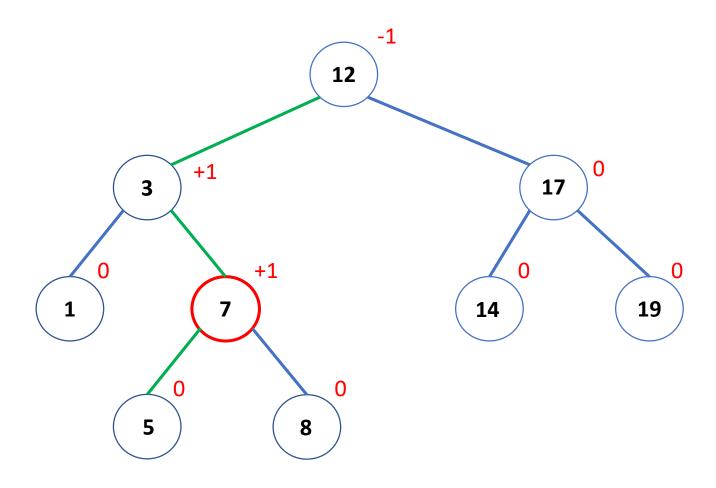


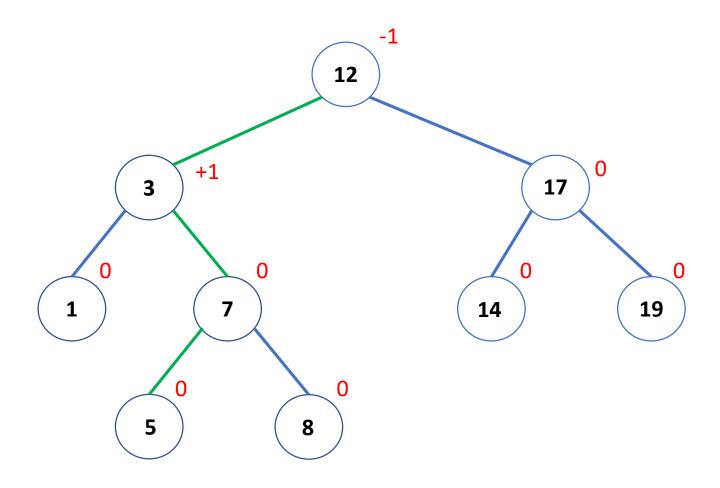


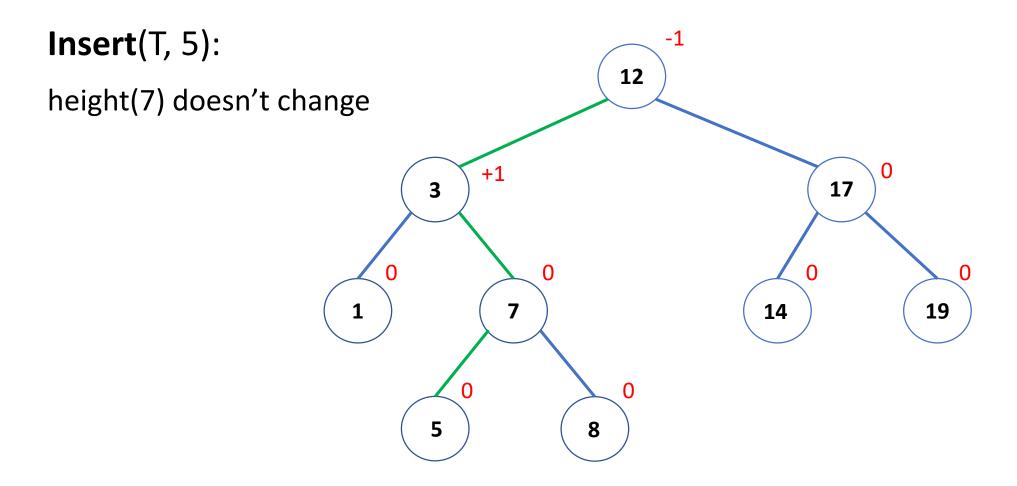


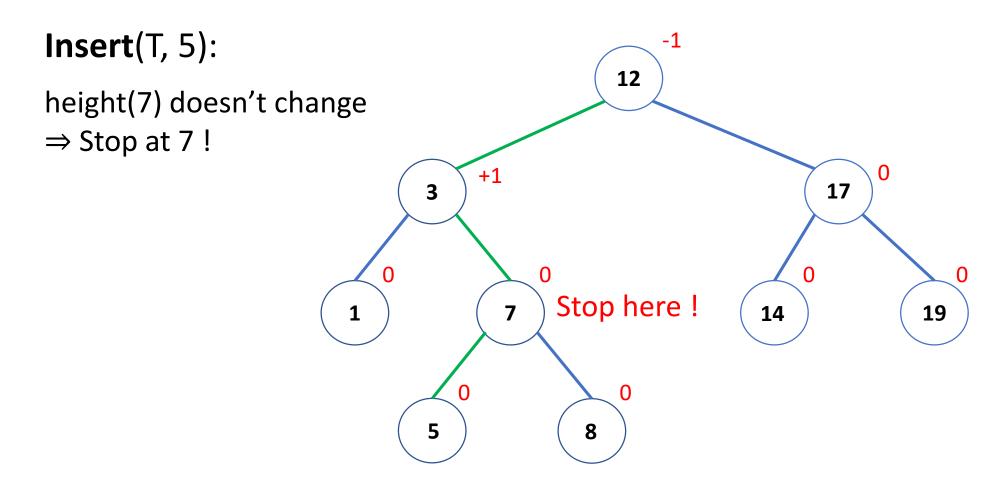


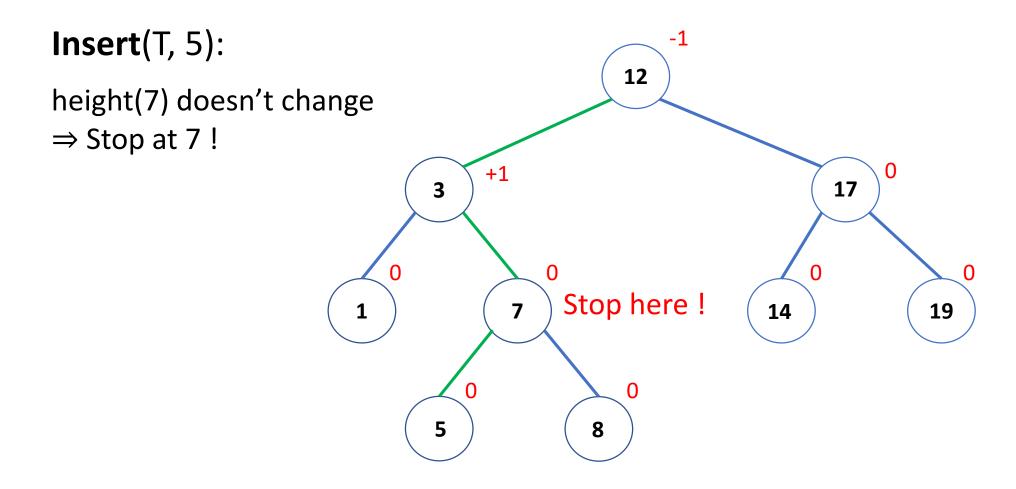




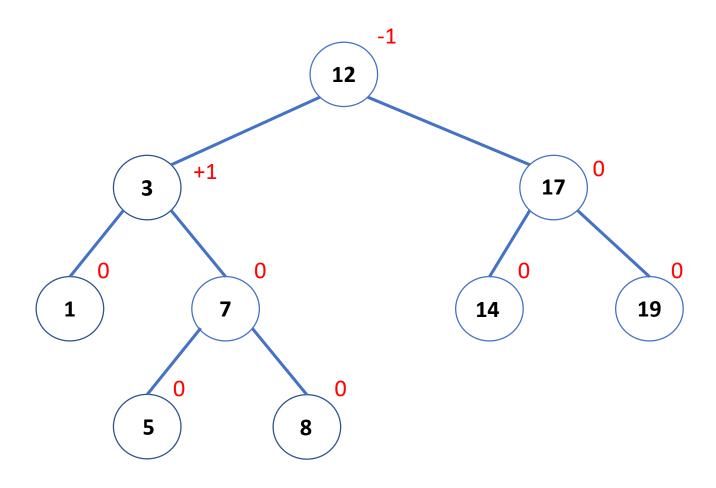


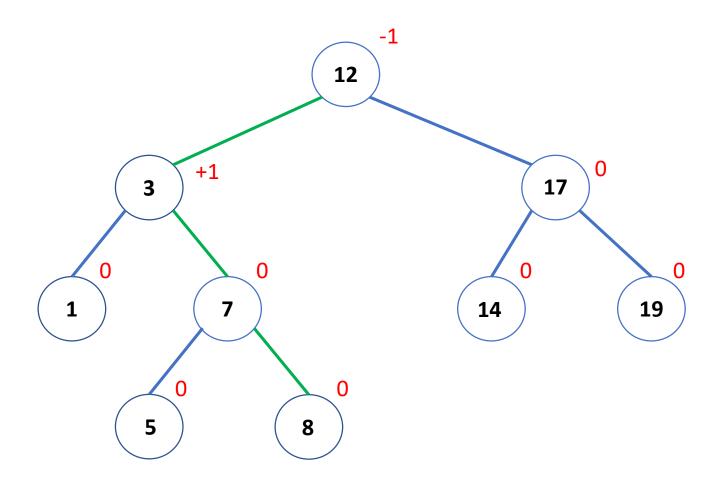


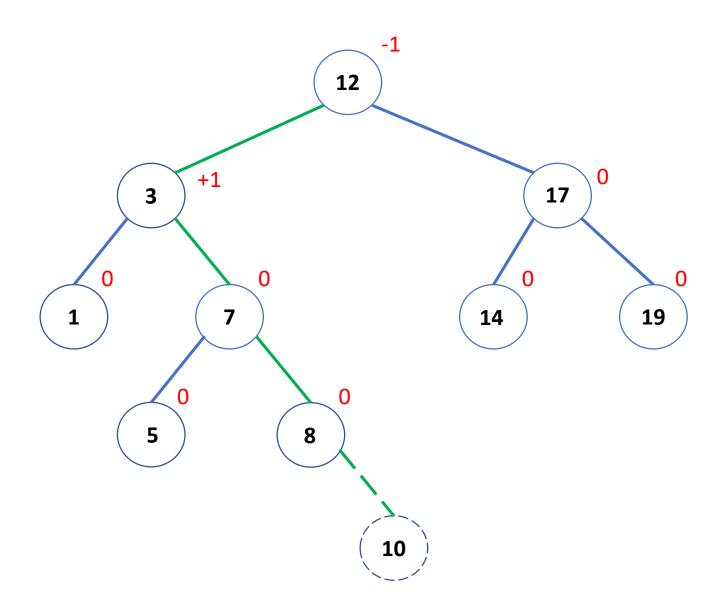


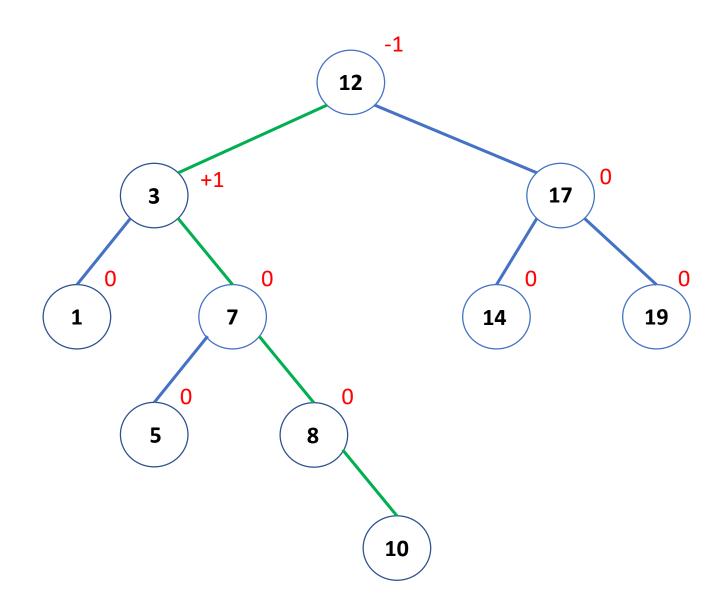


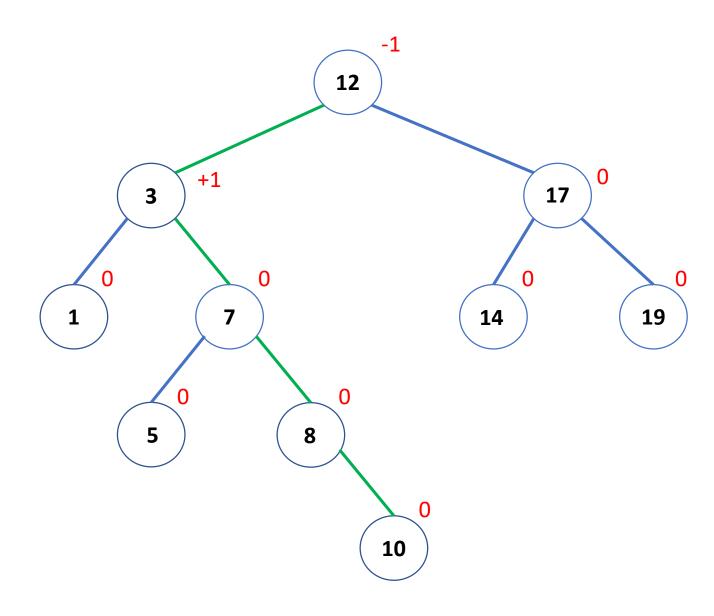
No rebalancing needed!

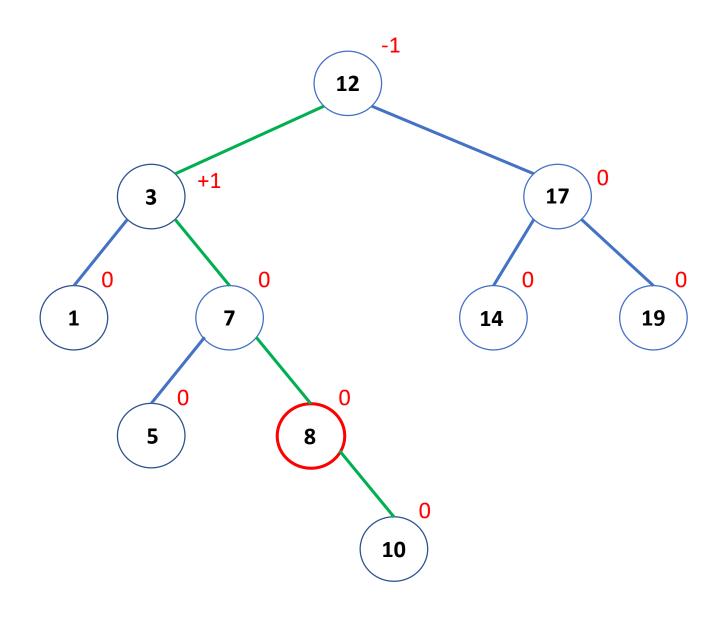


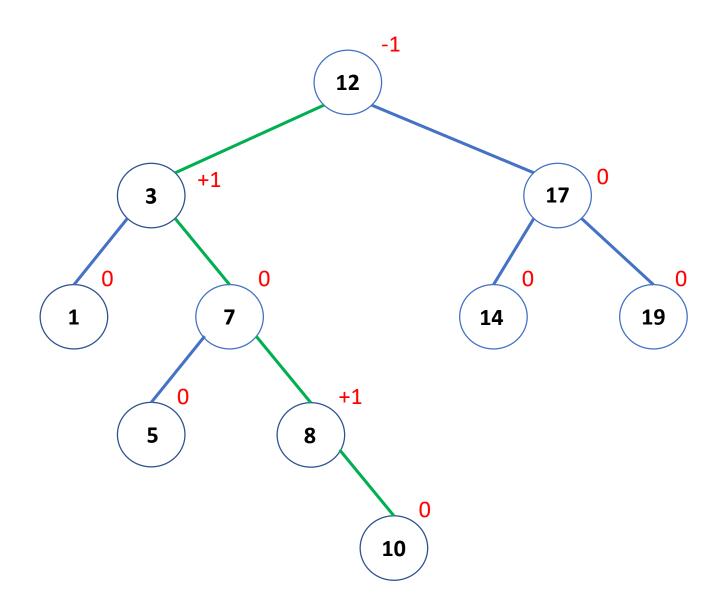


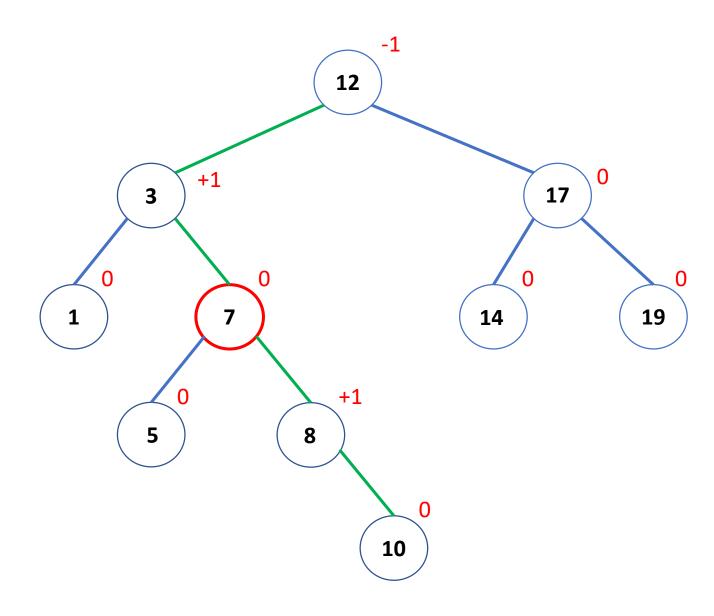


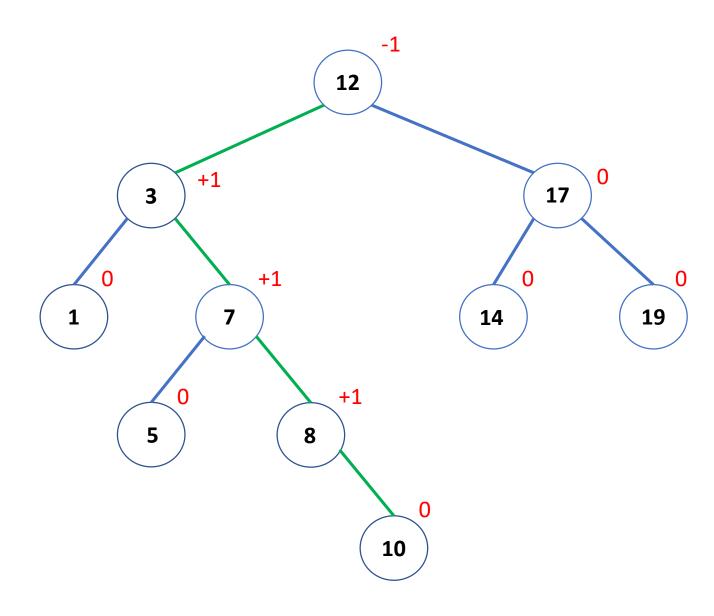


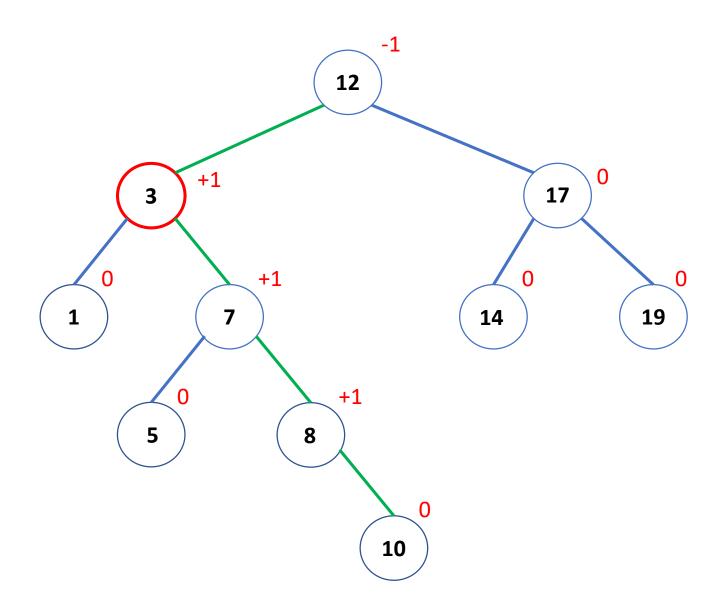


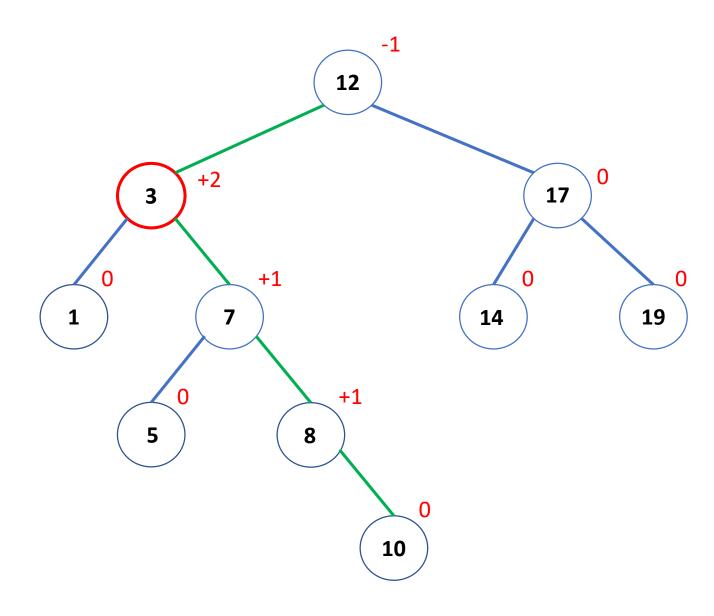


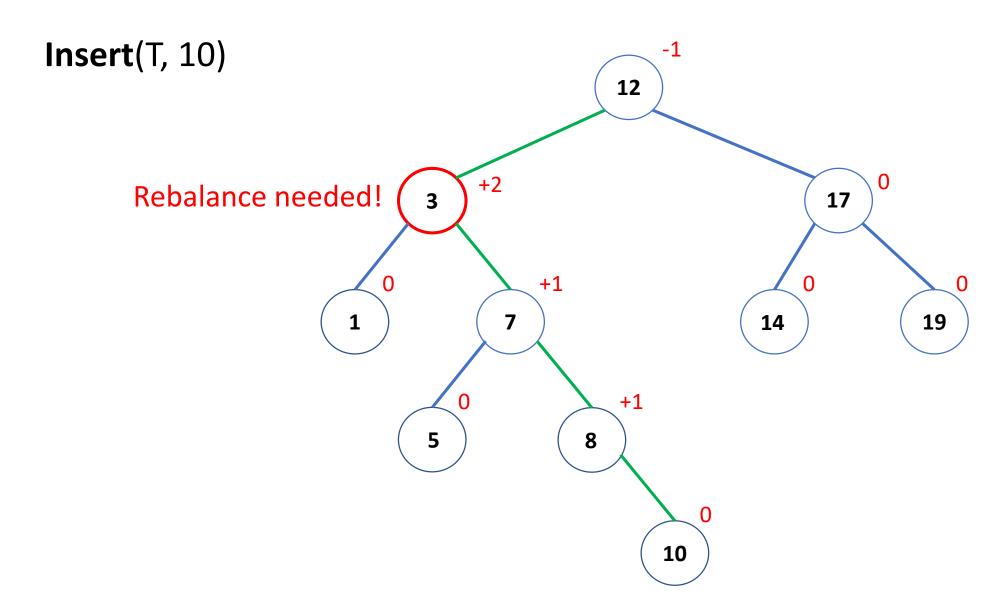


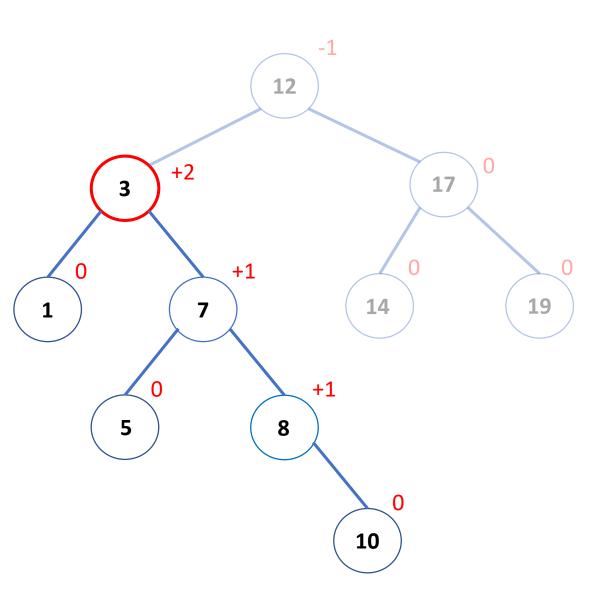


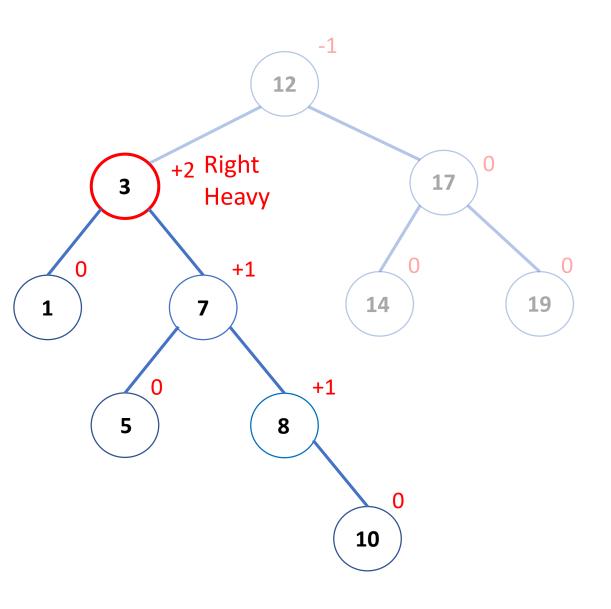


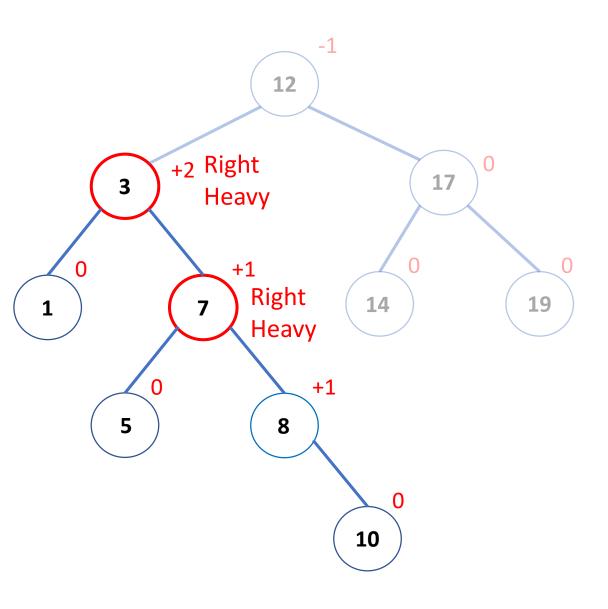


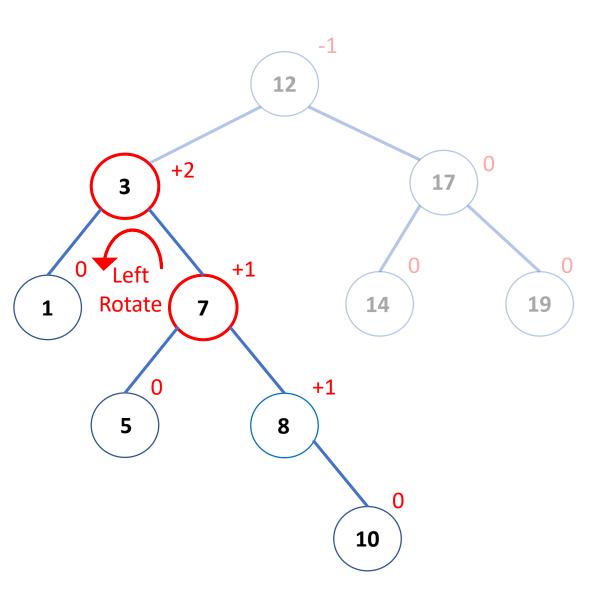


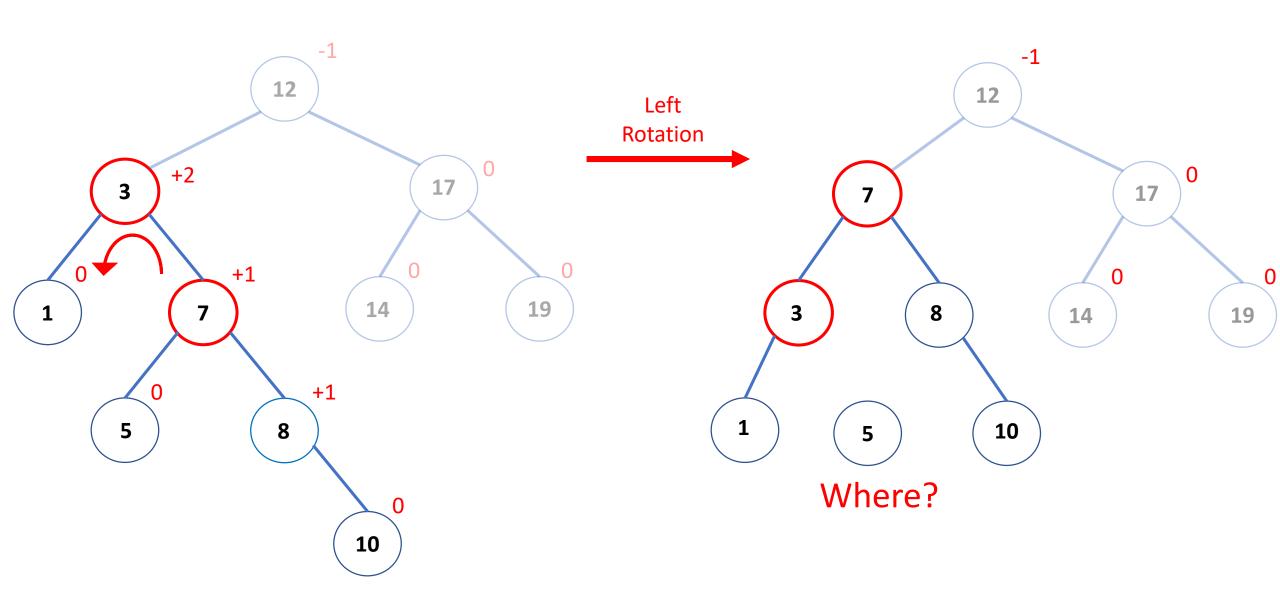


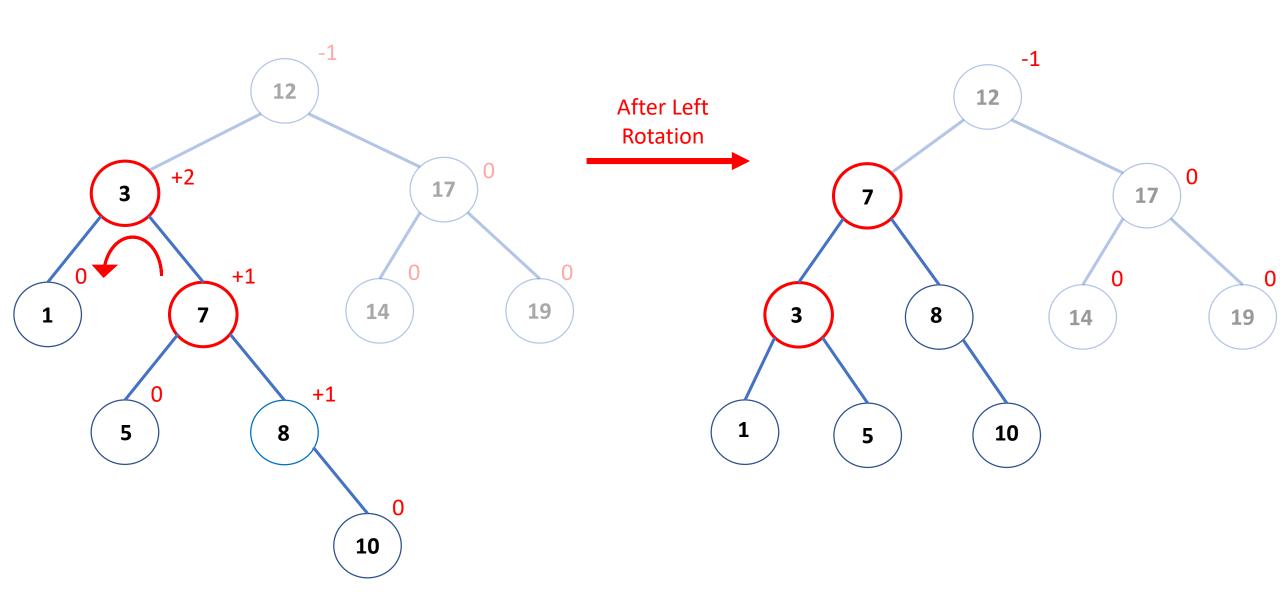


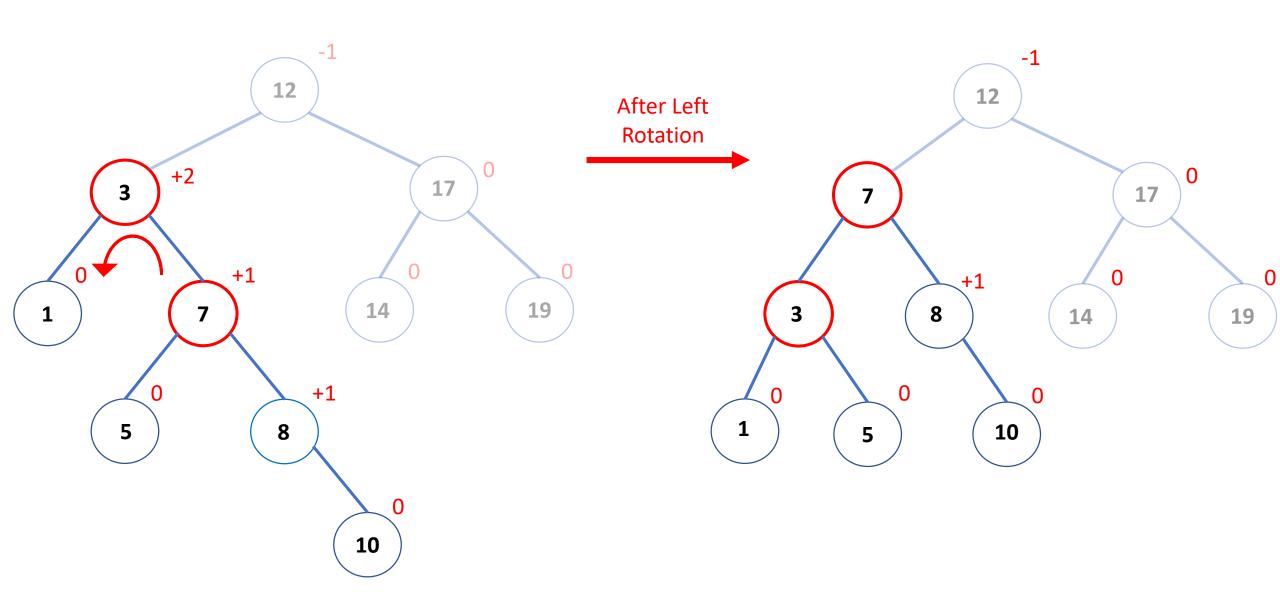


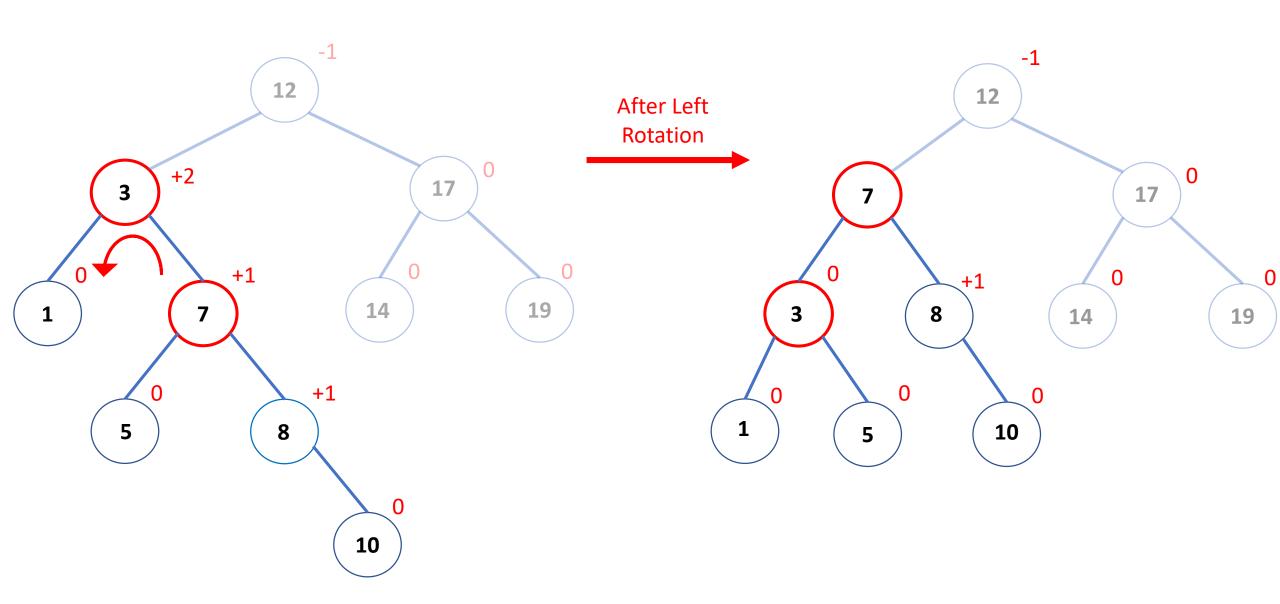


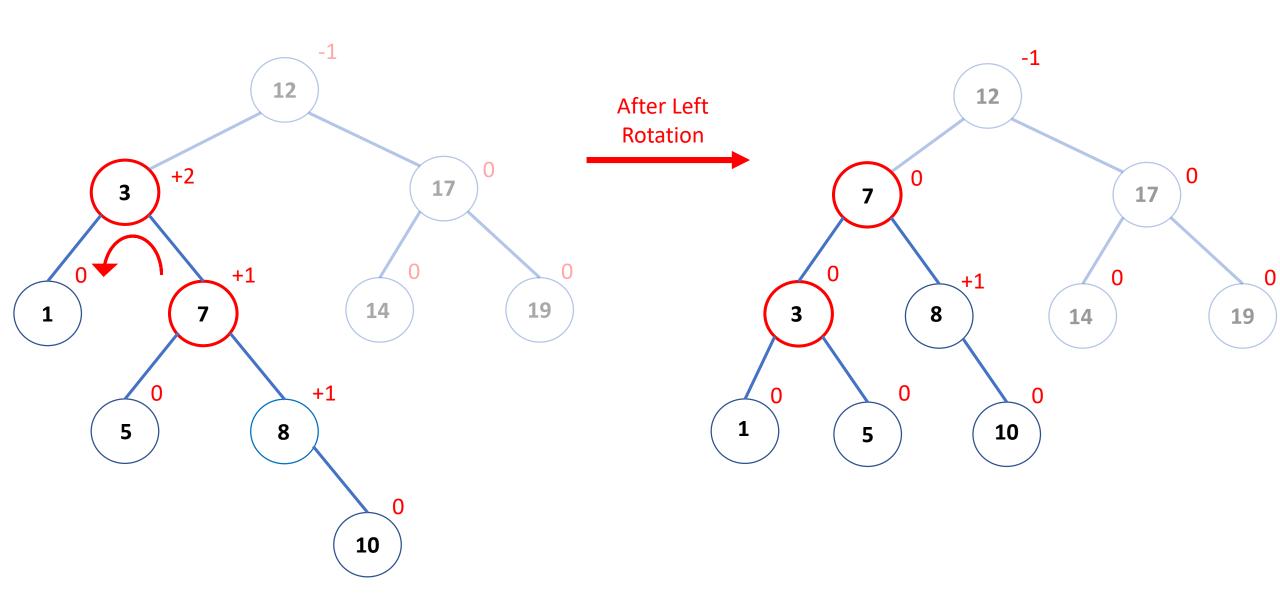


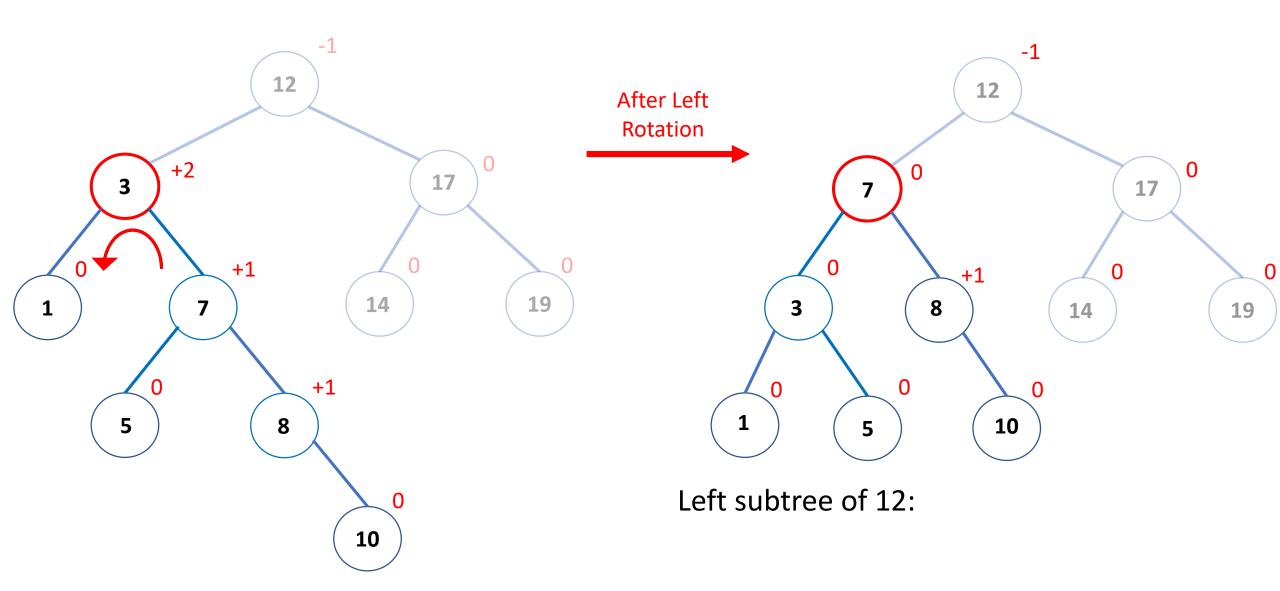


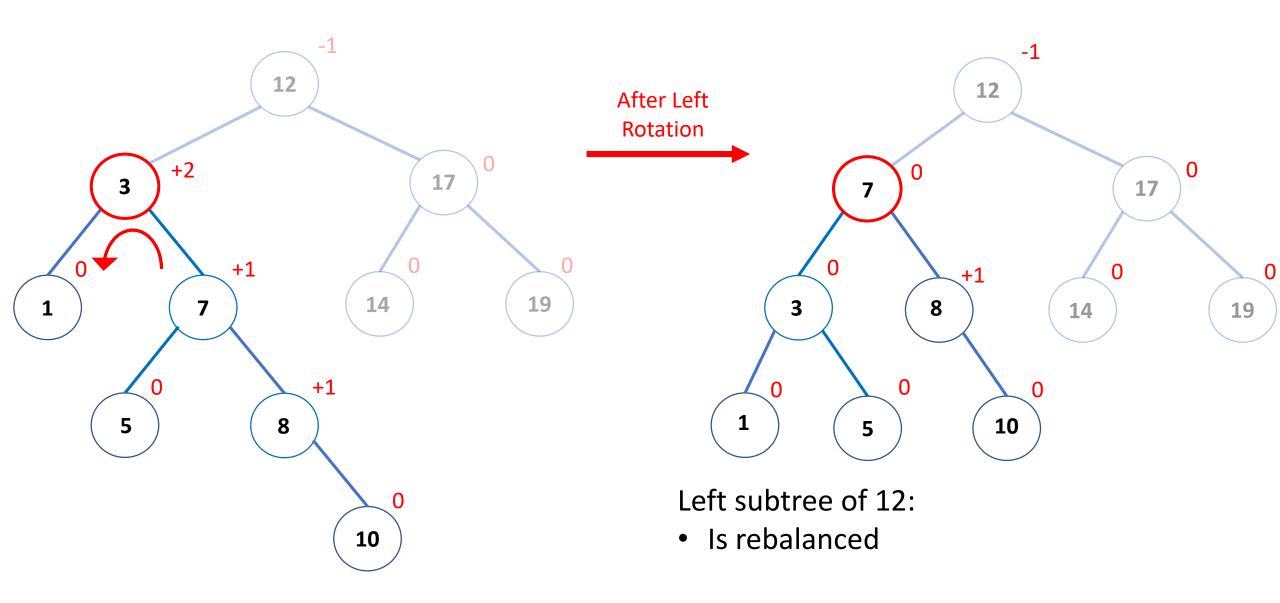


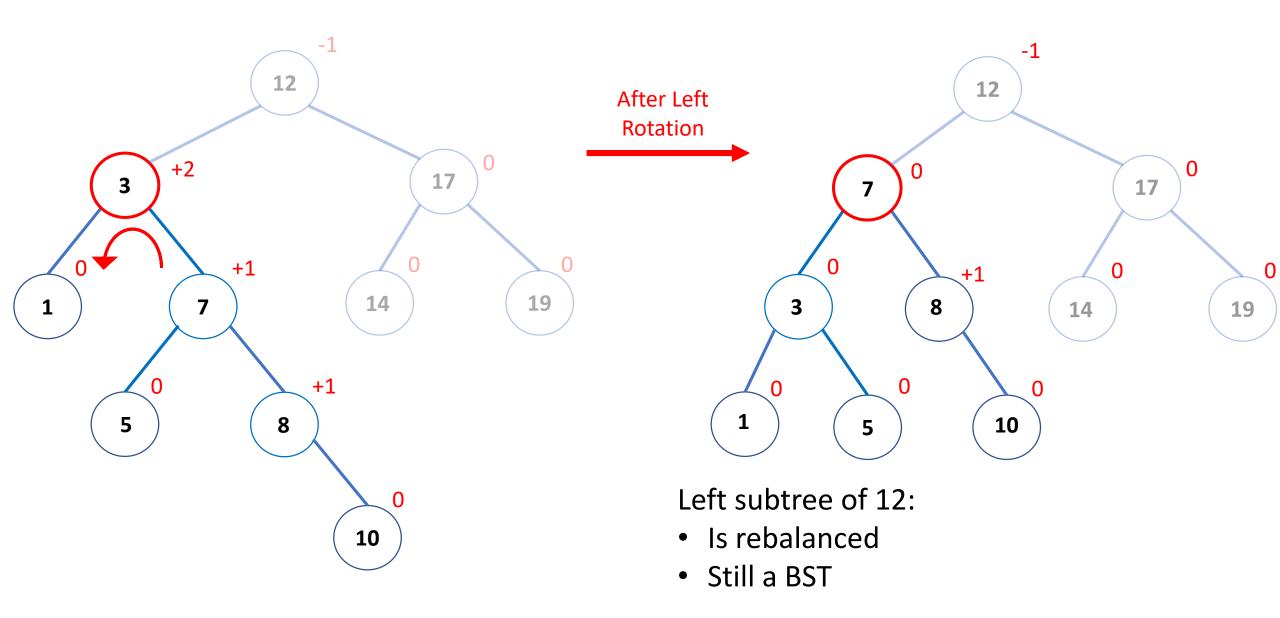


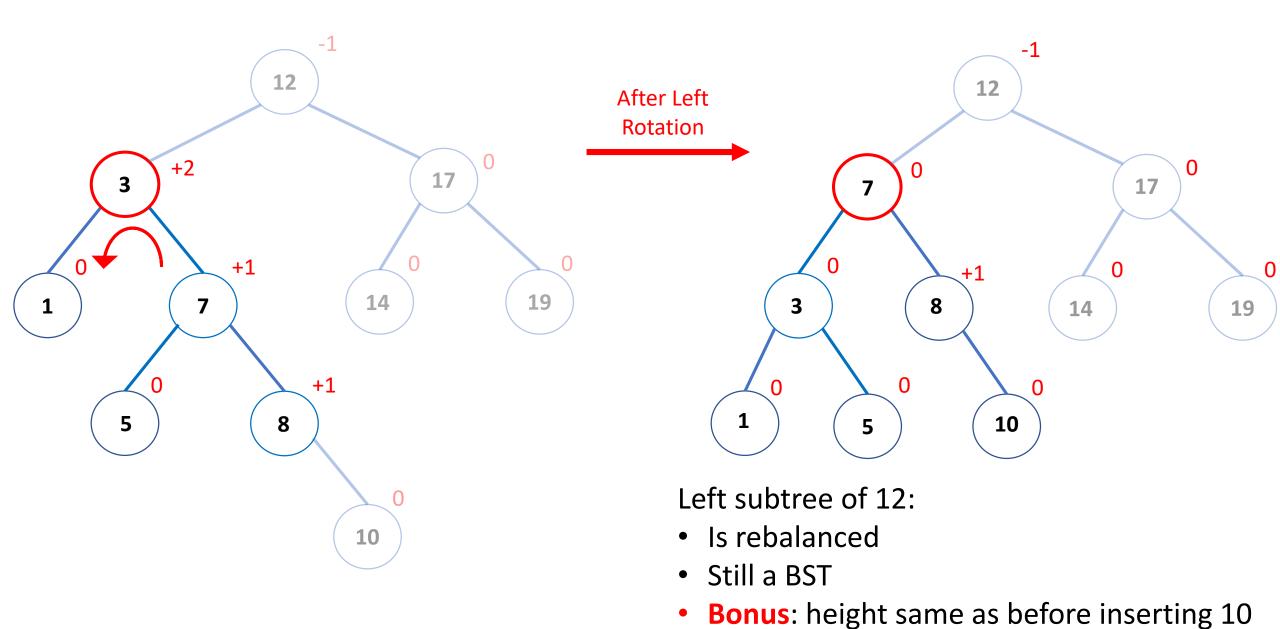


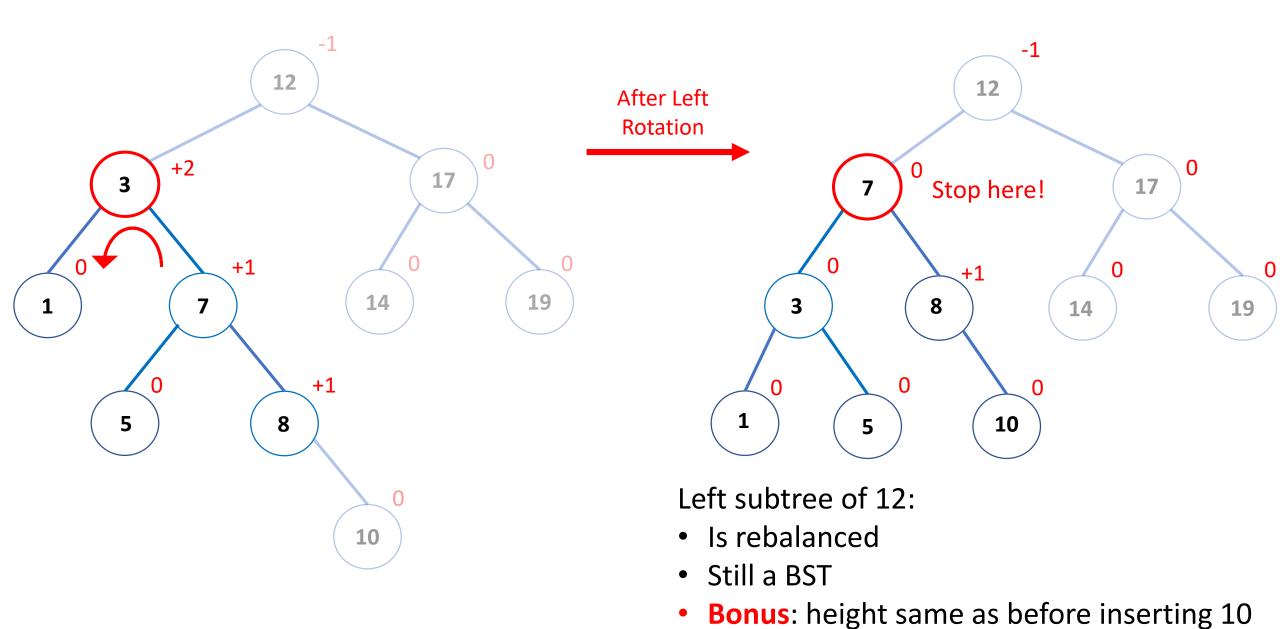


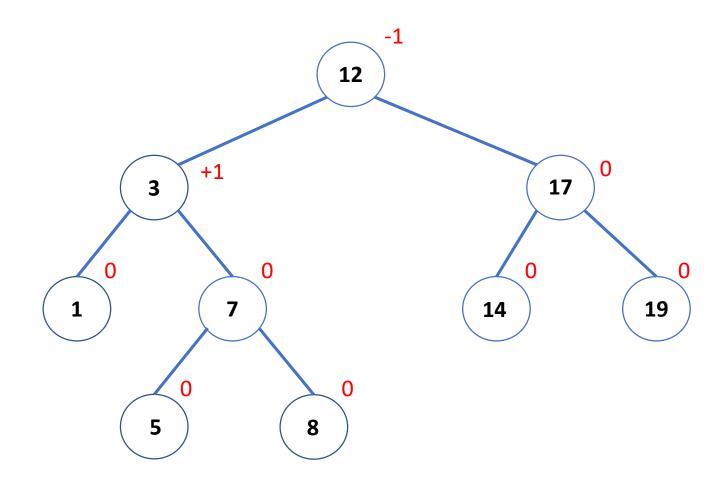


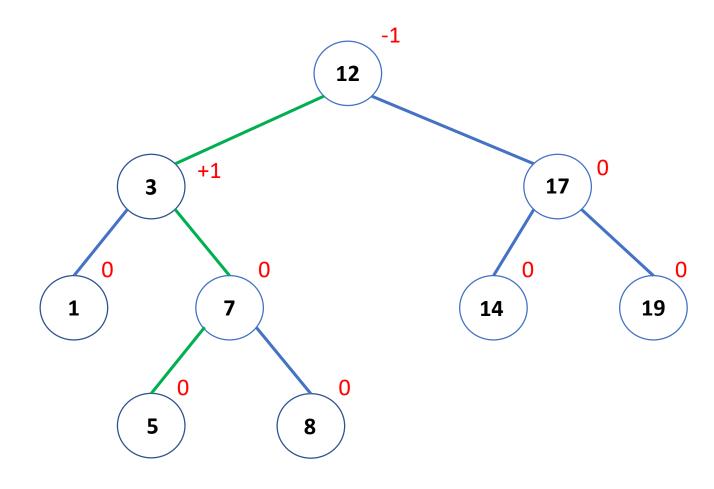


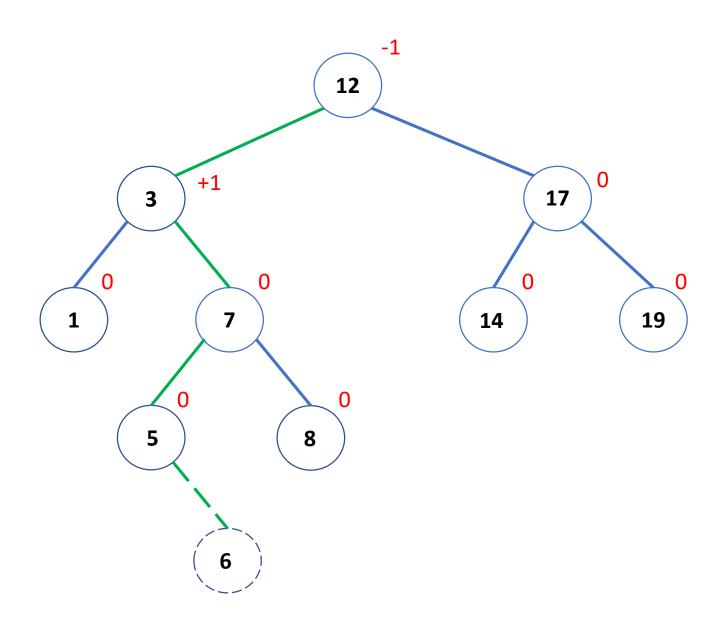


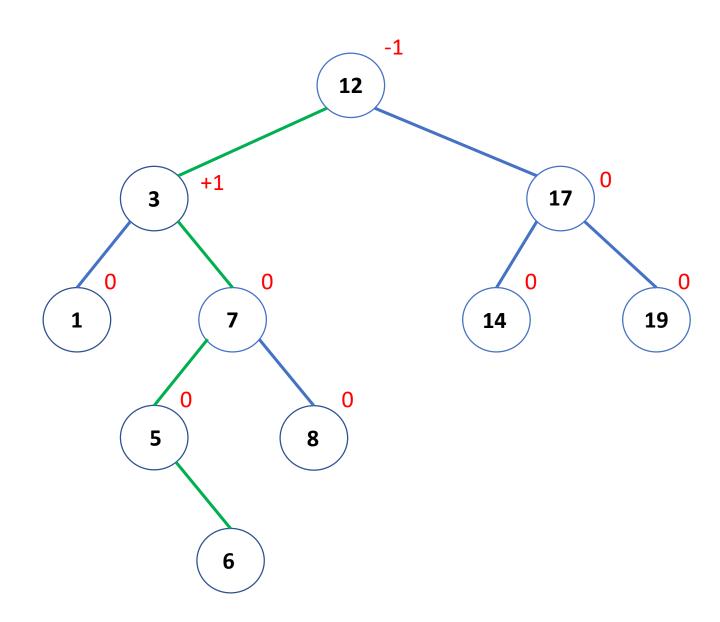


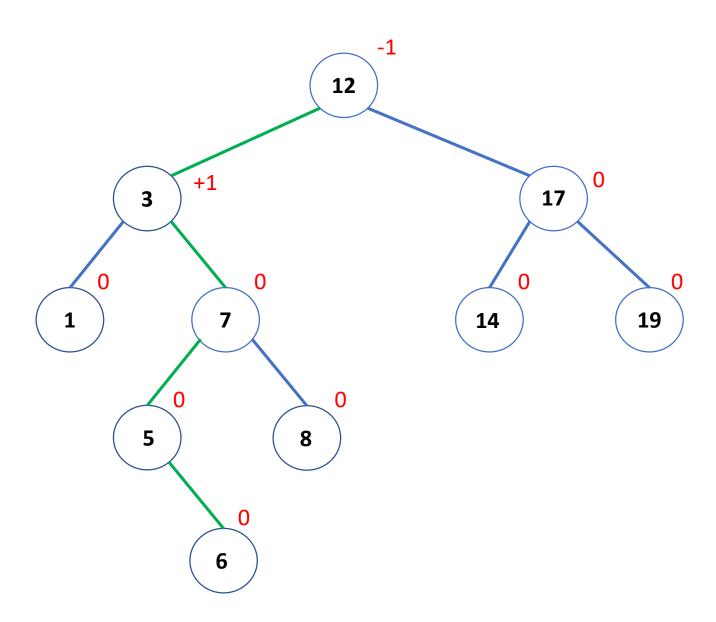


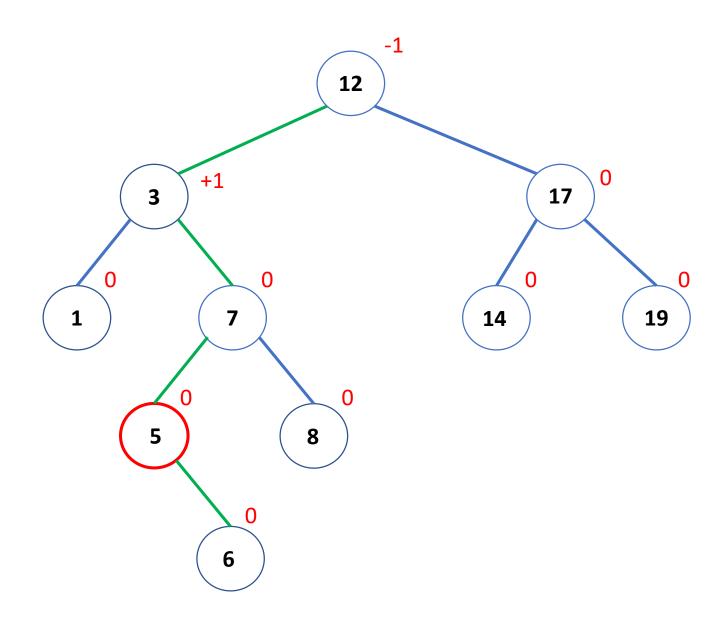


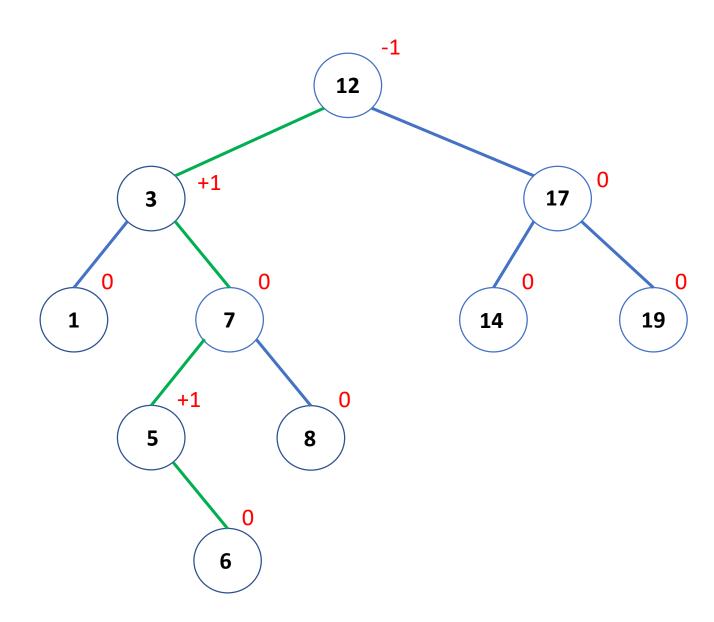


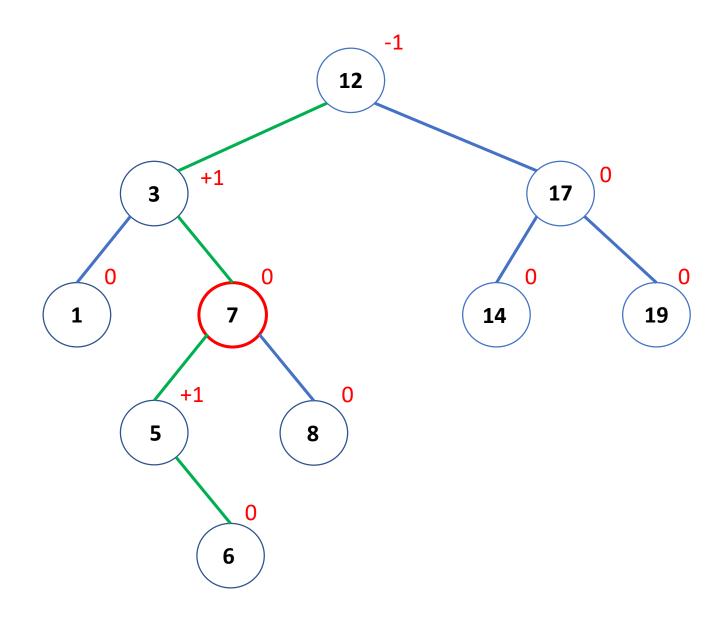


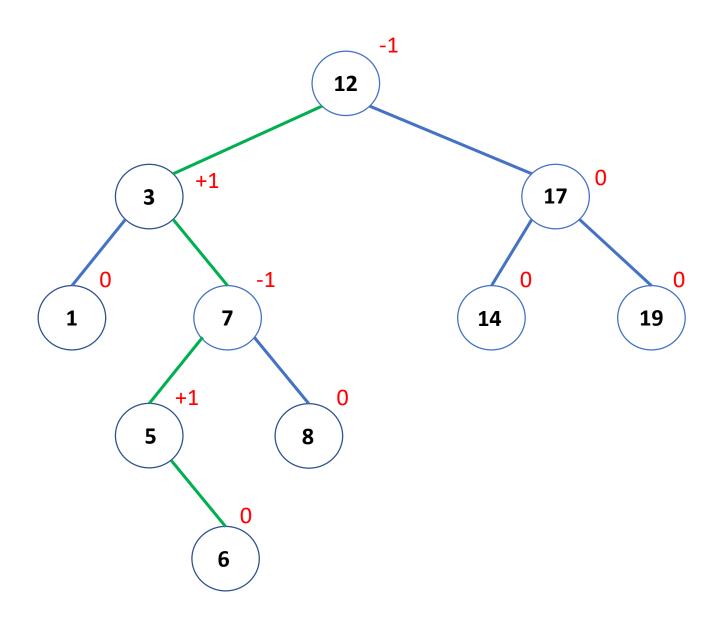


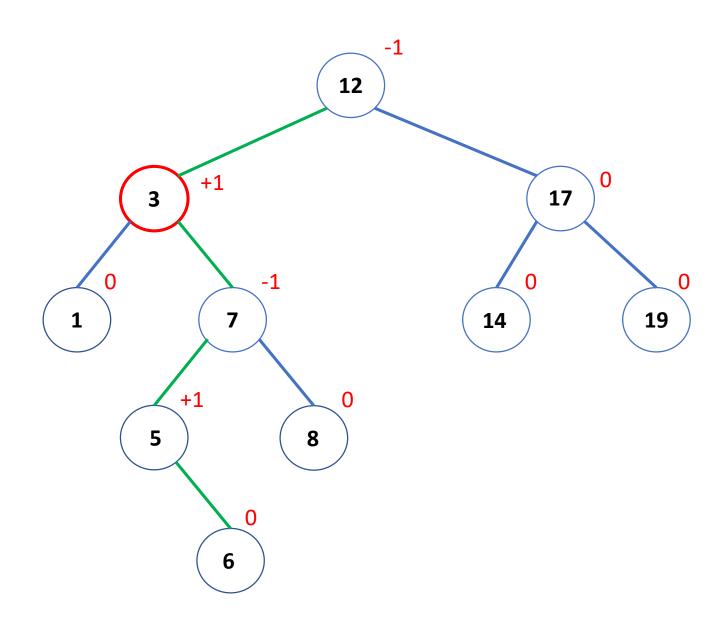


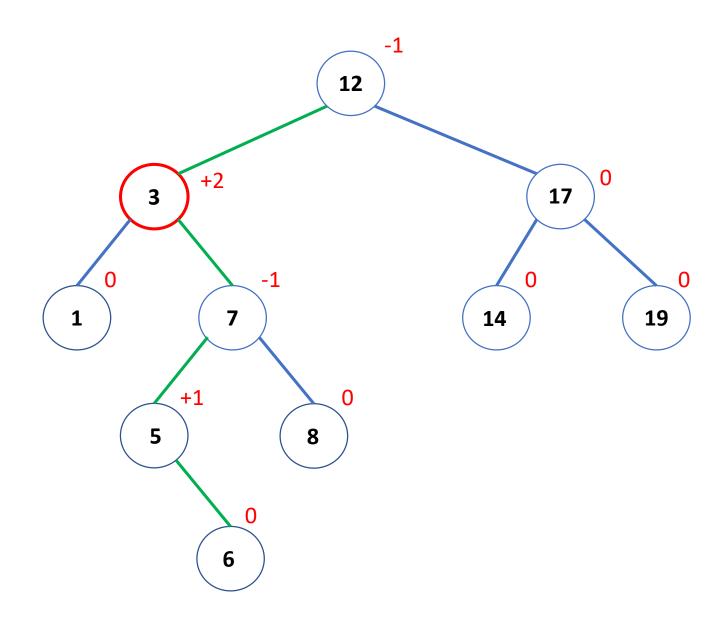


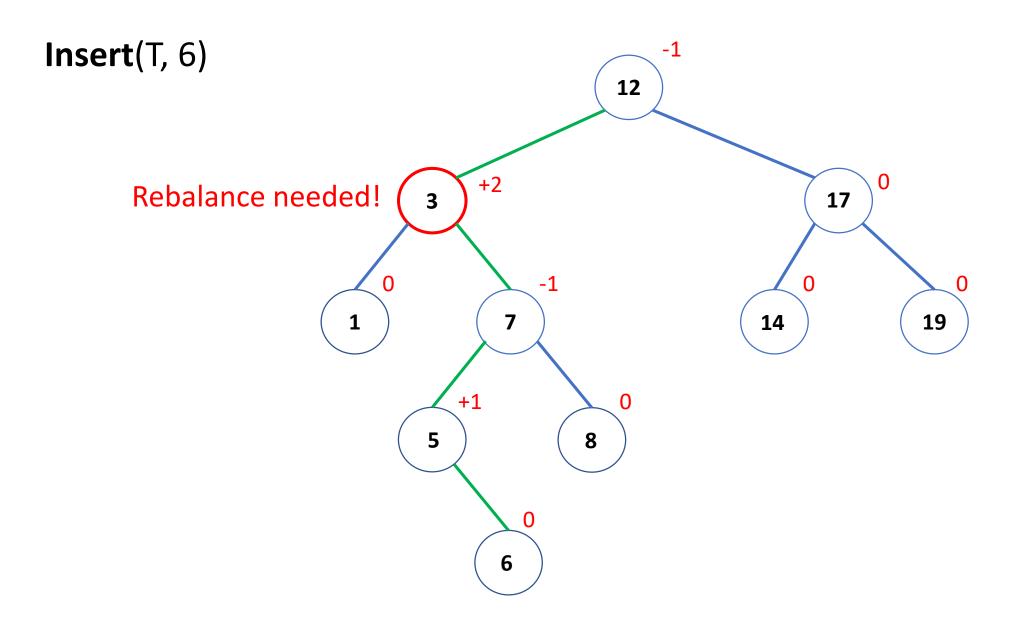


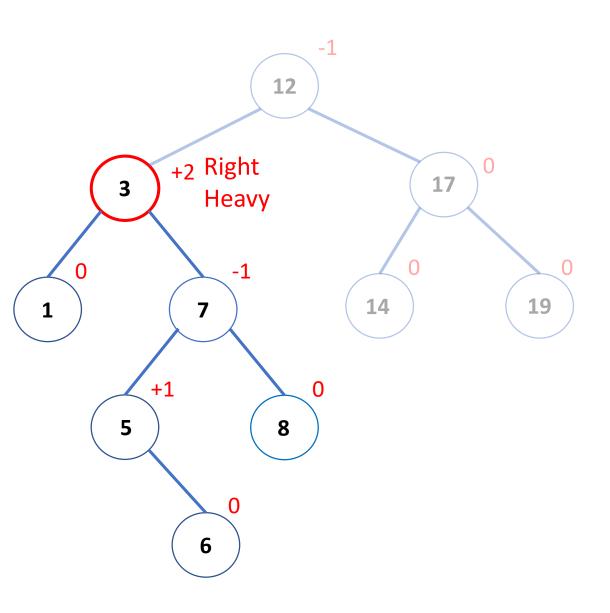


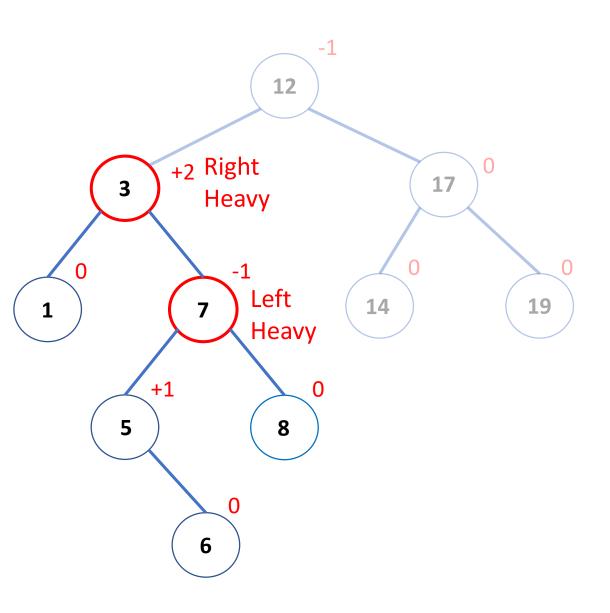


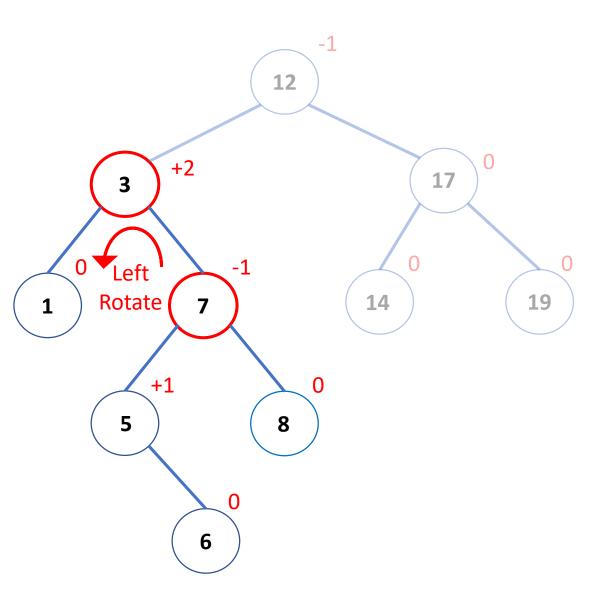


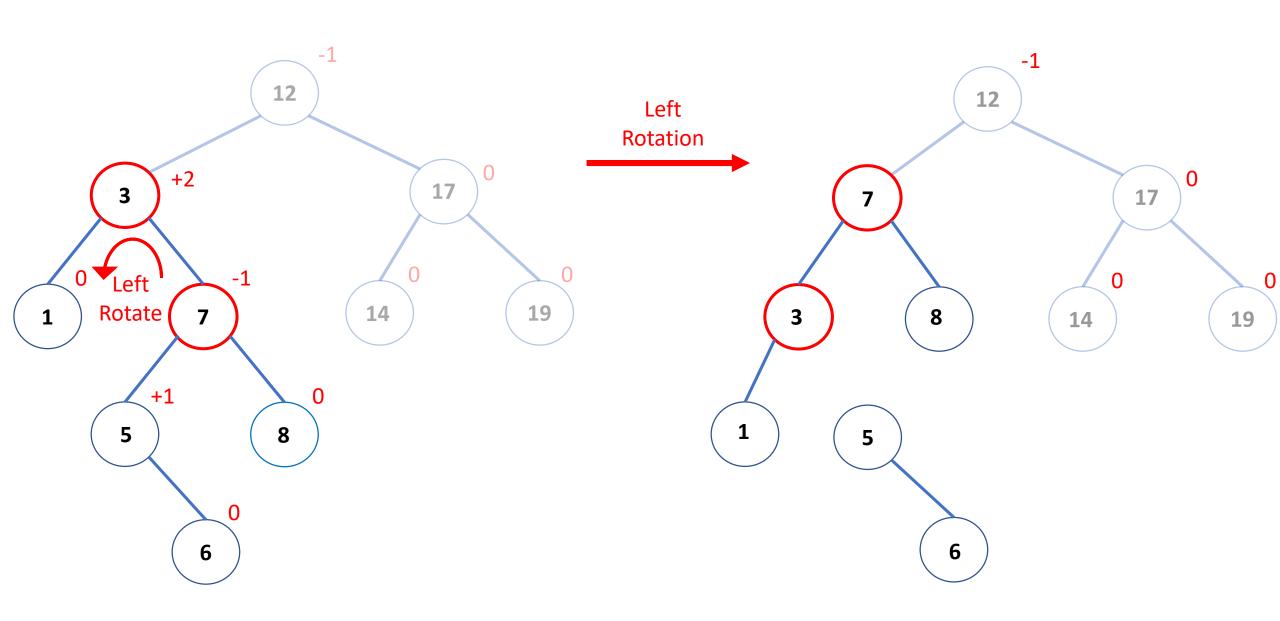


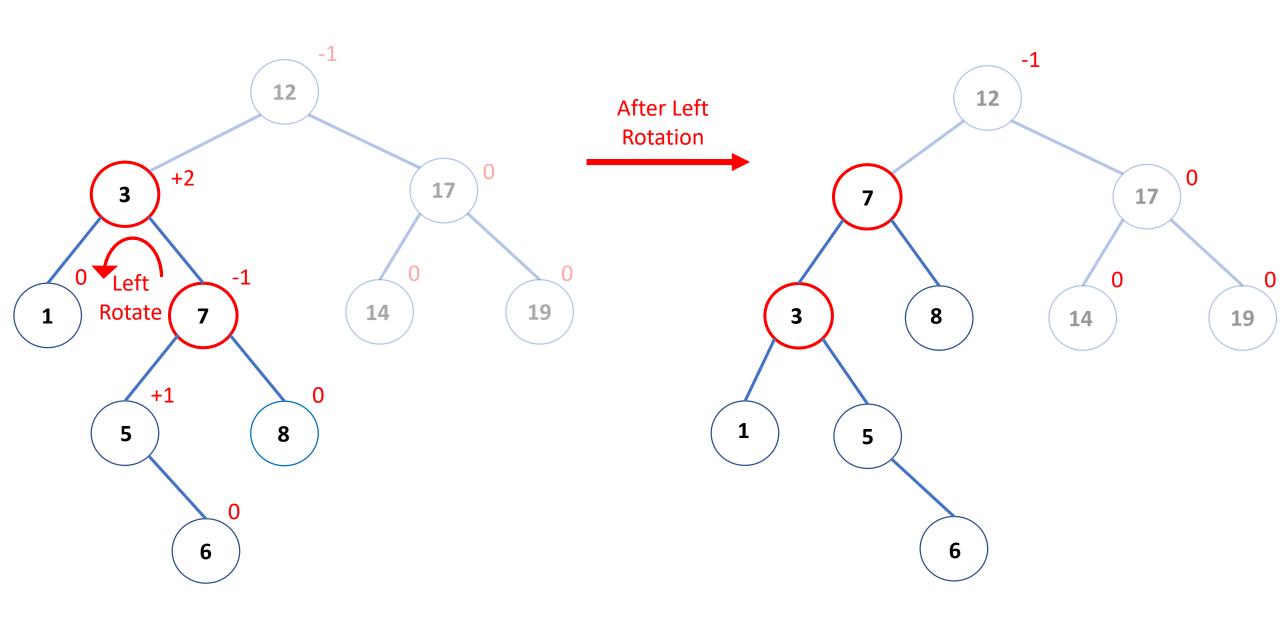


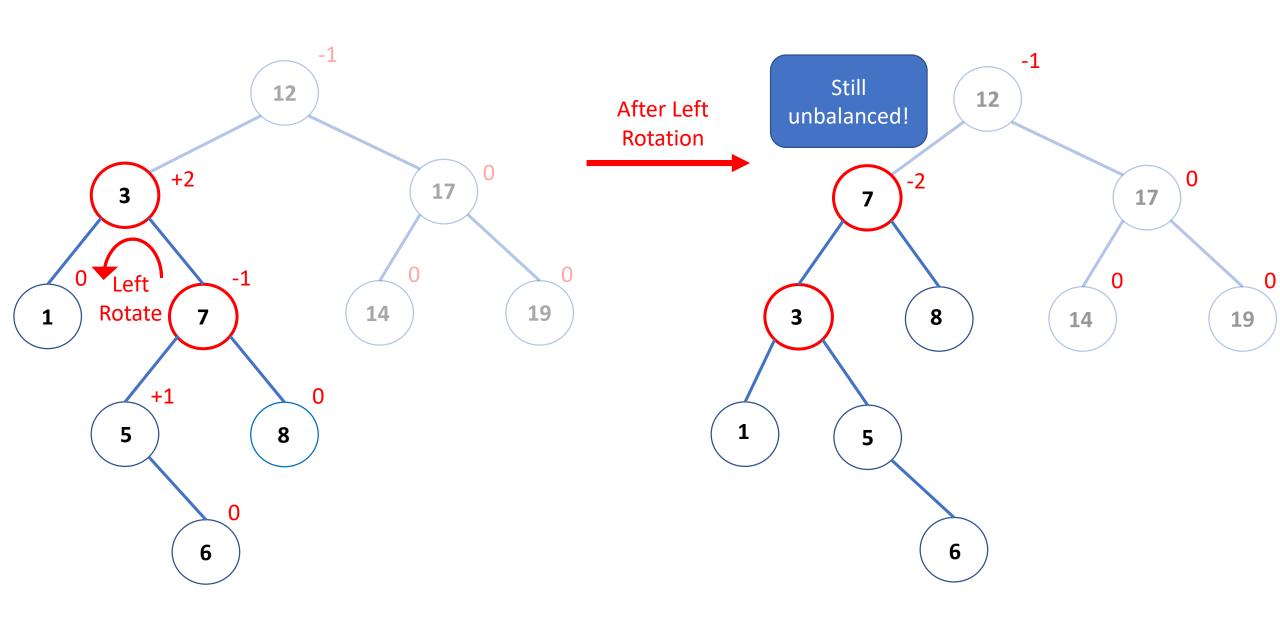


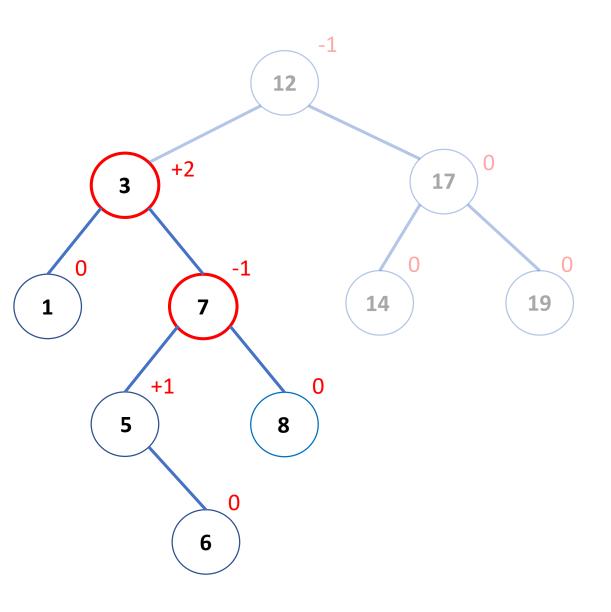


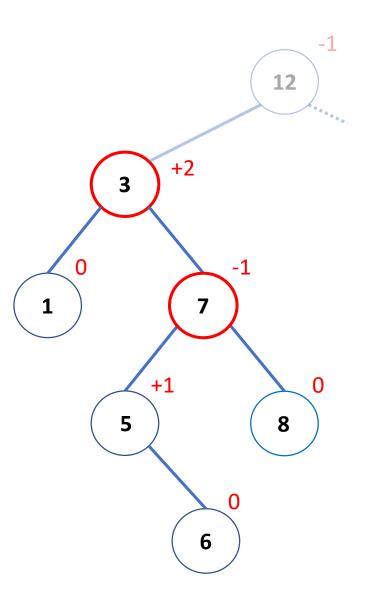


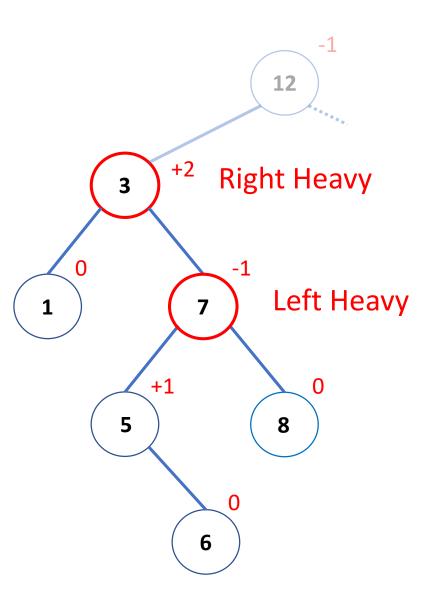




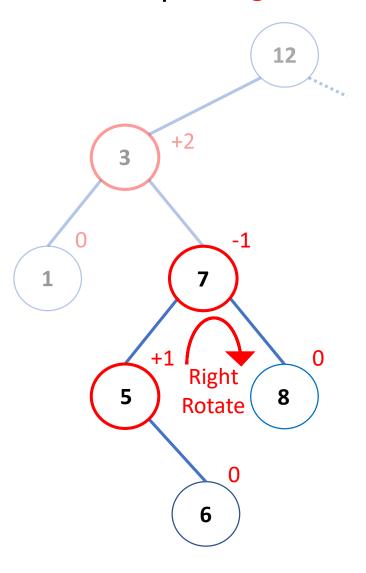




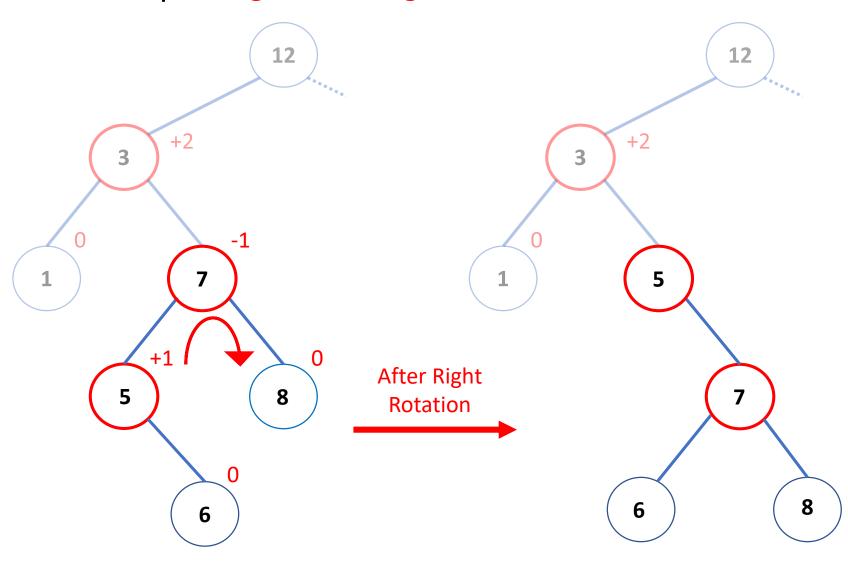




Step 1. Right rotate right subtree of 3



Step 1. Right rotate right subtree of 3



Step 1. Right rotate right subtree of 3 Step 2. Left rotate left subtree of 12 12 12 +2 Left Rotate After Right 8 Rotation 6 6

