

Walk in the Park

Ended

Edit

Description

Given a 2d array with 'U', 'D', 'L', 'R' representing up, down, left, right => count the number of steps you have taken before reaching an already visited cell or going outside the 2d array.



Note: The first step is always taken to cell located at (0,0).

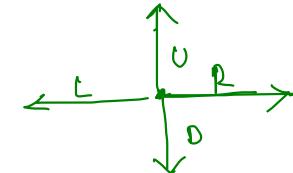


Input

The first line of the input contains T, the number of test cases.

The first line of each test case contains N, and M, the number of rows, and columns in the 2D Array.

The next lines of the test case, contains M characters each, denoting the values in the 2D array.



Constraints

$1 \leq T \leq 10$

$1 \leq N, M \leq 100$

Output

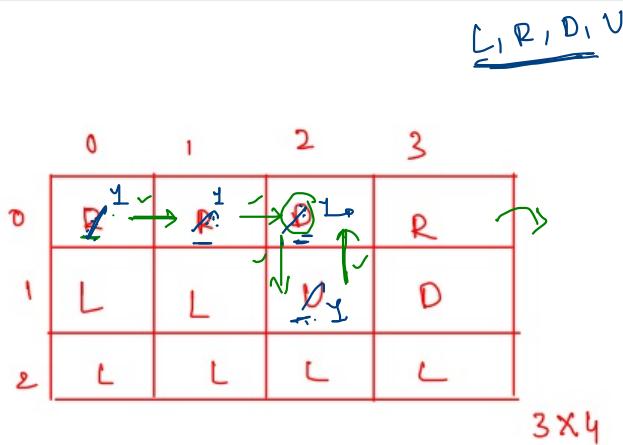
Print a single integer, denoting the number of steps, you take before stepping out of the 2D array, or visiting an already visited cell, on a new line.

Output

Print a single integer, denoting the number of steps, you take before stepping out of the 2D array, or visiting an already visited cell, on a new line.

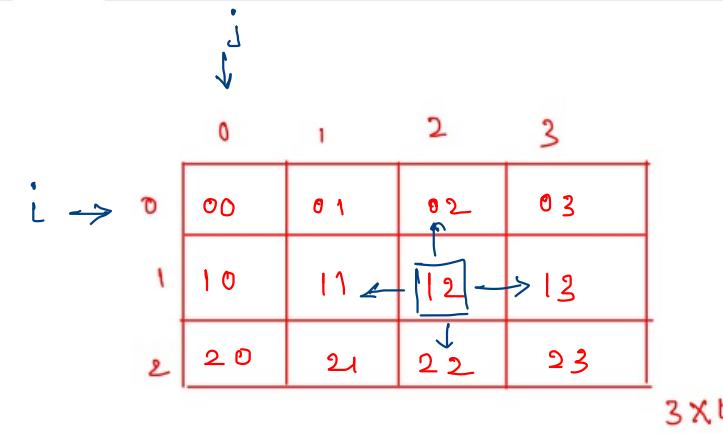
Sample Input 1

```
1  
3 4  
RRDR {  
LLUD }  
LLLL
```

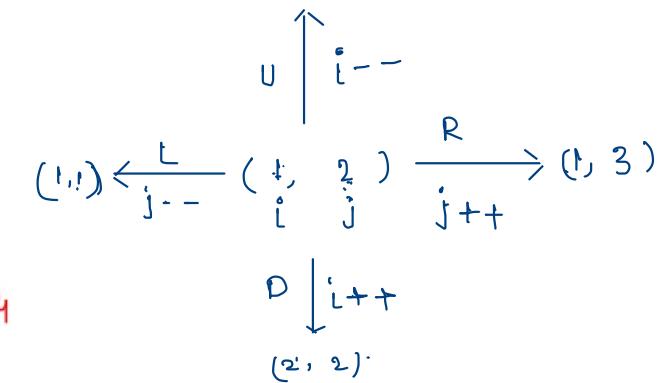


Sample Output 1

```
4
```



(0, 2)



$$\text{cond} = \emptyset \neq \mathbb{Z}\mathbb{Z}_4$$

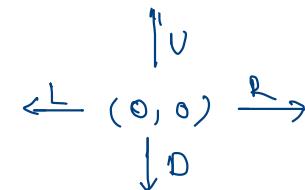
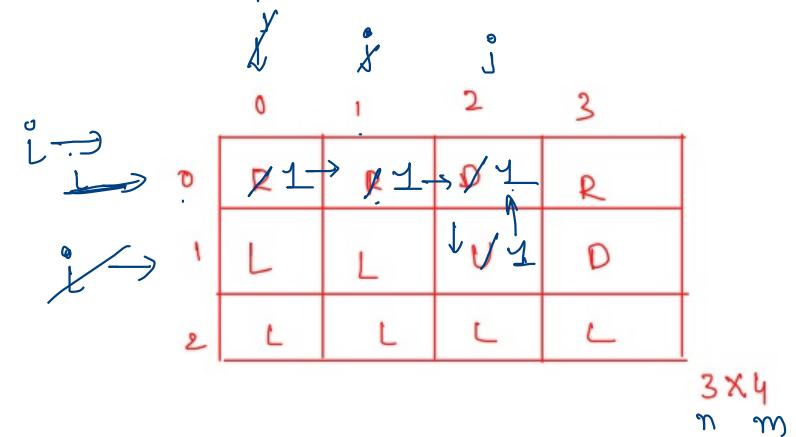
```

function countSteps(arr[][], n, m)
{
    i=0, j=0, count=0
    while(i>=0 && i<n && j>=0 && j<m) ✓
    {
        if(arr[i][j]==1)
            return count
        temp=arr[i][j]
        arr[i][j]=1
        if(temp=='L')
            j--
        else if(temp=='R')
            j++
        else if(temp=='U')
            i--
        else
            i++     $\Rightarrow$  D
        count++
    }
    return count
}

```

~~temp =~~

| | | | |
|---|--|--|--|
| R | | | |
| R | | | |
| D | | | |
| U | | | |



(0, 2)

$(\downarrow, \downarrow) \xleftarrow[\downarrow, \downarrow]{} (\uparrow, \frac{\downarrow}{\downarrow}) \xrightarrow[\downarrow, \downarrow]{R} (\uparrow, 3)$

x4

$D \downarrow \uparrow$
(2, 2)

Repeated and Missing

Ended

Edit

Description

Given an array of size n . Array elements are in the range from 1 to n . In the given array, one number from set $\{1, 2, \dots, n\}$ is missing and one number occurs twice in the array. Find these two numbers.

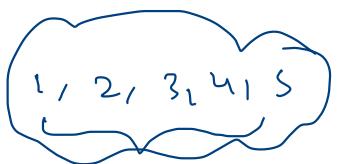
Print the missing one first and then the repeated one with a space character in between them.

Input

The first line of the input contains one integer t ($1 \leq t \leq 10$) — the number of test cases. Then t test cases follow.

The first line of each test case contains a single integer n ($2 \leq n \leq 100000$).

The second line of the test case contains n integers ($1 \leq A_i \leq n$).

$n=5 \rightarrow$ 

Output

For each test case, print the answer.

Sample Input 1

```
3
5
1 2 3 3 4
2
1 1
3
1 2 2
```

Sample Output 1

```
5 3
2 1
3 2
```

[1, 2, 3, 3, 4]

↳ 3 repeat
↳ 5 missing

arr

| 0 | 1 | 2 | 3 | 4 |
|---|---|---|----------|---|
| 1 | 2 | 3 | <u>3</u> | 4 |

$n=5$

Set: { 1, 2, 3,
• hasL) dup = 3 ✓

$n \cdot \log^2 n$
→ Sort +
log } duplicates

①

Step 1:- Set / Key + value
Obj \Rightarrow dup = 3 ✓

• Step 2:- $\text{sum_n} = \frac{n(n+1)}{2} \Rightarrow 1 + 2 + 3 + 4 + 5$

① Step 3:- $\text{sum_arr} = 1 + 2 + \underline{3} + 3 + 4$ { sum. of all elements of the given array)

• Step 4:- $\underbrace{\text{sum_arr}}_{\text{sum_arr}} = \text{sum_arr} - \text{dup} \Rightarrow 1 + 2 + 3 + 4$

• Step 5:- missing = $(\text{sum_n}) - (\text{sum_arr}) \Rightarrow 1 + 2 + 3 + 4 + 5 - 1 - 2 - 3 - 4 = 5$

Q(19)

✓ Stack & Queue

Stack

LIFO / FILO

comes
Last → first



comes
first → last

✓ push(α)

overflow

✓ pop(α)

⇒ Can I pass any
argument?

Peek()

underflow

Queue

FIFO

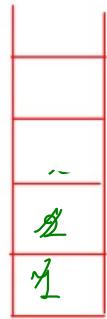
Enq(α)

Deq(α)



Which of the following permutations can be obtained in the output (in the same order) using a stack assuming that the input is the sequence 1, 2, 3, 4, 5 in that order?

- x A. 3, 4, 5, 1, 2 \Rightarrow
- ✓ B. 3, 4, 5, 2, 1
- x C. 1, 5, 2, 3, 4 \Rightarrow
- x D. 5, 4, 3, 1, 2 \Rightarrow



3 4 5 2 1

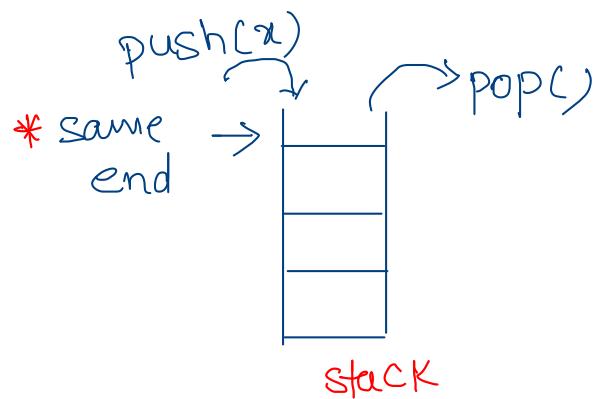
None

Stack [LIFO / FILO]

→ ① push(x) → overflow

② pop() → underflow

→ both operations, will happen
on the same end

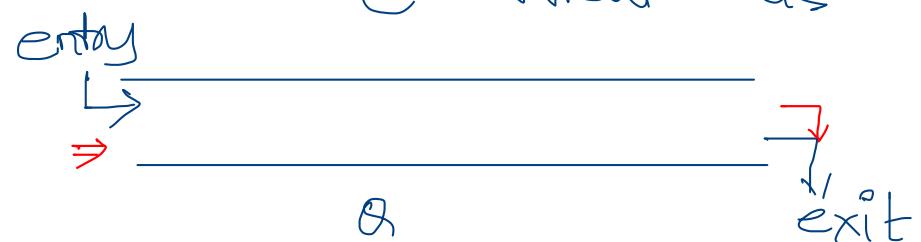


Queue (FIFO)

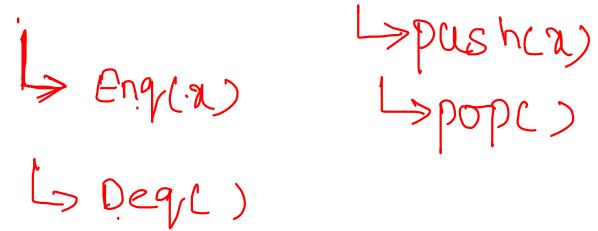
① Enq(x) → OF

② Deq() → UF

→ both operations will happen
@ different ends.



Implement Queue using Stacks



`Enq(x)`
{

`push(x) | pop()`

}

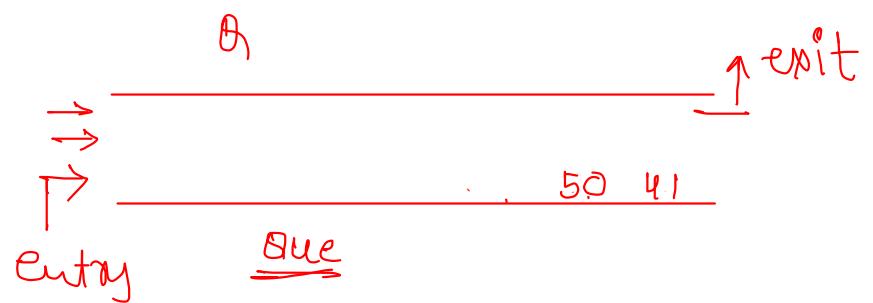
`Deq()`

{

`push(x) | pop()`

}

E(10), E(20), Dl₁₀, E(21), E(41), Dl₂₀, Dl₂₁, E(50), Dl₄₁



↑

2-ptx'-stuc

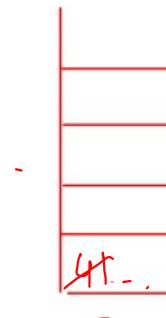
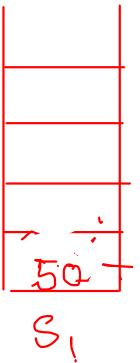
Eng(x) :-

$S_1 \cdot \text{push}(x)$

Dl(x) :-

1. [Move all ele's $S_1 \rightarrow S_2$] \Rightarrow when S_2 isEmpty

2. $S_2 \cdot \text{pop}()$ ✓

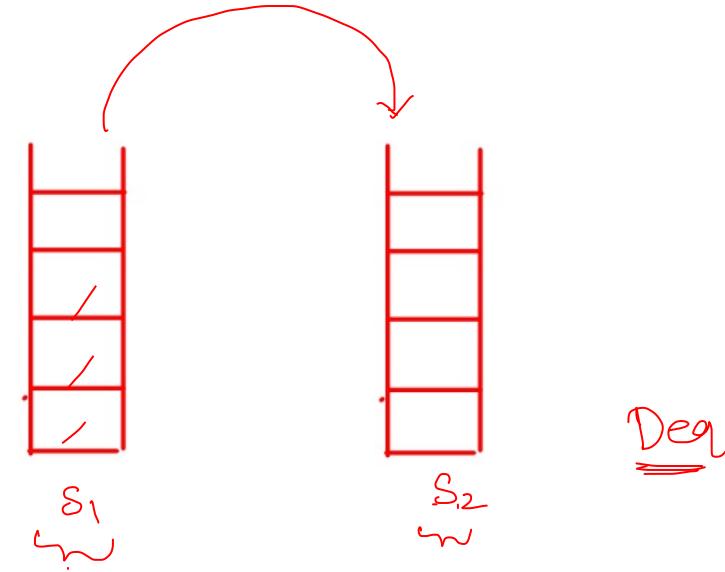


$\rightarrow 10$
 $\rightarrow 20$
 $\rightarrow 21$
 $\rightarrow 41$

PS-Code

```
void enqueue(Q,x)
{
    push(S1,x) ✓
}

dequeue(Q)
{
    →if(S2.isEmpty())
    {
        →if(S1.isEmpty())
        {
            print("Q is empty") ✓
            return
        }
        else
        {
            while(!S1.isEmpty())
            {
                x=S1.pop()
                S2.push(S2,x)
            }
        }
        x=S2.pop() ✓
        return x
    }
}
```



d(c)
4
B(L)

• Implement Stack using Queues

↳ push(x)

↳ pop()

↳ enq(x)
deq(c)

→ push(x)

{

Enq(x) | Deql

}

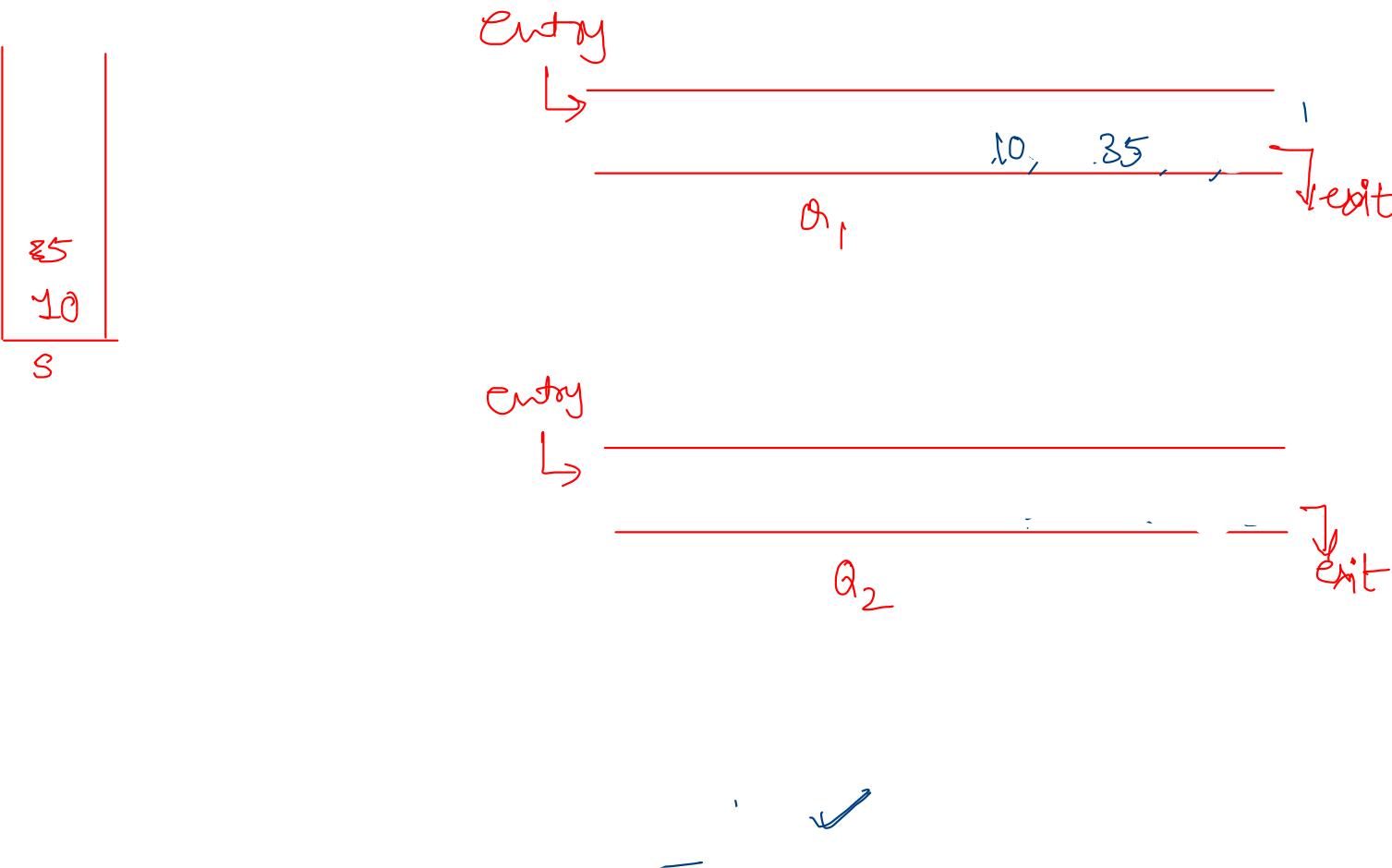
→ pop()

{

Enq(x) | Deql

}

$\checkmark \text{push}(10)$, $\checkmark \text{push}(20)$, $\checkmark \text{push}(30)$, $\checkmark \text{pop}()$, $\checkmark \text{pop}()$, $\checkmark \text{push}(35)$, $\checkmark \text{push}(40)$, $\checkmark \text{pop}()$



$\text{push}(x)$:-

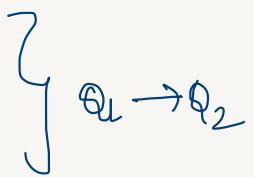
1. add to Q_2 ✓
2. $Q_1 \rightarrow Q_2$
3. Swap(Q_1, Q_2)

$\text{pop}()$:-

$Q_1 \cdot \text{deql}$)

Let Q1, Q2 be two Queues ✓
~~size=0~~

```
void push(x)
{
    1. Q2.enqueue(x) ✓
    2. while(!Q1.isEmpty())
        {
            Q2.enqueue(Q1.front())
            Q1.dequeue()
        }
    3. let temp be a Queue
        temp=Q1
        Q1=Q2
        Q2=temp
}
```



} swap(Q1, Q2)

→ pop()
{
 return Q1.dequeue() ✓
}

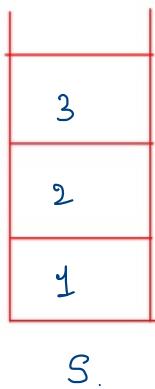


→ top()
{
 return Q1.front()
}

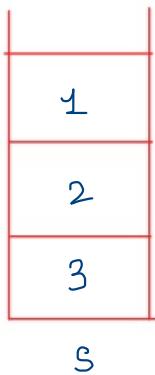
⇒ f, f

11:05

Reverse the contents of Stack



after reversing
?



2-ptor tech
+
Swap() } \Rightarrow array[] ✓

Ap: —

\rightarrow use another stack
(S_2)

S_1 : Input()

$O(n)TC$

S_2 : move all $S_1 \rightarrow S_2$

$O(n)SC$

✓ Declare a global stack called st

```
main()
{
    → add some elements to stack
    · reverse() // call the reverse function
}
reverse()
{
    → if(st.size>0)
    {
        1. x=st.peek()
        2. st.pop()
        3. reverse()
        4. insert(x)
    }
}
```

```
insert(x)
{
    1. if(st.isEmpty())
        st.push(x) ✓
    else
    {
        1. a=st.peek()
        2. st.pop() ✓
        3. insert(x)
        4. st.push(a) ✓
    }
}
```

$\log_2 n$

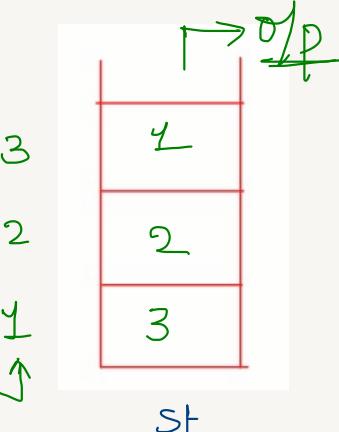
n

n

n^2
not sure

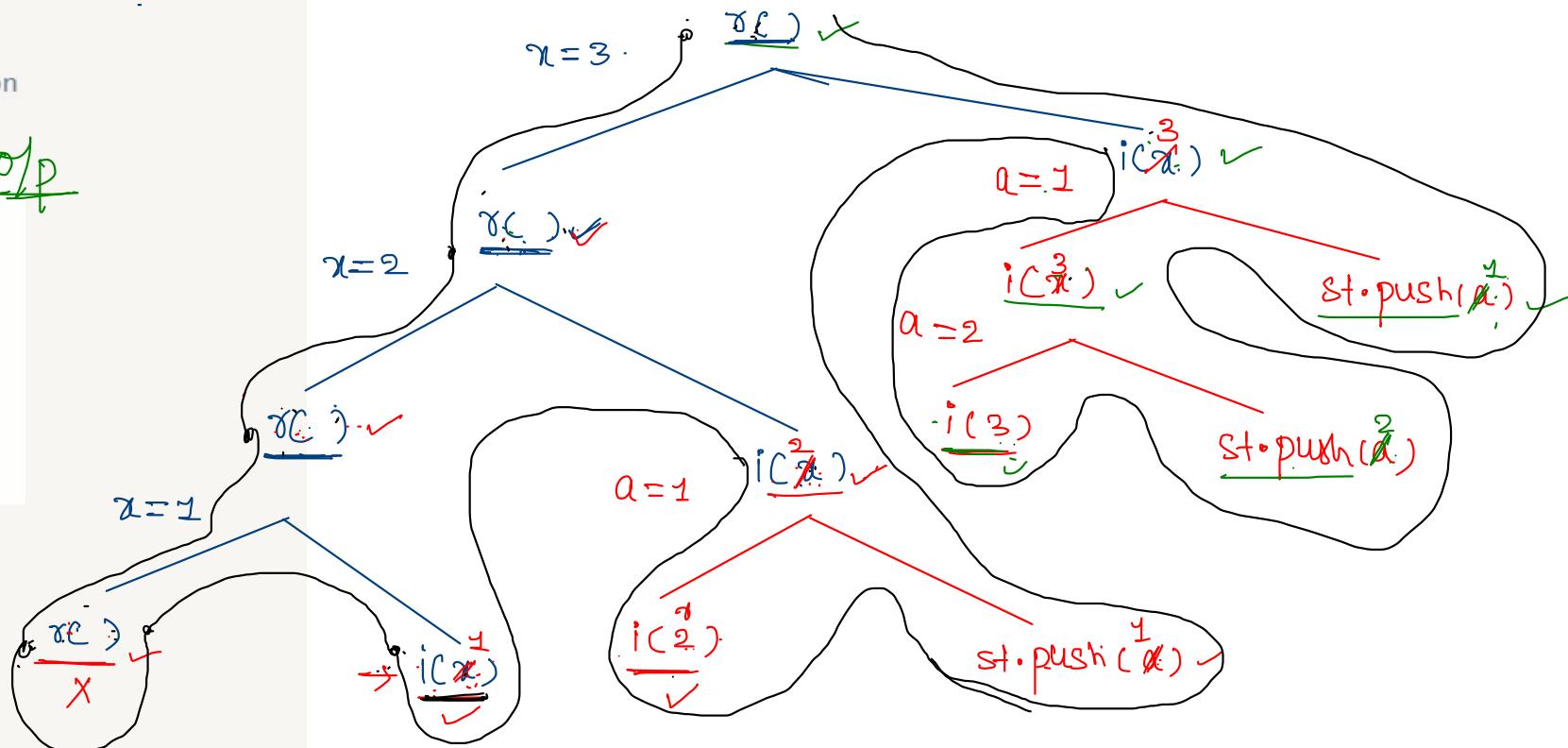
O(1)

peek() v/s pop()



given

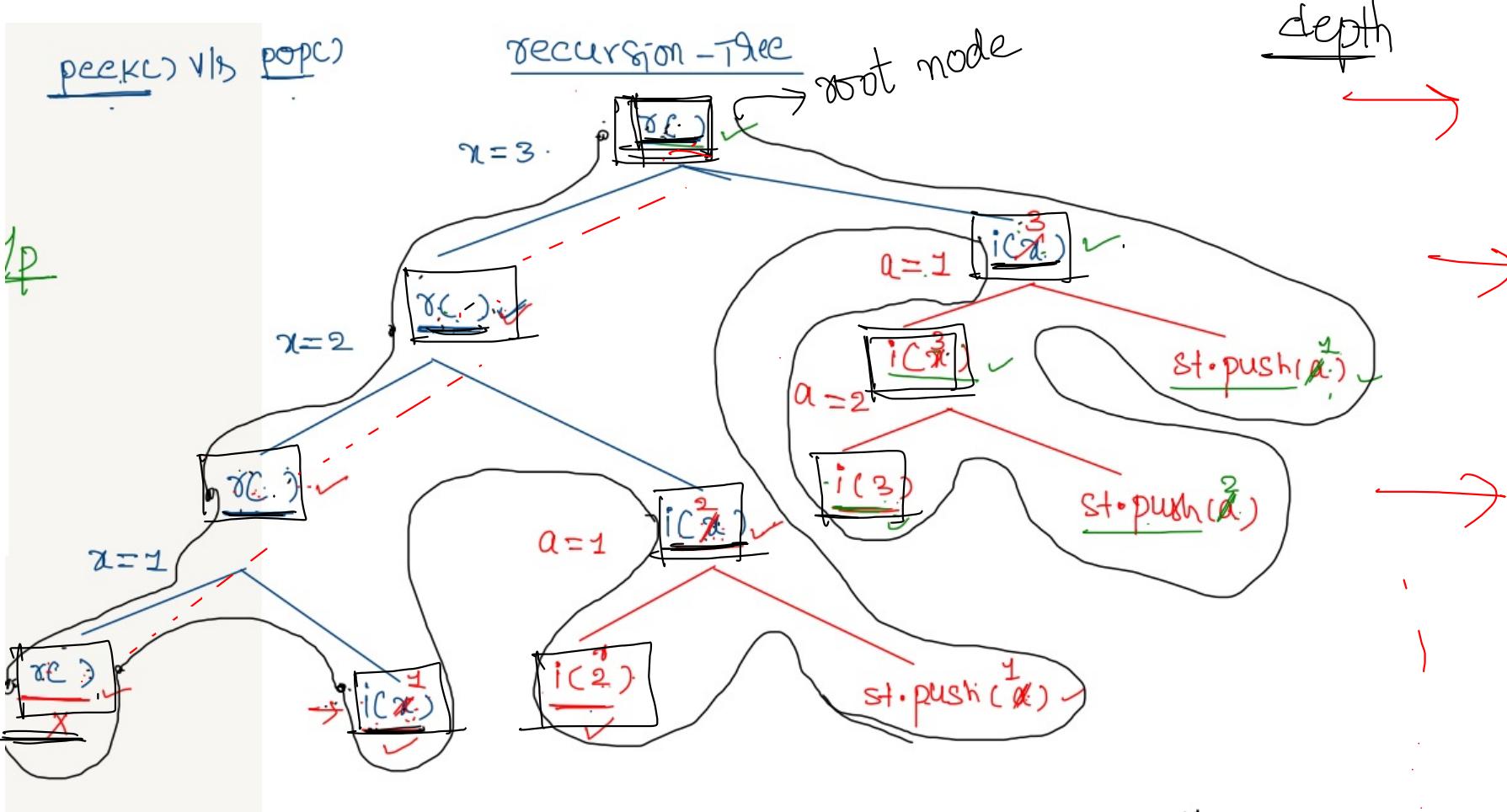
recursion-tree



$$\boxed{n+1} \rightarrow \begin{array}{l} i(1) \rightarrow 1 \\ i(2) \rightarrow 2 \\ i(3) \rightarrow 3 \end{array}$$

$$n + n^2 \Rightarrow \boxed{O(n^2)} \quad \underline{\underline{1+2+3}}$$

$$O(n^2) = \boxed{\frac{n(n+1)}{2}}$$



$3 \rightarrow$ edges [root \rightarrow last level]

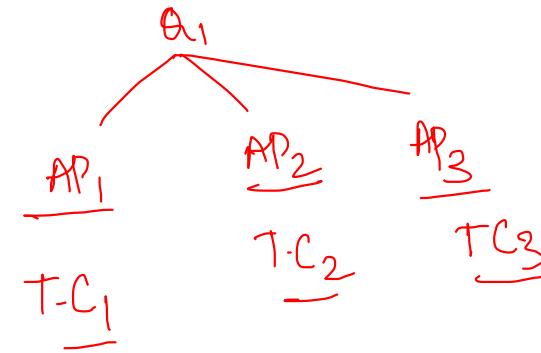
$$n \Rightarrow \underline{\mathcal{O}(n)} \text{ S.C} ; \quad \text{T.C: } \underline{\mathcal{O}(n^r)}$$

T.C
 # of funcalls

S.C
 Max depth

Time and Space Complexity

- 1) if code is given, then what is the T.C of that code
- 2) if two Time Complexities are given, then you should know which one is best



| | Notation | Name |
|--------|--|--|
| 1. | $O(1)$ | constant ✓ |
| 2. | $O(\log \log n)$ | double logarithmic ✓ |
| 3. | $O(\log n)$ | logarithmic ✓ |
| 4. | $O((\log n)^c)$ $c > 1$ $(\log^2 n)$ $(\log n)^3$ | polylogarithmic |
| 5. | $O(n^c)$ $0 < c < 1$ $n^{0.2}$, $n^{0.9}$ | fractional power |
| 6. | $O(n)$ | linear |
| 7. | $O(n \log^* n)$ | n log-star n |
| 8. | $O(n \log n) = O(\log n!)$ | linearithmic, loglinear, quasilinear, or " $n \log n$ " |
| 9. | $O(n^2)$ | quadratic |
| 10. | $O(n^c)$ $c > 2$ n^2 , n^3 , n^4 | polynomial or algebraic |
| X | $L_n[\alpha, c] = e^{(c+o(1))(\ln n)^\alpha (\ln \ln n)^{1-\alpha}}$ $0 < \alpha < 1$ | L-notation or sub-exponential |
| ** 11. | $O(c^n)$ $c > 1$ $\Rightarrow 2^n, 3^n, 4^n$ | exponential ✓ |
| ** 12. | $O(n!)$ ✓ | factorial |

$\Theta \Rightarrow$ Big-oh

$\log_2 n$

↑ Algo; (By-default)

$\rightarrow n^2, n^3, n^4. . .$ (poly)

$\rightarrow 2^n, 3^n, 4^n. . .$ (exp)

steps to decide which T.C is better?

- 1) cancel the common terms [if present]
- 2) apply log to the both T.C's given
- 3) put some very very larger values of n

then see which is having less values, then that is better [it will take less time, it means runs faster]

less time \Rightarrow better.

$$\underline{TC_1} \Rightarrow n^2$$

$$n * \underline{n}$$

let $n = 2$

$$\boxed{n = 2}$$

$\frac{1024}{2}$

$$\underline{TC_2} \Rightarrow n \cdot \log_2^n$$

$$n \cdot \log_2^n$$

$$\log_2^n$$

$$= \log_2^{1024} = 1024 \cdot \frac{\log_2 2}{1} = \underline{1024}$$

$TC_2 < TC_1 \therefore TC_2$ is better

$$TC_1 : \sqrt{n}$$

$$TC_2 : \underbrace{\log_2^n}_n$$

$$\sqrt{n} = n^{1/2}$$

$$\log_2^n$$

$$n = 2^{1024}$$

$$\log_2 \frac{1024}{2}$$

$$= \left[\frac{1024}{2} \right]^{1/2}$$

$$\frac{1024 \cdot \log_2 2}{2}$$

$$= \underline{\underline{\frac{512}{2}}}$$

$$\underline{\underline{1024}}$$

$TC_2 < TC_1 \Rightarrow TC_2$ is better

$$TC_1 : n \cdot \log_2^n$$

$$TC_2 : \sqrt{n} \cdot \log_2^n$$

$$n \cdot \cancel{\log_2^n}$$

$$\sqrt{n} \cdot \cancel{\log_2^n}$$

$$\text{let } n = \frac{n}{2^{1024}}$$

$$2^{1024}$$

$$\sqrt{n} = n^{1/2}$$

$$\frac{512}{2}^{1/2} = 2$$

$$TC_2 < TC_1 \Rightarrow TC_2 \text{ is better.}$$

①

Consider the following three functions.

$$\underline{f_1 = 10^n} \quad \underline{f_2 = n^{\log n}} \quad \underline{f_3 = n^{\sqrt{n}}}$$

Which one of the following options arranges the functions in the increasing order of asymptotic growth rate?

A

$$f_3, f_2, f_1$$

B

$$f_2, f_1, f_3$$

C

$$f_1, f_2, f_3$$

D

$$f_2, f_3, f_1$$

1st < 2nd < 3rd
Small Small Small

Consider the following functions from positive integers to real numbers:

$$10, \sqrt{n}, n, \log_2 n, \frac{100}{n}.$$

The CORRECT arrangement of the above functions in increasing order of asymptotic complexity is:

A

$$\log_2 n, \frac{100}{n}, 10, \sqrt{n}, n$$

B

$$\frac{100}{n}, 10, \log_2 n, \sqrt{n}, n$$

C

$$10, \frac{100}{n}, \sqrt{n}, \log_2 n, n$$

D

$$\frac{100}{n}, \log_2 n, 10, \sqrt{n}, n$$

Which of the given options provides the increasing order of asymptotic complexity of functions f_1 , f_2 , f_3 and f_4 ?

$$f_1(n) = 2^n;$$

$$f_2(n) = n^{3/2};$$

$$f_3(n) = n \log_2 n;$$

$$f_4(n) = n^{\log_2 n}$$

A

$$f_3, f_2, f_4, f_1$$

B

$$f_3, f_2, f_1, f_4$$

C

$$f_2, f_3, f_1, f_4$$

D

$$f_2, f_3, f_4, f_1$$