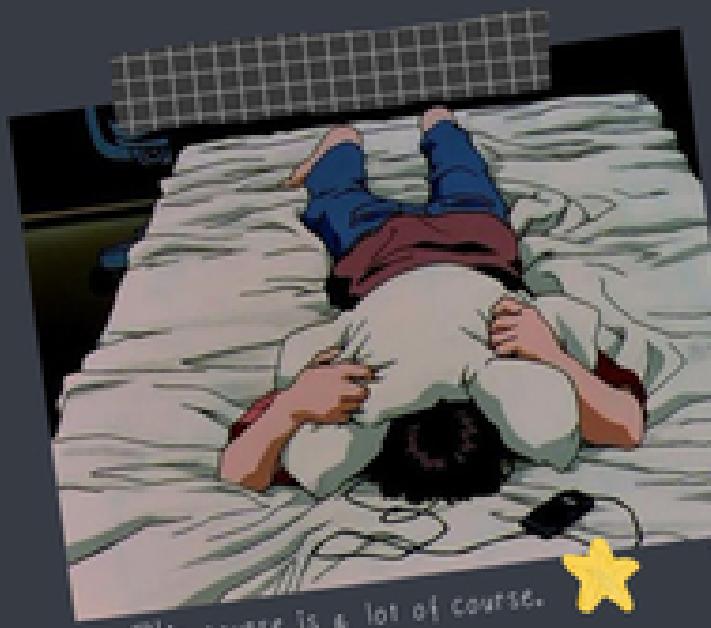


Cow Net



This course is a lot of course.

Term : 01

Fighting makes people



11 August 2020

Chapter 2 Application Layer Part 1

app architectures

2.1) Principles of network applications ↗ app requirements

2.2) Web and HTTP

2.3) FTP

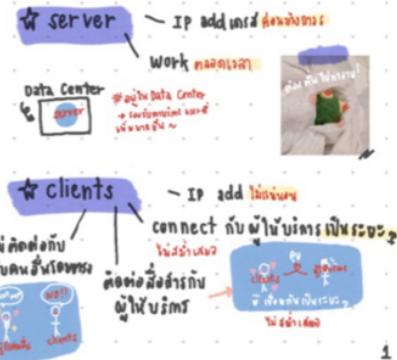
2.4) Electronic Mail ↗
 SMTP
 POP3

2.5) DNS

2.6) P2P applications

2.7) Socket programming with TCP

2.8) Socket programming with UDP



Some network apps ↗
 email, web
 remote login
etc. P2P file sharing } we call "app ภายนอก"

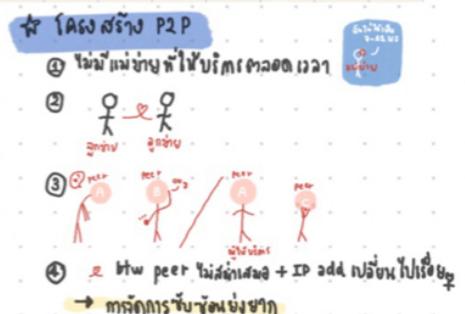
App ภายนอก layer ↗ layer ที่ take care Application that execute on that application layer
L take care end system

เต็มๆ

Creating a network app

2.1) Principles of network applications

Application architectures



(on all the time)

1. Client-server architectures (or) client model * always on host *

ஆகவே that we have 2 row of host 100% in side the network.

2. Peer-to-peer (P2P) ↗ what download ??

ஆக we have only 1 row in side the network * no always-on server *

3. Hybrid of client-server and P2P (on all the time)

• Skype : app ที่ working at instant message and also can call

is very popular 100% use hybrid system / connecting of client-server and P2P

• Instant messaging : chat btw 2 users is P2P

program running within a host

Processes communicating #Some processes it use for communication

Processes will be create a message and then message send over the network.
create by software , address, port

Sockets : the door of the processes.

when processes make the relation so the sockets will be create connecting
btw the application layer and the another layer ที่ต้องneed.

a lot of software we have the sockets served Ex. JAVA

Addressing processes (พิจารณาชื่อ)

App-layer protocol defines (พิจารณา)

What transport service does an app need?

/ ความต้องการของข้อมูล

1. Data loss : require some connecting that can lead to data loss ex. radio streaming

Ex music stop is likely data loss แต่ควรจะเป็น OK to you. ไม่ต้องการคุณกับเพื่อนแล้วเราต้องหาน้ำดื่ม

2. Timing : some app require low delay of data

Ex online game requires and ต้องการต่อเนื่องที่ต้องการต่อเนื่อง ± 2 s .. call "latency"

3. Throughput : to make sure that all from end-to-end connecting u have enough bandwidth to run
to all of data that u going to use in side the app

4. Security : very important that u have a make sure ว่า no one trap u data ไม่ถูก

Example แบบต่อไปนี้ (in slide)

very important in connecting service
(TCP)

Internet transport protocols services : have 2 types of word (TCP, UDP)

slow - TCP : ต้องใช้เวลาที่มากกว่า UDP

• connection-oriented ต้อง setup ต่อตัวกันก่อนที่จะเริ่มต้น

• reliable transport (ส่งข้อมูลที่ต้องการให้อ่านได้) btw. ผู้รับ - ผู้ส่ง

• flow control (ควบคุมการไหลของข้อมูล) ผู้ส่งไม่ส่งข้อมูลไปให้ผู้รับ

• congestion control (ควบคุมความคับคั่ง)

• does not provide (ส่งที่ไม่ได้แน่ ไม่เป็นไปได้) หมายความว่า

+ ความปะทะกันของข้อมูล

speed

↑ ต้องตอบสนองเร็วๆ กัน

UDP : ต้อง service ที่ can tolerate with
the data loss

• use AODV หรือ กันตัวเอง

faster than TCP

ไม่ต้องใช้ต่อระบบ TCP สักอย่าง เอาง่ายๆ ก็ได้
มีตัวต่อ เร็วๆ แต่ตัว อ่าน บันทึก ช้าๆ เช่น เฟรชช์ + วัสดุงาน
ชั้นนำ อาจต้องใช้เวลา

Example 例

2.2) Web and HTTP

Web and HTTP

web page คือ หน้าจอ file HTML มาก กันๆ

is web page is consists of objects

• the base HTML-file

• each object will be addressable by a URL

ที่อยู่ของ แต่ละ object ก็จะมี รอง

can be HTML file

can be JPEG image, etc.

or Java applet, audio file

HTTP ใช้ในการติดต่อเว็บไซต์กับเว็บเบราว์เซอร์ที่อยู่ในเครือข่าย WWW Server (World Wide web) โดยเอกสารจะอยู่ในรูปแบบที่เรียกว่า HTML และสำหรับลูกค้าที่ต้องการเข้าถึงเอกสารนั้น ต้องผ่านผู้ให้บริการอินเทอร์เน็ต อาจจะมีบุคคลที่รับผิดชอบในการจัดทำเอกสารให้กับผู้ให้บริการ ซึ่งผู้ให้บริการอินเทอร์เน็ตจะมีชื่อและที่อยู่ทางไปรษณีย์ รวมทั้งบุคคลที่รับผิดชอบในการจัดทำเอกสารให้กับผู้ให้บริการ

HTTP overview

: HTTP (Hypertext transfer protocol) is server/clients model

• so ในที่ HTML file all of web page obj will be อยู่บน web server

* Apache Web server =

- USES TCP = make รีเควส์ / server open the connecting send it over and then close
L uses port 80 # ต่อไปนี้
- HTTP is "stateless" (พังงาน)

connect ต่อ

HTTP connections

: have 2 kind of sending HTTP requests

การต่อไปต่อมาคือ

- Nonpersistent HTTP : At most one object is sent over a TCP connection.

L You hold the connecting and then send it until finish.

- Persistent HTTP : Multiple objects can be sent over single TCP connection btw client and server.
การเพื่อนร่วมเดินทาง

Nonpersistent HTTP

: at first (1a) HTTP initiates the connection btw host & server

(1b) server will be "accepts", and send accept notifying client

(2) Client receives the accept notifying from server start making a request message to contain URL

(3) Server receives the request message forms response message and then send over

(4) HTTP server closes TCP connection.

(5) HTTP client receives response message contain html file, displays html

(6) steps 1-5 repeated for each of 10 jpeg objects

when u not finish receiving the first object
u can not be receive to start receiving second
object # u should go one by one ต่อๆ

so it mean u have to do something

like this if 10 obj in this webpage

so it mean u have to repeat this one

10 time with nonpersistent

เวลาท่องกลับ
Non-Persistent HTTP : Response time is the same thing (33.5a)

- RTT = the round trip of request ≈ ④ ใช้ในการเดินทางของ package เล็กๆ จากลูกค้าไปยัง server และกลับ

Persistent HTTP

take slow

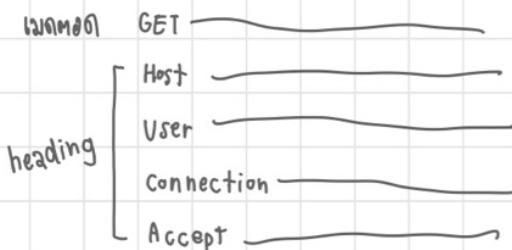
- Nonpersistent HTTP issues : have some problem because requires 2 RTTs per obj
- Persistent HTTP : leaves one connection open and server send ข้อมูลเพื่อ obj in one connection **# this is different**
so they use lit RTT and they return it more faster

↑ RTT 1 用来送所有 obj all

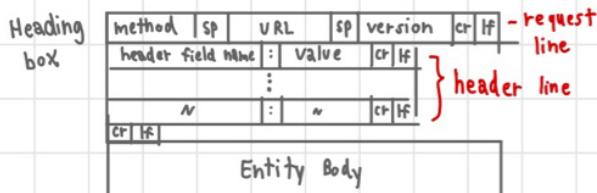
HTTP request message : 2 types of HTTP messages < request (ขอร้อง) response (ตอบกลับ)

• HTTP request message

- ASCII (human-readable format)



- N : general format



common this one
≈ ให้ดูได้ชัดเจน / clean กว่า

ส่งข้อมูลไปยัง server

Uploading from input

ผ่าน web ผ่านจากนี้มาในตัวอัตโนมัติ

/ หรือ input ที่ป้อนลงมาใน body

- Post method : use for upload data frome ภายนอก

- URL method : use get method (for URL request)

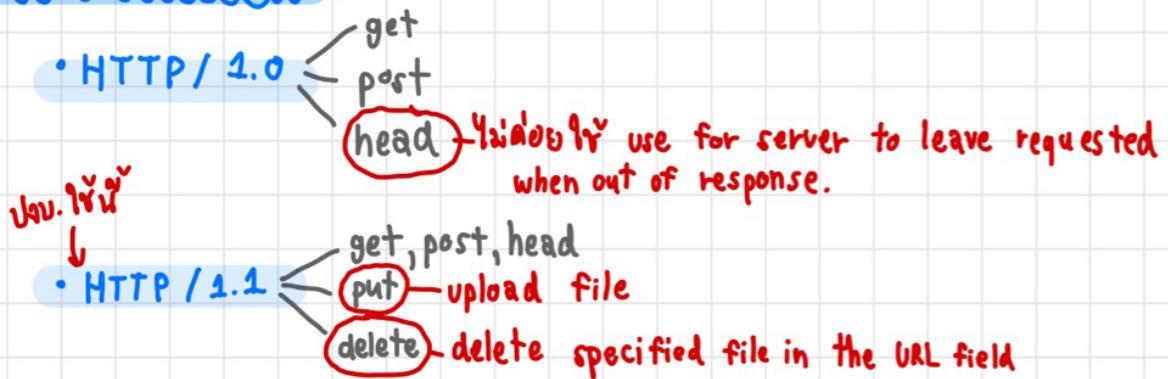
Ex) www.somesite.com/animalsearch?monkeys+banana

ข้อมูล input จะถูกต่อตัวกันเป็น field URL อยู่ที่ header

จะ request



Method type : different method



HTTP response message

HTTP / 1.1 200 OK
Date connection close
Content-Type Data :
Server :
Last-Modified :
Content-Length :
Content-Type :
HTML data data data data ...

* 200 OK = ສໍາເລັດ / ກຸກອ່າງ OK

HTTP response status codes

1. 200 = Ok // succeeded
2. 301 = Moved Permanently // obj moved new location
3. 400 = Bad Request // message not understand
4. 404 = Not found // file not found
5. 505 = HTTP Version Not Supported

// ລັບຂ່າຍ version ຂອງ HTTP ທີ່ server ບໍ່ໄດ້ support

Trying out HTTP (client side) for yourself

1. Telnet
- 2.
- 3.

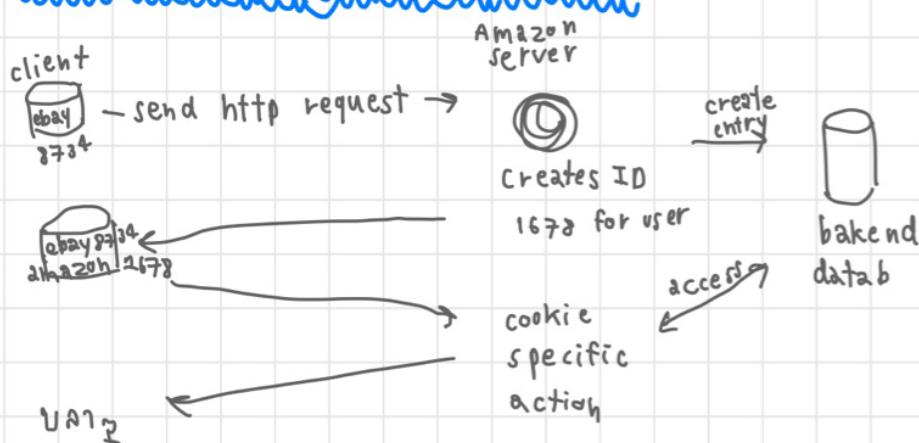
☞ ລົງທຶນດ້ວຍ
ຕັ້ງແຕ່ງ (P. 33)

User - server state : cookies

one of the important part of the website ຕີ່ຈະໄດ້ນຳໃຫຍ່
that website write down in to your machine ຈະ
ຄອນເຖິງ u person ດອລ information ex. login ກ່ອນ
ສ່ວນ cookie ມາເວລາເຂົ້າໃນມັກຈະໄດ້ໄໝ່ເຫັນ login ໃນນີ້ກ່ອງໄປແລ້ວ
read cookie ທີ່ນີ້ສ່ວນວ່າ (ປັບ???)

website ຂົນນີ້ໄດ້ໃຫຍ່ cookies

Cookies: keeping "state" (cont.)



cookies & privacy

- กรณี web รู้ชื่อของคุณ
- อย่างเช่น e-mail ฟ้า-หมาลูก

Cookies (cont.)

• What cookies can bring

information, authorization

shopping carts

recommendations

user session state (Web e-mail)

ເຫດຜົນເລີດ/ບໍລິສັດ

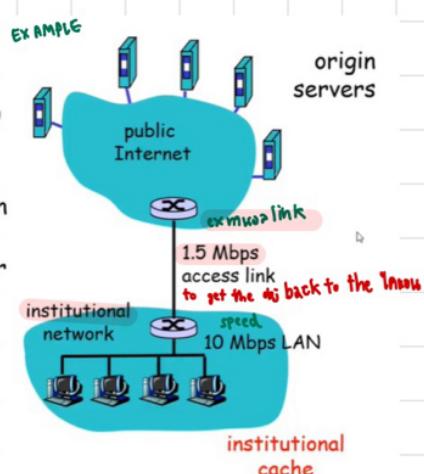
ໃນການ

Chapter 2 Application Layer Part 2

Web caches (proxy server) : working as the middle man (it in the middle) #at slide u see

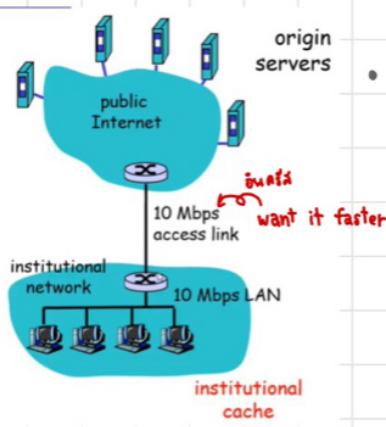
① Proxy server have information of the query it wait for the second or for another request that going to the same query so u don't have to go out that very single time that u request

More about web caching : Why we have to caching because reduce the time request
: reduce traffic

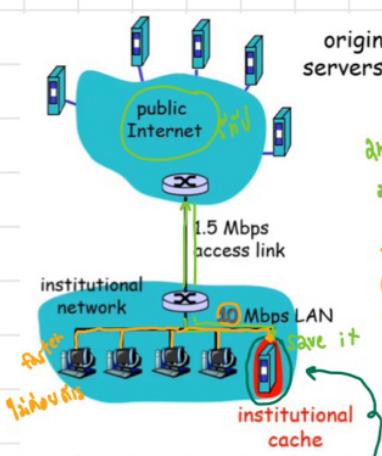


Caching example

- assumptions
- consequences



- possible solution consequences



and find what every they want
and save it in here

so every client asking the same
extending the content it always
go here

but if we have the proxy server

so it mean the proxy server will go out

Conditional GET

: GET method send to the proxy server
we have 2 situations

① When they get the request and the server find out
object isn't modified.

so it mean proxy know web content → ~~update~~ or update
or update you in update
it can over Adress to the client

but if it hot update proxy so go out and find new information and
save it to the proxy again.

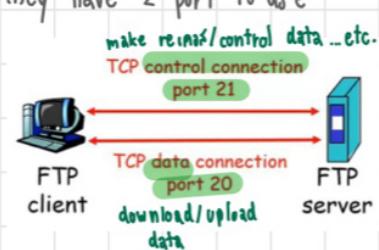
2.3) FTP is another app is very popular and useful in the network

FTP : the file transfer protocol

: is protocol app that running on this protocol it working for
upload file from client in to the server or download (up and down)

#ftp server : port 21

FTP: Separate control, data connections



: use TCP connection

FTP commands, responses

Sample commands:

- sent as ASCII text over control channel
- **USER** *username*
- **PASS** *password*
- **LIST** return list of files in current directory
- **RETR** *filename* retrieves (gets) file
- **STOR** *filename* stores (puts) file onto remote host

Sample return codes

- status code and phrase (as in HTTP)
- 331 Username OK, password required
- 125 data connection already open; transfer starting
- 425 Can't open data connection
- 452 Error writing file

: FTP use (ASCII text commands) = normal text commands

Ex. **USER** *username*, **PASS** *password*, **LIST** return list of files

: so when the client get return from the server

Ex. 331 Username OK, 425 / 452 is error message

we use at the time that we want to upload the big chunk of the file like a website etc.

2.4) Electronic Mail

- SMTP, POP3, IMAP

Electronic Mail (E-mail) : use **SMTP** protocol to send the e-mail

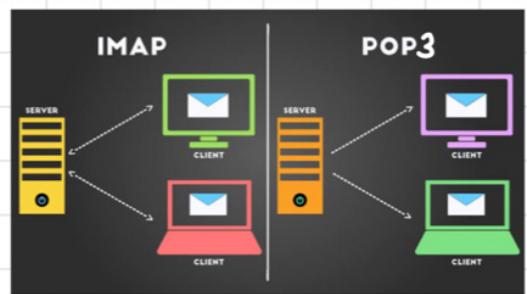
user agent =

agent =

we use agent \rightarrow e-mail and send to the server, mail server \rightarrow it will u email and then make relaid to the network to another mail server.

Mail Server

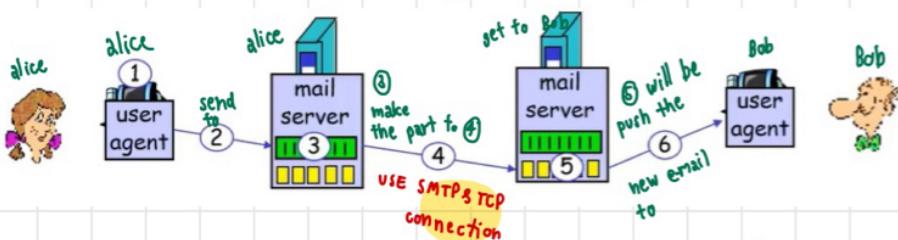
- mailbox
- message queue
- SMTP protocol



Electronic Mail: SMTP [RFC 2821]

- * use the ASCII text commands (same as MTP)
- * We have 3 phase of transfer
 - hands shaking
 - transfer of messages
 - closure
- * messages must be in 7-bit ASCII (is well ☺)
With 8-bit

Scenario: Alice sends message to Bob



Sample SMTP interaction & Ex. : start with hardchacking process

```
S: 220 hamburger.edu
C: HELO crepes.fr
S: 250 Hello crepes.fr, pleased to meet you
C: MAIL FROM: <alice@crepes.fr>
S: 250 alice@crepes.fr... Sender ok
C: RCPT TO: <bob@hamburger.edu>
S: 250 bob@hamburger.edu ... Recipient ok
C: DATA
S: 354 Enter mail, end with "." on a line by itself
C: Do you like ketchup?
C: How about pickles?
C: .
S: 250 Message accepted for delivery
C: QUIT
S: 221 hamburger.edu closing connection
```

— hello } hardchacking process
— hello back that the hamburger.edu is existinw

and then they making a relaid

— finish every thing

Try SMTP interaction for yourself

- * telnet servername 25
- * see 220 reply from server
- * enter HELO, MAIL FROM, RCPT TO, DATA, QUIT commands
above lets you send email without using email client
(reader)

ເລືອງດ້າຍຕົວເວັນ

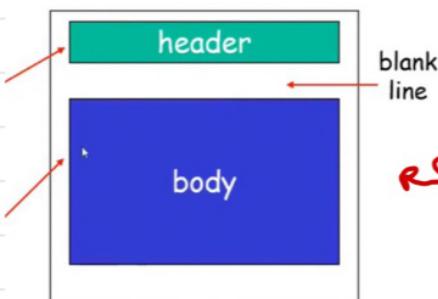
SMTP : final words

- SMTP uses persistent connections
- SMTP requires message (header & body) to be in 7-bit ASCII
- SMTP server uses CRLF, CRLF to determine end of message

Comparison with HTTP:

- HTTP: pull
- SMTP: push
- both have ASCII command/response interaction, status codes
- HTTP: each object encapsulated in its own response msg
- SMTP: multiple objects sent in multipart msg

Mail message format



re it same as HTTP

Message format: multimedia extensions

is the extension additional for e-mail to have more control

- * MIME : multimedia mail extension

and more info

Mail access protocols

is we have

- * POP : Post Office Protocol [RFC 1939]

User connect to mail server to download the e-mail back to the user

- * IMAP : Internet Mail Access Protocol [RFC 1730]

L it is another extenšun for ... they have the detail info

- * HTTP : gmail, Hotmail, Yahoo! Mail, etc. - free email

L it very popular that we use HTTP to working at the mail agent

so u don't have to run any mailagent to u pc @ e-mail

u can login g-mail, hotmail etc. use mail agent to connect with the mail server.

(more complex)
More features

manipulation of stored
msgs on server

POP3 protocol is protocol in making the request from the mail client to the mail server to receive mail from server and then end the list and end email that can receive.

POP3 (more) and IMAP

More about POP3

- Previous example uses "download and delete" mode.
- Bob cannot re-read e-mail if he changes client
- "Download-and-keep": copies of messages on different clients
- POP3 is stateless across sessions

IMAP

- Keep all messages in one place: the server
- Allows user to organize messages in folders
- IMAP keeps user state across sessions:
 - ❖ names of folders and mappings between message IDs and folder name

- can also folder, move, reorganize control by IMAP

Chapter 2 Application Layer Part 3 [End]

2.5) DNS

DNS: Domain Name System

it's like a name system for network ex. IP address

uniquely

- * DNS is when people very hard to remember only number so they invent mechanism that we ... name the num ...
- * What the DNS system is having distributed database so it mean we have database and we have many database working together

People: many identifiers:

- ❖ SSN, name, passport #

Internet hosts, routers:

- ❖ IP address (32 bit) - ^{host} used for addressing datagrams
- ❖ "name", e.g., ^{dns} www.yahoo.com - used by humans

Q: map between IP addresses and name ?

Domain Name System:

- **distributed database** implemented in hierarchy of many **name servers**
- **application-layer protocol** host, routers, name servers to communicate to **resolve** names (address/name translation)
 - ❖ note: core Internet function, implemented as application-layer protocol
 - ❖ complexity at network's "edge"

DNS

DNS services

- hostname to IP address translation
- host aliasing
 - ↳ Canonical, alias names
- mail server aliasing
- load distribution
 - ↳ replicated Web servers: set of IP addresses for one canonical name

Why not centralize DNS?

- single point of failure
- traffic volume
- distant centralized database
- maintenance

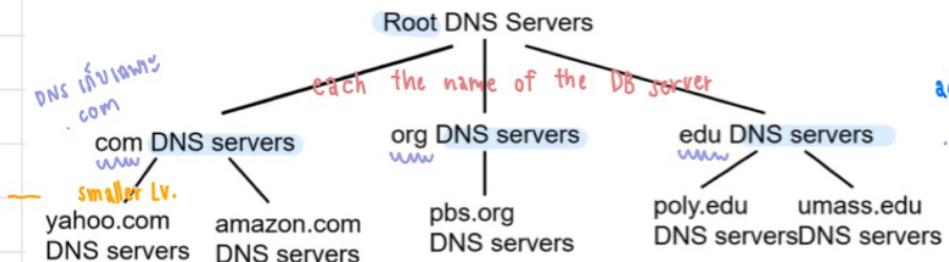
doesn't scale!

if we set the DNS server at a center it means when the server down all of the DNS server will be down

using many database would be better

it good if we have database sit in each country and they go step-to-step

Distributed, Hierarchical Database



Client wants IP for www.amazon.com; 1st approx:

- client queries a root server to find com DNS server
- client queries com DNS server to get amazon.com DNS server
- client queries amazon.com DNS server to get IP address for www.amazon.com

DNS : Root name servers

* it the biggest tree of the DNS tree

* it have many servers in the world

for example:



TLD and Authoritative Servers

- **Top-level domain (TLD) servers:** *edu, .com, .org, .net, .edu, .etc, and all top-level country domains .uk, .fr, .ca, .jp.*
- ❖ responsible for com, org, net, edu, etc, and all top-level country domains uk, fr, ca, jp.
- ❖ Network Solutions maintains servers for com TLD
- ❖ Educause for edu TLD
- **Authoritative DNS servers:** *(Low-Lv.)*
- ❖ organization's DNS servers, providing authoritative hostname to IP mappings for organization's servers (e.g., Web, mail).
- ❖ can be maintained by organization or service provider

* TLD = Top-LV.-domain

.edu .com .org .net .etc

comprovisos

/

.org

it is all maintained under .com or .net ... etc.

a big mode

ais @ anything that we provide internet connecting to client



Local Name Server

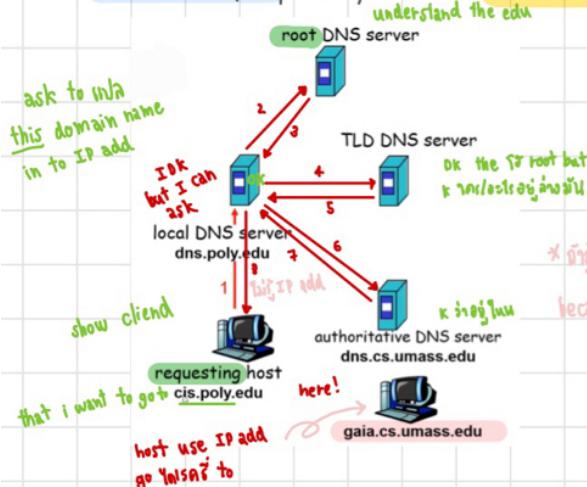
* so each host have to make the DNS query
so when every single time that the client inside the network want to connect to somewhere they have to send request to DNS server and DNS server send out the DNS query to upper layer of the DNS server

DNS name resolution example

* we have DNS transaction in 2 way

① iterative query : 1 way that we make the DNS query to get the answer

understand the edu



* you only name it mean isn't in the network

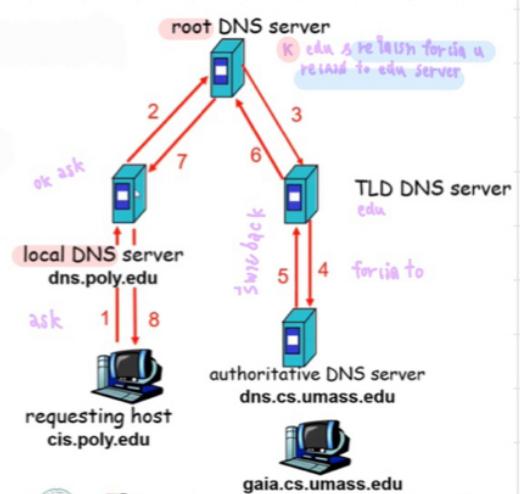
(will show in window)

because the network only is know the IP add

② recursive query : this's another way to refind the DNS transaction

* Q & A : it more heavy load more than the last one or not?

: yes, it more heavy load.

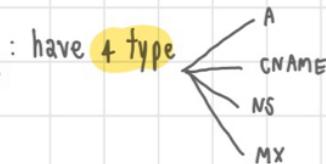


DNS: caching and updating records

- once (any) name server learns mapping, it **caches** mapping
 - ❖ cache entries timeout (disappear) after some time
 - ❖ TLD servers typically cached in local name servers
 - Thus root name servers not often visited
- update/notify mechanisms under design by IETF
 - ❖ RFC 2136
 - ❖ <http://www.ietf.org/html.charters/dnsind-charter.html>

សំនួរពីខ្លួន : Aj. ខ្លួន Aj តាមអាជីវកម្មបានបង្កើតបាន
ដូច Aj ធ្វើវាត្រូវបានគិតជាបាន

DNS records



it look like RR format: (name, value, type, ttl)

* Type A -> name is hostname
value is IP address

* Type NS -> name is domain name
value is hostname

that it under the this
domain name

sub domain / child domain under main domain

when u want to alias the wrong
name in to short name.

* Type CNAME -> name is alias name
value is canonical name

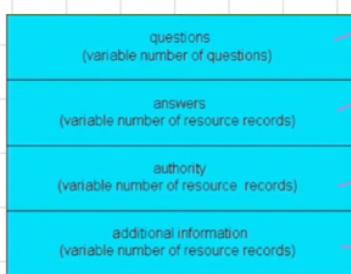
* Type MX - value is name of mailserver

DNS protocol, messages

= in each messages of the DNS it will look like this when it send out to get the transaction

msg header

- identification : 16 bit #
- flags :
 - query or reply
 - recursion desired
 - recursion available
 - reply is authoritative



Name, type fields for a query

RRs in response to query

records authoritative

helpful info

header →

data itself @
the msg is a

| 16 bit | have identifier |
|-------------------------|---|
| identification | flags |
| number of questions | number of answer RRs |
| number of authority RRs | number of additional RRs |
| body | questions (variable number of questions) |
| body | answers (variable number of resource records) |
| body | authority (variable number of resource records) |
| body | additional information (variable number of resource records) |

12 bytes

Inserting records into DNS

= u register new name with new IP add

the name of the domain name it have to be unique: have to be a new one not the same name

must be unique with the IP add — have to be unique IP add
unique

2.6) P2P applications

Pure P2P architecture

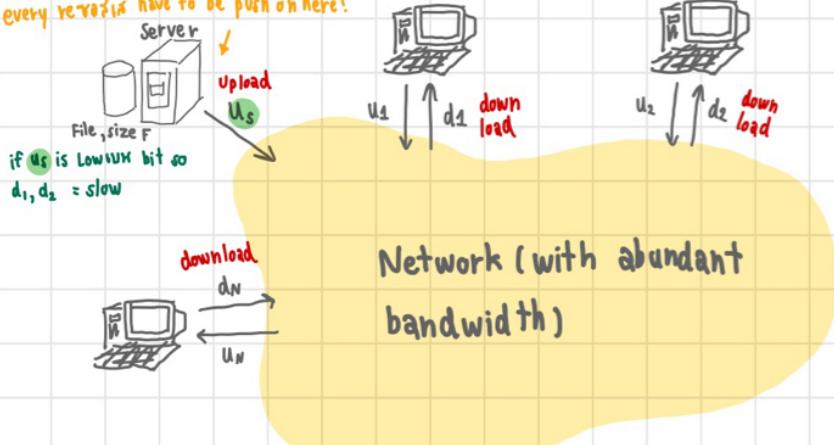
= is the ~~functionality~~ that we have a row in the network

Client connects to Client. So not always on server inside the network
and end system connects to end system

- * 3 topics
 - File distribution
 - Searching for information
 - Case study : skype (~~even using~~)

File Distribution : Server-Client vs P2P

every request have to be pushed on here!



not limit U_S with someone else

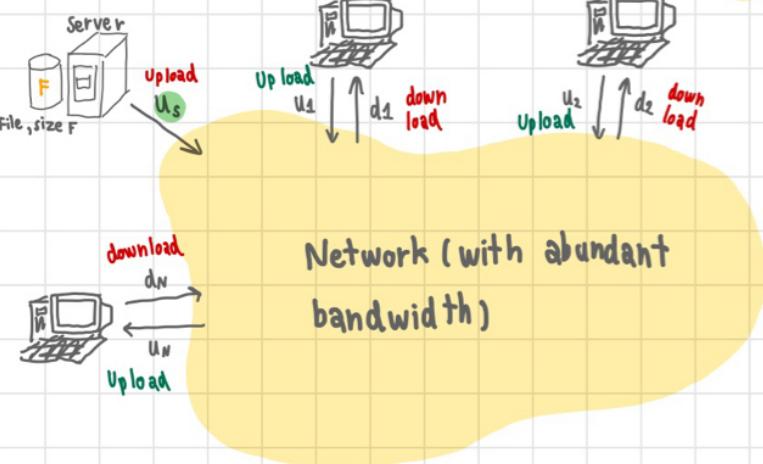
- * U_S : server upload bandwidth
- * U_i : peer i upload
- * d_i : download

to make all names' download fast

File distribution time : server-client

: many ways of upload but download only U_S

(U₁, U₂, U_N)



// Working Part calculate

File distribution time : P2P

= we have upload & download equally.

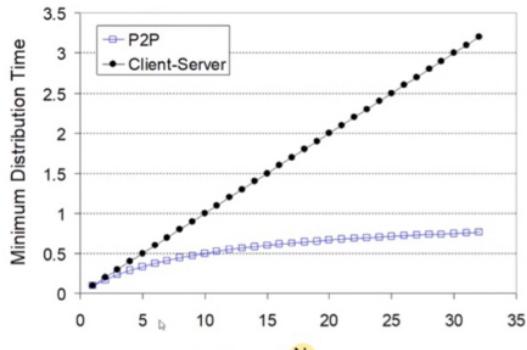
U₁, U₂, U_N
* helping upload for each other so when download d₁ may be connect to U_{1,2} or b₂ connect to d_n or d_{4,5} so it
is good distribution to help database calm out from the network so it faster because the U_S + Σ U_i

must than this one !!

Server-client vs. P2P : example

Client upload rate = u , $F/u = 1$ hour, $u_s = 10u$, $d_{\min} \geq u_s$

(worst/good)



is item for number of client that connect → that make request for download

: Client-server take more time when n increase

if N increase the time will be more

more client = more upload reqs

because when client get info that download it. They

also upload what they have in to the network

not just download.

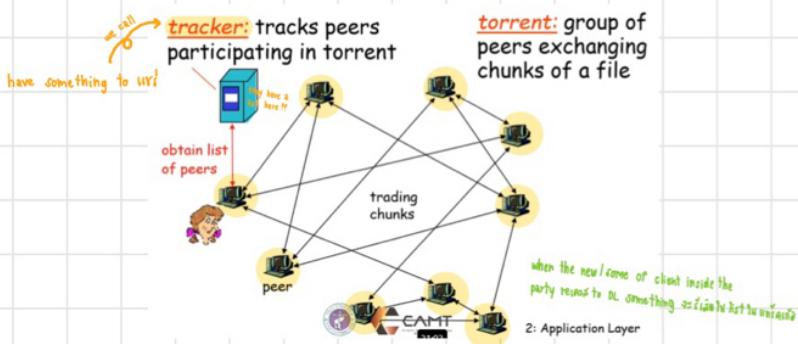
* When we have the file info, it take less the same time because client don't have the file but when they have more client it will be faster because each client also help upload the file info that they already have to another one who come later as well.

more client = more file download ← must better than

this ?

File distribution : BitTorrent

* at first bitTorrent have to know who has the file so it mean we have to create a list of file / info in each client that who has it and how many what they have



BitTorrent (1)

: ภูมิคุณในเงินล่วงหน้าที่เรียกว่า chunks

size : 256 KB

* so if ① ask ② for a chunks ③...

if we have 1000 chunks as a one file info ... person ② have chunks can provide info to the 1 who get DL

bit ② have id for chunks of data

so to ④ first

1) if answer
j) ④ go to another
info ④
info ④

so ④

they don't have a file info

so when the person DL in the chunks in the same G

if another person ask to the chunks this person already

have → it will be send to the chunks that the person here already have to the new guy as well.

in the same time ④

BitTorrent (2) : we have

- * Pulling Chunks (we call a "DL chunks")
- * Sending Chunks

P2P: searching for information : we have something each

- * file sharing (like e-mule system)

when client runs e-mule PG initiating file it want to use file upload folder and e-mule get a list to another p that can see u UL folder info and then they can get from u UL folder — very easy system to use

run winfile kinda in bittorrent system run on e-mule

* Instant messaging : it also the serio of P2P

L it working in the same thing since it's server working at the index ex. when u want to go with someone the server tell IP add of that friend member that u want to go and then u just go P2P to them.

P2P: centralized index

{ (พัฒนา)

P2P: problems with centralized directory

2.7) Socket programming with TCP

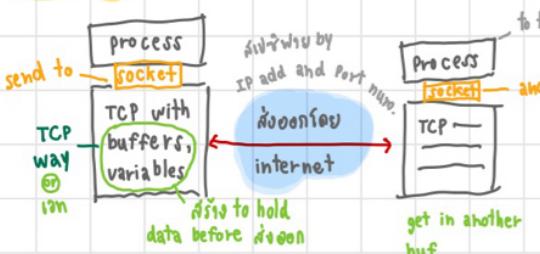
Socket programming is the PG we do to send / send a message / something info over the network

အောက်ပါတဲ့ we need socket API → to use with

L use java API or by UNIX (socket)

Socket-programming using TCP

* Socket = door btw app and under need a protocol . We have 2 kind serial



TCP
TCP
} transmission

TCP > connection establish
cu have to retrans connect / hold connection until the end of transaction

— UCP = u don't have retrans anything
and send out it side away on the network

Socket programming with TCP

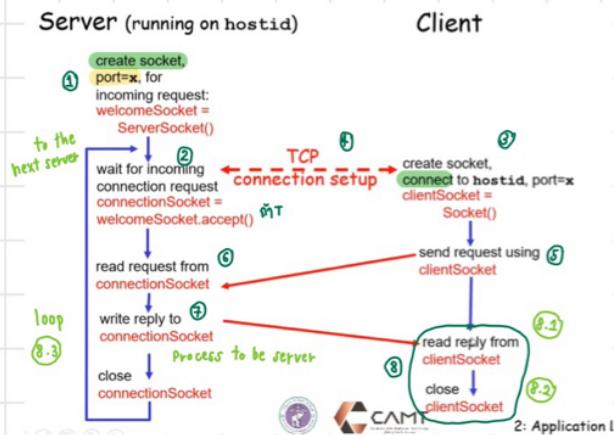
* Client must contact server (it work on client server model & server model also on running), run socket that welcomes client contact

* Client contacts server by :

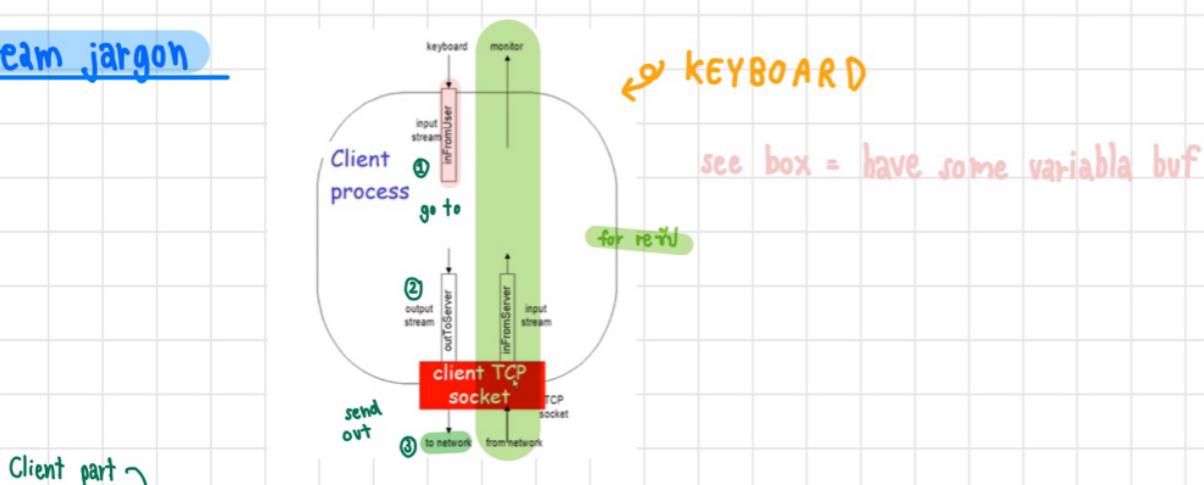
- have to do contact to the server by own client TCP socket & own server socket it must be the same port (with specifying IP address, port num have to be the same)
- When client own socket → establishes connection when server return confirm

* start សំវារណ៍នៃ server និង then close the connection

Client/server socket interaction : TCP



Stream jargon



Example : Java client (TCP)

JAVA server (TCP)

2.8 Socket programming with UDP

Socket programming with UDP

: ត្រូវយក TCP ឡើងទៅក្នុង connection btw server & client

: UDP no connection btw client & server

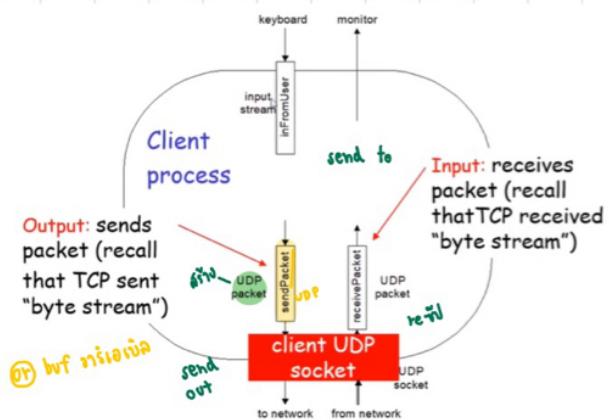
: no handshaking

Server (running on hostid)

Client



Example: Java client (UDP)



Example: Java client (UDP) : code źródłowy na slajdzie

Java server (UDP)

Chapter 2: Summary

* our study of network app now complete! *

- application architectures
 - ❖ client-server
 - ❖ P2P
 - ❖ hybrid
- application service requirements:
 - ❖ reliability, bandwidth, delay
- Internet transport service model
 - ❖ connection-oriented, reliable: TCP
 - ❖ unreliable, datagram: UDP

- specific protocols:
 - ❖ HTTP
 - ❖ FTP
 - ❖ SMTP, POP, IMAP
 - ❖ DNS
 - ❖ P2P: BitTorrent, Skype
- socket programming

Most importantly: learned about protocols

- typical request/reply message exchange:
 - ❖ client requests info or service
 - ❖ server responds with data, status code
- message formats:
 - ❖ headers: fields giving info about data
 - ❖ data: info being communicated

Important themes:

- control vs. data msgs
 - ❖ in-band, out-of-band
- centralized vs. decentralized
- stateless vs. stateful
- reliable vs. unreliable msg transfer
- "complexity at network edge"

Chapter 3 Transport Layer Part 1

2.1) Principles of network applications

logical

L differ

is mean end to end to ser

fisic

→ to software

CHAPTER 3 : p1

* segments

* transport layer = comm in the real life.

cave end-end

* network layer = how to inv way to send?

split it out

↑

Demultiplexing

* inv on m... control by _____

VDP = ~~data~~ ~~information~~ ~~control~~ check

↓
need ~~for~~

some kind

SP = ~~sequence~~ = PN
DP = ~~data~~ = PDU.

read before connect = no handshaking.

* flipped bits = bits have 0,1 but it ref 1

data now

sum num.

* check sum

| | | |
|----|----|----|
| 17 | 20 | 34 |
| 5 | 12 | 11 |
| 6 | | 34 |

+ all = 34 → ~~error~~ -

if mod data = 0 = ~~no error~~

not 0 = have error in data