

CSC3200 Final Review

Data Structures and Advanced Programming

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① Introduction

② C++ Basics

③ Data Structures

④ Application

⑤ Q & A

Self-Introduction



Information

- **Name** Chaoyi Sun
- **Year** Second-year Undergraduate
- **Major** Computer Science

Achievements

- **Silver Medal**, 50th ICPC Asia Regional Contest (Wuhan)
- **Silver Medal**, 50th ICPC Asia Regional Contest (Nanjing)
- **Bronze Medal**, 11th CCPC National Contest (Chongqing)

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C++ Keywords

alignas	alignof	and	and_eq	asm	auto
bitand	bitor	bool	break	case	catch
char	char8_t	char16_t	char32_t	class	compl
concept	const	constexpr	constexpr	constinit	const_cast
continue	co_await	co_return	co_yield	decltype	default
delete	do	double	dynamic_cast	else	enum
explicit	export	extern	false	float	for
friend	goto	if	inline	int	long
mutable	namespace	new	noexcept	not	not_eq
nullptr	operator	or	or_eq	private	protected
public	register	reinterpret_cast	requires	return	short
signed	sizeof	static	static_assert	static_cast	struct
switch	template	this	thread_local	throw	true
try	typedef	typeid	typename	union	unsigned
using	virtual	void	volatile	wchar_t	while
xor	xor_eq				

Operator Precedence

Prec.	Operator	Assoc.
1	() [] -> . :: ++ -- (postfix)	L→R
2	! ~ ++ -- (prefix) - + * &	R→L
3	->* .*	L→R
4	* / %	L→R
5	+ -	L→R
6	<< >>	L→R
7	< <= > >=	L→R
8	== !=	L→R
9	&	L→R
10	~	L→R
11		L→R
12	&&	L→R
13		L→R
14	? :	R→L
15	= += -= *= /= %= &= ^= = <= >=	R→L
16	,	L→R

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Memory

Memory Region	Characteristics
Stack	Automatic stack frames created for each method call to hold local variables
Heap	Dynamic memory pool for manual allocation and lifetime control.
Dynamic Allocation	[manual] malloc/free
Global/Static	[Auto] Global variables/constants that persist throughout the lifetime of the program.
Constant	[Auto] constants/read-only

Example

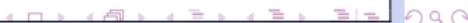
```
1 void f1 ()  
2 {  
3     int *p = new int [5];  
4     int variable = 10;  
5 }
```

```
1 0000000000001149 <_Z2f1v>:a  
2     1149:   f3 0f 1e fa          endbr64  
3     114d:   55                push    %rbp  
4     114e:   48 89 e5          mov     %rsp,%rbp  
5     1151:   48 83 ec 10          sub    $0x10,%rsp  
6     1155:   bf 14 00 00 00          mov    $0x14,%edi  
7     115a:   e8 f1 fe ff ff          call   1050 <_Znam@plt>  
8     115f:   48 89 45 f8          mov    %rax,-0x8(%rbp)  
9     1163:   c7 45 f4 0a 00 00 00          movl   $0xa,-0xc(%rbp)  
10    116a:   90                nop  
11    116b:   c9                leave  
12    116c:   c3                ret
```

Example (Cont.)

```
1 void f1 ()  
2 {  
3     int *p = new int[5];  
4     delete [] p;  
5 }
```

```
1 0000000000001149 <_Z2f1v>:  
2 1169:   f3 0f 1e fa          endbr64  
3 116d:   55                  push    %rbp  
4 116e:   48 89 e5            mov     %rsp,%rbp  
5 1171:   48 83 ec 10          sub    $0x10,%rsp  
6 1175:   bf 14 00 00 00        mov    $0x14,%edi  
7 117a:   e8 e1 fe ff ff      call   1060 <_Znam@plt>  
8 117f:   48 89 45 f8          mov    %rax,-0x8(%rbp)  
9 1183:   48 83 7d f8 00        cmpq   $0x0,-0x8(%rbp)  
10 1188:  74 0c                je     1196 <_Z2f1v+0x2d>  
11 118a:   48 8b 45 f8          mov    -0x8(%rbp),%rax  
12 118e:   48 89 c7            mov    %rax,%rdi  
13 1191:   e8 da fe ff ff      call   1070 <_ZdaPv@plt>  
14 1196:   90                  nop  
15 1197:   c9                  leave  
16 1198:   c3                  ret
```



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Pointer / Reference

```
1 int vari = 3100;
2 int *ptr = &vari; // pointer
3 int &ref = vari; // reference
4 *ptr = 3200;
```

Differences

- **Pointer:** Can be reassigned, can be null, has separate memory.
- **Reference:** Must be initialized, cannot be reassigned, alias to existing variable.

Pointer / Reference

```
1 int vari = 3100;
2 int *ptr = &vari; // pointer
3 int &ref = vari; // reference
4 *ptr = 3200;
```

Differences

- **Pointer:** Can be reassigned, can be null, has separate memory.
- **Reference:** Must be initialized, cannot be reassigned, alias to existing variable.

The meaning of *

- `int *ptr = &vari;` Type specifier in declarations (store address).
- `*ptr = 3200;` Dereference operator in expressions.

Multi-level Pointer

```
1 #include <cstdio>
2 void calc (int a,int b,int **f1,int **f2)
3 {
4     **f1 = a + b;**f2 = a - b;
5     int *temp = *f1;*f1 = *f2;*f2 = temp;
6 }
7 int main ()
8 {
9     int a = 5,b = 3,x,y;
10    int *px = &x,*py = &y;
11    calc (a,b,&px,&py);
12    printf ("%d %d\n",x,y);
13    return 0;
14 }
```

Multi-level Pointer

```
1 #include <cstdio>
2 void calc (int a,int b,int **f1,int **f2)
3 {
4     **f1 = a + b;**f2 = a - b;
5     int *temp = *f1;*f1 = *f2;*f2 = temp;
6 }
7 int main ()
8 {
9     int a = 5,b = 3,x,y;
10    int *px = &x,*py = &y;
11    calc (a,b,&px,&py);
12    printf ("%d %d\n",x,y);
13    return 0;
14 }
```

Answer

1 8 2

Swap

```
1 #include <cstdio>
2 void swap1 (int *x,int *y) {int tmp = *x;*x = *y;*y = tmp;}
3 void swap2 (int *x,int *y) {int *tmp = x;x = y;y = tmp;}
4 void swap3 (int **x,int **y) {int *tmp = *x;*x = *y;*y = tmp;}
5 void swap4 (int *&x,int *&y) {int *tmp = x;x = y;y = tmp;}
6 int main ()
7 {
8     int a = 5,b = 3;
9     swap1 (&a,&b);
10    printf ("%d%d\n",a,b);
11    a = 5,b = 3;
12    swap2 (&a,&b);
13    printf ("%d%d\n",a,b);
14    a = 5,b = 3;
15    int *pa = &a,*pb = &b;
16    swap3 (&pa,&pb);
17    printf ("%d%d%d%d\n",a,b,*pa,*pb);
18    a = 5,b = 3;
19    pa = &a;pb = &b;
20    swap4 (pa,pb);
21    printf ("%d%d%d%d\n",a,b,*pa,*pb);
22    return 0;
23 }
```



Swap (cont.)

Answer

- 1 3 5
- 2 5 3
- 3 5 3 3 5
- 4 5 3 3 5

Hint

Drawing memory diagrams can greatly assist in solving such problems!

Array

```
1 int arr[5] = {0,1,3,4,7};  
2 *(arr + 3) = 5;  
3 for (int i = 0;i < 5;++i) printf ("The %d-th: %d\n",i,*(arr + i));
```

```
1 The 0-th: 0  
2 The 1-th: 1  
3 The 2-th: 3  
4 The 3-th: 5  
5 The 4-th: 7
```

Array (Cont.)

```
1 #include <cstdio>
2 /* declaration here */
3 int main ()
4 {
5     char arr[2][3] = {{'a','b','c'},{'e','f','g'}};
6     calc (arr);
7     return 0;
8 }
```

Which of the following function declarations are correct?

- A) void calc (char (*arr)[3]);
- B) void calc (char arr[2][3]);
- C) void calc (char **arr);
- D) void calc (char arr[] [3]);

Array (Cont.)

```
1 #include <iostream>
2 /* declaration here */
3 int main ()
4 {
5     char arr[2][3] = {{'a','b','c'}, {'e','f','g'}};
6     calc (arr);
7     return 0;
8 }
```

Which of the following function declarations are correct?

- A) void calc (char (*arr)[3]);
- B) void calc (char arr[2][3]);
- C) void calc (char **arr);
- D) void calc (char arr[] [3]);

Answer

1 ABD

Pointer and Const

- **Const pointer** Pointer itself cannot be changed.

```
int* const p1;
```

- **Pointer to const** Pointed value cannot be changed.

```
int const *p1; or const int *p1;
```

- **Const pointer to const** Both cannot be changed.

```
const int* const p1;
```

const int const *p1; is wrong due to redundant const qualifier.

Increment/Decrement

```
1 #include <cstdio>
2 using namespace std;
3 int main ()
4 {
5     int a[] = {3,1,4,1,5,9};
6     int *p = a + 1,*q = a + 4;
7     int x = *++p,y = *q--;
8     int z = ++*p,w = (*q)--;
9     int t = (*p++) + (--q);
10    int u = *p,v = *q;
11    for (int i = 0;i < 6;++i) printf ("%d",a[i]);
12    printf ("\n%d%d%d%d%d%d\n",x,y,z,w,t,u,v);
13    return 0;
14 }
```

Increment/Decrement

```
1 #include <cstdio>
2 using namespace std;
3 int main ()
4 {
5     int a[] = {3,1,4,1,5,9};
6     int *p = a + 1,*q = a + 4;
7     int x = *++p,y = *q--;
8     int z = ++*p,w = (*q)--;
9     int t = (*p++) + (--q);
10    int u = *p,v = *q;
11    for (int i = 0;i < 6;++i) printf ("%d",a[i]);
12    printf ("\n%d%d%d%d%d%d\n",x,y,z,w,t,u,v);
13    return 0;
14 }
```

Answer

```
1 315059
2 45511005
```

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Vector

```
1 int main ()
2 {
3     vector <int> lst = {2,0,2,5,1,2,1,4};
4     int tot = 0,flag = 0,n = lst.size ();
5     for (auto it = lst.begin ();it != lst.end ();++it)
6     {
7         if (flag) tot += lst.front ();
8         else tot += *it;
9         flag = 1 - flag;
10    }
11    printf ("%d,%d\n",tot,tot / n);
12
13 }
```

Vector

```
1 int main ()
2 {
3     vector <int> lst = {2,0,2,5,1,2,1,4};
4     int tot = 0,flag = 0,n = lst.size ();
5     for (auto it = lst.begin ();it != lst.end ();++it)
6     {
7         if (flag) tot += lst.front ();
8         else tot += *it;
9         flag = 1 - flag;
10    }
11    printf ("%d,%d\n",tot,tot / n);
12
13 }
```

Answer

1 14,1

Vector (Cont.)

resize (), reserve ()

```
1 vector <int> lst = {2,0,2,5};  
2 lst.reserve (8);  
3 printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
4 lst.reserve (6);  
5 printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
6 lst.resize (10,3);  
7 printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
8 lst.resize (3,5);  
9 printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());
```

Vector (Cont.)

resize (), reserve ()

```
1 vector <int> lst = {2,0,2,5};  
2 lst.reserve (8);  
3 printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
4 lst.reserve (6);  
5 printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
6 lst.resize (10,3);  
7 printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());  
8 lst.resize (3,5);  
9 printf ("%ld,%ld,%d\n",lst.size (),lst.capacity (),lst.back ());
```

Answer

```
1 4,8,5  
2 4,8,5  
3 10,10,3  
4 3,10,2
```

Stack/Queue

```
1 #include <stack>
2 #include <cstdio>
3 class CSC
4 {
5     private:
6         std::stack <int> s1,s2;
7         void tran () {while (!s1.empty ()) s2.push (s1.top ()),s1.pop ();}
8     public:
9         CSC () = default;
10        void add (int val) {s1.push (val);}
11        int del ()
12        {
13            if (s2.empty ())
14            {
15                if (s1.empty ()) return -1;
16                tran ();
17            }
18            int val = s2.top ();s2.pop ();
19            return val;
20        }
21    };
```

Stack/Queue (Cont.)

```
1 int main ()
2 {
3     CSC arr;
4     for (int i = 1;i <= 5;++i) arr.add (i);
5     for (int i = 0;i <= 5;++i)
6     {
7         printf ("%d",arr.del ());
8         if (i != 5) printf (",");
9     }
10    printf ("\n");
11    return 0;
12 }
```

Stack/Queue (Cont.)

```
1 int main ()
2 {
3     CSC arr;
4     for (int i = 1;i <= 5;++i) arr.add (i);
5     for (int i = 0;i <= 5;++i)
6     {
7         printf ("%d",arr.del ());
8         if (i != 5) printf (",");
9     }
10    printf ("\n");
11    return 0;
12 }
```

Answer

1 1,2,3,4,5,-1

Stack/Queue (Cont.)

Sliding Window

```
1 void calc ()
2 {
3     int q[1005] = {0}, num[] = {1,2,7,3,8,5,2,9};
4     int k = 3, n = 8, head = 0, tail = -1;
5     for (int i = 0; i < n; ++i)
6     {
7         while (head <= tail && q[head] <= i - k) ++head;
8         while (head <= tail && num[q[tail]] <= num[i]) --tail;
9         q[++tail] = i;
10        if (i >= k - 1) printf ("%d", num[q[head]]);
11    }
12    printf ("\n");
13 }
```

Stack/Queue (Cont.)

Sliding Window

```
1 void calc ()
2 {
3     int q[1005] = {0}, num[] = {1,2,7,3,8,5,2,9};
4     int k = 3, n = 8, head = 0, tail = -1;
5     for (int i = 0; i < n; ++i)
6     {
7         while (head <= tail && q[head] <= i - k) ++head;
8         while (head <= tail && num[q[tail]] <= num[i]) --tail;
9         q[++tail] = i;
10        if (i >= k - 1) printf ("%d", num[q[head]]);
11    }
12    printf ("\n");
13 }
```

Answer (Hint: It also can be implemented with std::deque.)

1 778889

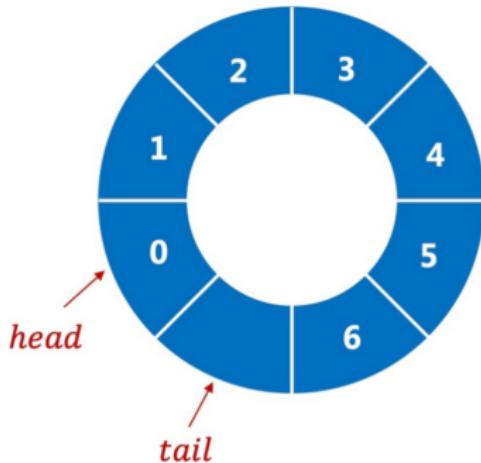


Stack/Queue (Cont.)

Circular Queue

Empty head == tail;

Full (tail + 1) % capacity == front;



Link list

Version 1

```
1 LinkNode* func (LinkNode* head)
2 {
3     LinkNode* pre = nullptr;
4     LinkNode* cur = head;
5     LinkNode* nxt = nullptr;
6     while (cur != nullptr)
7     {
8         nxt = cur -> next;
9         cur -> next = pre;
10        pre = cur; cur = nxt;
11    }
12    return pre;
13 }
```

Link list (Cont.)

Version 2

```
1 LinkNode* func (LinkNode* head)
2 {
3     if (head == nullptr || head -> next == nullptr) return head;
4     LinkNode* new_head = func (head -> next);
5     head -> next -> next = head;
6     head -> next = nullptr;
7     return new_head;
8 }
```

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Map

Hashmap

```
1 #define TABLE_SIZE 7
2 struct HashTable {int table[TABLE_SIZE];};
3 int hashFunction (int key) {return key % TABLE_SIZE;}
4 void insert(int table[],int key)
5 {
6     int id = hashFunction (key);
7     while (table[id] != -1) id = (id + 1) % TABLE_SIZE;
8     table[id] = key;
9 }
```

Map (Cont.)

1. The following keys are inserted in order: $\{7, 14, 21, 28, 35, 42\}$. After inserting all keys, what is the average successful search length (average number of comparisons to find an existing key)?

- A) 2.0
- B) 2.5
- C) 3.0
- D) 3.5

2. In the worst-case scenario for linear probing, inserting n elements into a table of size n requires?

- A) $O(n)$ time per insertion
- B) $O(\log n)$ time per insertion
- C) $O(n \log n)$ total time for all insertions
- D) $O(n^2)$ total time for all insertions

Map (Cont.)

1. The following keys are inserted in order: $\{7, 14, 21, 28, 35, 42\}$. After inserting all keys, what is the average successful search length (average number of comparisons to find an existing key)?
 - A) 2.0
 - B) 2.5
 - C) 3.0
 - D) 3.5
2. In the worst-case scenario for linear probing, inserting n elements into a table of size n requires?
 - A) $O(n)$ time per insertion
 - B) $O(\log n)$ time per insertion
 - C) $O(n \log n)$ total time for all insertions
 - D) $O(n^2)$ total time for all insertions

Answer

1

D ; AD



Set

Bitset

```
1 const int RANGE_SIZE = 128;
2 const int BITS_PER_BYTE = 8;
3 const int BITS_PER_LONG = BITS_PER_BYTE * sizeof(long);
4 const int CVEC_WORDS = (RANGE_SIZE + BITS_PER_LONG - 1) / BITS_PER_LONG;
5 struct BitSet
6 {
7     unsigned long words[CVEC_WORDS];
8     BitSet ()
9     {
10         for (int i = 0;i < CVEC_WORDS;++i) words[i] = 0;
11     }
12 };
13 unsigned long createMask (int k)
14 {
15     unsigned long x = 1L;
16     return x << k % BITS_PER_LONG;
17 }
```

Set (Cont.)

```
1 // check if the k-th characters in ASCII table in the Set
2 bool inSet(BitSet &cv, int k)
3 {
4     if (k < 0 || k >= RANGE_SIZE) return;
5     return cv.words[k / BITS_PER_LONG] & createMask (k);
6 }
7 // set the k-th character in ASCII table to the set
8 void setBit(BitSet &cv, int k)
9 {
10    if (k < 0 || k >= RANGE_SIZE) return;
11    cv.words[k / BITS_PER_LONG] |= createMask (k);
12 }
13 // remove the k-th characters in ASCII Table from the set
14 void remove(BitSet &cv, int k)
15 {
16     if (k < 0 || k >= RANGE_SIZE) return;
17     cv.words[k / BITS_PER_LONG] &= ~createMask (k);
18 }
```

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Binary Tree

Three Tree Traversal Methods

- **Preorder Traversal** Root → Left → Right
- **Inorder Traversal** Left → Root → Right
- **Postorder Traversal** Left → Right → Root

Question

Preorder traversal 1, 2, 4, 8, 5, 9, 3, 6, 10, 7

Inorder traversal 8, 4, 2, 9, 5, 1, 6, 10, 3, 7

Based on the given preorder and inorder traversals, what is the postorder traversal?

Binary Tree

Three Tree Traversal Methods

- **Preorder Traversal** Root → Left → Right
- **Inorder Traversal** Left → Root → Right
- **Postorder Traversal** Left → Right → Root

Question

Preorder traversal 1, 2, 4, 8, 5, 9, 3, 6, 10, 7

Inorder traversal 8, 4, 2, 9, 5, 1, 6, 10, 3, 7

Based on the given preorder and inorder traversals, what is the postorder traversal?

Answer

8, 4, 9, 5, 2, 10, 6, 7, 3, 1

Binary Tree (Cont.)

```
1 int find (int x)
2 {
3     for (int i = 0;i < n;++i)
4         if (in[i] == x) return i;
5     return -1;
6 }
7 void getpost (int sl,int sr,int fl,int fr)
8 {
9     if (sl > sr || fl > fr) return;
10    int k = find (_1_);
11    getpost (_2_);getpost (_3_);
12    printf ("%d\n",pre[sl]);
13 }
14 void solve () {getpost (0,n - 1,0,n - 1);}
```

Binary Tree (Cont.)

```
1 int find (int x)
2 {
3     for (int i = 0;i < n;++i)
4         if (in[i] == x) return i;
5     return -1;
6 }
7 void getpost (int sl,int sr,int fl,int fr)
8 {
9     if (sl > sr || fl > fr) return;
10    int k = find (_1_);
11    getpost (_2_);getpost (_3_);
12    printf ("%d\n",pre[sl]);
13 }
14 void solve () {getpost (0,n - 1,0,n - 1);}
```

Answer

```
1 pre[sl]
2 sl + 1,sl + k - fl,fl,k - 1
3 sl + k - fl + 1,sr,k + 1,fr
```

BST

```
1 struct TreeNode
2 {
3     int key,sz,cnt;
4     TreeNode *left,*right;
5     TreeNode (int val) :
6         key (val), sz (1), cnt (1), left (nullptr), right (nullptr) {}
7 };
```

```
1 bool search (TreeNode* u,int val)
2 {
3     if (u == nullptr) return false;
4     if (u -> key == val) return true;
5     else if (val < u -> key) return search (rt -> left,val);
6     else return search (rt -> right,val);
7 }
```

BST (Cont.)

```
1 int get_sz (TreeNode* u) {return u == nullptr ? 0 : u -> size;}
2 TreeNode* insert (TreeNode* u, int val)
3 {
4     if (u == nullptr) return new TreeNode (val);
5     if (val < u -> key) u -> left = insert (u -> left, val);
6     else if (val > u -> key) u -> right = insert (u -> right, val);
7     else ++(u -> cnt);
8     u -> sz = u -> cnt + get_sz (u -> left) + get_sz (u -> right);
9     return u;
10 }
```

```
1 TreeNode* findMinNode (TreeNode* u)
2 {
3     while (u -> left != nullptr) u = u -> left;
4     return u;
5 }
```

BST (Cont.)

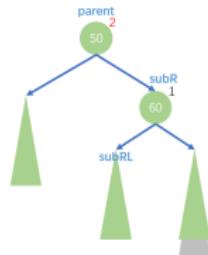
```
1  TreeNode* remove (TreeNode* u,int val)
2  {
3      if (u == nullptr) return u;
4      if (val < u -> key) u -> left = remove (u -> left,val);
5      else if (val > u -> key) u -> right = remove (u -> right,val);
6      else
7      {
8          if (u -> cnt > 1) {--(rt -> cnt);return;}
9          if (u -> left != nullptr && u -> right != nullptr)
10         {
11             TreeNode* nxt = findMinNode (u -> right);
12             u -> key = nxt -> key;u -> cnt = nxt -> cnt;nxt -> cnt = 1;
13             u -> right = remove (u -> right,nxt -> key);
14         }
15     else
16     {
17         TreeNode *tmp = u;
18         u = (u -> left != nullptr) ? u -> left : u -> right;
19         delete tmp;return u;
20     }
21 }
22 u -> sz = u -> cnt + get_sz (u -> left) + get_sz (u -> right);
23 return u;
24 }
```

BST (Cont.)

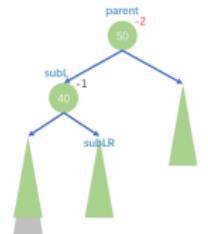
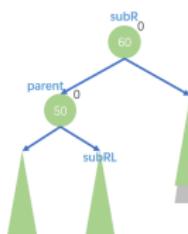
```
1 int queryRank (TreeNode* u,int val)
2 {
3     if (u == nullptr) return;
4     if (v == u -> key) return get_sz (u -> left) + 1;
5     else if (v < u -> key) return queryRank (u -> left,val);
6     return queryRank (u -> right,val) + get_sz (u -> left) + u -> cnt;
7 }
```

```
1 int querykth (TreeNode* u,int k)
2 {
3     if (u == nullptr) return -1;
4     if (u -> left)
5     {
6         if (u -> left -> sz >= k) return querykth (u -> left,k);
7         if (u -> left -> sz + u -> cnt >= k) return u -> key;
8     }
9     else if (k <= u -> cnt) return u -> cnt;
10    return querykth (u -> right,k - get_sz (u -> left) - u -> cnt);
11 }
```

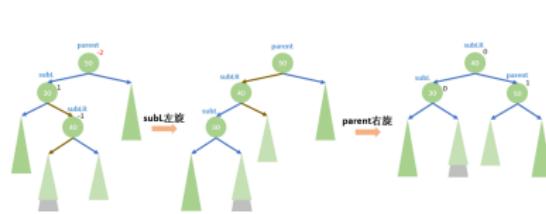
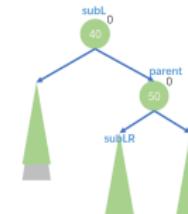
AVL Tree



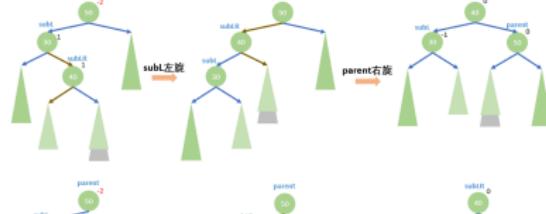
左旋



右旋



subL 左旋



subL 右旋

AVL Tree (Cont.)

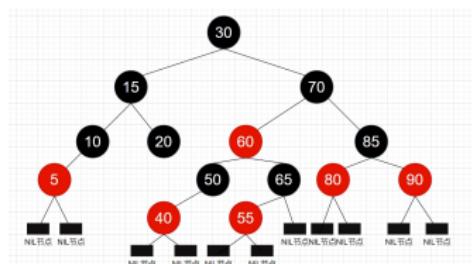
```
1  avl_node* R_rotate (avl_node* u)
2  {
3      avl_node* x = u -> L; avl_node* y = x -> R;
4      x -> R = u; u -> L = y;
5      x -> fat = u -> fat;
6      u -> fat = x;
7      if (y != nullptr) y -> fat = u;
8      u -> upd (); x -> upd ();
9      return x;
10 }
11 avl_node* L_rotate (avl_node* u)
12 {
13     avl_node* x = u -> R; avl_node* y = x -> L;
14     x -> L = u; u -> R = y;
15     x -> fat = u -> fat;
16     u -> fat = x;
17     if (y != nullptr) y -> fat = u;
18     u -> upd (); x -> upd ();
19     return x;
20 }
```

AVL Tree (Cont.)

```
1  avl_node* LR_rotate (avl_node* u) {u -> L = L_rotate (u -> L);return R_rotate (u);} 
2  avl_node* RL_rotate (avl_node* u) {u -> R = R_rotate (u -> R);return L_rotate (u);}
3  int height (avl_node* u) const {return u ? u -> dep : -1;}
4  int calc (avl_node* u) const {return !u ? 0 : height (u -> L) - height (u -> R);}
5  avl_node* balance (avl_node* u)
6  {
7      if (u == nullptr) return nullptr;
8      int fac = calc (u);
9      if (fac > 1)
10     {
11         if (calc (u -> L) >= 0) return R_rotate (u);
12         u -> L = L_rotate (u -> L);
13         return R_rotate (u);
14     }
15     else if (fac < -1)
16     {
17         if (calc (u -> R) <= 0) return L_rotate (u);
18         u -> R = R_rotate (u -> R);
19         return L_rotate (u);
20     }
21 }
22 }
```

Red-Black Tree

Key Characteristics of RB Trees



- Binary Search Tree Property
- Color Rules
 - Every node is either **red** or **black**
 - Root is always **black**
 - All leaves (NIL nodes) are **black**
- Red Node Constraint
 - No two consecutive red nodes
- Black Height Property
 - Every path from root to leaf has the same number of black nodes
 - Ensures tree remains balanced

Red-Black Tree (Cont.)

Theorem

For a red-black tree with n nodes, its height h satisfies $h \leq 2 \log_2(n + 1)$.

Proof

Let bh be the black height of a RB tree.

Then we have $n \geq 2^{bh} - 1$ and $h \leq 2bh$.

So we can conclude $h \leq 2bh \leq \log_2(n + 1)$.

Heap

Heap Construction

$$\begin{aligned}T(n) &= 1 \times (H - 1) + 2 \times (H - 2) + \cdots + 2^{H-1} \times 0 \\&= 2^H - H - 1 = O(n)\end{aligned}$$

```
1 void down (int *a,int n,int i)
2 {
3     int pos = i,l = 2 * i + 1,r = 2 * i + 2;
4     if (l < n && a[l] > a[pos]) pos = l;
5     if (r < n && a[r] > a[pos]) pos = r;
6     if (pos != i) swap (a[i],a[pos]),down (a,n,pos);
7 }
8 void build () //Default: max-heap
9 {
10    for (int i = a.size () / 2 - 1;i >= 0;--i) down (a,a.size (),i);
11 }
```

Heap (Cont.)

Insert

```
1 void up (int *a,int id)
2 {
3     int fa = (id - 1) / 2;
4     while (id > 0 && a[id] > a[fa])
5         {swap (a[id],a[fa]);id = fa;fa = (id - 1) / 2;}
6 }
7 void insert (int val)
8 {
9     if (n == _MAX_SIZE) return;
10    a[n] = val; ++n; up (a,n);
11 }
```

Delete

```
1 void del (int *a,int n)
2 {
3     if (n <= 0) return;
4     swap (a[0],a[n]); --n; down (a,n,0);
5 }
```

Heap (Cont.)

Q1. Given the keyword sequence 5, 8, 12, 19, 28, 20, 15, 22 is a **min-heap**, after inserting keyword 3 and performing heap adjustment, the resulting min-heap is:

- A) 3, 5, 12, 8, 28, 20, 15, 22, 19
- B) 3, 5, 12, 19, 20, 15, 22, 8, 28
- C) 3, 8, 12, 5, 20, 15, 22, 28, 19
- D) 3, 12, 5, 8, 28, 20, 15, 22, 19

Heap (Cont.)

Q1. Given the keyword sequence 5, 8, 12, 19, 28, 20, 15, 22 is a **min-heap**, after inserting keyword 3 and performing heap adjustment, the resulting min-heap is:

- A) 3, 5, 12, 8, 28, 20, 15, 22, 19
- B) 3, 5, 12, 19, 20, 15, 22, 8, 28
- C) 3, 8, 12, 5, 20, 15, 22, 28, 19
- D) 3, 12, 5, 8, 28, 20, 15, 22, 19

Answer

1

A

Heap (Cont.)

Q2. Which of the following statements about a **max-heap** (containing at least 2 elements) is/are correct?

- I. A heap can be viewed as a **complete binary tree**.
 - II. A heap can be stored using **sequential (array-based) storage**.
 - III. A heap can be viewed as a **binary search tree**.
 - IV. The **second largest element** in the heap must be in the level directly below the root.
- A) I and II only
- B) II and III only
- C) I, II, and IV only
- D) I, III, and IV only

Heap (Cont.)

Q2. Which of the following statements about a **max-heap** (containing at least 2 elements) is/are correct?

- I. A heap can be viewed as a **complete binary tree**.
 - II. A heap can be stored using **sequential (array-based) storage**.
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- A) I and II only
- B) II and III only
- C) I, II, and IV only
- D) I, III, and IV only

Answer

1

c

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Complexity

Complexity	Name	Example
$O(1)$	Constant	Array access
$O(\log n)$	Logarithmic	Binary search
$O(n)$	Linear	Linear search
$O(n \log n)$	Log-linear	Merge sort
$O(n^2)$	Quadratic	Bubble sort
$O(n^3)$	Cubic	Matrix multiplication
$O(2^n)$	Exponential	Subset enumeration

Growth Rates:

$$O(1) \leq O(\log n) \leq O(n) \leq O(n \log n) \leq O(n^2) \leq O(n^3) \leq O(2^n)$$

Complexity (Cont.)

Time Complexity Analysis

```
1 for (int i = 1;i <= n;++i)
2     for (int j = 1;j <= n / 2;++j) ++cnt;
3 for (int i = 1;i <= n;i *= 2)
4     for (int j = 1;j <= n;++j) ++cnt;
5 for (int i = 1;i <= n;i *= 2)
6     for (int j = 1;j <= i;++j) ++cnt;
7 for (int i = 1;i <= n;++i)
8     for (int j = 1;j <= n / i;++j) ++cnt;
```

Complexity (Cont.)

Time Complexity Analysis

```
1 for (int i = 1;i <= n;++i)
2     for (int j = 1;j <= n / 2;++j) ++cnt;
3 for (int i = 1;i <= n;i *= 2)
4     for (int j = 1;j <= n;++j) ++cnt;
5 for (int i = 1;i <= n;i *= 2)
6     for (int j = 1;j <= i;++j) ++cnt;
7 for (int i = 1;i <= n;++i)
8     for (int j = 1;j <= n / i;++j) ++cnt;
```

Answer

$O(n^2)$, $O(n \log n)$, $O(n)$, $O(n \log n)$

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Sort

```
1 void Bubble_Sort () {
2     for (int i = 0; i < n; ++i)
3         for (int j = 0; j < n - i - 1; ++j)
4             if (a[j] > a[j + 1]) swap (a[j], a[j + 1]);
5 }
6 void Selection_Sort () {
7     for (int i = 0; i < n; ++i)
8     {
9         int minIndex = i;
10        for (int j = i + 1; j < n; ++j)
11            if (a[j] < a[minIndex]) minIndex = j;
12        swap (a[minIndex], a[i]);
13    }
14 }
15 void Insertion_Sort () {
16     for (int i = 1; i < n; ++i)
17     {
18         int key = a[i], j = i - 1;
19         while (j >= 0 && a[j] > key) a[j + 1] = a[j], --j;
20         a[j + 1] = key;
21     }
22 }
```



Sort (Cont.)

```
1 template <typename T>
2 void Merge_Sort (T a[],int l,int r)
3 {
4     if (l >= r) return;
5     int mid = (l + r) / 2;
6     Merge_Sort (a,l,mid);Merge_Sort (a,mid + 1,r);
7     vector <T> lst;
8     int posl = l, posr = mid + 1;
9     while (posl <= mid && posr <= r)
10    {
11        if (a[posl] <= a[posr]) lst.push_back (a[posl++]);
12        else lst.push_back (a[posr++]);
13    }
14    while (posl <= mid) lst.push_back (a[posl++]);
15    while (posr <= r) lst.push_back (a[posr++]);
16    for (int i = l;i <= r;++i) a[i] = lst[i - 1];
17 }
18 Merge_Sort (a,0,n - 1);
```

Sort (Cont.)

```
1 template <typename T>
2 int partition (T a[],int l,int r)
3 {
4     int pivot = a[l];
5     while (l < r)
6     {
7         while (l < r && pivot <= a[r]) --r;
8         a[l] = a[r];
9         while (l < r && a[l] <= pivot) ++l;
10        a[r] = a[l];
11    }
12    a[l] = pivot;
13    return l;
14}
15 template <typename T>
16 void Quick_Sort (T a[],int l,int r)
17 {
18     if (l >= r) return;
19     int pivot = partition (a,l,r);
20     Quick_Sort (a,l,pivot - 1);Quick_Sort (a,pivot + 1,r);
21 }
22 Quick_Sort (a,0,n - 1);
```

Sort (Cont.)

The Complexity of Different Sort Algorithms

Algorithm	Best	Average	Worst
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Insertion Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Bubble Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$
Heap Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
BST Tree Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$

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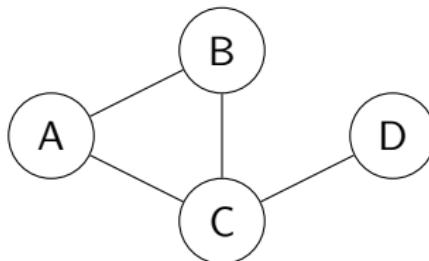
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Graph



Adjacency List : $O(|V| + |E|)$.

A: → B → C → ×

B: → A → C → ×

C: → A → B → D → ×

D: → C → ×

Adjacency Matrix : $O(|V|^2)$.

$$\begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

DFS/BFS

Visualization

Algorithm 1 DFS

```
procedure SEARCH( $u$ )
    mark  $u$  as vis
    for each neighbor  $v$  of  $u$  do
        if  $v$  is not vis then
            SERACH( $v$ )
        end if
    end for
end procedure
```

Algorithm 2 BFS

```
procedure SEARCH( $source$ )
     $Q \leftarrow$  new Queue()
     $Q.enqueue(source)$ 
    while  $Q$  not empty do
         $u \leftarrow Q.dequeue()$ 
        for each neighbor  $v$  of  $u$  do
            if  $v$  is not vis then
                mark  $v$  as vis
                 $Q.enqueue(v)$ 
            end if
        end for
    end while
end procedure
```

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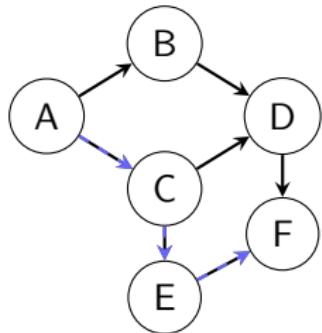
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Unweighted Shortest-Path

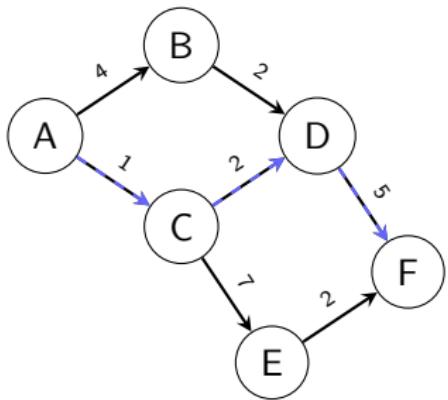
Algorithm 3 0-1 BFS

```
procedure CALC(source, graph)
    Q  $\leftarrow$  new Queue()
    dist  $\leftarrow$  new Dictionary()
    Q.enqueue(source)
    dist[source]  $\leftarrow$  0
    while Q is not empty do
        u  $\leftarrow$  Q.dequeue()
        for each neighbor v of u in graph do
            if v is not in dist then
                dist[v]  $\leftarrow$  dist[u] + 1
                Q.enqueue(v)
            end if
        end for
    end while
    return dist
end procedure
```



Weighted Shortest-Path

Algorithm 4 Dijkstra



```
procedure CALC(source, graph)
    dist  $\leftarrow$  array of size  $|V|$ , initialized to  $\infty$ 
    vis  $\leftarrow$  boolean array of size  $|V|$ , initialized to false
    dist[source]  $\leftarrow 0$ 
    for i  $\leftarrow 1$  to  $|V|$  do
        u  $\leftarrow$  vertex with minimum dist not vis
        vis[u]  $\leftarrow$  true
        for each neighbor v of u with weight w do
            if not vis[v] and dist[u] + w  $<$  dist[v] then
                dist[v]  $\leftarrow$  dist[u] + w
            end if
        end for
    end for
    return dist
end procedure
```

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Algorithm 5 Topological Sort

```
procedure SEARCH( $G(V, E)$ )
     $indeg[|V|] \leftarrow 0$ ,  $q \leftarrow \emptyset$ ,  $ord \leftarrow \emptyset$ 
    for  $(u, v) \in E$  do  $indeg[v] \leftarrow indeg[v] + 1$ 
    end for
    for  $v \in V$  do
        if  $indeg[v] = 0$  then  $q.push(v)$ 
        end if
    end for
    while  $q \neq \emptyset$  do
         $u \leftarrow q.pop()$ ,  $ord.push(u)$ 
        for  $v \in adj[u]$  do
             $indeg[v] \leftarrow indeg[v] - 1$ 
            if  $indeg[v] = 0$  then  $q.push(v)$ 
            end if
        end for
    end while
    if  $|ord| \neq |V|$  then return  $\emptyset$ 
    end if
    return  $ord$ 
end procedure
```

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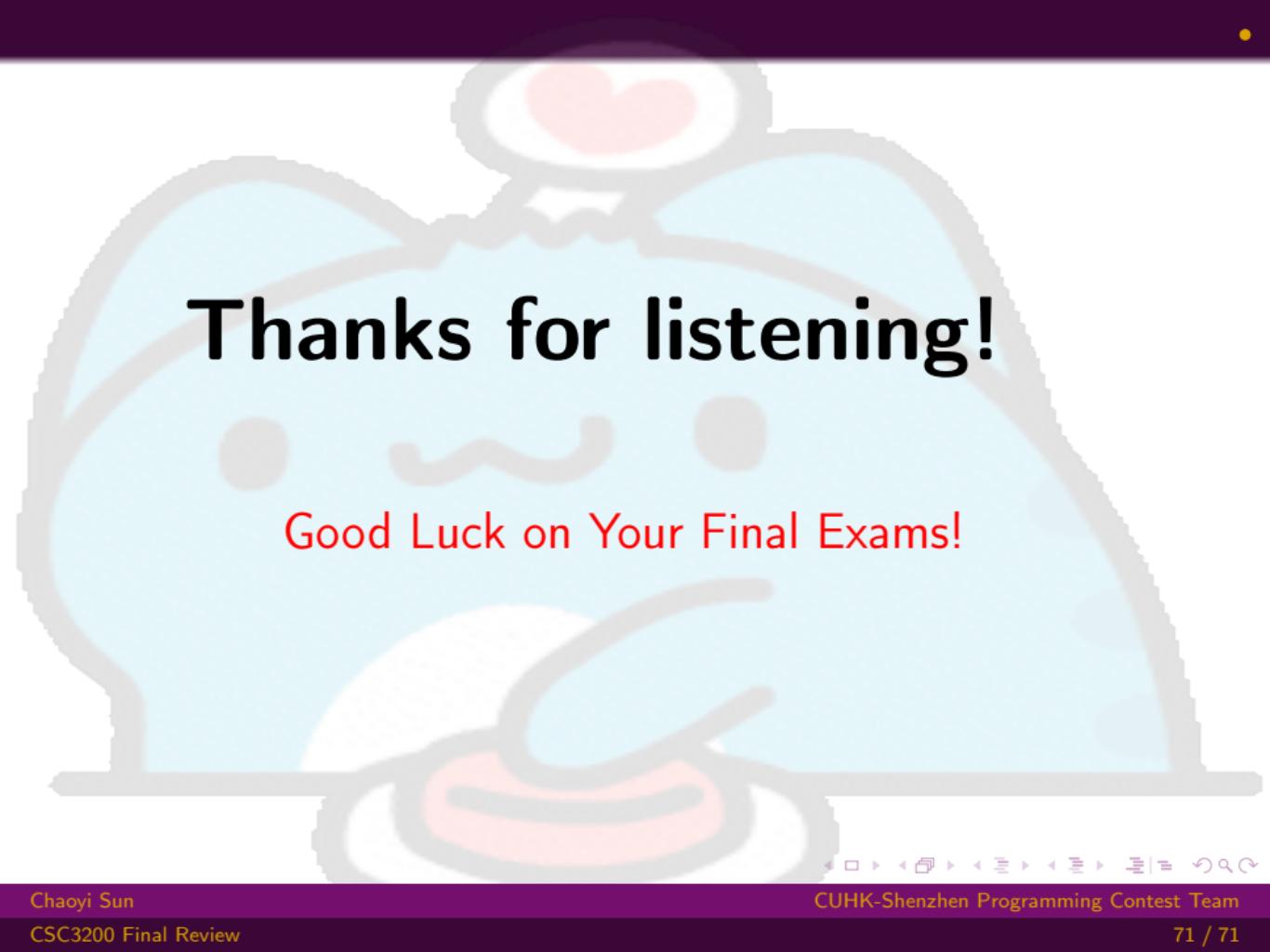
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Thanks for listening!

Good Luck on Your Final Exams!