

# Chapter 5 Curve Fitting

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In this chapter, we will introduce the approximation theory involves **two general types of problems**.

- Approximation Problem of a Function: to find a 'simple' type of function, such as polynomial, that can be used to determine approximate values of the given functions.
- Curve Fitting Problem: to find the **"best" function** in a certain class to fit given data.

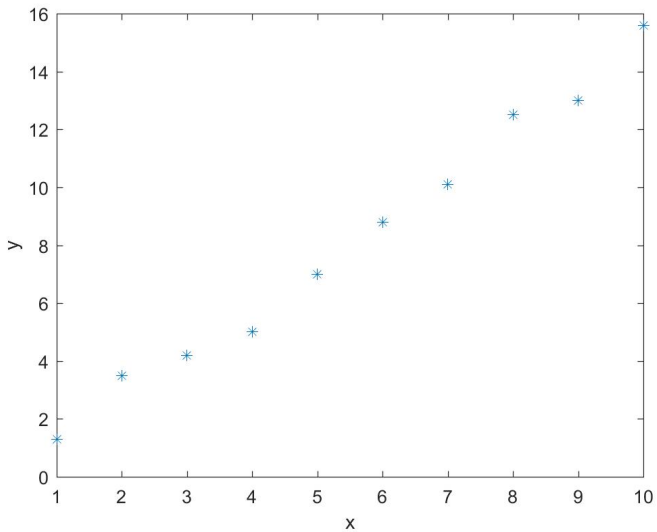
## 5.1 Discrete Least Squares Approximation

### Example 1:

- Consider the problem of estimating the values of a function at nontabulated points, given the experimental data in Table 8.1.

$x_i$	$y_i$	$x_i$	$y_i$
1	1.3	6	8.8
2	3.5	7	10.1
3	4.2	8	12.5
4	5.0	9	13.0
5	7.0	10	15.6

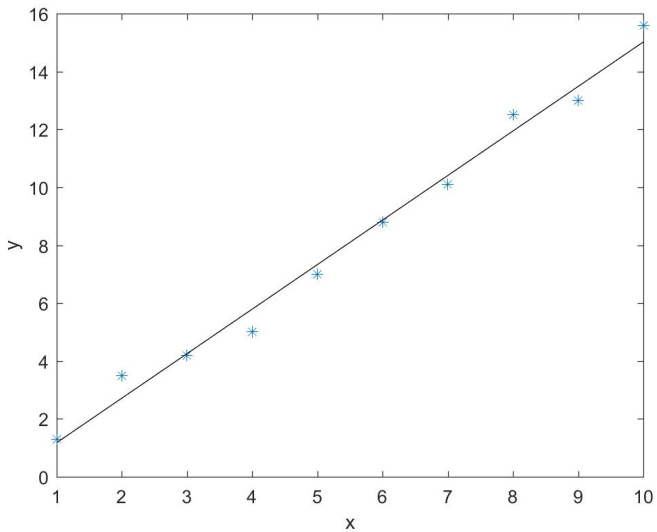
- Using these given data to make a graph, to view the relationship between  $x$  and  $y$ .
- Conclusion: It seems to be linear.



We fit these data with the linear polynomial, by using Matlab7.0 commands:

```
x=[1,2,3,4,5,6,7,8,9,10];  
y=[1.3,3.5,4.2,5.0,7.0,8.8,10.1,12.5,13.0,15.6];  
z=polyfit(x,y,1)  
z=  
1.5382 -0.3600  
Y=1.5382*x-0.36;  
plot(x,y,'*',x,Y,'r')
```

The result can be seen in the following graph.



Let  $a_1x_i + a_0$  denote the  $i$ th value on the approximating line and  $y_i$  be the  $i$ th given  $y$ -value.

## I. Minimax Rule

- The problem of finding the equation of the best linear approximation in the absolute sense requires that values of  $a_0$  and  $a_1$  be found to minimize

$$E_1 = \min_{a_0, a_1} \max_{1 \leq i \leq 10} \{|y_i - (a_1x_i + a_0)|\}.$$

- This is commonly called a **minimax** problem and cannot be handled by elementary techniques.

## II. Absolute Deviation Rule

- Another approach to determining the best linear approximation involves finding Values of  $a_0$  and  $a_1$  to minimize

$$E_2 = \min_{a_0, a_1} \sum_{i=1}^{10} |y_i - (a_1 x_i + a_0)|.$$

- This quantity is called the **absolute deviation**.



- To minimize this function of two variables, we need to set its partial derivatives to zero.
- That is we need to find  $a_0$  and  $a_1$  with

$$0 = \frac{\partial}{\partial a_0} \sum_{i=1}^{10} |y_i - (a_1 x_i + a_0)|,$$

and

$$0 = \frac{\partial}{\partial a_1} \sum_{i=1}^{10} |y_i - (a_1 x_i + a_0)|.$$

- The difficulty is that the absolute-value function is not differentiable at zero, and we may not be able to find solutions to this pair of equations.

### III. Least Square Rule

- The least squares approach to this problem involves determining the best approximating line when the error involved is the sum of the squares of the differences between the  $y$ -values on the approximating line and the given  $y$ -values.
- Hence, constants  $a_0$  and  $a_1$  must be found that minimize the least squares error:

$$E = \min_{a_0, a_1} \sum_{i=1}^{10} [y_i - (a_1 x_i + a_0)]^2$$

# 直线拟合的一般形式

- The least squares method is the most convenient procedure for determining best linear approximations, but there are also important theoretical considerations that favor it.
- The **general problem of fitting the best least squares line** to a collection of data  $\{(x_i, y_i)\}_{i=1}^m$  involves minimizing the total error,

$$E \equiv \min_{a_0, a_1} \sum_{i=1}^m [y_i - (a_1 x_i + a_0)]^2$$

# 法方程或正则方程

- For a minimum to occur, we need

$$0 = \frac{\partial}{\partial a_0} \sum_{i=1}^m [y_i - (a_1 x_i + a_0)]^2 = \sum_{i=1}^m 2(y_i - a_1 x_i - a_0)(-1),$$

and

$$0 = \frac{\partial}{\partial a_1} \sum_{i=1}^m [y_i - (a_1 x_i + a_0)]^2 = \sum_{i=1}^m 2(y_i - a_1 x_i - a_0)(-x_i).$$

- These equations simplify to the **normal equations**:

$$\begin{cases} a_0 \cdot m + a_1 \sum_{i=1}^m x_i &= \sum_{i=1}^m y_i \\ a_0 \sum_{i=1}^m x_i + a_1 \sum_{i=1}^m x_i^2 &= \sum_{i=1}^m x_i y_i. \end{cases}$$

To solve the equations, we get the solution

$$a_0 = \frac{\sum_{i=1}^m x_i^2 \sum_{i=1}^m y_i - \sum_{i=1}^m x_i y_i \sum_{i=1}^m x_i}{m \left( \sum_{i=1}^m x_i^2 \right) - \left( \sum_{i=1}^m x_i \right)^2}$$

and

$$a_1 = \frac{m \sum_{i=1}^m x_i y_i - \sum_{i=1}^m x_i \sum_{i=1}^m y_i}{m \left( \sum_{i=1}^m x_i^2 \right) - \left( \sum_{i=1}^m x_i \right)^2}$$

## The General Form of Discrete Least Square Rule

- The general problem of approximating a set of data:

$$\{(x_i, y_i) | i = 1, 2, \dots, m\},$$

with an algebraic polynomial

$$P_n(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$$

of degree  $n < m - 1$ .

- The General Form of Discrete Least Square Rule:

$$\begin{aligned} \min_{a_0, a_1, \dots, a_n} E &= \sum_{i=1}^m (y_i - P_n(x_i))^2 \\ &= \sum_{i=1}^m \left( y_i - \sum_{k=0}^n a_k x_i^k \right)^2. \end{aligned}$$

- To find the suitable parameters  $a_0, a_1, \dots, a_n$ , such that  $E$  gets to be minimized.
- Let

$$0 = \frac{\partial E}{\partial a_j} = -2 \sum_{i=1}^m y_i x_i^j + 2 \sum_{k=0}^n a_k \sum_{i=1}^m x_i^{j+k}.$$

for each  $j = 0, 1, \dots, n$ .

- This gives  $n + 1$  **normal equations** in the  $n + 1$  unknown parameters  $a_j, j = 0, 1, \dots, n$ .

$$\sum_{k=0}^n a_k \sum_{i=1}^m x_i^{j+k} = \sum_{i=1}^m y_i x_i^j,$$

for each  $j = 0, 1, \dots, n$ .

- Let

$$\mathbf{R} = \begin{bmatrix} x_1^n & x_1^{n-1} & \cdots & x_1 & 1 \\ x_2^n & x_2^{n-1} & \ddots & x_2 & 1 \\ x_3^n & x_3^{n-1} & \ddots & x_3 & 1 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ x_m^n & x_m^{n-1} & \cdots & x_m & 1 \end{bmatrix}, \mathbf{a} = \begin{bmatrix} a_n \\ a_{n-1} \\ \vdots \\ a_1 \\ a_0 \end{bmatrix}, \mathbf{y} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_{m-1} \\ y_m \end{bmatrix}$$

- Then the above normal equations can be written as

$$\mathbf{R}^T \mathbf{R} \mathbf{a} = \mathbf{R}^T \mathbf{y}.$$

- **Note that:** These normal equations have a unique solution provided that the  $x_i$  are distinct.



## 一些可简化为直线拟合的非线性拟合问题

(1) 幂函数:  $y = \alpha x^\beta$  可化为

$$\ln y = \ln \alpha + \beta \ln x.$$

(2) 指数曲线:  $y = \alpha e^{\beta x}$  可化为

$$\ln y = \ln \alpha + \beta x.$$

(3) 对数曲线:  $y = \ln bx$  可化为

$$e^y = bx.$$

(4) 双曲线(单支):  $y = \frac{a}{x} + b$  可化为

$$y = a \frac{1}{x} + b.$$

## 5.2 Orthogonal Polynomials and Least Square Approximation—正交多项式及其最小二乘逼近

- Suppose  $f \in C[a, b]$  and  $P_n(x)$  is a polynomial of degree at most  $n$  with form:

$$P_n(x) = a_n x^n + a_{n-1} x^{n-1} + \cdots + a_1 x + a_0 = \sum_{k=0}^n a_k x^k.$$

- To determine a least squares approximating polynomial  $P_n(x)$ , define

$$\begin{aligned} E &\equiv E(a_0, a_1, \cdots, a_n) = \int_a^b [f(x) - P_n(x)]^2 dx \\ &= \int_a^b \left( f(x) - \sum_{k=0}^n a_k x^k \right)^2 dx. \end{aligned}$$

- Finding real coefficients  $a_0, a_1, \dots, a_n$  so that

$$\begin{aligned} \min_{a_0, a_1, \dots, a_n} E(a_0, a_1, \dots, a_n) &= \int_a^b [f(x) - P_n(x)]^2 dx \\ &= \int_a^b \left( f(x) - \sum_{k=0}^n a_k x^k \right)^2 dx. \\ &= \int_a^b [f(x)]^2 dx - 2 \sum_{k=0}^n a_k \int_a^b x^k f(x) dx + \int_a^b \left( \sum_{k=0}^n a_k x^k \right)^2 dx, \end{aligned}$$

- Let

$$\frac{\partial E}{\partial a_j} = 0, \quad j = 0, 1, \dots, n.$$

- we have normal equations for  $a_0, a_1, \dots, a_n$ :

$$\frac{\partial E}{\partial a_j} = -2 \int_a^b x^j f(x) dx + 2 \sum_{k=0}^n a_k \int_a^b x^{j+k} dx, \quad j = 0, 1, \dots, n.$$

- To find  $P_n(x)$ , the  $(n + 1)$  linear **normal equations**

$$\sum_{k=0}^n a_k \int_a^b x^{j+k} dx = \int_a^b x^j f(x) dx, \quad j = 0, 1, \dots, n.$$

- Rewrite it in linear system of equations

$$\begin{aligned} a_0 \int_a^b 1 dx + a_1 \int_a^b x dx + \dots + a_n \int_a^b x^n dx &= \int_a^b f(x) dx \\ a_0 \int_a^b x dx + a_1 \int_a^b x^2 dx + \dots + a_n \int_a^b x^{n+1} dx &= \int_a^b x f(x) dx \\ &\vdots \\ a_0 \int_a^b x^n dx + a_1 \int_a^b x^{n+1} dx + \dots + a_n \int_a^b x^{2n} dx &= \int_a^b x^n f(x) dx \end{aligned}$$

- **Note that:** The normal equations always have a unique solution provided  $f \in C[a, b]$ .

- The coefficients in the linear system are of the form

$$\int_a^b x^{j+k} dx = \frac{b^{j+k+1} - a^{j+k+1}}{j+k+1},$$

for each  $j, k = 0, 1, 2, \dots, n$  and in the right side are of the form

$$\int_a^b x^j f(x) dx, \text{ for } j = 0, 1, 2, \dots, n.$$

- The matrix in the linear system is known as a **Hilbert matrix**.

## Remarks:

- 1 The linear system does not have an easily computed numerical solution.
- 2 The calculations that were performed in obtaining the best  $n$ th-degree polynomial,  $P_n(x)$ , do not lessen the amount of work required to obtain  $P_{n+1}(x)$ , the polynomial of next higher degree.

- To consider the computational efficiency, a different technique of least squares approximations will now be considered.
- To facilitate the discussion, we need some new concepts.

### Definition 5.1

- The set of functions  $\{\phi_0, \phi_1, \dots, \phi_n\}$  is said to be **linearly independent** on  $[a, b]$  if, whenever

$$c_0\phi_0(x) + c_1\phi_1(x) + \dots + c_n\phi_n(x) = 0,$$

for all  $x \in [a, b]$ , we have  $c_0 = c_1 = \dots = c_n = 0$ .

- Otherwise the set of functions is said to be **linearly dependent**.

## Theorem 5.2

If  $\phi_j(x)$  is a polynomial of degree  $j$ , for each  $j = 0, 1, \dots, n$ , then  $\{\phi_0, \phi_1, \dots, \phi_n\}$  is linearly independent on any interval  $[a, b]$ .

### Proof:

- Suppose  $c_0, c_1, \dots, c_n$  are real numbers for which

$$P(x) = c_0\phi_0(x) + c_1\phi_1(x) + \dots + c_n\phi_n(x) = 0,$$

for all  $x \in [a, b]$ .

- The polynomial  $P(x)$  vanishes on  $[a, b]$ , so it must be the zero polynomial, and the coefficients of all the powers of  $x$  are zero.
- particular, the coefficient of  $x^n$  is zero.



- Since  $c_n \phi_n(x)$  is the only term in  $P(x)$  that contains  $x_n$ , we must have  $c_n = 0$  and

$$P(x) = \sum_{j=0}^{n-1} c_j \phi_j(x).$$

- With same idea above, since the only term that contains a power of  $x^{n-1}$  is  $c_{n-1} \phi_{n-1}(x)$ , so this term must also be zero and

$$P(x) = \sum_{j=0}^{n-2} c_j \phi_j(x).$$

- With a similar manner, the remaining constants  $c_{n-2}, c_{n-3}, \dots, c_0$  are all zero, which implies that  $\{\phi_0, \phi_1, \dots, \phi_n\}$  is linearly independent. ■■

**Notation:** Let  $\Pi_n$  be the **set of all polynomials of degree at most  $n$ .**

### Theorem 5.3:

If  $\{\phi_0(x), \phi_1(x), \dots, \phi_n(x)\}$  is a collection of linearly independent polynomials in  $\Pi_n$ , then any polynomial in  $\Pi_n$  can be written uniquely as a linear combination of  $\phi_0(x), \phi_1(x), \dots, \phi_n(x)$ . ■

## Definition 5.4

An integrable function  $w(x)$  is called a **weight function** on the interval  $I$ , if  $w(x) \geq 0$ , for all  $x \in I$ , but  $w(x) \neq 0$  on any subinterval of  $I$ .

Suppose  $\{\phi_0(x), \phi_1(x), \dots, \phi_n(x)\}$  is a set of linearly independent functions on  $[a, b]$ ,  $w(x)$  is a weight function for  $[a, b]$ , and, for  $f \in C[a, b]$ , a linear combination

$$P(x) = \sum_{k=0}^n a_k \phi_k(x).$$

is sought to minimize the error

$$\begin{aligned} & E(a_0, a_1, \dots, a_n) \\ &= \int_a^b w(x) \left[ f(x) - \sum_{k=0}^n a_k \phi_k(x) \right]^2 dx. \end{aligned} \quad (1)$$

- This problem reduces to the situation considered at the beginning of this section in the special case when  $w(x) \equiv 1$  and  $\phi_k(x) = x^k$ , for each  $k = 0, 1, \dots, n$ .
- The **normal equations** associated with this problem are derived from the fact that for each  $j = 0, 1, \dots, n$ ,

$$0 = \frac{\partial E}{\partial a_j} = 2 \int_a^b w(x) \left[ f(x) - \sum_{k=0}^n a_k \phi_k(x) \right] \phi_j(x) dx.$$

- The system of normal equations can be written

$$\sum_{k=0}^n a_k \int_a^b w(x) \phi_k(x) \phi_j(x) dx = \int_a^b w(x) f(x) \phi_j(x) dx$$

for each  $j = 0, 1, \dots, n$ .

We rewrite it as linear system form:

$$\begin{aligned} & a_0 \int_a^b w(x) \phi_0(x) \phi_0(x) dx + a_1 \int_a^b w(x) \phi_0(x) \phi_1(x) dx + \dots \\ & + a_n \int_a^b w(x) \phi_0(x) \phi_n(x) dx = \int_a^b w(x) f(x) \phi_0(x) dx \\ & a_0 \int_a^b w(x) \phi_1(x) \phi_0(x) dx + a_1 \int_a^b w(x) \phi_1(x) \phi_1(x) dx + \dots \\ & + a_n \int_a^b w(x) \phi_1(x) \phi_n(x) dx = \int_a^b w(x) f(x) \phi_n(x) dx \\ & \dots\dots\dots \\ & a_0 \int_a^b w(x) \phi_n(x) \phi_0(x) dx + a_1 \int_a^b w(x) \phi_n(x) \phi_1(x) dx + \dots \\ & + a_n \int_a^b w(x) \phi_n(x) \phi_n(x) dx = \int_a^b w(x) f(x) \phi_n(x) dx \end{aligned}$$

If the functions  $\phi_0, \phi_1, \dots, \phi_n$  can be chosen so that

$$\int_a^b w(x) \phi_k(x) \phi_j(x) dx = \begin{cases} 0, & \text{when } j \neq k; \\ \alpha_j > 0, & \text{when } j = k. \end{cases} \quad (2)$$

then the normal equations reduce to

$$\int_a^b w(x) f(x) \phi_j(x) dx = a_j \int_a^b w(x) [\phi_j(x)]^2 dx = a_j \alpha_j$$

for each  $j = 0, 1, \dots, n$ , and easily solved to give

$$a_j = \frac{1}{\alpha_j} \int_a^b w(x) f(x) \phi_j(x) dx$$

## Definition 5.5

$\phi_0, \phi_1, \dots, \phi_n$  is said to be an **orthogonal set of functions** for the interval  $[a, b]$  with respect to the weight function  $w$  if

$$\int_a^b w(x) \phi_j(x) \phi_k(x) dx = \begin{cases} 0, & \text{when } j \neq k; \\ \alpha_k > 0, & \text{when } j = k. \end{cases}$$

If, in addition,  $\alpha_k = 1$  for each  $k = 0, 1, 2, \dots, n$ , the set is said to be **orthonormal**.



## Theorem 5.6

If  $\phi_0, \phi_1, \dots, \phi_n$  is an orthogonal set of functions on an interval  $[a, b]$  with respect to the weight function  $w$ , then the least squares approximation to  $f$  on  $[a, b]$  with respect to  $w$  is

$$P(x) = \sum_{k=0}^n a_k \phi_k(x).$$

where for each  $k = 0, 1, 2, \dots, n$ ,

$$a_k = \frac{\int_a^b w(x) \phi_k(x) f(x) dx}{\int_a^b w(x) [\phi_k(x)]^2 dx} = \frac{1}{\alpha_k} \int_a^b w(x) \phi_k(x) f(x) dx.$$

## Theorem 5.7 (Gram-Schmidt Orthogonalize Process)

The set of polynomial functions  $\{\phi_0, \phi_1, \dots, \phi_n\}$  defined in the following way is orthogonal on  $[a, b]$  with respect to the weight function  $w$ .

$$\phi_0(x) = 1, \phi_1(x) = x - B_1, \text{ for each } x \text{ in } [a, b],$$

where

$$B_1 = \frac{\int_a^b xw(x)[\phi_0(x)]^2 dx}{\int_a^b w(x)[\phi_0(x)]^2 dx}$$

and when  $k \geq 2$ ,

$$\phi_k(x) = (x - B_k)\phi_{k-1}(x) - C_k\phi_{k-2}(x), \text{ for each } x \text{ in } [a, b],$$

where

$$B_k = \frac{\int_a^b xw(x)[\phi_{k-1}(x)]^2 dx}{\int_a^b w(x)[\phi_{k-1}(x)]^2 dx}$$

Theorem 8.7 is the known Gram-Schmidt process, it gives a method that how to construct orthogonal polynomials on  $[a, b]$  with respect to a weight function  $w$ .

### Corollary 8.8

For any  $n > 0$ , the set of polynomial functions  $\{\phi_0, \phi_1, \dots, \phi_n\}$  given in Theorem 8.7 is linearly independent on  $[a, b]$  and

$$\int_a^b w(x) \phi_n(x) Q_k(x) dx = 0,$$

for any polynomial  $Q_k(x)$  of degree  $k < n$ .

# Proof:

- Since  $\phi_n(x)$  is a polynomial of degree  $n$ , Theorem 8.2 implies that  $\{\phi_0(x), \phi_1(x), \dots, \phi_n(x)\}$  is a linearly independent set.
- Let  $Q_k(x)$  be a polynomial of degree  $k$ . By Theorem 8.3 there exist numbers  $c_0, c_1, \dots, c_k$  such that

$$Q_k(x) = \sum_{j=0}^k c_j \phi_j(x).$$

# Proof:

Thus,

$$\begin{aligned}\int_a^b w(x) Q_k(x) \phi_n(x) dx &= \sum_{j=0}^k c_j \int_a^b w(x) \phi_j(x) \phi_n(x) dx \\ &= \sum_{j=0}^k c_j \cdot 0 = 0,\end{aligned}$$

Since  $\phi_n(x)$  is orthogonal to  $\phi_j(x)$  for each  $j = 0, 1, \dots, k$ . ■■■

## Example:

The set of **Legendre Polynomial** on  $[-1,1]$  with respect to weight function  $w(x) = 1$ .

Using the method given in theorem 8.7, we can easily give the set of Legendre Polynomial:

$$P_0(x) = 1,$$

$$P_1(x) = x,$$

$$P_2(x) = x^2 - \frac{1}{3},$$

$$P_3(x) = x^3 - \frac{3}{5}x$$

$$\vdots$$

**Note that:** the Legendre Polynomials were ever mentioned in section 4.7, where their roots were used as the nodes in Gaussian Quadrature.

## 5.3 Chebyshev Polynomials and Economization(压缩) of Power Series

- **Chebyshev Polynomials** (Chebyshev 多项式):

$$T_n(x) = \cos[n \arccos x], \quad n = 0, 1, 2, \dots$$

in  $[-1, 1]$ .

- Is  $T_n(x)$  a **polynomial** in  $x \in [-1, 1]$ ?
- Are Chebyshev polynomials orthogonal to each other?

First we show that  $T_n(x)$  is a **polynomial** in  $x$ .

- We note that by definition

$$T_0(x) = \cos 0 = 1,$$

and

$$T_1(x) = \cos[\arccos x] = x.$$

- When  $n > 1$ , we introduce the substitution

$$\theta = \arccos x$$

to change this equation to

$$T_n(\theta(x)) = T_n(\theta) = \cos(n\theta), \text{ where } \theta \in [0, \pi].$$



- A recurrence relation is derived by noting that

$$T_{n+1}(\theta) = \cos[(n+1)\theta] = \cos(n\theta) \cos \theta - \sin(n\theta) \sin \theta$$

and

$$T_{n-1}(\theta) = \cos[(n-1)\theta] = \cos(n\theta) \cos \theta + \sin(n\theta) \sin \theta.$$

- Adding these two equations, gives

$$T_{n+1}(\theta) + T_{n-1}(\theta) = 2 \cos(n\theta) \cos \theta.$$

- Note that

$$\cos \theta = \cos(\arccos x) = x,$$

and

$$\cos(n\theta) = \cos(n \arccos x) = T_n(x),$$

- So we have for each  $n \geq 1$ ,

$$T_{n+1}(x) = 2xT_n(x) - T_{n-1}(x).$$

- Since  $T_0(x) = 1$ ,  $T_1(x) = x$ , the recurrence relation implies that  $T_n(x)$  is a polynomial of degree  $n$  with **leading coefficient**  $2^{n-1}$ , when  $n \geq 1$ .
- **The Chebyshev polynomials** are

$$T_0(x) = 1,$$

$$T_1(x) = x,$$

$$T_2(x) = 2xT_1(x) - T_0(x) = 2x^2 - 1,$$

$$T_3(x) = 2xT_2(x) - T_1(x) = 4x^3 - 3x,$$

$$T_4(x) = 2xT_3(x) - T_2(x) = 8x^4 - 8x^2 + 1,$$

$$T_5(x) = 16x^5 - 20x^3 + 5x,$$

$$T_6(x) = 32x^6 - 48x^4 + 18x^2 - 1,$$

$\vdots$

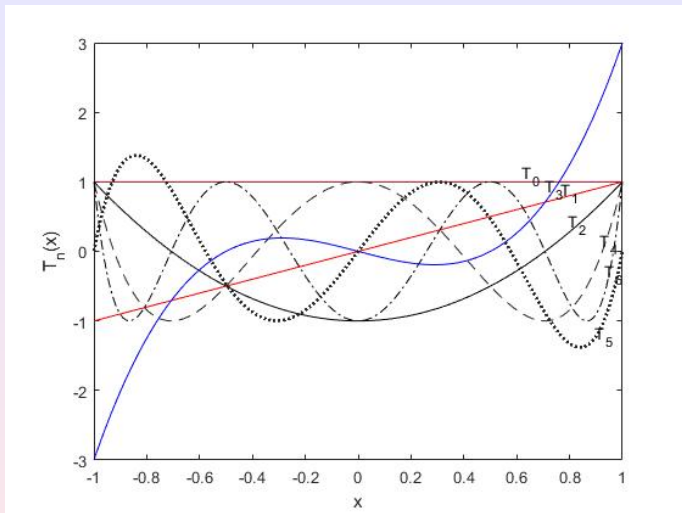


Figure: Chebyshev Polynomials

- Second, we show the **orthogonality of the Chebyshev polynomials** with respect to the weight function

$$w(x) = (1 - x^2)^{-1/2}$$

- That is we need to show that for any  $n \neq m$

$$\int_{-1}^1 w(x) T_n(x) T_m(x) dx = 0, \forall n \neq m$$

# Proof of orthogonality of Chebyshev Polynomials

- Considering

$$\begin{aligned} & \int_{-1}^1 w(x) T_n(x) T_m(x) dx \\ &= \int_{-1}^1 \frac{\cos(n \arccos x) \cos(m \arccos x)}{\sqrt{1-x^2}} dx \end{aligned}$$

- Reintroducing the substitution

$$\theta = \arccos x$$

gives

$$d\theta = -\frac{1}{\sqrt{1-x^2}} dx$$

- Thus

$$\begin{aligned}\int_{-1}^1 \frac{T_n(x) T_m(x)}{\sqrt{1-x^2}} dx &= - \int_{\pi}^0 \cos(n\theta) \cos(m\theta) d\theta \\ &= \int_0^{\pi} \cos(n\theta) \cos(m\theta) d\theta\end{aligned}$$

- Since

$$\cos(n\theta) \cos(m\theta) = \frac{1}{2} [\cos((n+m)\theta) + \cos((n-m)\theta)],$$

- If  $n \neq m$ , we have

$$\begin{aligned} & \int_{-1}^1 \frac{T_n(x) T_m(x)}{\sqrt{1-x^2}} dx \\ &= \frac{1}{2} \int_0^\pi \cos((n+m)\theta) d\theta + \frac{1}{2} \int_0^\pi \cos(n-m)\theta d\theta \\ &= \left[ \frac{1}{2(n+m)} \sin((n+m)\theta) + \frac{1}{2(n-m)} \sin((n-m)\theta) \right]_0^\pi \\ &= 0 \end{aligned}$$

- If  $n = m$ , with a similar technique, we have

$$\int_{-1}^1 \frac{[T_n(x)]^2}{\sqrt{1-x^2}} dx = \frac{\pi}{2}, \quad \text{for each } n \geq 1$$

# Remarks:

- The Chebyshev polynomials are used to minimize approximation error.
- We will see how they are used to solve two problems of this type:
  - ① An **optimal placing of interpolating points** to minimize the error in Lagrange interpolation;
  - ② A means of reducing the degree of an approximating polynomial with minimal loss of accuracy.



# Zeros of Chebyshev polynomial $T_n(x)$

## Theorem 5.9

- The Chebyshev polynomial  $T_n(x)$  of degree  $n \geq 1$  has  $n$  simple zeros in  $[-1, 1]$  at

$$\bar{x}_k = \cos\left(\frac{2k-1}{2n}\pi\right), \text{ for each } k = 1, 2, \dots, n.$$

- Moreover,  $T_n(x)$  assumes its absolute extrema (极值) at

$$\bar{x}'_k = \cos\left(\frac{k\pi}{n}\right) \text{ with } T_n(\bar{x}'_k) = (-1)^k,$$

for each  $k = 0, 1, \dots, n$ .

# Proof of Theorem 5.9

- If we let

$$\bar{x}_k = \cos\left(\frac{2k-1}{2n}\pi\right), \text{ for each } k = 1, 2, \dots, n.$$

then

$$\begin{aligned} T_n(\bar{x}_k) &= \cos(n \arccos \bar{x}_k) \\ &= \cos\left(n \arccos\left(\cos\left(\frac{2k-1}{2n}\pi\right)\right)\right) \\ &= \cos\left(\frac{2k-1}{2}\pi\right) = 0. \end{aligned}$$

- This means that each  $\bar{x}_k$  is a distinct zero of  $T_n$ .
- Since  $T_n(x)$  is a polynomial of degree  $n$ , all zeros of  $T_n(x)$  must be of this form.

To show the second part, first note that

$$T'_n(x) = \frac{d}{dx}[\cos(n \arccos x)] = \frac{n \sin(n \arccos x)}{\sqrt{1-x^2}},$$

and that, when  $k = 1, 2, \dots, n-1$ .

$$T'_n(\bar{x}'_k) = \frac{n \sin\left(n \arccos\left(\cos\left(\frac{k\pi}{n}\right)\right)\right)}{\sqrt{1 - \left[\cos\left(\frac{k\pi}{n}\right)\right]^2}} = \frac{n \sin(k\pi)}{\sin\left(\frac{k\pi}{n}\right)} = 0$$

- Since  $T_n(x)$  is a polynomial of degree  $n$ , its derivative  $T'_n(x)$  is a polynomial of degree  $(n - 1)$ .
- All the zeros of  $T'_n(x)$  occur at these  $n - 1$  points  $\bar{x}'_k, k = 1, 2, \dots, n - 1$ .
- The only other possibilities for extrema of  $T_n(x)$  occur at the endpoints of the interval  $[-1, 1]$ ; that is, at  $\bar{x}'_0 = -1$  and at  $\bar{x}'_n = 1$ .
- Since for any  $k = 0, 1, \dots, n$ , we have

$$\begin{aligned} T_n(\bar{x}'_k) &= \cos\left(n \arccos\left(\cos\left(\frac{k\pi}{n}\right)\right)\right) \\ &= \cos(k\pi) = (-1)^k, \end{aligned}$$

a maximum occurs at each even value of  $k$  and a minimum at each odd value. ■■■

# The Monic(首项系数为1) Chebyshev Polynomial

- The monic polynomials are the ones with leading coefficient 1
- The monic Chebyshev polynomials  $\tilde{T}_n(x)$  are derived from the Chebyshev polynomial  $T_n(x)$  by dividing by the leading coefficient  $2^{n-1}$ .
- Hence,

$$\tilde{T}_0(x) = 1 \quad \text{and} \quad \tilde{T}_n(x) = \frac{1}{2^{n-1}} T_n(x),$$

for each  $n \geq 1$

The recurrence relationship satisfied by the Chebyshev polynomials implies that

$$\tilde{T}_0(x) = 1,$$

$$\tilde{T}_1(x) = \frac{1}{2^0} T_1(x) = x,$$

$$\tilde{T}_2(x) = x \tilde{T}_1(x) - \frac{1}{2} \tilde{T}_0(x) = x^2 - \frac{1}{2}$$

$$\tilde{T}_{n+1}(x) = x \tilde{T}_n(x) - \frac{1}{4} \tilde{T}_{n-1}(x), n \geq 2$$

# Properties of $\tilde{T}_n(x)$ :

1. The zeros of  $\tilde{T}_n(x)$  occur at

$$\bar{x}_k = \cos\left(\frac{2k-1}{2n}\pi\right), \text{ for each } k = 1, 2, \dots, n.$$

2. The extreme values of  $\tilde{T}_n(x)$ , for  $n \geq 1$ , occur at

$$\bar{x}'_k = \cos\left(\frac{k\pi}{n}\right), \text{ with } \tilde{T}_n(\bar{x}'_k) = \frac{(-1)^k}{2^{n-1}} \quad (3)$$

for each  $k = 0, 1, 2, \dots, n$ .

Let  $\tilde{\Pi}_n$  denote the **set of all monic polynomials of degree  $n$** .

The relation expressed in Eq. (3) leads to an important minimization property that distinguishes  $\tilde{T}_n(x)$  from the other members of  $\tilde{\Pi}_n$ .

### Theorem 5.10

The polynomials of the form  $\tilde{T}_n(x)$ , when  $n \geq 1$ , have the property that

$$\frac{1}{2^{n-1}} = \max_{x \in [-1,1]} |\tilde{T}_n(x)| \leq \max_{x \in [-1,1]} |P_n(x)|,$$

for all  $P_n(x) \in \tilde{\Pi}_n$ .

Moreover, equality can occur only if  $P_n \equiv \tilde{T}_n$ .



# Proof of Theorem 5.10(反证法):

- Suppose that  $P_n(x) \in \tilde{\Pi}_n$  and

$$\max_{x \in [-1,1]} |P_n(x)| \leq \frac{1}{2^{n-1}} = \max_{x \in [-1,1]} |\tilde{T}_n(x)|.$$

- Let  $Q = \tilde{T}_n - P_n$ .
- Since  $\tilde{T}_n(x)$  and  $P_n(x)$  are both monic polynomials of degree  $n$ ,  $Q(x)$  is a polynomial of degree at most  $(n-1)$ .
- Moreover, at the extreme points of  $\tilde{T}_n(x)$ ,

$$Q(\bar{x}'_k) = \tilde{T}_n(\bar{x}'_k) - P_n(\bar{x}'_k) = \frac{(-1)^k}{2^{n-1}} - P_n(\bar{x}'_k).$$

- Since

$$|P_n(\bar{x}'_k)| \leq \frac{1}{2^{n-1}}, \text{ for each } k = 0, 1, \dots, n$$

we have

$$Q(\bar{x}'_k) \leq 0, \text{ when } k \text{ is odd}$$

and

$$Q(\bar{x}'_k) \geq 0, \text{ when } k \text{ is even.}$$

- Since  $Q$  is continuous, the Intermediate Value Theorem implies that  $Q(x)$  has at least one zero between  $\bar{x}'_j$  and  $\bar{x}'_{j+1}$ , for each  $j = 0, 1, \dots, n-1$ .
- Thus  $Q$  has at least  $n$  zeros in the interval  $[-1, 1]$ .
- But the degree of  $Q(x)$  is less than  $n$ , so  $Q \equiv 0$ , this implies that  $P_n \equiv \tilde{T}_n$ . ■■■

# Application I. Error Estimation for Lagrange Interpolation

- Suppose that  $x_0, x_1, x_2, \dots, x_n$  are distinct points in the interval  $[-1, 1]$
- $P(x)$  is the Lagrange interpolating polynomial of degree  $n$
- If  $f \in C^{n+1}[-1, 1]$ , then, for each  $x \in [-1, 1]$ , a number  $\xi(x)$  exists in  $(-1, 1)$  with

$$f(x) - P(x) = \frac{f^{(n+1)}(\xi(x))}{(n+1)!} (x - x_0)(x - x_1) \cdots (x - x_n),$$

- Since  $\xi(x)$  just show existence, and we don't know where it is, thus the bound of  $|f^{(n+1)}(\xi(x))|$  can't be done. So the left problem is to minimize the quantity

$$|(x - x_0)(x - x_1) \cdots (x - x_n)|$$

throughout the interval  $[-1, 1]$ .

- If we choose the nodes of  $x_0, x_1, \dots, x_n$  for Lagrange Interpolation as the zeros of Chebyshev polynomial  $T_{n+1}(x)$ , then

$$(x - x_0)(x - x_1) \cdots (x - x_n) = \tilde{T}_{n+1}(x).$$

- The maximum value of

$$\max_{-1 \leq x \leq 1} |(x - x_0)(x - x_1) \cdots (x - x_n)|$$

is smallest when  $x_k$  is chosen to be the  $(k + 1)$ st zeros of  $\tilde{T}_{n+1}$ , for each  $k = 0, 1, \dots, n$

- That is, when  $x_k$  is

$$\bar{x}_{k+1} = \cos\left(\frac{2k + 1}{2(n + 1)}\pi\right).$$

- Since  $\max_{x \in [-1,1]} |\tilde{T}_{n+1}(x)| = \frac{1}{2^n}$ , this also implies that

$$\begin{aligned} \frac{1}{2^n} &= \max_{x \in [-1,1]} |(x - \bar{x}_1)(x - \bar{x}_2) \cdots (x - \bar{x}_{n+1})| \\ &\leq \max_{x \in [-1,1]} |(x - x_0)(x - x_1) \cdots (x - x_n)|, \end{aligned}$$

for any choice of  $x_0, x_1, \dots, x_n$  in the interval  $[-1,1]$ .

## Application II. To Reduce the Degree of an Approximating Polynomial with a Minimal Loss of Accuracy.

- Consider approximating an arbitrary  $n$ th-degree polynomial

$$P_n(x) = a_n x^n + a_{n-1} x^{n-1} + \cdots + a_1 x + a_0$$

on  $[-1, 1]$  with a polynomial  $P_{n-1}(x)$  of degree at most  $n - 1$ .

- The object is to choose  $P_{n-1}(x)$  in  $\Pi_{n-1}$ , so that

$$\max_{x \in [-1, 1]} |P_n(x) - P_{n-1}(x)|$$

is as small as possible.

- We first note that  $(P_n(x) - P_{n-1}(x))/a_n$  is a monic polynomial of degree  $n$ , so applying Theorem 5.10 gives

$$\max_{x \in [-1,1]} \left| \frac{1}{a_n} (P_n(x) - P_{n-1}(x)) \right| \geq \frac{1}{2^{n-1}}.$$

- Equality occurs precisely when

$$\frac{1}{a_n} (P_n(x) - P_{n-1}(x)) = \tilde{T}_n(x).$$

- This means that we should choose

$$P_{n-1}(x) = P_n(x) - a_n \tilde{T}_n(x),$$

- With this choice we have the minimum value of

$$\begin{aligned} & \max_{x \in [-1,1]} |(P_n(x) - P_{n-1}(x))| \\ &= |a_n| \max_{x \in [-1,1]} \left| \frac{1}{a_n} (P_n(x) - P_{n-1}(x)) \right| \\ &= \frac{|a_n|}{2^{n-1}}. \end{aligned}$$



## Corollary 5.11

If  $P(x)$  is the interpolating polynomial of degree at most  $n$  with nodes at the roots of  $T_n(x)$ , then

$$\begin{aligned} & \max_{x \in [-1,1]} |f(x) - P(x)| \\ & \leq \frac{1}{2^n(n+1)!} \max_{x \in [-1,1]} |f^{(n+1)}(x)|, \end{aligned}$$

for each  $f \in C^{n+1}[-1, 1]$ .

# Notes:

For the case of a general closed interval  $[a, b]$ , we can use the change of variables

$$\tilde{x} = \frac{1}{2}[(b - a)x + a + b]$$

to transform the numbers  $\bar{x}_k$  in the interval  $[-1, 1]$  into the corresponding number  $\tilde{x}_k$  in the interval  $[a, b]$ .

## 5.4 Trigonometric Polynomial Approximation(三角多项式逼近)

- Using series of sine and cosine functions to represent arbitrary functions began in the 1750s with the study of the motion of a vibrating string(弦振动).
- In the early part of the 19th century, Jean Baptiste Joseph Fourier used these series to study the flow of heat and developed quite a complete theory of the subject.
- How to construct an function to approximate periodic function?

# Observation

- Define a set of functions as following

$$\begin{aligned}\phi_0(x) &= 1/2 \\ \phi_k(x) &= \cos kx, \quad k = 1, 2, \dots, n \\ \phi_{n+k}(x) &= \sin kx, \quad k = 1, 2, \dots, n-1,\end{aligned}$$

- Then for each positive integer  $n$ , the set of functions

$$\{\phi_0, \phi_1, \dots, \phi_{2n-1}\},$$

is an **orthogonal set** on  $[-\pi, \pi]$  with respect to weighted function  $w(x) \equiv 1$ .

# Orthogonality

- This **orthogonality** follows from the fact that, for every integer  $j$ , the integrals of  $\sin jx$  and  $\cos jx$  over  $[-\pi, \pi]$  are 0, that is

$$\int_{-\pi}^{\pi} \sin(jx) \cos(jx) dx = 0,$$

- we can rewrite products of sine and cosine functions as sums by using the **three trigonometric identities**

$$\sin x \sin y = \frac{1}{2} [\cos(x - y) - \cos(x + y)],$$

$$\cos x \cos y = \frac{1}{2} [\cos(x - y) + \cos(x + y)],$$

$$\sin x \cos y = \frac{1}{2} [\sin(x - y) + \sin(x + y)].$$

# Construction of Fourier series

- Let  $\mathfrak{I}_n$  denote the set of all linear combinations of the functions  $\{\phi_0, \phi_1, \dots, \phi_{2n-1}\}$ .
- This set is called the **set of trigonometric polynomials** of degree less than or equal to  $n$   
(Notes: Some sources also include an additional function  $\phi_{2n}(x) = \sin nx$  in the set.)
- For a function  $f \in C[-\pi, \pi]$ , we want to find the continuous least squares approximation by functions in  $\mathfrak{I}_n$  in the form

$$S_n(x) = \frac{a_0}{2} + a_n \cos nx + \sum_{k=1}^{n-1} (a_k \cos kx + b_k \sin kx).$$

- Since the set of functions  $\{\phi_0, \phi_1, \dots, \phi_{2n-1}\}$  is orthogonal on  $[-\pi, \pi]$  with respect to  $w(x) \equiv 1$
- It follows from Theorem 5.6, that the appropriate selection of coefficients is

$$a_k = \frac{1}{\pi} \int_{-\pi}^{\pi} f(x) \cos kx dx, \quad k = 0, 1, \dots, n$$
$$b_k = \frac{1}{\pi} \int_{-\pi}^{\pi} f(x) \sin kx dx, \quad k = 1, 2, \dots, n - 1.$$

- $\lim_{n \rightarrow \infty} S_n(x)$  is called the **Fourier series** of  $f$ .

Fourier series are used to describe the solution of various ordinary and partial-differential equations that occur in physical situations.

# Discrete Least Square Approximation in the Sense of Trigonometric Polynomials

- Suppose that a collection of  $2m$  paired data points

$$\{(x_j, y_j)\}_{j=0}^{2m-1}$$

is given, with equally spaced points  $\{x_j\}_{j=0}^{2m-1}$  in a closed interval  $[a, b]$ .

- For convenience, we assume that the interval is  $[-\pi, \pi]$ , so,

$$x_j = -\pi + \left(\frac{j}{m}\right)\pi, \quad j = 0, 1, \dots, 2m-1. \quad (4)$$

- Note that if it is not  $[-\pi, \pi]$ , a simple linear transformation could be used to translate the data into this form.



- The goal in the discrete case is to determine the trigonometric polynomial  $S_n(x)$  in  $\mathfrak{T}_n$  that will minimize

$$E(S_n) = \sum_{j=0}^{2m-1} [y_j - S_n(x_j)]^2.$$

- Choosing the constants  $a_0, a_1, \dots, a_n; b_1, b_2, \dots, b_{n-1}$ , so that

$$\begin{aligned} & \min_{a_0, a_1, \dots, a_n; b_1, b_2, \dots, b_{n-1}} E(S_n) & (5) \\ = & \sum_{j=0}^{2m-1} \left[ y_j - \left( \frac{a_0}{2} + a_n \cos nx_j + \sum_{k=1}^{n-1} (a_k \cos kx_j + b_k \sin kx_j) \right) \right]^2 \end{aligned}$$

- The determination of the constants is simplified by the fact that the set

$$\{\phi_0, \phi_1, \dots, \phi_{2n-1}\}$$

is orthogonal with respect to summation over the equally spaced points  $\{x_j\}_{j=0}^{2m-1}$  in  $[-\pi, \pi]$ .

- By this we mean that for each  $k \neq l$ ,

$$\sum_{j=0}^{2m-1} \phi_k(x_j) \phi_l(x_j) = 0.$$

## Lemma 5.12

- If the integer  $r$  is not a multiple of  $2m$ , then

$$\sum_{j=0}^{2m-1} \cos rx_j = 0, \text{ and } \sum_{j=0}^{2m-1} \sin rx_j = 0$$

- Moreover, if  $r$  is not a multiple of  $m$ , then

$$\sum_{j=0}^{2m-1} (\cos rx_j)^2 = m, \text{ and } \sum_{j=0}^{2m-1} (\sin rx_j)^2 = m.$$

# Proof:

- Euler's Formula states that if  $i^2 = -1$ , then for every real number  $z$ , we have

$$e^{iz} = \cos z + i \sin z.$$

- Applying this result gives

$$\begin{aligned} \sum_{j=0}^{2m-1} \cos rx_j + i \sum_{j=0}^{2m-1} \sin rx_j &= \sum_{j=0}^{2m-1} (\cos rx_j + i \sin rx_j) \\ &= \sum_{j=0}^{2m-1} e^{irx_j} \end{aligned}$$

- But

$$e^{irx_j} = e^{ir(-\pi + j\pi/m)} = e^{-ir\pi} \cdot e^{ir\frac{j\pi}{m}},$$

so

$$\sum_{j=0}^{2m-1} \cos rx_j + i \sum_{j=0}^{2m-1} \sin rx_j = e^{-ir\pi} \sum_{j=0}^{2m-1} e^{ir\frac{j\pi}{m}}.$$

- Since  $\sum_{j=0}^{2m-1} e^{ir\frac{j\pi}{m}}$  is a geometric series with first term 1 and ratio  $e^{ir\frac{\pi}{m}} \neq 1$ , we have

$$\sum_{j=0}^{2m-1} e^{ir\frac{j\pi}{m}} = \frac{1 - (e^{ir\frac{\pi}{m}})^{2m}}{1 - e^{ir\frac{\pi}{m}}} = \frac{1 - e^{2ir\pi}}{1 - e^{ir\frac{\pi}{m}}}$$

- But  $e^{2ir\pi} = \cos 2r\pi + i \sin 2r\pi = 1$ , so  $1 - e^{2ir\pi} = 0$  and

$$\sum_{j=0}^{2m-1} \cos rx_j + i \sum_{j=0}^{2m-1} \sin rx_j = e^{-ir\pi} \sum_{j=0}^{2m-1} e^{ir\frac{j\pi}{m}} = 0$$

- This implies that both

$$\sum_{j=0}^{2m-1} \cos rx_j = 0, \quad \text{and} \quad \sum_{j=0}^{2m-1} \sin rx_j = 0$$

- If  $r$  is not a multiple of  $m$ , these sums imply that

$$\begin{aligned}\sum_{j=0}^{2m-1} (\cos rx_j)^2 &= \sum_{j=0}^{2m-1} \frac{1}{2} (1 + \cos 2rx_j) \\ &= \frac{1}{2} \left[ \sum_{j=0}^{2m-1} 1 + \sum_{j=0}^{2m-1} \cos 2rx_j \right] \\ &= \frac{1}{2} (2m + 0) = m\end{aligned}$$

- Similarly, that

$$\sum_{j=0}^{2m-1} (\sin rx_j)^2 = m. \quad \blacksquare \blacksquare \blacksquare.$$

- Now let's show the orthogonality of the set

$$\{\phi_0, \phi_1, \dots, \phi_{2n-1}\},$$

which means that for  $k \neq l$ , we have

$$\sum_{j=0}^{2m-1} \phi_k(x_j) \phi_l(x_j) = 0.$$

- Consider, for example, the case

$$\sum_{j=0}^{2m-1} \phi_k(x_j) \phi_{n+l}(x_j) = \sum_{j=0}^{2m-1} (\cos kx_j)(\sin lx_j).$$

- Since

$$\cos kx_j \sin lx_j = \frac{1}{2}[\sin(l+k)x_j + \sin(l-k)x_j]$$

and  $(l+k)$  and  $(l-k)$  are both integers that are not multiples of  $2m$

- By Lemma 5.12, implies that

$$\begin{aligned} & \sum_{j=0}^{2m-1} (\cos kx_j)(\sin lx_j) \\ &= \frac{1}{2} \left[ \sum_{j=0}^{2m-1} \sin(l+k)x_j + \sum_{j=0}^{2m-1} \sin(l-k)x_j \right] \\ &= \frac{1}{2}(0+0) = 0. \end{aligned}$$

- This technique is used to show that the orthogonality condition is satisfied for any pairs of the functions and is used to produce the following result.



## Theorem 5.13

The constants in the summation

$$S_n(x) = \frac{a_0}{2} + a_n \cos nx + \sum_{k=1}^{n-1} (a_k \cos kx + b_k \sin kx)$$

that minimize the least squares sum

$$E(a_0, a_1, \dots, a_n; b_1, b_2, \dots, b_{n-1}) = \sum_{j=0}^{2m-1} (y_j - S_n(x_j))^2$$

where

$$a_k = \frac{1}{m} \sum_{j=0}^{2m-1} y_j \cos kx_j, k = 0, 1, \dots, n$$
$$b_k = \frac{1}{m} \sum_{j=0}^{2m-1} y_j \sin kx_j, k = 1, 2, \dots, n.$$



