

Algorithm-6

—— Matrix-Multiply[Dynamic.Programming]

A. Problem Description

Matrix chain multiplication is an optimization problem which is that can be solved using *dynamic programming*. Given a sequence of matrices, we want to find the most efficient way to *multiply these matrices* together. The problem is not actually to perform the multiplications, but merely to decide in which order to perform the multiplications.

We have many options because matrix multiplication is *associative*. In other words, no matter how we parenthesize the product, the result will be the same. For example, if we had four matrices A, B, C, and D, we would have:

$$(ABC)D = (AB)(CD) = A(BCD) = A(BC)D = \dots$$

However, the order in which we parenthesize the product affects the number of simple arithmetic operations needed to compute the product, or the efficiency. For example,

suppose A is a 10 x 30 matrix, B is a 30 x 5 matrix, and C is a 5 x 60 matrix. Then,

$$(AB)C = (10 \times 30 \times 5) + (10 \times 5 \times 60) = 1500 + 3000 = 4500$$

operations

$$A(BC) = (30 \times 5 \times 60) + (10 \times 30 \times 60) = 9000 + 18000 = 27000$$

operations.

B. Description of algorithm

Let $m[i][j]$ be the minimum number of scalar multiplications needed to compute the matrix $A[i : j]$.

$$m[i][j] = \begin{cases} 0 & i = j \\ \min_{i \leq k < j} \{m[i][k] + m[k+1][j] + p_{i-1}p_kp_j\} & i < j \end{cases}$$

MATRIX-CHAIN(p)

$n = p.length - 1$

let $m[1..n, 1..n]$ and $s[1..n-1, 2..n]$ be new tables

for $I = 1$ to n

$m[i, j] = 0$

for $l = 2$ to n

for $i = 1$ to $n - l + 1$

$j = i + l - 1$

$m[i, j] = \inf$

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    for k = i to j - 1
         $q = m[i, k] + p_{i-1}p_kp_j m[k+1, j] +$ 
        if  $q < m[i, j]$ 
             $m[i, j] = q$ 
             $s[i, j] = k$ 

    return m and s

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C. Time Complexity

$T = O()$ n^3

D. Code[Python]

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#!/usr/bin/python
# Filename: MatrixChain.py

def MatrixChain(p, n, m, s):
    for i in range(1, n + 1):
        m[i][i] = 0;
    for r in range(2, n + 1):
        for i in range(1, n - r + 2):
            j = i + r - 1
            m[i][j] = m[i + 1][j] + p[i - 1] * p[i] * p[j]
            s[i][j] = i
            for k in range(i + 1, j):
                t = m[i][k] + m[k + 1][j] + p[i - 1] * p[k] * p[j]
                if t < m[i][j]:
                    m[i][j] = t
                    s[i][j] = k

def Traceback(i, j, s):
    if i == j:
        return
    Traceback(i, s[i][j], s)
    Traceback(s[i][j] + 1, j, s)

```

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print 'Multiply A', i, ',', s[i][j],  
print 'and A', s[i][j] + 1, ',', j
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