

Assignment 4 Answer

Q1. 总分10分

- 状态表

A_n	B_n	input: x	A_{n+1}	B_{n+1}
0	0	0	0	1
0	0	1	0	0
0	1	0	1	1
0	1	1	1	0
1	0	0	1	1
1	0	1	1	0
1	1	0	0	0
1	1	1	1	1

当 $A_n B_n = 01$ 与 $A_n B_n = 10$ 时, 会得到相同的 $A_{n+1} B_{n+1}$, 因此简化后得到的状态表如下:

A_n	B_n	input: x	A_{n+1}	B_{n+1}
0	0	0	0	1
0	0	1	0	0
0	1	0	1	1
0	1	1	0	1
1	1	0	0	0
1	1	1	1	1

我们有 $A(n+1) = T \oplus A(n)$, 因此 $T = A(n) \oplus A(n+1)$, 状态表如下:

A_n	B_n	input: x	A_{n+1}	B_{n+1}	T_A	T_B
0	0	0	0	1	0	1
0	0	1	0	0	0	0
0	1	0	1	1	1	0
0	1	1	0	1	0	0
1	0	0	x	x	x	x
1	0	1	x	x	x	x
1	1	0	0	0	1	1
1	1	1	1	1	0	0

- K-map化简并得出激励方程

A \ Bx	00	01	11	10
0	0	0	0	1
1	x	x	0	1

A \ Bx	00	01	11	10
0	1	0	0	0
1	x	x	0	1

得出激励方程：

$$T_A = Bx' + Ax'$$

$$T_B = B'x' + Ax'$$

Q2. 总分10分

- 获得表达式

- 真值表：

Q(t)	G	In	Q(t+1)
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	0
1	1	1	1

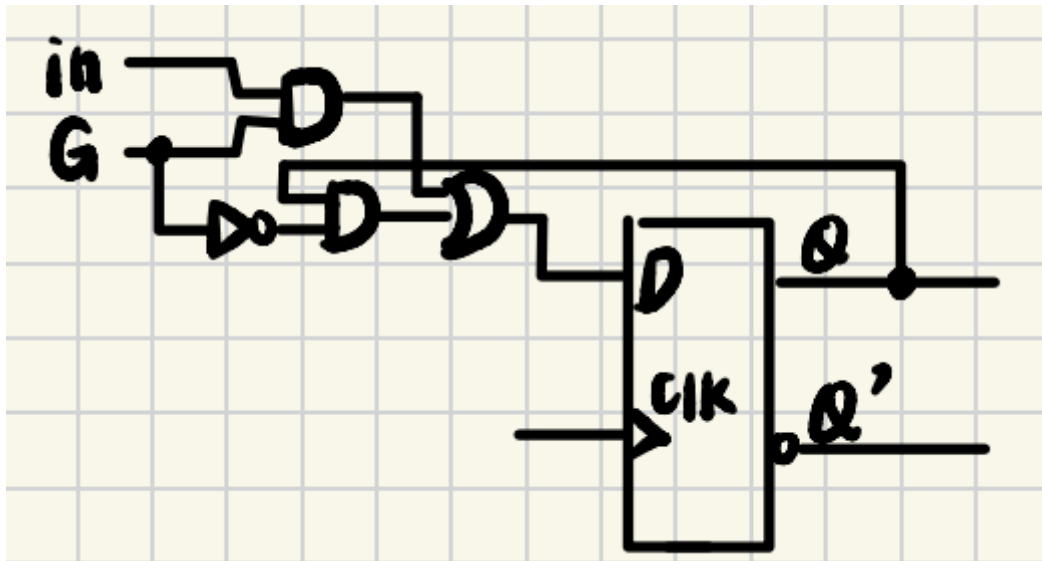
- K-map：

$Q(t+1) :$

$Q(t) \backslash G_{in}$	00	01	11	10
0	0	0	1	0
1	1	1	1	0

得出激励方程: $Q(t+1) = G'Q(t) + GIn$

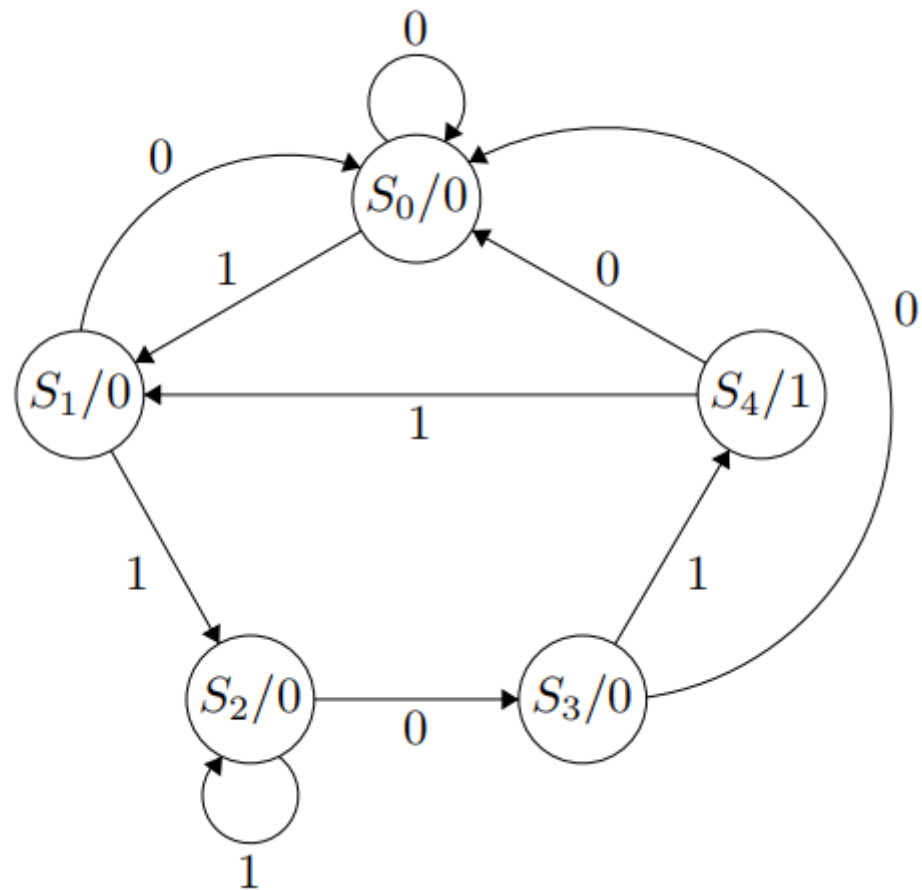
- 电路图



Q3. 总分25分

$$S_0 = \text{None}, S_1 = 1, S_2 = 11, S_3 = 110, S_4 = 1101$$

- 状态图



- 状态表

Present State	Next State		Output	
	$x = 0$	$x = 1$	$x = 0$	$x = 1$
S_0	S_0	S_1	0	0
S_1	S_0	S_2	0	0
S_2	S_3	S_2	0	0
S_3	S_0	S_4	0	0
S_4	S_0	S_1	1	1

- 状态编码 (State Encoding/Assignment)
 $S_0 = 000, S_1 = 001, S_2 = 010, S_3 = 011, S_4 = 100$
- 状态表和JK触发器输入

A	B	C	x	$A(t+1)$	$B(t+1)$	$C(t+1)$	J_A	K_A	J_B	K_B	J_C	K_C
0	0	0	0	0	0	0	0	X	0	X	0	X
0	0	0	1	0	0	1	0	X	0	X	1	X
0	0	1	0	0	0	0	0	X	0	X	X	1
0	0	1	1	0	1	0	0	X	1	X	X	1
0	1	0	0	0	1	1	0	X	X	0	1	X
0	1	0	1	0	1	0	0	X	X	0	0	X
0	1	1	0	0	0	0	0	X	X	1	X	1
0	1	1	1	1	0	0	1	X	X	1	X	1
1	0	0	0	0	0	0	X	1	0	X	0	X
1	0	0	1	0	0	1	X	1	0	X	1	X

A	B	C	x	$A(t+1)$	$B(t+1)$	$C(t+1)$	J_A	K_A	J_B	K_B	J_C	K_C
1	0	1	X	X	X	X	X	X	X	X	X	X
1	1	X	X	X	X	X	X	X	X	X	X	X

注意：ABC=101、110、111(unused state)可以填也可以不填,不在状态表中做要求，但需要在K-map中的需要体现出是X

- 推导激励方程

		cd			
		00	01	11	10
ab	00	0	0	0	0
	01	0	0	1	0
	11	X	X	X	X
	10	X	X	X	X

		cd			
		00	01	11	10
ab	00	X	X	X	X
	01	X	X	X	X
	11	X	X	X	X
	10	1	1	X	X

$$J_A = BCx, \quad K_A = 1$$

		cd			
		00	01	11	10
ab	00	0	0	1	0
	01	X	X	X	X
	11	X	X	X	X
	10	0	0	X	X

		cd			
		00	01	11	10
ab	00	X	X	X	X
	01	0	0	1	1
	11	X	X	X	X
	10	X	X	X	X

$$J_B = Cx, \quad K_B = C$$

		cd			
		00	01	11	10
ab	00	0	1	X	X
	01	1	0	X	X
	11	X	X	X	X
	10	0	1	X	X

		cd			
		00	01	11	10
ab	00	X	X	1	1
	01	X	X	1	1
	11	X	X	X	X
	10	X	X	X	X

$$J_C = B \oplus x, \quad K_C = 1,$$

Q4. 总分25分

$7 \leq 2^n - 1, n \geq 3$, 至少需要3个FF以生成 1011110

- $n = 3$ 时, 状态表如下:

Clock	Sequence Z	FF's Output		
		Q_2	Q_1	Q_0
↑	0	1	0	1
↑	1	0	1	0
↑	1	1	0	1
↑	1	1	1	0
↑	1	1	1	1
↑	0	1	1	1
↑	1	0	1	1

存在相同 Q_2, Q_1, Q_0 但输出结果不同的情况, 因此3个FF不足以生成序列 1011110

- $n = 4$ 时, 状态表如下:

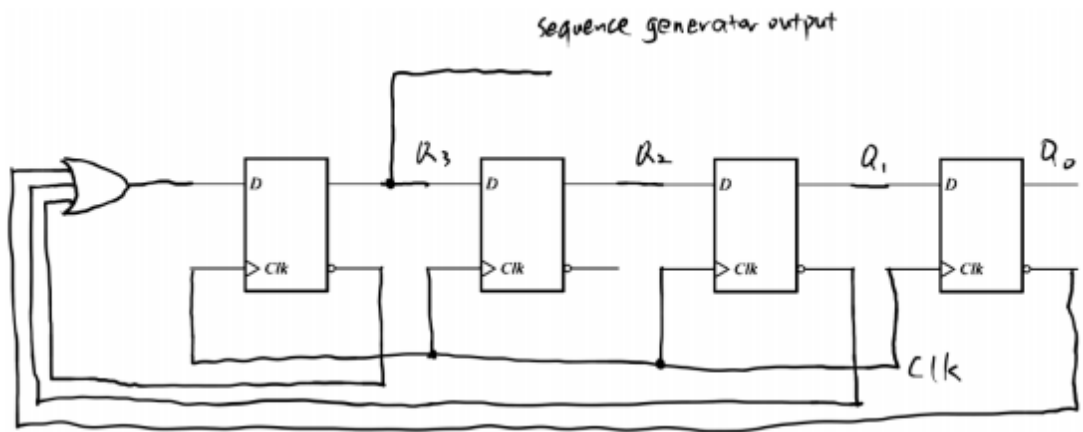
Clock	Sequence Z	FF's Output			
		Q_3	Q_2	Q_1	Q_0
↑	0	1	0	1	1
↑	1	0	1	0	1
↑	1	1	0	1	0
↑	1	1	1	0	1
↑	1	1	1	1	0
↑	0	1	1	1	1
↑	1	0	1	1	1

- 推导方程

$$Z = Q_0' + Q_1' + Q_3'$$

cd ab	00	01	11	10
00	X	X	X	X
01	X	1	1	X
11	X	1	0	1
10	X	X	0	1

- 电路图



Q5. 总分30分

- 状态表

Present State			Next State					
A	B	C	A	B	C	T_A	T_B	T_C
0	0	0	0	0	1	0	0	1
0	0	1	0	1	1	0	1	0
0	1	0	X	X	X	X	X	X
0	1	1	1	1	1	1	0	0
1	0	0	0	0	0	1	0	0
1	0	1	X	X	X	X	X	X
1	1	0	1	0	0	0	1	0
1	1	1	1	1	0	0	0	1

- 推导方程

a	bc			
	00	01	11	10
0	0	0	1	X
1	1	X	0	0

a	bc			
	00	01	11	10
0	0	1	0	X
1	0	X	0	1

a	bc			
	00	01	11	10
0	1	0	0	X
1	0	X	1	0

$$T_A = AB' + A'B = A \oplus B$$

$$T_B = BC' + B'C = B \oplus C$$

$$T_C = AC + A'C' = A \oplus C'$$

- 将don't care conditions情况进行修改，修改合理即可：

1) 修改一：

设定 010 和 101 下一状态均为 000：

Present State			Next State			Flip-Flop Inputs		
A	B	C	A	B	C	T_A	T_B	T_C
0	0	0	0	0	1	0	0	1
0	0	1	0	1	1	0	1	0
0	1	0	0	0	0	0	1	0
0	1	1	1	1	1	1	0	0
1	0	0	0	0	0	1	0	0
1	0	1	0	0	0	1	0	1
1	1	0	1	0	0	0	1	0
1	1	1	1	1	0	0	0	1

推导方程：

$$T_A = AB' + A'BC$$

$$T_B = BC' + A'B'C$$

$$T_C = AC + A'B'C'$$

T_A

A \ bc	00	01	11	10
0	0	1	1	0
1	1	1	0	0

$T_A = AB' + A'B$

T_B

A \ bc	00	01	11	10
0	0	1	1	0
1	1	1	1	0

$T_B = B' + A'B'C$

T_C

A \ bc	00	01	11	10
0	1	0	0	0
1	0	0	1	0

$T_C = AC + A'B'C'$

2) 修改二:

设定 010 的下一状态为 100, 101 的下一状态为 011

Present State			Next State					
A	B	C	A	B	C	T_A	T_B	T_C
0	1	0	1	0	0	1	1	0
1	0	1	0	1	1	1	1	0

推导方程:

		bc			
		00	01	11	10
a	0	0	0	1	1
	1	1	1	0	0

		bc			
		00	01	11	10
a	0	0	1	0	1
	1	0	1	0	1

		bc			
		00	01	11	10
a	0	1	0	0	0
	1	0	0	1	0

The final equation is

$$T_A = AB' + A'B = A \oplus B$$

$$T_B = BC' + B'C = B \oplus C$$

$$T_C = ABC + A'B'C'$$