Principles of Database Systems (CS307)

Lecture 14: Transaction

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- Most contents are from slides made by Stéphane Faroult and the authors of Database System Concepts (7th Edition).
- Their original slides have been modified to adapt to the schedule of CS307 at SUSTech.

Transaction in Real Life

- "An exchange of goods for money"
  - A series of steps
  - All or nothing

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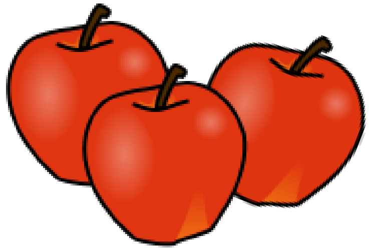

Transaction in Computer

- A transaction is a unit of program execution that accesses and possibly updates various data items
  - A classical example in database: money transfer

E.g., transaction to transfer CNY ¥50 from account A to account B:

1. **read**(A)
2. $A := A - 50$
3. **write**(A)
4. **read**(B)
5. $B := B + 50$
6. **write**(B)

An Example of Transactions in PostgreSQL

- BEGIN, COMMIT, ROLLBACK

```
begin;  -- Start a transaction

update people_1 set num_movies = 50000 where peopleid = 1;

select * from people_1 where peopleid = 1;

delete from people_1 where peopleid > 100 and peopleid < 200;

commit;  -- start executing all the queries above
-- or "rollback;", which means to revoke the operationso of all the queries
```

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6. **write**(B)

  - Two main issues to deal with:
    - Failures of various kinds, such as hardware failures and system crashes
    - Concurrent execution of multiple transactions

Requirements in Transactions

- Atomicity Requirement
  - If the transaction ***fails*** after step 3 and before step 6, money will be "lost" leading to an inconsistent database state
    - Failure could be due to software or hardware
  - The system should ensure that <u>updates of a partially executed transaction are not reflected in the database</u>

E.g., transaction to transfer CNY ¥50 from account A to account B:

```
1. read(A)
2. A := A – 50
3. write(A)
4. read(B)
5. B := B + 50
6. write(B)
```

Requirements in Transactions

- Durability Requirement
  - Once the user has been notified that the transaction has completed (i.e., the transfer of the ¥50 has taken place), the updates to the database by the transaction **must persist** even if there are software or hardware failures.

Requirements in Transactions

- Consistency Requirement
  - Explicitly specified integrity constraints such as primary keys and foreign keys
  - Implicit integrity constraints
    - e.g., sum of balances of all accounts
      - In the example: The sum of A and B is unchanged by the execution of the transaction

E.g., transaction to transfer CNY ¥50 from account A to account B:

```
1. read(A)
2. A := A – 50
3. write(A)
4. read(B)
5. B := B + 50
6. write(B)
```

Requirements in Transactions

- Isolation Requirement
  - If between steps 3 and 6, another transaction **T2** is allowed to access the partially updated database, it will see an inconsistent database
    - The sum A + B will be less than it should be

```
T1                              T2
1.  read(A)
2.  A := A – 50
3.  write(A)
                                   read(A), read(B), print(A+B)

4.  read(B)
5.  B := B + 50
6.  write(B)
```

  - Isolation can be ensured trivially by running transactions serially, that is, one after the other
    - However, executing multiple transactions concurrently has significant benefits

ACID Properties

- A transaction is a unit of program execution that accesses and possibly updates various data items
  - To preserve the integrity of data the database system must ensure:

Atomicity: Either <u>all operations</u> of the transaction are properly reflected in the database, or <u>none</u> are

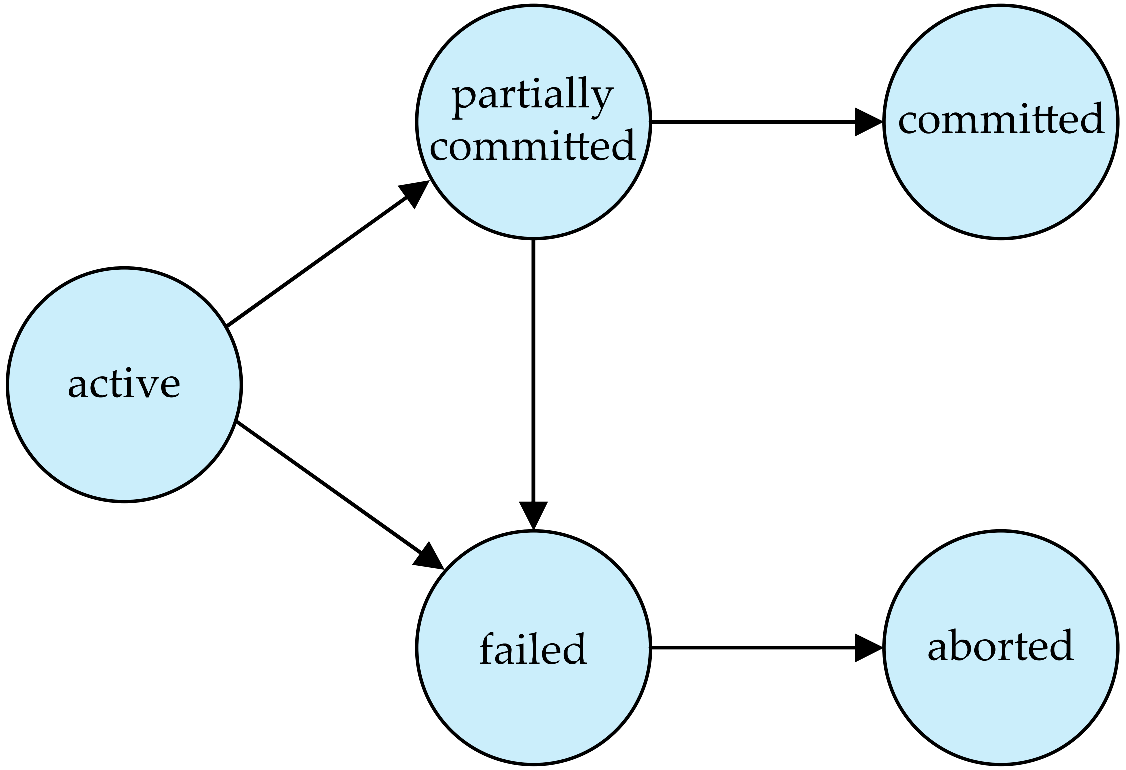
Consistency: Execution of a transaction in isolation preserves the <u>consistency of the database</u>.

Isolation: Although multiple transactions may execute concurrently, each transaction must be <u>unaware of other concurrently executing transactions</u>. Intermediate transaction results must be hidden from other concurrently executed transactions.

- That is, for every pair of transactions T_i and T_j, it appears to T_i that either T_j, finished execution before T_i started, or T_j started execution after T_i finished.

Durability: After a transaction completes successfully, the <u>changes</u> it has made to the database <u>persist</u>, even if there are system failures.

Transaction State



- Active
  - The initial state; the transaction stays in this state while it is executing
- Partially committed
  - After the final statement has been executed.
- Failed
  - After the discovery that normal execution can no longer proceed.
- Aborted – after the transaction has been rolled back and the database restored to its state prior to the start of the transaction. Two options after it has been aborted:
  - Restart the transaction
    - Can be done only if no internal logical error
  - Kill the transaction
- Committed
  - After successful completion.

Concurrent Executions

- Multiple transactions are allowed to run concurrently in the system. Advantages are:
  - Increased processor and disk utilization, leading to better transaction throughput
    - E.g., one transaction can be using the CPU while another is reading from or writing to the disk
  - Reduced average response time for transactions
    - Short transactions do not need to wait behind long ones
- Concurrency control schemes – mechanisms to achieve isolation
  - That is, to control the interaction among the concurrent transactions in order to prevent them from destroying the consistency of the database

Schedules

- **Schedule** – a sequences of instructions that specify the chronological order in which instructions of concurrent transactions are executed
  - A schedule for a set of transactions must consist of all instructions of those transactions
  - Must preserve the order in which the instructions appear in each individual transaction

- A transaction that successfully completes its execution will have a commit instructions as the last statement
  - By default, transaction assumed to execute commit instruction as its last step

- A transaction that fails to successfully complete its execution will have an abort instruction as the last statement

Schedule 1

- Let T_1 transfer CNY ¥50 from A to B, and T_2 transfer 10% of the balance from A to B
  - A **serial schedule** in which T_1 is followed by T_2 :

T_1	T_2
read (A)	
$A := A - 50$	
write (A)	
read (B)	
$B := B + 50$	
write (B)	
commit	
	read (A)
	$temp := A * 0.1$
	$A := A - temp$
	write (A)
	read (B)
	$B := B + temp$
	write (B)
	commit

Schedule 2

- A **serial schedule** where T_2 is followed by T_1

T_1	T_2
	read (A) $temp$:= A * 0.1 A := A - $temp$ write (A) read (B) B := B + $temp$ write (B) commit
read (A) A := A – 50 write (A) read (B) B := B + 50 write (B) commit	

Schedule 3

- Let T_1 and T_2 be the transactions defined previously
  - The following schedule is not a serial schedule, but it is *equivalent* to Schedule 1
    - In Schedules 1, 2 and 3, the sum A + B is preserved.

T_1	T_2
read (A) $A := A - 50$ write (A)	
	read (A) $temp := A * 0.1$ $A := A - temp$ write (A)
read (B) $B := B + 50$ write (B) commit	
	read (B) $B := B + temp$ write (B) commit

Schedule 4

- The following concurrent schedule does not preserve the value of (A + B)

T_1	T_2
read (A)	
$A := A - 50$	
	read (A)
	$temp := A * 0.1$
	$A := A - temp$
	write (A)
	read (B)
write (A)	
read (B)	
$B := B + 50$	
write (B)	
commit	
	$B := B + temp$
	write (B)
	commit

Serializability

- Basic Assumption:
  - Each transaction preserves database consistency
  - Thus, serial execution of a set of transactions preserves database consistency

- A (possibly concurrent) schedule is **serializable** if it is equivalent to a **serial schedule**
  - Different forms of schedule equivalence give rise to the notions of:
    - 1. Conflict serializability
    - 2. * View serializability

Simplified View of Transactions

- We ignore operations other than **read** and **write** instructions

- We assume that transactions may perform arbitrary computations on data in local buffers in between reads and writes.

- Our simplified schedules consist of only **read** and **write** instructions.

Conflicting Instructions

- Instructions I_i and I_j , of transactions T_i and T_j respectively, **conflict** if and only if there exists some item Q accessed by both I_i and I_j, and at least one of these instructions wrote Q.
  - 1. I_i = **read**(Q), I_j = **read**(Q). I_i and I_j don't conflict
  - 2. I_i = **read**(Q), I_j = **write**(Q). They conflict.
  - 3. I_i = **write**(Q), I_j = **read**(Q). They conflict
  - 4. I_i = **write**(Q), I_j = **write**(Q). They conflict

- Intuitively, a conflict between I_i and I_j forces a (logical) temporal order between them.
  - If I_i and I_j are consecutive in a schedule and they do not conflict, their results would remain the same even if they had been interchanged in the schedule.

Conflict Serializability

- If a schedule *S* can be transformed into a schedule *S'* by a series of swaps of non-conflicting instructions, we say that *S* and *S'* are **conflict equivalent**

- We say that a schedule *S* is **conflict serializable** if it is conflict equivalent to a serial schedule

Conflict Serializability

- Schedule 3 can be transformed into Schedule 6, a serial schedule where T_2 follows T_1
  - … by series of swaps of non-conflicting instructions
  - Therefore, Schedule 3 is *conflict serializable*.

Operations on different data

- … and hence swappable in temporal order

T_1	T_2
read (A)	
write (A)	
	read (A)
	write (A)
read (B)	
write (B)	
	read (B)
	write (B)

Schedule 3

T_1	T_2
read (A)	
write (A)	
read (B)	
write (B)	
	read (A)
	write (A)
	read (B)
	write (B)

Schedule 6

Conflict Serializability

- Example of a schedule that is not conflict serializable:

T_3	T_4
read (Q)	
	write (Q)
write (Q)	

  - We are unable to swap instructions in the above schedule to obtain either the serial schedule < T_3, T_4 >, or the serial schedule < T_4, T_3 >.

* View Serializability

- Let *S* and *S'* be two schedules with the same set of transactions. *S* and *S'* are **view equivalent** if the following three conditions are met, for each data item *Q*,
  - If in schedule S, transaction T_i reads the initial value of *Q*, then in schedule *S'* also transaction T_i must read the initial value of *Q*.
  - If in schedule S transaction T_i executes **read**(*Q*), and that value was produced by transaction T_j (if any), then in schedule *S'* also transaction T_i must read the value of *Q* that was produced by the same **write**(Q) operation of transaction T_j .
  - The transaction (if any) that performs the final **write**(*Q*) operation in schedule *S* must also perform the final **write**(*Q*) operation in schedule *S'*.
- As can be seen, view equivalence is also based purely on **reads** and **writes** alone.

* View Serializability

- A schedule S is **view serializable** if it is view-equivalent to a serial schedule
- Every conflict serializable schedule is also view serializable

- Below is a schedule which is view-serializable but *not* conflict serializable

Two "blind writes" in T27 and T28
- Since the written values were not used anywhere else

T_{27}	T_{28}	T_{29}
read (Q)		
	write (Q)	
write (Q)		
		write (Q)

Overwrites values from T27 and T28
- … and hence, swapping write(Q) in T27 and T28 will not affect the resulting value of Q

- What serial schedule is above equivalent to?
- Every view-serializable schedule that is not conflict serializable has **blind writes**
  - **Blind write:** Write operations <u>without</u> further reading

Testing for Serializability

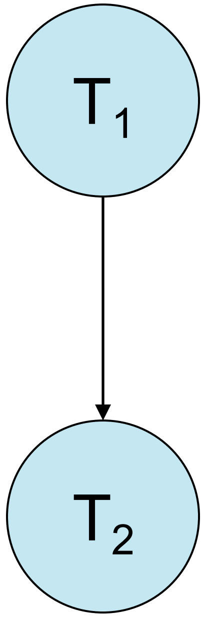
- Consider some schedule of a set of transactions $T_1, T_2, \ldots, T_n$
- Precedence graph
  - A directed graph where the vertices are the transactions (names of the transactions)
  - We draw an arc from T_i to T_j if the two transactions **conflict**
    - which means, in the schedule S, T_i **must** appear earlier than T_j
  - We may label the arc by the item that was accessed.

Conflict – At least one of the following situations exists for a data item Q:

- T_i: write(Q) -> T_j: read(Q)
- T_i: read(Q) -> T_j: write(Q)
- T_i: write(Q) -> T_j: write(Q)

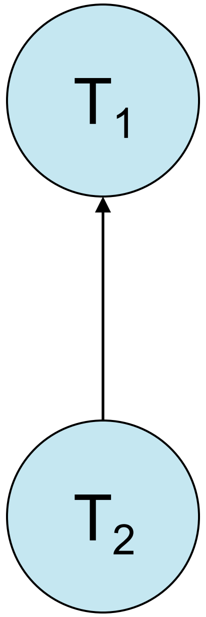
Testing for Serializability

T_1	T_2
read (A)	
$A := A - 50$	
write (A)	
read (B)	
$B := B + 50$	
write (B)	
commit	
	read (A)
	$temp := A * 0.1$
	$A := A - temp$
	write (A)
	read (B)
	$B := B + temp$
	write (B)
	commit



Schedule 1

T_1	T_2
	read (A)
	$temp := A * 0.1$
	$A := A - temp$
	write (A)
	read (B)
	$B := B + temp$
	write (B)
	commit
read (A)	
$A := A - 50$	
write (A)	
read (B)	
$B := B + 50$	
write (B)	
commit	

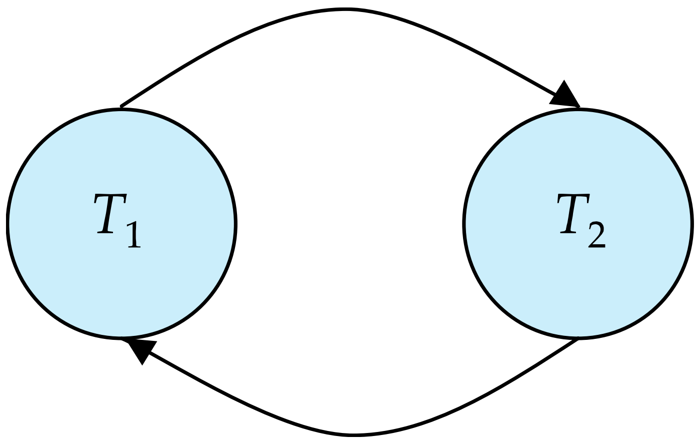


Schedule 2

Testing for Serializability

T_1	T_2
read (A)	
$A := A - 50$	
	read (A)
	$temp := A * 0.1$
	$A := A - temp$
	write (A)
	read (B)
write (A)	
read (B)	
$B := B + 50$	
write (B)	
commit	
	$B := B + temp$
	write (B)
	commit

Schedule 4

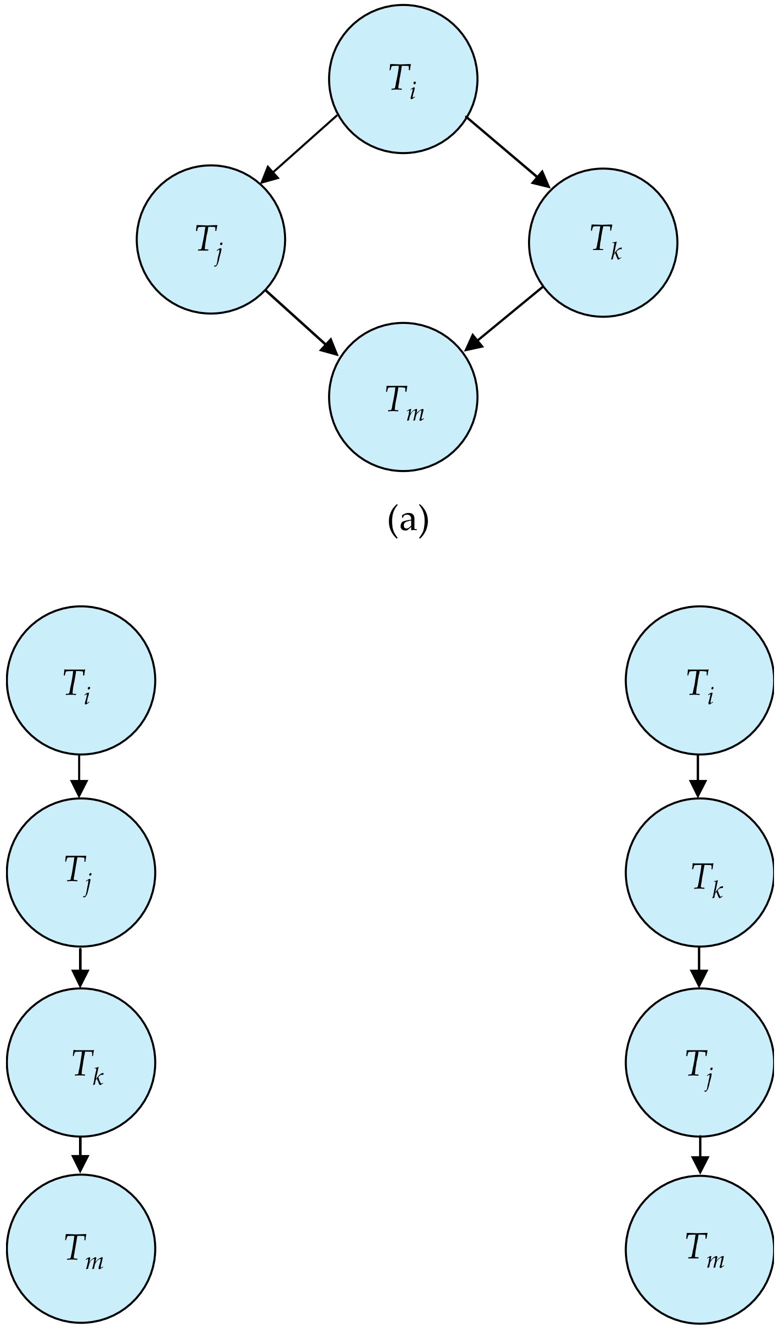


Testing for Serializability

- A schedule is <u>conflict serializable</u> if and only if its precedence graph is acyclic

Cycle-detection: Cycle-detection algorithms exist which take n^2 time, where n is the number of vertices in the graph.

- If the precedence graph is acyclic, the <u>serializability order</u> can be obtained by a *topological sorting* of the graph
  - E.g., The topological order of (a) can be (b) and (c)



(b) (c)

Recoverable Schedules

- Need to address the effect of transaction failures on concurrently running transactions
- **Recoverable schedule** — if a transaction T_j reads a data item previously written by a transaction T_i, then the commit operation of T_i appears before the commit operation of T_j.
- The following schedule is not recoverable

T_8	T_9
read (A)	
write (A)	
	read (A)
	commit
read (B)	

  - If T_8 should abort, T_9 would have read (and possibly shown to the user) an inconsistent database state. Hence, database must ensure that schedules are recoverable.

Weak Levels of Consistency

- Some applications are willing to live with weak levels of consistency, allowing schedules that are not serializable
  - E.g., a read-only transaction that wants to get an approximate total balance of all accounts
    - Such transactions do not need to be serializable with respect to other transactions
  - Purpose: Trade-off between accuracy and performance

Levels of Consistency (in SQL-92)

- Serializable (Strongest)
  - Default
- Repeatable read — only committed records to be read.
  - Repeated reads of same record must return same value.
  - However, a transaction may not be serializable – it may find some records inserted by a transaction but not find others.
- Read committed — only committed records can be read.
  - Successive reads of record may return different (but committed) values.
- Read uncommitted (Weakest) — even uncommitted records may be read.

Levels of Consistency

- Lower degrees of consistency can be useful for gathering approximate information about the database

- Warning: some database systems do not ensure serializable schedules by default
  - E.g., Oracle (and PostgreSQL prior to version 9) by default support a level of consistency called snapshot isolation (not part of the SQL standard)

- Warning 2: All SQL-92 consistency levels infer that dirty writes are prohibited
  - Dirty write - when one transaction overwrites a value that has previously been written by another still in-flight transaction