



UNSW
SYDNEY

COMP9020

Foundations of Computer Science

Lecture 10: Induction

Administrivia

- Assignment 1 due on Friday at 12 noon
 - Late penalty: 5 marks per 24 hours or part thereof for maximum 5 days.
 - Contact me if you need an extension – include zld and length of requested extension
- Assignment 2 released next week
- Feedback:
 - More in-class quizzes
 - More offline consultations
 - Quiz solutions

Outline

Motivation

Basic Induction

Variations on Basic Induction

Structural Induction

Feedback

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Recursive datatypes

Describe arbitrarily large objects in a finite way

Recursive functions

Define behaviour for these objects in a finite way

Induction

Reason about these objects in a finite way

Example

Recall the recursive program:

Example

Summing the first n natural numbers:

```
sum( $n$ ):  
  if( $n = 0$ ): 0  
  else:  $n + \text{sum}(n - 1)$ 
```

Another attempt:

Example

```
sum2( $n$ ):  
  return  $n * (n + 1) / 2$ 
```

Induction proof **guarantees** that these programs will behave the same.

Inductive Reasoning

Suppose we would like to reach a conclusion of the form

$P(x)$ for all x (of some type)

Inductive reasoning (as understood in philosophy) proceeds from examples.

E.g. From “This swan is white, that swan is white, in fact every swan I have seen so far is white”

Conclude: “Every Swan is white”

NB

This may be a good way to discover hypotheses.

But it is not a valid principle of reasoning!

Mathematical induction is a variant that is valid.

Mathematical Induction

Mathematical Induction is based not just on a set of examples, but also a rule for deriving new cases of $P(x)$ from cases for which P is known to hold.

General structure of reasoning by mathematical induction:

Base Case [B]: $P(a_1), P(a_2), \dots, P(a_n)$ for some small set of examples $a_1 \dots a_n$ (often $n = 1$)

Inductive Step [I]: A general rule showing that if $P(x)$ holds for some cases $x = x_1, \dots, x_k$ then $P(y)$ holds for some new case y , constructed in some way from x_1, \dots, x_k .

Conclusion: Starting with $a_1 \dots a_n$ and repeatedly applying the construction of y from existing values, we can eventually construct all values in the domain of interest.

Induction proof structure

Let $P(x)$ be the proposition that ...

We will show that $P(x)$ holds for all x by induction on x .

Base case: $x = \dots$:

- $P(x)$: ...
-
- so $P(x)$ holds.

[Repeat for all base cases]

Inductive case: $P(x)$ implies $P(y)$

- Assume $P(x)$ holds. That is,
- We will show $P(y)$ holds.
- ...
- So $P(x)$ implies $P(y)$.

[Repeat for all inductive cases]

Conclusion

We have shown $P(\dots)$, and for all x , $P(x)$ implies $P(y)$. Therefore, by induction, $P(x)$ holds for all x .

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Basic induction

Basic induction is the general principle applied to the natural numbers.

Goal: Show $P(n)$ holds for all $n \in \mathbb{N}$.

Approach: Show that:

Base case (B): $P(0)$ holds; and

Inductive case (I): If $P(k)$ holds then $P(k + 1)$ holds.

Conclusion (C): $P(n)$ holds for all $n \in \mathbb{N}$.

Example

Recall the recursive program:

Example

Summing the first n natural numbers:

```
sum( $n$ ):  
  if( $n = 0$ ): 0  
  else:  $n + \text{sum}(n - 1)$ 
```

Another attempt:

Example

```
sum2( $n$ ):  
  return  $n * (n + 1) / 2$ 
```

Induction proof **guarantees** that these programs will behave the same.

Example

Let $P(n)$ be the proposition that: $\sum_{i=0}^n i = \frac{n(n+1)}{2}$.

We will show that $P(n)$ holds for all $n \in \mathbb{N}$ by induction on n .

[B] Base case: $n = 0$:

$$\left(\sum_{i=0}^0 i \right) = 0 = \frac{0(0+1)}{2}.$$

So $P(0)$ holds.

[I] Inductive case: $P(k) \Rightarrow P(k+1)$:

That is,

$$\sum_{i=0}^k i = \frac{k(k+1)}{2} \Rightarrow \sum_{i=0}^{k+1} i = \frac{(k+1)(k+2)}{2}$$

(proof?)

[C] Conclusion: We have $P(0)$ is true, and $P(k)$ implies $P(k+1)$.
Therefore, by induction, $P(n)$ holds for all $n \in \mathbb{N}$.

Proof of inductive step.

Assume $P(k)$ holds. That is:

$$\left(\sum_{i=0}^k i \right) = \frac{k(k+1)}{2} + (k+1). \quad (\text{IH})$$

Then:

$$\begin{aligned} \sum_{i=0}^{k+1} i &= \left(\sum_{i=0}^k i \right) + (k+1) \\ &= \frac{k(k+1)}{2} + (k+1) \quad (\text{by the inductive hypothesis}) \\ &= \frac{k(k+1) + 2(k+1)}{2} \\ &= \frac{(k+1)(k+2)}{2} \end{aligned}$$

Therefore $P(k)$ implies $P(k+1)$. □

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Variations

There are many variants of basic induction that may be more useful in certain circumstances. For example:

- 1 Induction from m upwards
- 2 Induction steps > 1
- 3 Strong induction
- 4 Backward induction
- 5 Forward-backward induction
- 6 Structural induction

Induction From m Upwards

If

$$[B] \quad P(m)$$

$$[I] \quad \forall k \geq m (P(k) \rightarrow P(k + 1))$$

then

$$[C] \quad \forall n \geq m (P(n))$$

Example

Theorem. For all $n \geq 1$, the number $8^n - 2^n$ is divisible by 6.

[B] $8^1 - 2^1$ is divisible by 6

[I] if $8^k - 2^k$ is divisible by 6, then so is $8^{k+1} - 2^{k+1}$, for all $k \geq 1$

Prove [I] using the “trick” to rewrite 8^{k+1} as $8 \cdot (8^k - 2^k + 2^k)$
which allows you to apply the IH on $8^k - 2^k$

Induction Steps $\ell > 1$

If

[B] $P(m)$

[I] $P(k) \rightarrow P(k + \ell)$ for all $k \geq m$

then

[C] $P(n)$ for every ℓ 'th $n \geq m$

Example

Every 4th Fibonacci number is divisible by 3.

[B] $F_4 = 3$ is divisible by 3

[I] if $3 \mid F_k$, then $3 \mid F_{k+4}$, for all $k \geq 4$

Prove [I] by rewriting F_{k+4} in such a way that you can apply the IH on F_k

Strong Induction

This is a version in which the inductive hypothesis is stronger.
Rather than using the fact that $P(k)$ holds for a single value, we use *all* values up to k .

If

[B] $P(m)$

[I] $[P(m) \wedge P(m+1) \wedge \dots \wedge P(k)] \rightarrow P(k+1)$ for all $k \geq m$

then

[C] $P(n)$, for all $n \geq m$

Example

Claim: All integers ≥ 2 can be written as a product of primes.

[B] 2 is a product of primes

[I] If all x with $2 \leq x \leq k$ can be written as a product of primes,
then $k + 1$ can be written as a product of primes, for all $k \geq 2$

Proof for [I]?

Negative Integers, Backward Induction

NB

Induction can be conducted over any subset of \mathbb{Z} with least element. Thus m can be negative; eg. base case $m = -10^6$.

NB

One can apply induction in the 'opposite' direction $p(m) \rightarrow p(m - 1)$. It means considering the integers with the opposite ordering where the next number after n is $n - 1$. Such induction would be used to prove some $p(n)$ for all $n \leq m$.

NB

Sometimes one needs to reason about all integers \mathbb{Z} . This requires two separate simple induction proofs: one for \mathbb{N} , another for $-\mathbb{N}$. They both would start from some initial values, which could be the same, e.g. zero. Then the first proof would proceed through positive integers; the second proof through negative integers.

Forward-Backward Induction

Idea

To prove $P(n)$ for all $n \geq k_0$

- verify $P(k_0)$
- prove $P(k_i)$ for infinitely many $k_0 < k_1 < k_2 < k_3 < \dots$
- fill the gaps

$$P(k_1) \rightarrow P(k_1 - 1) \rightarrow P(k_1 - 2) \rightarrow \dots \rightarrow P(k_0 + 1)$$

$$P(k_2) \rightarrow P(k_2 - 1) \rightarrow P(k_2 - 2) \rightarrow \dots \rightarrow P(k_1 + 1)$$

.....

NB

This form of induction is extremely important for the analysis of algorithms.

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Structural Induction

Basic induction allows us to assert properties over **all natural numbers**. The induction scheme (layout) uses the recursive definition of \mathbb{N} .

The induction scheme can be applied not only to natural numbers (and integers) but to any partially ordered set in general – especially those defined recursively.

The basic approach is always the same — we need to verify that

- **[B]** the property holds for all minimal objects — objects that have no predecessors; they are usually very simple objects allowing immediate verification
- **[I]** for any given object, if the property in question holds for *all* its predecessors ('smaller' objects) then it holds for the object itself

Example: Induction on Σ^*

Recall definition of Σ^* :

$$\lambda \in \Sigma^*$$

If $w \in \Sigma^*$ then $aw \in \Sigma^*$ for all $a \in \Sigma$

Structural induction on Σ^* :

Goal: Show $P(w)$ holds for all $w \in \Sigma^*$.

Approach: Show that:

Base case (B): $P(\lambda)$ holds; and

Inductive case (I): If $P(w)$ holds then $P(aw)$ holds for all $a \in \Sigma$.

Example: Induction on Σ^*

Recall:

Formal definition of Σ^* :

$$\lambda \in \Sigma^*$$

If $w \in \Sigma^*$ then $aw \in \Sigma^*$ for all $a \in \Sigma$

Formal definition of concatenation:

(concat.B) $\lambda v = v$

(concat.I) $(aw)v = a(wv)$

Formal definition of length:

(length.B) $\text{length}(\lambda) = 0$

(length.I) $\text{length}(aw) = 1 + \text{length}(w)$

Prove:

$$\text{length}(wv) = \text{length}(w) + \text{length}(v)$$

Example: Induction on Σ^*

Let $P(w)$ be the proposition that, for all $v \in \Sigma^*$:

$$\text{length}(wv) = \text{length}(w) + \text{length}(v).$$

We will show that $P(w)$ holds for all $w \in \Sigma^*$ by **structural induction on w** .

Proof:

Base case ($w = \lambda$):

$$\begin{aligned}\text{length}(\lambda v) &= \text{length}(v) && (\text{concat.B}) \\ &= 0 + \text{length}(v) \\ &= \text{length}(w) + \text{length}(v) && (\text{length.B})\end{aligned}$$

Example: Induction on Σ^*

Proof cont'd:

Inductive case ($w = aw'$): Assume that $P(w')$ holds. That is, for all $v \in \Sigma^*$:

$$(IH): \quad \text{length}(w'v) = \text{length}(w') + \text{length}(v).$$

Then, for all $a \in \Sigma$, we have:

$$\begin{aligned} \text{length}((aw')v) &= \text{length}(a(w'v)) && (\text{concat.l}) \\ &= 1 + \text{length}(w'v) && (\text{length.l}) \\ &= 1 + \text{length}(w') + \text{length}(v) && (IH) \\ &= \text{length}(aw') + \text{length}(v) && (\text{length.l}) \end{aligned}$$

So $P(aw')$ holds.

We have $P(\lambda)$ and for all $w' \in \Sigma^*$ and $a \in \Sigma$: $P(w') \rightarrow P(aw')$.
Hence $P(w)$ holds for all $w \in \Sigma^*$.

Example 2: Induction on Σ^*

Recall $\text{append} : \Sigma^* \times \Sigma \rightarrow \Sigma^*$ defined as:

- $\text{append}(\lambda, x) = x$
- $\text{append}(aw, x) = a(\text{append}(w, x))$

Prove:

For all $w, v \in \Sigma^*$ and $x \in \Sigma$:

$$\text{append}(wv, x) = w(\text{append}(v, x))$$

Example 2: Induction on Σ^*

Theorem

For all $w, v \in \Sigma^$ and $x \in \Sigma$: $\text{append}(wv, x) = w(\text{append}(v, x))$.*

Proof: By induction on w ...

[B]	$\text{append}(\lambda v, x) = \text{append}(v, x)$	(concat.B)
[I]	$\text{append}((aw)v, x) = \text{append}(a(wv), x)$	(concat.I)
	$= a \text{ append}(wv, x)$	(append.I)
	$= a (w \text{ append}(v, x))$	(IH)
	$= (aw) \text{ append}(v, x)$	(concat.I)

Example 3: Induction on Σ^*

Define $\text{rev} : \Sigma^* \rightarrow \Sigma^*$:

(rev.B) $\text{rev}(\lambda) = \lambda$,

(rev.I) $\text{rev}(a \cdot w) = \text{append}(\text{reverse}(w), a)$

Example 3: Induction on Σ^*

Theorem

For all $w, v \in \Sigma^$, $\text{reverse}(wv) = \text{reverse}(v) \cdot \text{reverse}(w)$.*

Proof: By induction on w ...

$$\begin{aligned} \text{[B]} \quad \text{rev}(\lambda v) &= \text{rev}(v) && (\text{concat.B}) \\ &= \text{rev}(v)\lambda && (*) \\ &= \text{rev}(v)\text{rev}(\lambda) && (\text{rev.B}) \end{aligned}$$

$$\begin{aligned} \text{[I]} \quad \text{rev}((aw')v) &= \text{rev}(a(w'v)) && (\text{concat.I}) \\ &= \text{append}(\text{rev}(w'v), a) && (\text{rev.I}) \\ &= \text{append}(\text{rev}(v)\text{rev}(w'), a) && (\text{IH}) \\ &= \text{rev}(v)\text{append}(\text{rev}(w'), a) && (\text{Example 2}) \\ &= \text{rev}(v)\text{rev}(aw') && (\text{rev.I}) \end{aligned}$$

Example 4: Induction on more complex structures

Recall expressions in the Proof assistant:

- (B) $A, B, \dots, Z, a, b, \dots, z$ are expressions
- (B) \emptyset and \mathcal{U} are expressions
- (R) If E is an expression then so is (E) and E^c
- (R) If E_1 and E_2 are expressions then:
 - $(E_1 \cup E_2)$,
 - $(E_1 \cap E_2)$,
 - $(E_1 \setminus E_2)$, and
 - $(E_1 \oplus E_2)$ are expressions.

Theorem

In any valid expression, the number of (equals the number of)

Proof: By induction on the structure of E ...

Exercise

Exercise

RW: 4.4.2 Define $s_1 = 1$ and $s_{n+1} = \frac{1}{1+s_n}$ for $n \geq 1$

Then $s_1 = 1$, $s_2 = \frac{1}{2}$, $s_3 = \frac{2}{3}$, $s_4 = \frac{3}{5}$, $s_5 = \frac{5}{8}$, \dots

The numbers in numerator and denominator remind one of the Fibonacci sequence.

Prove by induction that

$$s_n = \frac{\text{FIB}(n)}{\text{FIB}(n+1)}$$

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Weekly Feedback

I would appreciate any comments/suggestions/requests you have on this week's lectures.



<https://forms.office.com/r/xKKrxYMRn9>