

SYNTHETIC GUARDED DOMAIN THEORY AND CLASSIFYING TOPOI

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1. INTRODUCTION

This is an internal analysis on synthetic guarded domain theory [1]. Let \mathbf{W} be the theory of filters on the poset $\omega = (\mathbf{N}, \leq) = \{0 < 1 < \dots\}$. For convenience, we will in fact use the poset $\omega_\perp = \{\perp < 0 < 1 < \dots\}$. Explicitly, \mathbf{W} has propositional constants p_\perp, p_0, p_1, \dots , with axioms:

- (1) $p_\perp \vdash \perp$;
- (2) For all $n \leq m \in \omega_\perp$, $p_n \vdash p_m$;
- (3) $\top \vdash \bigvee_{n \in \omega} p_n$.

We work internally in the classifying topos $\mathbf{Set}[\mathbf{W}] \simeq \mathbf{Psh}(\omega)$. We will use \mathbf{Prop} to denote the type of propositions. We will also assume a univalent type universe \mathbf{Type} . We will use ω_\perp to denote the poset internalised in $\mathbf{Set}[\mathbf{W}]$ as well. For any $n : \omega_\perp$, we use $\llbracket n \rrbracket$ to denote the representable proposition.

Axiom (QC). For any $n, m : \omega_\perp$,

$$\llbracket n \rrbracket \rightarrow \llbracket m \rrbracket = \begin{cases} \llbracket m \rrbracket & m < n \\ 1 & \text{otherwise} \end{cases}$$

Axiom (GR). Every proposition is definable: $\forall p : \mathbf{Prop}. \exists n : \omega_\perp. p = \llbracket n \rrbracket$.

Remark 1.1. Notice that (GR) plus (QC) implies the sequence of propositions is a model of \mathbf{W} . By (GR), $\emptyset = \llbracket n \rrbracket$ for some $n : \omega$, and by (QC) $\llbracket \perp \rrbracket \rightarrow \llbracket n \rrbracket = 1$, which implies $\llbracket \perp \rrbracket \rightarrow \emptyset$, thus $\llbracket \perp \rrbracket = \emptyset$. By (GR) again, $\exists n : \omega_\perp. 1 = \llbracket n \rrbracket$, while $1 = \llbracket n \rrbracket$ is equivalent to $\llbracket n \rrbracket$, thus this implies $\bigvee_{n : \omega_\perp} \llbracket n \rrbracket$.

As an easy consequence, we can show the generic model is indeed generic in the following sense:

Lemma 1.1. For any $n : \omega$, $\neg\neg \llbracket n \rrbracket$.

Proof. By (QC), given $n : \omega$, $\llbracket n \rrbracket \rightarrow \llbracket \perp \rrbracket = \llbracket \perp \rrbracket = \emptyset$. Equivalently, $\neg\neg \llbracket n \rrbracket$. \square

2. INTERNAL LATER MODALITY ON PROPOSITIONS

We can construct an internal later modality from the data of a certain endo-function on ω_\perp :

Construction 2.1 (Later modality on propositions). Let $\theta : \omega_\perp \rightarrow \omega_\perp$ be a monotone, non-decreasing function. Then the θ -later modality $\square_\theta : \mathbf{Prop} \rightarrow \mathbf{Prop}$ takes $p : \mathbf{Prop}$ to $\llbracket \theta n \rrbracket$, where $p = \llbracket n \rrbracket$ by (GR).

We first show this construction is well-defined.

Proposition 2.1. $\square_\theta : \mathbf{Prop} \rightarrow \mathbf{Prop}$ is well-defined, i.e. we have a diagram

$$\begin{array}{ccc} \omega_\perp & \xrightarrow{\theta} & \omega_\perp \\ \downarrow & & \downarrow \\ \mathbf{Prop} & \dashrightarrow_{\square_\theta} & \mathbf{Prop} \end{array}$$

Proof. Suppose $p = \llbracket n \rrbracket = \llbracket m \rrbracket$ for some $n, m : \omega_\perp$, and we need to show $\llbracket \theta n \rrbracket = \llbracket \theta m \rrbracket$ as well. We may assume $m < n$, since the case $m = n$ is trivial. By (QC), $\llbracket n \rrbracket \rightarrow \llbracket m \rrbracket$ is equivalent to $\llbracket m \rrbracket$, thus $\llbracket m \rrbracket$ holds. Since θ is non-decreasing, by (QC) $\llbracket \theta m \rrbracket$ holds as well, which by monotonicity of θ , $\llbracket \theta n \rrbracket$ holds as well. Thus this shows $\llbracket \theta n \rrbracket = \llbracket \theta m \rrbracket$. \square

From now on, fix a monotone and non-decreasing function $\theta : \omega_\perp \rightarrow \omega_\perp$.

Proposition 2.2. \square_θ is pointed.

Proof. For any $p : \text{Prop}$, by (GR) let $p = \llbracket n \rrbracket$. By construction, $\square_\theta p = \llbracket \theta n \rrbracket$, which by θ being non-decreasing we have $p \rightarrow \square_\theta p$. \square

Proposition 2.3. \square_θ is left exact.

Proof. For 1, by (GR) let $1 = \llbracket n \rrbracket$. Then since θ is non-decreasing, $\square_\theta 1 = \llbracket \theta n \rrbracket$ holds as well. Preserving binary meets is easy. \square

Theorem 2.1. If θ is strictly increasing, then \square_θ satisfies Löb induction,

$$\forall p : \text{Prop}. (\square_\theta p \rightarrow p) \rightarrow p.$$

Proof. Take $p : \text{Prop}$, and by (GR) let $p = \llbracket n \rrbracket$. Assume $\square_\theta p \rightarrow p$, viz. $\llbracket \theta n \rrbracket \rightarrow \llbracket n \rrbracket$. Since θ is strictly increasing, by (QC) this implication is equivalent to $\llbracket n \rrbracket$, viz. p , holds. \square

3. INTERNAL LATER MODALITY ON TYPES

For any type $X : \text{Type}$, we need a way to relate the function spaces $X^{\llbracket n \rrbracket}$ for $n : \omega_\perp$ to X itself. It turns out we can write any type as a coend:

Theorem 3.1. For any $X : \text{Type}$, it is isomorphic to the coend

$$X \cong \int^{n : \omega_\perp} \llbracket n \rrbracket \times X^{\llbracket n \rrbracket}.$$

Proof. This is a consequence of (GR), where we can assume $1 = \llbracket n \rrbracket$ for some $n : \omega_\perp$. This way, the coend becomes constant, thus is isomorphic to X . \square

Construction 3.1 (Later modality on types). Let $\theta : \omega_\perp \rightarrow \omega_\perp$ be a monotone, non-decreasing map on ω_\perp . The θ -later modality $\square_\theta : \text{Type} \rightarrow \text{Type}$ is constructed as follows: For any $X : \text{Type}$,

$$\square_\theta X := \int^{n : \omega_\perp} \llbracket \theta n \rrbracket \times X^{\llbracket n \rrbracket}.$$

Again, let us fix a monotone and non-decreasing map $\theta : \omega_\perp \rightarrow \omega_\perp$.

Proposition 3.1. The modality $\square_\theta : \text{Type} \rightarrow \text{Type}$ given in Construction 3.1 restricts to the modality $\square_\theta : \text{Prop} \rightarrow \text{Prop}$ on propositions given in Construction 2.1.

Proof. By (GR), it suffices to show $\square_\theta \llbracket n \rrbracket$, according to Construction 3.1, will be equal to $\llbracket \theta n \rrbracket$ for all $n : \omega_\perp$. By construction and (QC), we have

$$\begin{aligned} \square_\theta \llbracket n \rrbracket &\equiv \int^{m : \omega_\perp} \llbracket \theta m \rrbracket \times \llbracket n \rrbracket^{\llbracket m \rrbracket} \\ &= \bigvee_{m : \omega_\perp} \llbracket \theta m \rrbracket \wedge (\llbracket m \rrbracket \rightarrow \llbracket n \rrbracket) \\ &= \bigvee_{m \leq n} \llbracket \theta m \rrbracket \vee \bigvee_{m > n} \llbracket \theta m \rrbracket \wedge \llbracket n \rrbracket \\ &= \llbracket \theta n \rrbracket \end{aligned}$$

The last step holds by the fact that θ is monotone and non-decreasing, thus $\llbracket \theta m \rrbracket \wedge \llbracket n \rrbracket = \llbracket n \rrbracket$ for all $m > n$. \square

Proposition 3.2. \square_θ is left exact.

Proof. \square_θ preserves 1 by Proposition 3.1 and Proposition 2.3. For pullback, suppose we have $f : Y \rightarrow X$ and $g : Z \rightarrow X$. We then have

$$\begin{aligned} \square_\theta(Y \times_X Y) &\equiv \int^{n:\omega_1} [\theta n] \times (Y \times_X Z)^{[n]} \\ &\cong \int^{n:\omega_1} [\theta n] \times Y^{[n]} \times_{X^{[n]}} Z^{[n]} \\ &\cong \square_\theta Y \times_{\square_\theta X} \square_\theta Z \end{aligned}$$

The final step uses the fact that finite limits commutes with the filtered colimit. \square

Proposition 3.3. \square_θ is well-pointed.

Proof. To construct the unit, by (GR) we may assume $1 = [n]$, thus by the universal property there is a comparison map

$$X \xrightarrow{\cong} [n] \times X^{[n]} \longrightarrow [\theta n] \times X^{[n]} \longrightarrow \square_\theta X.$$

This is well-defined due to the fact that $\square_\theta X$ is a coend, and naturality is evident. Checking well-pointedness is routine. \square

Corollary 3.1. If θ is strictly increasing, then $\square_\theta : \text{Type} \rightarrow \text{Type}$ supports guarded recursion fixed-point construction.

Proof. This follows from Proposition 3.3, Proposition 3.1, and [2]. \square

4. A NON-LOCALIC MODEL OF GUARDED RECURSION

Consider \mathbf{Inj} , the category of finite sets with injections between them. Notice that this category has pullback but not a terminal object. There is an evident functor

$$1 + - : \mathbf{Inj} \rightarrow \mathbf{Inj},$$

which induces an adjunction

$$\begin{array}{ccc} \mathbf{Psh}(\mathbf{Inj}) & \begin{array}{c} \xleftarrow{L} \\ \perp \\ \xrightarrow{R} \end{array} & \mathbf{Psh}(\mathbf{Inj}) \end{array}$$

where L is given by precomposition with $1 + -$ and R is obtained via right Kan extension. Since L is also continuous, this gives us a geometric morphism

$$R : \mathbf{Psh}(\mathbf{Inj}) \rightarrow \mathbf{Psh}(\mathbf{Inj}).$$

Proposition 4.1. We have a self-indexed path family

$$\begin{array}{ccc} \mathbf{Psh}(\mathbf{Inj}) & \begin{array}{c} \xrightarrow{R} \\ \Downarrow \theta \\ \xrightarrow{R} \end{array} & \mathbf{Psh}(\mathbf{Inj}) \end{array}$$

Proof. It suffices to provide a natural transformation $\theta : L \Rightarrow \text{id}$. But such a natural transformation is induced by the natural inclusion $\text{id} \Rightarrow 1 + -$. \square

Theorem 4.1. The self-indexing family in Proposition 4.1 is contractive, which means the internal modality associated to this self-indexing family satisfies Löb induction.

Proof. To show contractivity, it suffices to observe that for any representable $\mathbb{1}_n$, there exists some k that $L^k \mathbb{1}_n \cong \emptyset$. Now we first show that

$$L \mathbb{1}_n \cong n \cdot \mathbb{1}_{n-1}.$$

For any $i \in n$, let \hat{n}^i denote $n - \{i\}$. By construction we have

$$L \mathbb{1}_n(m) \cong \mathbf{Inj}(m+1, n) \cong \sum_{i:n} \mathbf{Inj}(m, \hat{n}^i) \cong n \cdot \mathbb{1}_{n-1}(m).$$

It is worth explicitly specifying the last isomorphism: given a pair $(i, f : m \rightarrowtail n - 1)$, the induced injection $(i, f) : m + 1 \rightarrowtail n$ is given by

$$(i, f)(j) = \begin{cases} f(j) & j < m \wedge f(j) < i \\ f(j) + 1 & j < m \wedge f(j) \geq i \\ i & j = m \end{cases}$$

This isomorphism is natural: Given any $f : l \rightarrowtail m$, we verify

$$\begin{array}{ccc} n \times \mathbf{Inj}(m, n - 1) & \xrightarrow{\cong} & \mathbf{Inj}(m + 1, n) \\ \downarrow n \times (- \circ f) & & \downarrow - \circ (f + 1) \\ n \times \mathbf{Inj}(l, n - 1) & \xrightarrow{\cong} & \mathbf{Inj}(l + 1, n) \end{array}$$

For any pair (i, g) with $i \in m$ and $g : m \rightarrowtail n - 1$, and for $j \in l + 1$ we have

$$\begin{aligned} (i, g)((f + 1)(j)) &= \begin{cases} gf(j) & j < l \wedge gf(j) < i \\ gf(j) + 1 & j < l \wedge gf(j) \geq i \\ i & j = l \end{cases} \\ &= (i, gf)(j) \end{aligned}$$

This shows that the isomorphism is natural, thus $L \circ_n \cong n \cdot \circ_{n-1}$. This way, since L preserves colimits, it follows that

$$L^{n+1} \circ_n \cong \emptyset,$$

which completes the proof. \square

5. ANOTHER INTERESTING EXAMPLE

Consider $\mathbf{Psh}(\mathbf{Fin})$, where \mathbf{Fin} is the category of finite sets, and consider the functor

$$1 + - : \mathbf{Fin} \rightarrow \mathbf{Fin},$$

and consider the functor $L : \mathbf{Psh}(\mathbf{Fin}) \rightarrow \mathbf{Psh}(\mathbf{Fin})$ given by

$$LX(n) \cong X(n + 1).$$

L is cocontinuous, thus in particular has a right adjoint $R : \mathbf{Psh}(\mathbf{Fin}) \rightarrow \mathbf{Psh}(\mathbf{Fin})$. To compute the right adjoint, it suffices to compute what L does on representables.

Lemma 5.1. *For any $n \in \mathbf{Fin}$, we have*

$$L \circ_n \cong n \cdot \circ_n,$$

where $n \cdot -$ denotes the tensor of the finite set n with the presheaf \circ_n .

Proof. By construction, for any $m \in \mathbf{Fin}$,

$$(L \circ_n)(m) \cong n^{m+1} \cong n \times n^m \cong n \times \circ_n(m).$$

This shows that $L \circ_n \cong n \cdot \circ_n$. \square

Proposition 5.1. *For any $X \in \mathbf{Psh}(\mathbf{Fin})$, we have*

$$(RX)(n) \cong X(n)^n,$$

where for any $f : n \rightarrow m$, the action on f takes $a \in X(m)^m$ to the following composite,

$$\begin{array}{ccc} n & \dashrightarrow^{af} & X(n) \\ f \downarrow & & \uparrow - \circ f \\ m & \xrightarrow{a} & X(m) \end{array}$$

Proof. This follows from Lemma 5.1, where by construction

$$(RX)(n) \cong \mathbf{Psh}(\mathbf{Fin})(L \downarrow n, X) \cong \mathbf{Psh}(\mathbf{Fin})(n \cdot \downarrow_n, X) \cong X(n)^n.$$

The action on morphisms in \mathbf{Fin} follows from routine computation. \square

Proposition 5.2. *R is well-pointed.*

Proof. The point is easy to provide, where for any $X \in \mathbf{Psh}(\mathbf{Fin})$, the unit $\eta_X : X \rightarrow RX$ is given by the constant function for any $n \in \mathbf{Fin}$,

$$\eta_X := c_- : X(n) \rightarrow X(n)^n.$$

Naturality is evident. For well-pointedness, this is routine to check. \square

Theorem 5.1. *$R : \mathbf{Psh}(\mathbf{Fin}) \rightarrow \mathbf{Psh}(\mathbf{Fin})$ restricts to identity $\text{id} : \Omega \rightarrow \Omega$ on propositions.*

Proof. For any $n \in \mathbf{Fin}$, let $p \in \Omega(n)$ be a sieve on n in \mathbf{Fin} . Over \downarrow_n , the fibred modality R acting on p is given by the pullback

$$\begin{array}{ccc} \square \downarrow_n p & \longrightarrow & Rp \\ \downarrow & \lrcorner & \downarrow \\ \downarrow_n & \xrightarrow{\eta} & R \downarrow_n \end{array}$$

By construction, for any $f : m \rightarrow n$, $f \in \square \downarrow_n p$ iff $\square \downarrow_m f \in \square \downarrow_n p$. \square

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