

# Queues

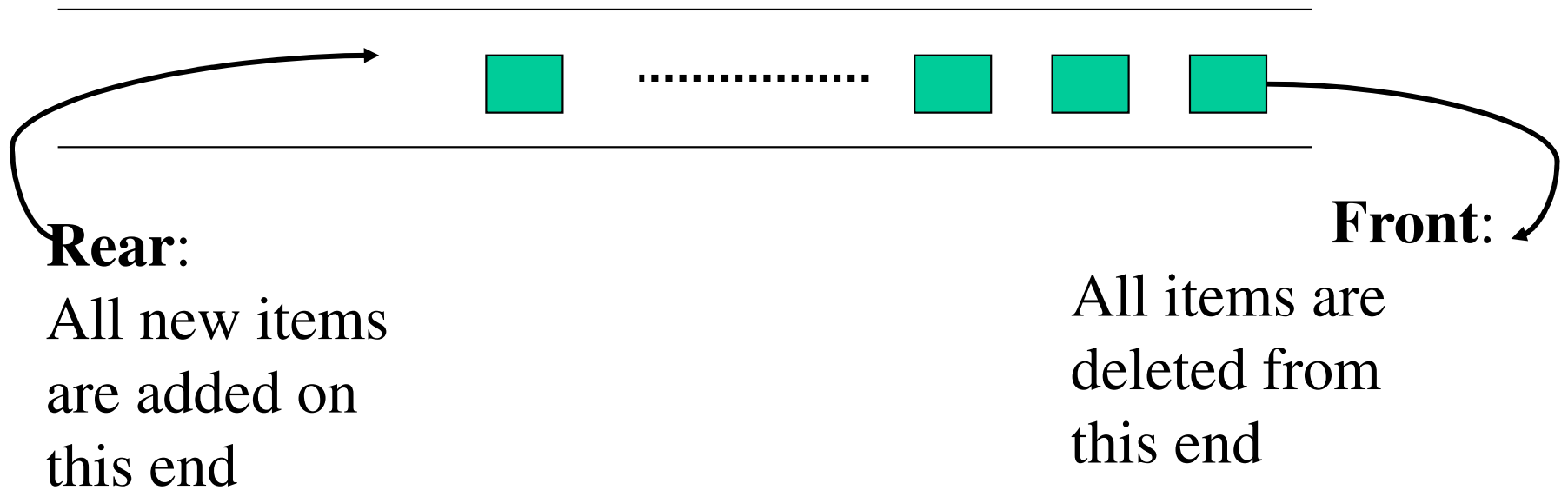
# Definition of a Queue

- A **queue** is a data structure that models/enforces the **first-come first-serve** order, or equivalently the **first-in first-out** (FIFO) order.
- The element that is inserted first into the queue is deleted first
  - The element that is inserted last is deleted last.
- A **waiting line** is a good real-life example of a queue.

# Queues

- Linear list.
- One end is called **front (head)**.
- Other end is called **rear (tail)**.
- Insertions are done at the **rear (tail)** only.
- Deletions are made from the **front(head)** only.

# A Graphic Model of a Queue



# Examples of Queues

- An electronic mailbox is a queue
  - The ordering is chronological (by arrival time)
- A waiting line in a store, at a service counter, on a one-lane road
- Equal-priority processes waiting to run on a processor in a computer system

# Example

Rear

Front

			4	1	3
--	--	--	---	---	---

enqueue(1);

		1	4	1	3
--	--	---	---	---	---

enqueue(5);

	5	1	4	1	3
--	---	---	---	---	---

dequeue();

	5	1	4	1	
--	---	---	---	---	--

dequeue();

	5	1	4		
--	---	---	---	--	--

dequeue();

	5	1			
--	---	---	--	--	--

Given the following Queue, how will it change when we apply the given operations?

# Queue ADT

## instances

ordered list of elements; one end is the front; the other is the rear;

## operations

`empty()`: Return true if queue is empty, return false otherwise

`size()`: Return the number of elements in the queue

`front()`: Return the front element of queue

`rear()`: Return the rear element of queue

`dequeue()`: Remove an element from the queue

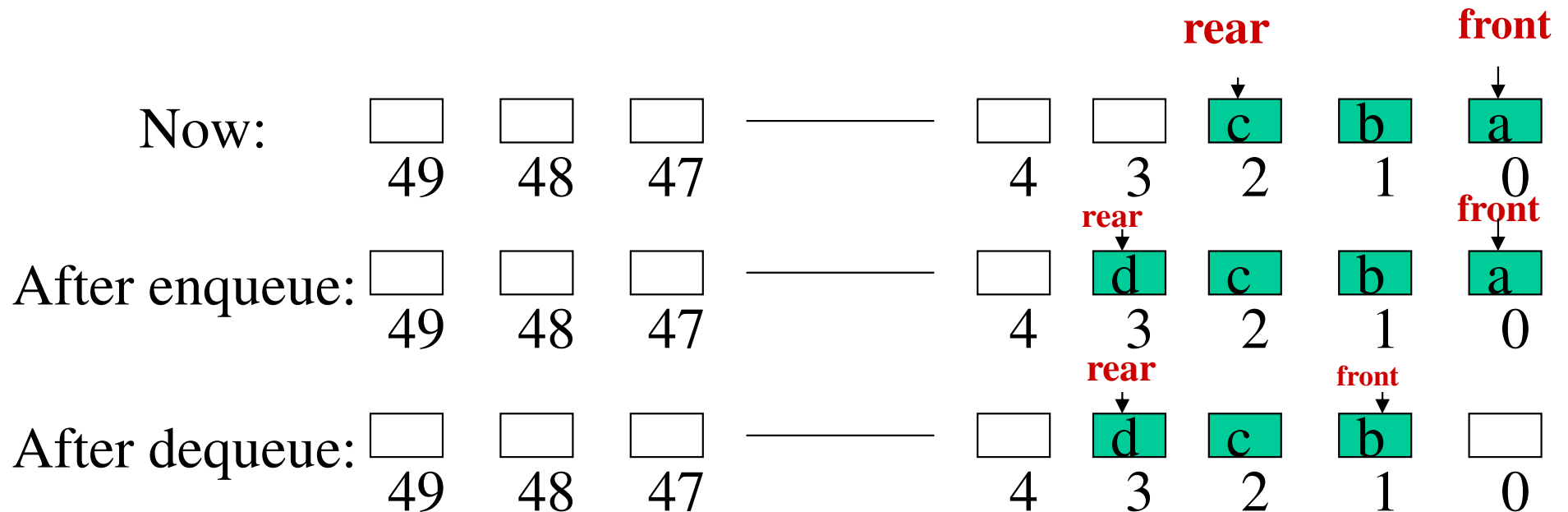
`enqueue(x)`: Add element x to the queue

It is also possible to represent Queues using

1. Array-based representation
2. Linked representation

# How front and rear Change

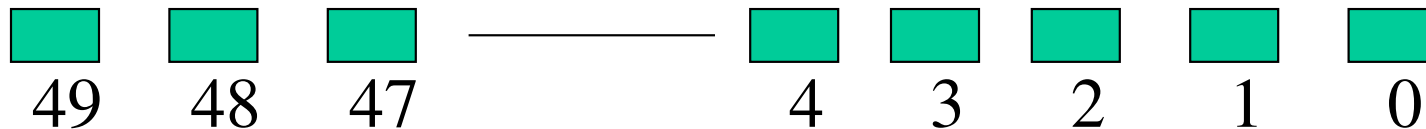
- **front** increases by 1 after each dequeue( )
- **rear** increases by 1 after each enqueue( )



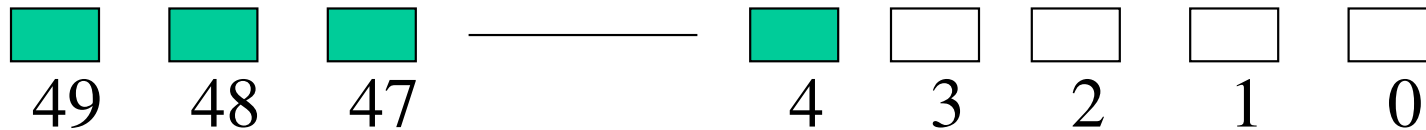


## False-Overflow Issue First

- Suppose 50 calls to enqueue have been made, so now the queue array is full



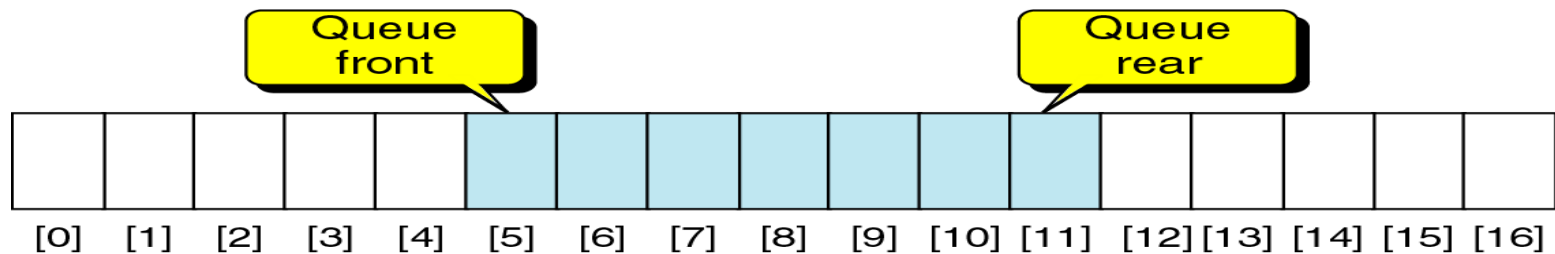
- Assume 4 calls to dequeue( ) are made



- Assume a call to enqueue( ) is made now.
- The rear part have no space, but the front has 4 unused spaces; if never used, they are wasted.

# Circular Queue

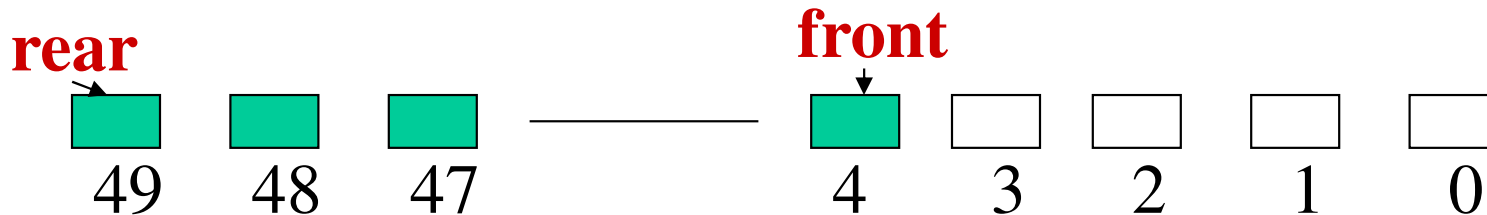
- Use Linear Array to implement a queue.



- **Waste of memory: The deleted elements can not be re-used.**
- Solution: to use circular queue.
- Two implementations:
  - Using  $n-1$  space.
  - Using  $n$  space + full tag

## Solution: A Circular Queue

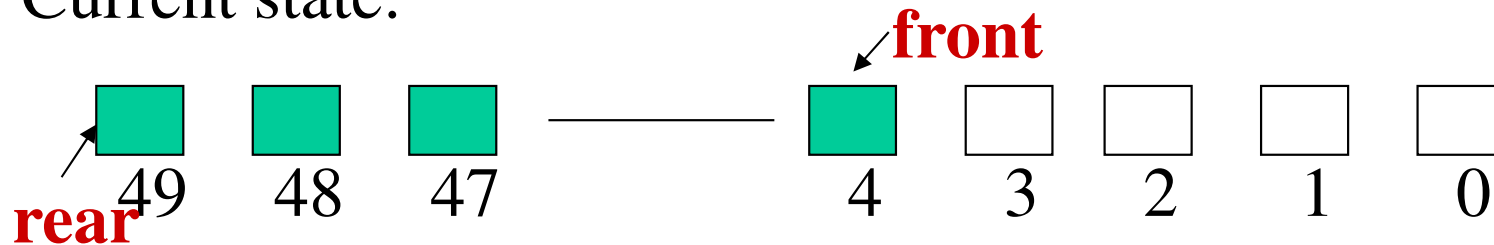
- Allow the front (and the rear) to be moving targets
- When the rear end fills up and front part of the array has empty slots, new insertions should go into the front end



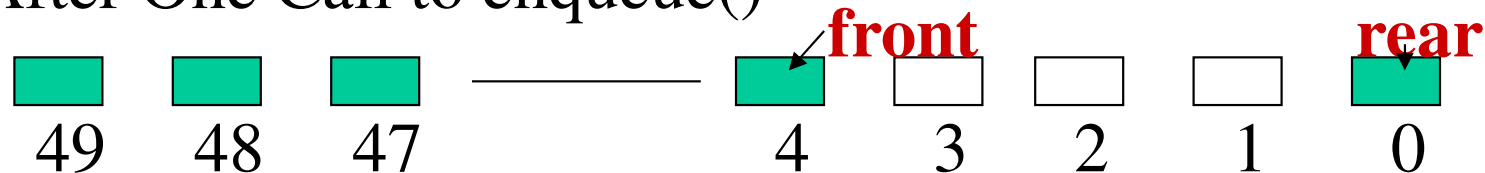
- Next insertion goes into slot 0, and rear tracks it. The insertion after that goes into a lot 1, etc.

# Illustration of Circular Queues

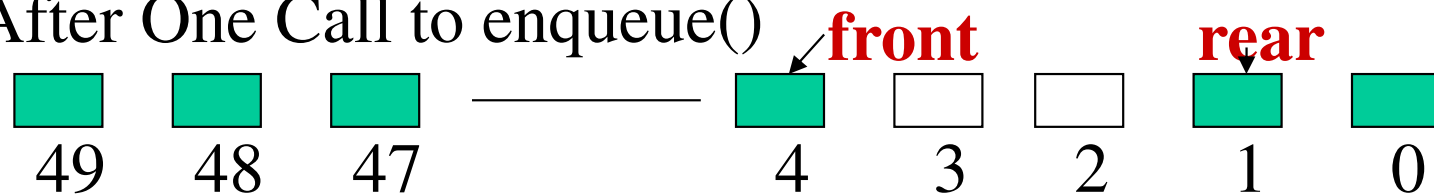
- Current state:



- After One Call to enqueue()



- After One Call to enqueue()



# Numeric for Circular Queues

- **front** increases by (1 modulo capacity) after each dequeue( ):

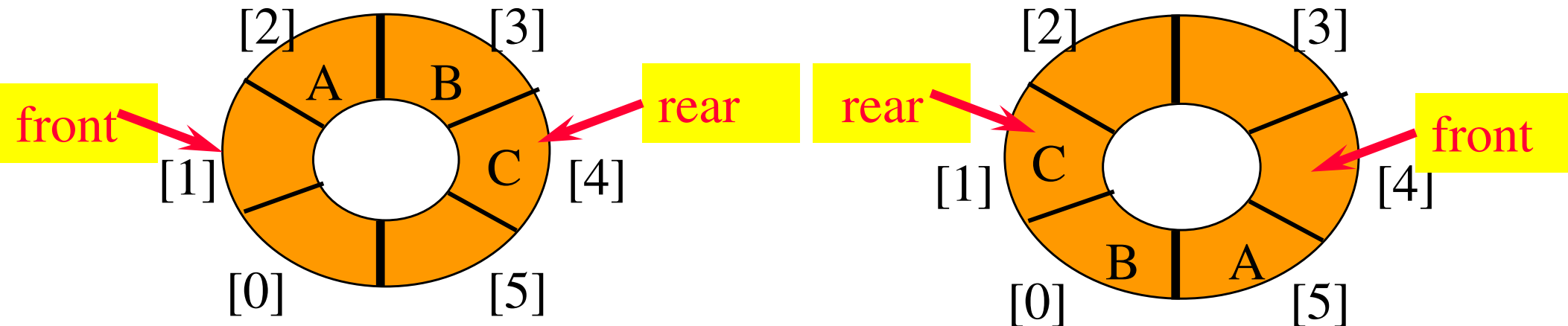
$$\text{front} = (\text{front} + 1) \% \text{capacity};$$

- **rear** increases by (1 modulo capacity) after each enqueue( ):

$$\text{rear} = (\text{rear} + 1) \% \text{capacity};$$

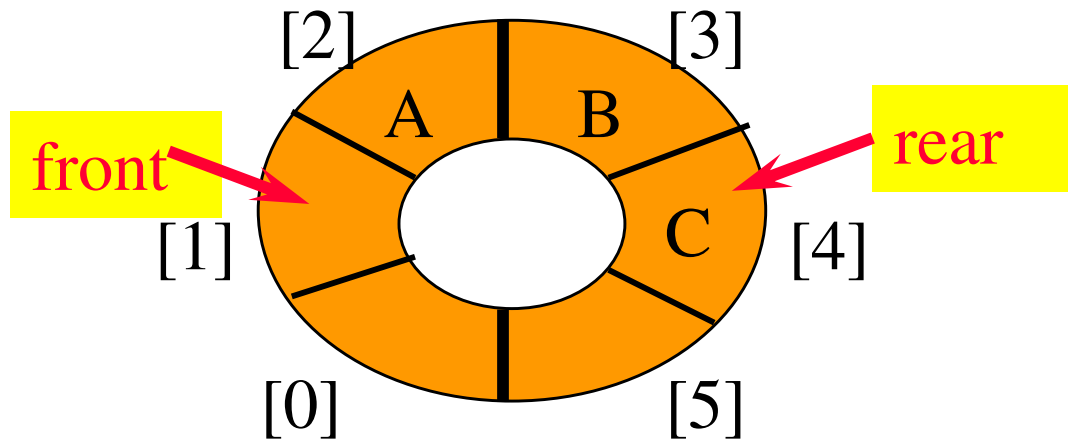
## Yet another Illustration of Circular Queues(Using n-1 space)

- Use integer variables **front** and **rear**.
  - **front** is one position counterclockwise from first element
  - **rear** gives position of last element



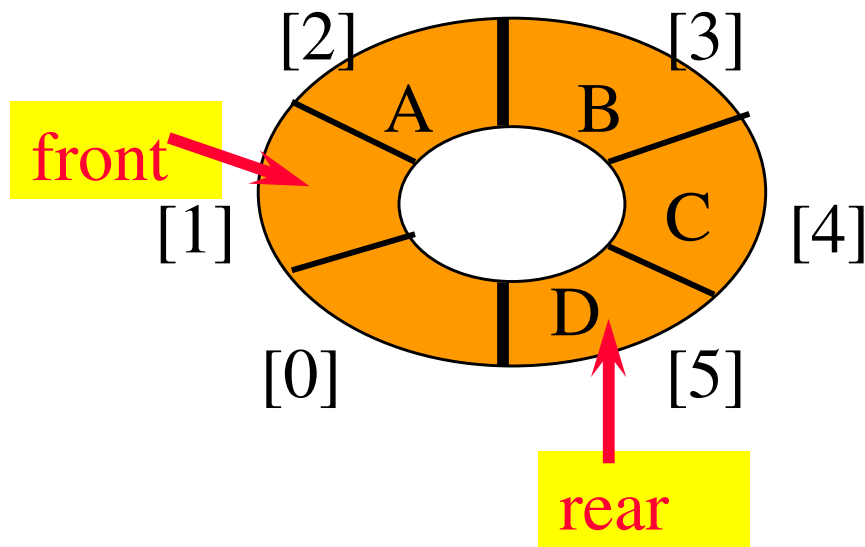
# Add An Element

- Move **rear** one clockwise.



# Add An Element

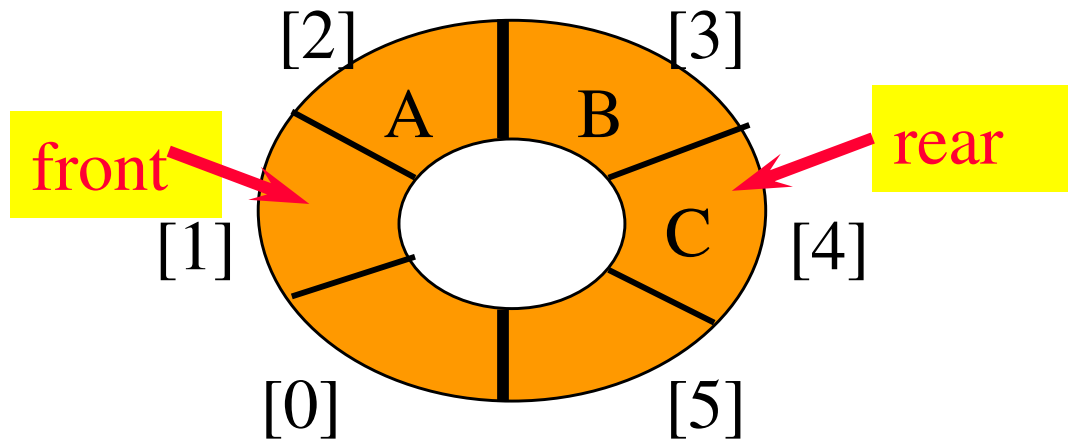
- Move **rear** one clockwise.
- Then put into **queue[rear]**.





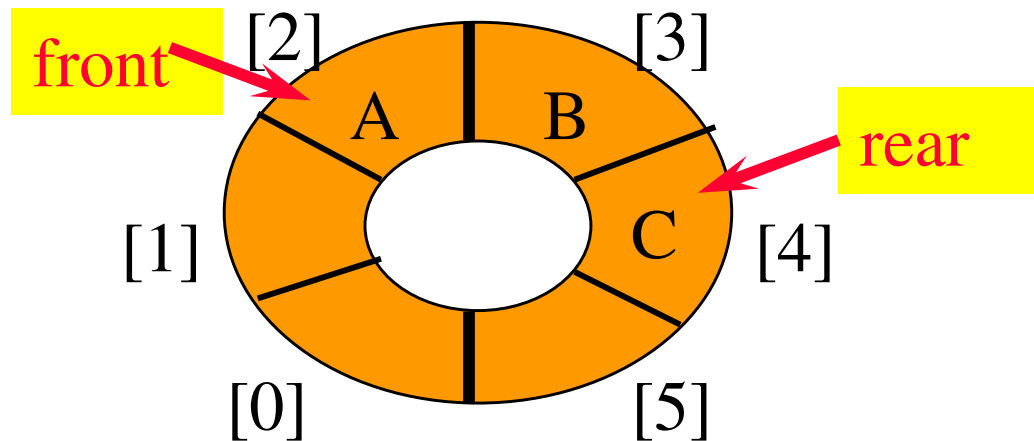
# Remove An Element

- Move **front** one clockwise.



# Remove An Element

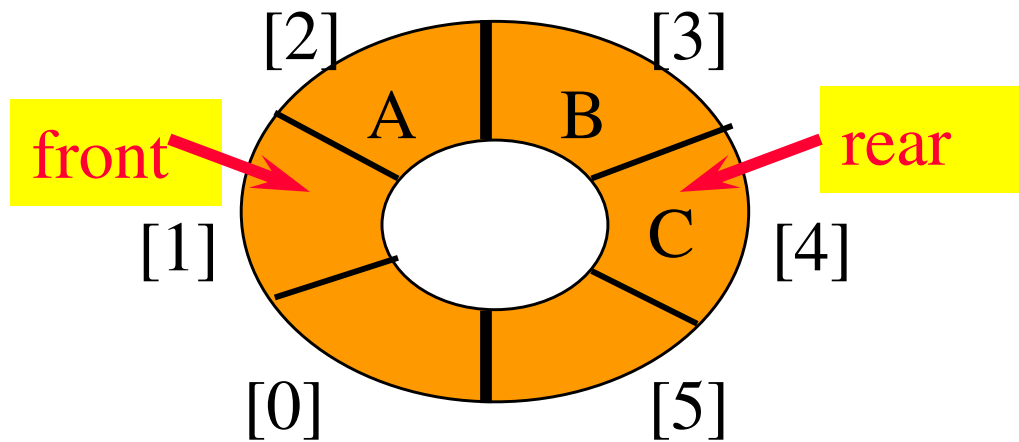
- Move **front** one clockwise.
- Then extract from **queue[front]**.



## Moving rear Clockwise

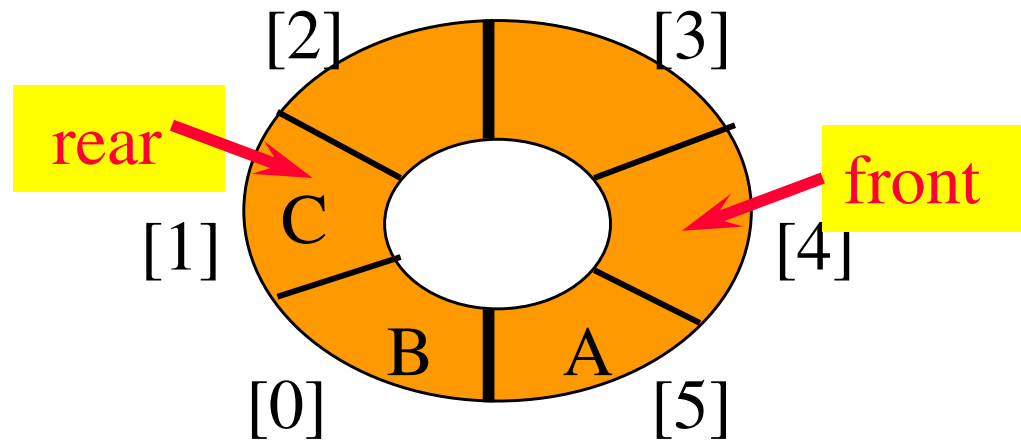
- `rear++;`

`if (rear == queue.length) rear = 0;`

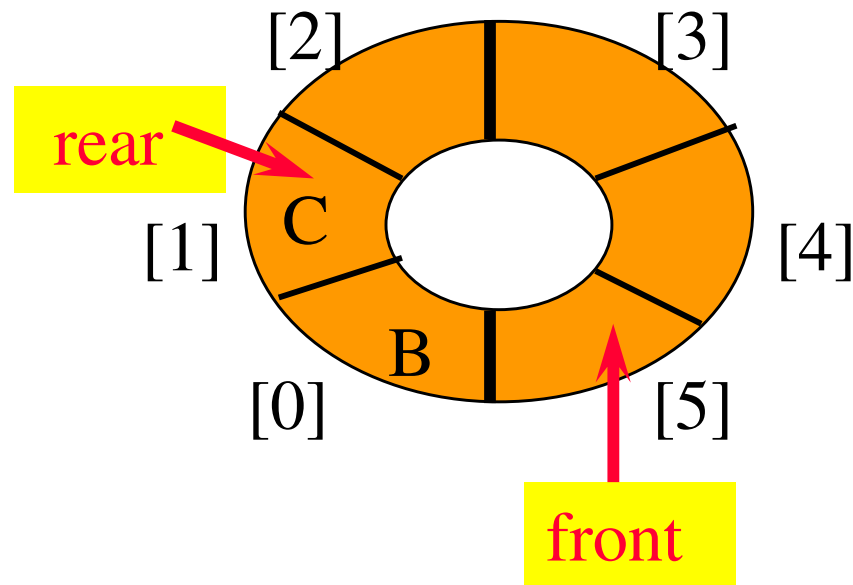


- `rear = (rear + 1) % queue.length;`

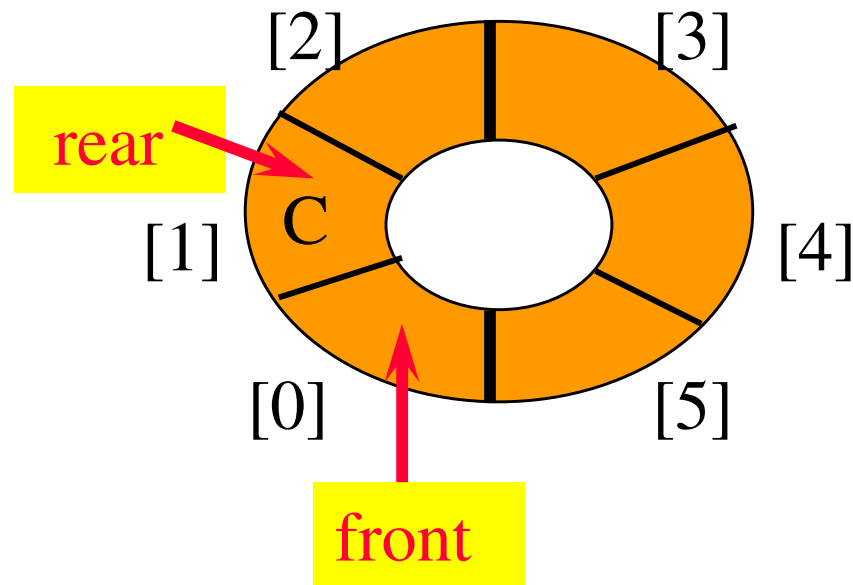
# Empty Queue



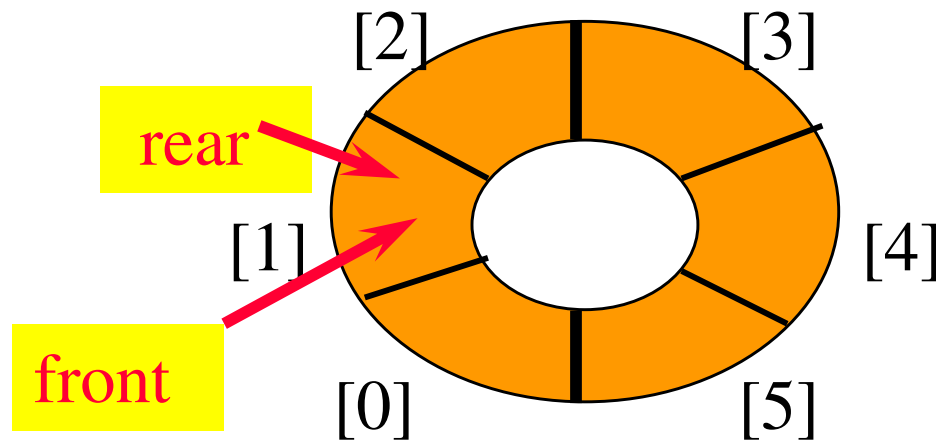
# Empty Queue



# Empty Queue

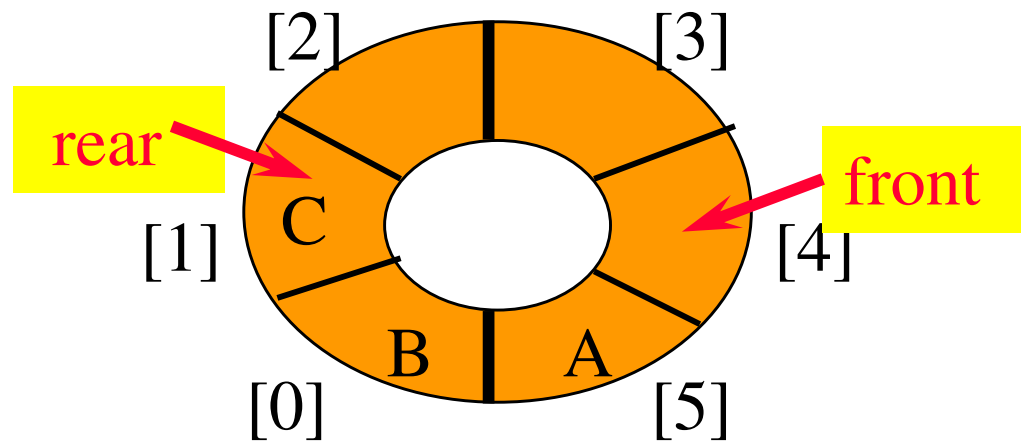


# Empty Queue



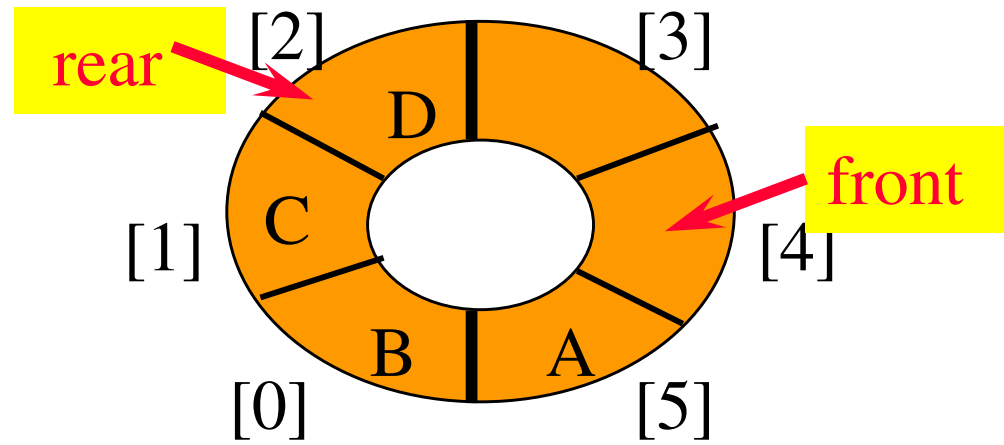
- When a series of removes causes the queue to become empty, **front = rear**.
- When a queue is constructed, it is empty.
- So initialize **front = rear = 0**.

# Full Queue

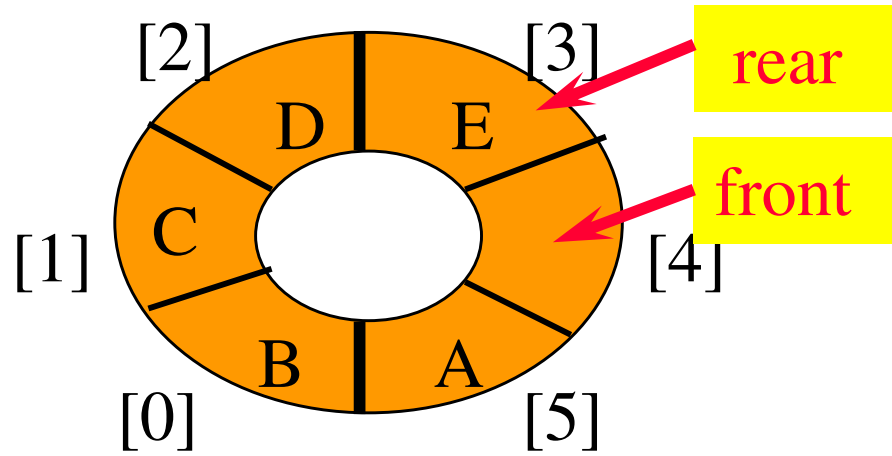




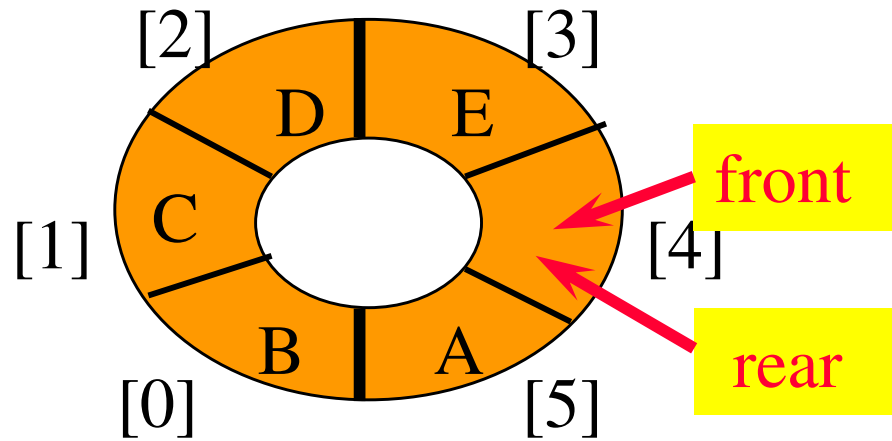
# Full Queue



# Full Queue



## Full Queue



- When a series of enqueue causes the queue to become full,  $\text{front} = \text{rear}$ .
- So, we cannot distinguish between a full queue and an empty queue!

# Implementation of Circular queue with (n-1) space used

## ■ Create(Q)

Q: Array[0...n-1]

front = rear = 0 //initialize

## ■ Enqueue(item, Q) $\Rightarrow$ Queue

```
{ rear = (rear+1) mod n; //rear moves forward;
```

```
if( rear == front){
```

```
    QueueFull;    // Queue is full.
```

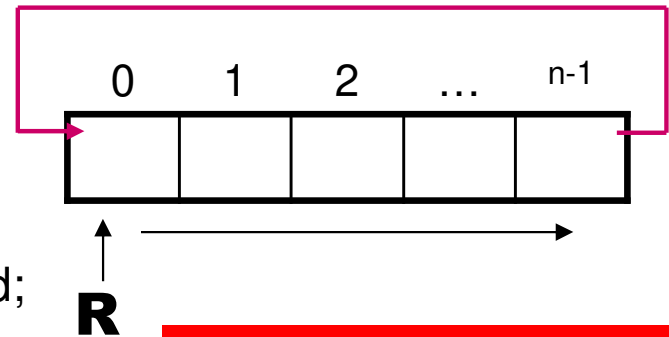
```
    rear = (rear-1) mod n; // rear back to the previous position;
```

```
}
```

```
else
```

```
    Q[rear]=item;
```

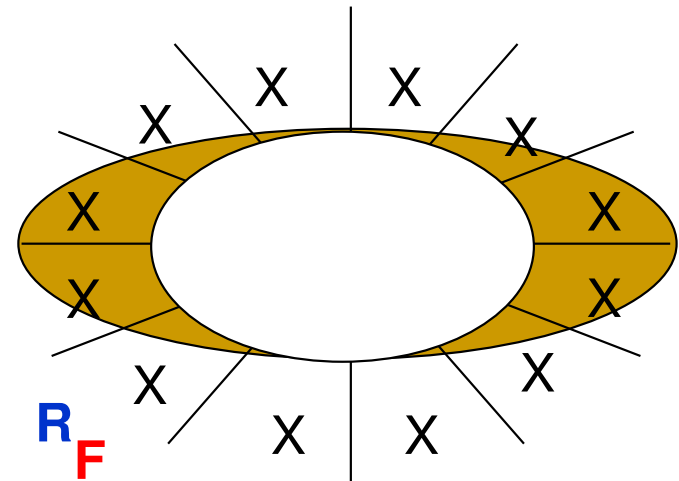
```
}
```



$\Rightarrow \text{Rear} = (\text{Rear}+1) \bmod n$

# Implementation of Circular queue with (n-1) space used

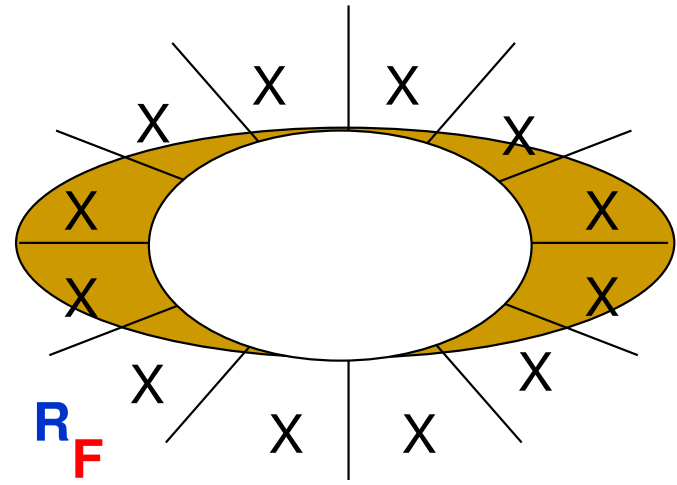
- Dequeue(Q)  $\Rightarrow$  item  
{  
  if( front==rear)  
    QueueEmpty;  
  else{  
    front = (front+1) mod n;  
    item = Q[front];  
  }  
}
- Note: only (n-1) space used;



# Implementation of Circular Queue with n space used

- A parameter “Tag” is introduced to help to make sure the queue is Empty or Full:
  - Boolean
  - If Tag = True, combined with other conditions => queue is Full
  - If Tag = False, combined with other conditions => queue is Null
  - **“Tag” can determine the states of the queue solely!**

# Implementation of Circular Queue with n space used



# Implementation of Circular Queue with n space used

- **Create(Q)**

Q: Array[0...n-1]

int front = rear = 0

Boolean Tag = False (0)

- **Enqueue(item, Q)  $\Rightarrow$  Queue**

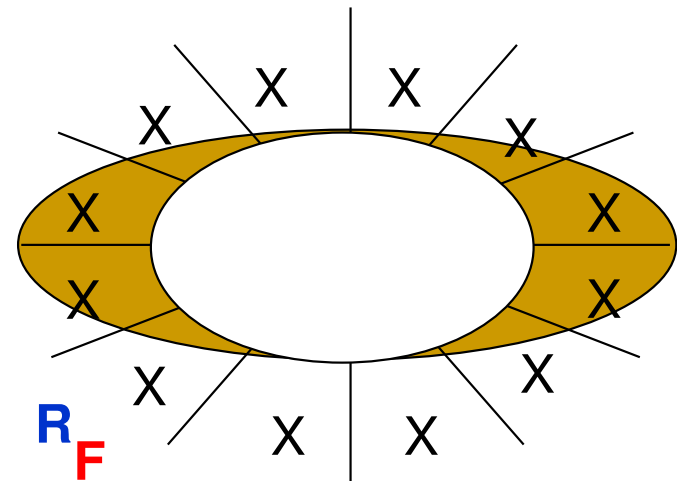
```
{  
    if (rear == front && Tag == True)  
        Overflow message: QueueFull;  
    else {  
        rear = (rear+1) mod n; //rear moves forward;  
        Q[rear]=item;  
        if (rear==front)  
            Tag=True;  
    }  
}
```



# Implementation of Circular Queue with n space used

## ■ Dequeue(Q) $\Rightarrow$ item

```
{  
    if (front==rear && Tag==False)  
        Underflow message: QueueEmpty;  
    else {  
        front = (front+1) mod n;  
        item = Q[front];  
        if (front==rear)  
            Tag=False;  
    }  
}
```



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Slides and figures have been collected from various publicly available Internet sources for preparing the lecture slides of IT2001 course. I acknowledge and thank all the original authors for their contribution to prepare the content.

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