

INF1004 L2 Linear and Affine Equation Systems

Alexander Binder

SIT

February 25, 2023

- 1 Recap
- 2 Linear equation Systems
- 3 Affine equation systems
- 4 Solving affine systems with transformations on rows
- 5 Extending a set of k independent vectors to a basis of the whole space \mathbb{R}^d

- is the product of vector norms times the cosine of the angle between the two vectors
- a similarity measure

Inner product

Let be $u \in \mathbb{R}^d$, $v \in \mathbb{R}^d$, **Definition:** Then the inner product between u and v is defined as:

$$u \cdot v = \sum_{k=1}^d u_k v_k$$

$$[1, -3, 5, 2] \cdot [2, 3, 4, -1] = 1 * 2 + (-3 * 3) + 5 * 4 + 2 * (-1) = 11$$

Interpretation of the inner product

It holds for the inner product defined above that:

$$u \cdot v = \|u\|_2 \|v\|_2 \cos(\angle(u, v))$$

it is the product of the euclidean length of u , of v and the cosine of the angle between these two vectors.

The canonical inner product defines the euclidean norm via:

$$\|v\|_2 = \sqrt{v \cdot v}$$

$$\|[1, -3, 5, 2]\|_2 = \sqrt{1 + (-3)^2 + 5^2 + 2^2} = \sqrt{39}$$

- A way to measure length of vectors
- Euclidean norm is well known.

$$\|v\|_2 = \left(\sum_{k=1}^d v_k^2 \right)^{1/2}$$

Other norms are in use, as well:

$$\|v\|_p = \left(\sum_{k=1}^d |v_k|^p \right)^{1/p}$$

Difference to the Euclidean norm: Replace 2 by p , use $|v_k|$ for a vector component

Properties of any norm

- $\|0\| = 0$
- $\|\lambda v\| = |\lambda| \|v\|$
- $\|v + u\| \leq \|v\| + \|u\|$
- a norm induces a distance measure $d(\cdot, \cdot)$ via $d(u, v) = \|u - v\|$

Projecting onto a vector and removing the direction of a vector

Project a vector x onto a vector v :

$$x_{\parallel v} = \frac{x \cdot v}{v \cdot v} v = \left(x \cdot \frac{v}{\|v\|_2} \right) \frac{v}{\|v\|_2}$$

Remove from a vector x a vector v :

$$x_{\perp v} = x - x_{\parallel v}$$

Definition of a vector space

The vector space spanned by a set of vectors $w^{\{0\}}, \dots, w^{\{k-1\}}$ is the set of all their linear combinations using all possible real numbers $a_0, \dots, a_{k-1} \in \mathbb{R}$

$V = \{w \text{ such that}$

$$w = \sum_{r=0}^k a_r w^{\{r\}}, a_0, \dots, a_{k-1} \in \mathbb{R}\}$$

The dimension of this set is the largest number of independent vectors which we can obtain from $w^{\{0\}}, \dots, w^{\{k-1\}}$.

Given a matrix of shape (n, d) and a vector of dimensionality d , the multiplication Ax of them (with x being on the right hand side) is defined as a vector of length n such that

$$\begin{aligned} Ax &= \begin{bmatrix} a_{0,0} & \dots & a_{0,d-1} \\ a_{1,0} & \dots & a_{1,d-1} \\ \vdots & \ddots & \vdots \\ a_{n-1,0} & \dots & a_{n-1,d-1} \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \\ \dots \\ x_{d-1} \end{bmatrix} \\ &= \begin{bmatrix} \begin{bmatrix} a_{0,0} & \dots & a_{0,d-1} \end{bmatrix} \cdot x \\ \begin{bmatrix} a_{1,0} & \dots & a_{1,d-1} \end{bmatrix} \cdot x \\ \vdots \\ \begin{bmatrix} a_{n-1,0} & \dots & a_{n-1,d-1} \end{bmatrix} \cdot x \end{bmatrix} \end{aligned}$$

Thus, Ax is a vector and the k -th component of vector Ax is given as an inner product

$$\begin{aligned}(Ax)_k &= A[k, :] \cdot x \\ &= (a_{k,0} \dots a_{k,d-1}) \cdot x \\ &= \sum_{r=0}^{d-1} a_{k,r} x_r\end{aligned}$$

between the k -th row of A and vector x . The matrix multiplication with a vector from the right is only defined, if the number of columns in A is equal to the dimensionality of the vector x .

Given a matrix of shape (n, d) and a vector of dimensionality n , the multiplication $x^\top A$ of them (with x^\top being on the left hand side) is defined as a vector of length d such that

$$\begin{aligned}
 x^\top A &= \begin{bmatrix} x_0, x_1, \dots, x_{n-1} \end{bmatrix} \begin{bmatrix} a_{0,0} & \dots & a_{0,d-1} \\ a_{1,0} & \dots & a_{1,d-1} \\ \vdots & \ddots & \vdots \\ a_{n-1,0} & \dots & a_{n-1,d-1} \end{bmatrix} \\
 &= \begin{bmatrix} x \cdot \begin{bmatrix} a_{0,0} \\ \vdots \\ a_{n-1,0} \end{bmatrix} & x \cdot \begin{bmatrix} a_{0,1} \\ \vdots \\ a_{n-1,1} \end{bmatrix} & \dots & x \cdot \begin{bmatrix} a_{0,d-1} \\ \vdots \\ a_{n-1,d-1} \end{bmatrix} \end{bmatrix}
 \end{aligned}$$

Thus, $x^\top A$ is a vector and the k -th component of vector $x^\top A$ is given as an inner product

$$\begin{aligned}(x^\top A)_k &= x \cdot A[:, k] \\ &= (a_{0,k} \dots a_{n-1,k}) \cdot x \\ &= \sum_{r=0}^{n-1} x_r a_{r,k}\end{aligned}$$

Consequence:

inner products as matrix vector multiplication

The inner product $x \cdot y$ of two column-shaped vectors can be written in matrix-vector multiplication notation as

$$x \cdot y = [x_0, x_1, \dots, x_{d-1}] \begin{bmatrix} y_0 \\ y_1 \\ \vdots \\ y_{d-1} \end{bmatrix} = x^\top y$$

where x^\top is the transpose of a vector or matrix x .

The transpose was used here to convert the column-shaped vector into a row-shaped vector.

Given a matrix A of shape (n, d) and a matrix B of shape (d, f) , their multiplication AB is defined as a matrix of shape (n, f) and its component $(AB)_{i,k}$ at row i and column k is given as

$$(AB)_{i,k} = A_{(i,:)} \cdot B_{(:,k)} = \sum_{r=1}^d A_{i,r} B_{rk}$$

as an inner product between the i -th row of the left matrix and the k -th column of the right matrix.

Important: the number or dimensions in the second axis of A must be equal to the number or dimensions in the first axis of B . Otherwise matrix multiplication is not possible.

Therefore a matrix- matrix multiplication is a matrix consisting of inner products:

$$AB = \begin{bmatrix} A_{(0,:)} \cdot B_{(:,0)} & A_{(0,:)} \cdot B_{(:,1)} & A_{(0,:)} \cdot B_{(:,2)} & \dots & A_{(0,:)} \cdot B_{(:,f-1)} \\ A_{(1,:)} \cdot B_{(:,0)} & A_{(1,:)} \cdot B_{(:,1)} & A_{(1,:)} \cdot B_{(:,2)} & \dots & A_{(1,:)} \cdot B_{(:,f-1)} \\ A_{(2,:)} \cdot B_{(:,0)} & A_{(2,:)} \cdot B_{(:,1)} & A_{(2,:)} \cdot B_{(:,2)} & \dots & A_{(2,:)} \cdot B_{(:,f-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ A_{(n-1,:)} \cdot B_{(:,0)} & A_{(n-1,:)} \cdot B_{(:,1)} & A_{(n-1,:)} \cdot B_{(:,2)} & \dots & A_{(n-1,:)} \cdot B_{(:,f-1)} \end{bmatrix}$$

Memorizing

- Matrix-Matrix multiplication results in a matrix, if shapes are permissible:
 $(n, d)(d, f) \rightarrow (n, f)$.
- The component $(AB)_{i,k}$ at row i and column k is given as the inner product between row i of the left matrix and column k of the right matrix.
That also tells you which axes must match in dimensionality: left matrix - the number of columns = the dimensionality of the second axis. right matrix - the number of rows = the dimensionality of the first axis.

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & -1 \\ -1 & 0 & 6 \end{bmatrix} \begin{bmatrix} 5 & -1 & -2 \\ 3 & 9 & -3 \\ -2 & 1 & 2 \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & -1 \\ -1 & 0 & 6 \end{bmatrix} \begin{bmatrix} 5 & -1 & -2 \\ 3 & 9 & -3 \\ -2 & 1 & 2 \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & -1 \\ -1 & 0 & 6 \end{bmatrix} \begin{bmatrix} 5 & -1 & -2 \\ 3 & 9 & -3 \\ -2 & 1 & 2 \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & -1 \\ -1 & 0 & 6 \end{bmatrix} \begin{bmatrix} 5 & -1 & -2 \\ 3 & 9 & -3 \\ -2 & 1 & 2 \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & -1 \\ -1 & 0 & 6 \end{bmatrix} \begin{bmatrix} 5 & -1 & -2 \\ 3 & 9 & -3 \\ -2 & 1 & 2 \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & -1 \\ -1 & 0 & 6 \end{bmatrix} \begin{bmatrix} 5 & -1 & -2 \\ 3 & 9 & -3 \\ -2 & 1 & 2 \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & -1 \\ -1 & 0 & 6 \end{bmatrix} \begin{bmatrix} 5 & -1 & -2 \\ 3 & 9 & -3 \\ -2 & 1 & 2 \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & -1 \\ -1 & 0 & 6 \end{bmatrix} \begin{bmatrix} 5 & -1 & -2 \\ 3 & 9 & -3 \\ -2 & 1 & 2 \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 4 & -1 \\ -1 & 0 & 6 \end{bmatrix} \begin{bmatrix} 5 & -1 & -2 \\ 3 & 9 & -3 \\ -2 & 1 & 2 \end{bmatrix} = \begin{bmatrix} * & * & * \\ * & * & * \\ * & * & * \end{bmatrix}$$

Learning goals

- Linear and affine equation systems (short)
- Solving Linear and affine equation systems using Gauss-Jordan algorithm
- Identifying the solutions of affine equation systems: no solution, a single solution, an affine space of solutions, based on the outcome of the Gauss-Jordan algorithm
- Extending a set of vectors to a basis of the whole space (by applying Gram-Schmid)

like a python try: except: clause

- ⦿ if you understand the explanation, it is best as it allows you to remember things easier, because you can deduct them again without memorizing too much – however this happens 10% all the time only
- ⦿ if you do not understand the explanation, take a look at it after a sleep. The brain processes things from the previous days (cf. nightmares!)
- ⦿ if you still do not understand the explanation, search for others
- ⦿ if you still do not understand the explanation after trying n different of them, it is okay ... not everyone can be a Messi, a van Gogh, a supermodel, a gifted cook, same holds for math skills, then be pragmatical: memorize the rules how to apply it

- 1 Recap
- 2 Linear equation Systems
- 3 Affine equation systems
- 4 Solving affine systems with transformations on rows
- 5 Extending a set of k independent vectors to a basis of the whole space \mathbb{R}^d

Linear equation system

A set of n linear equations which are to be solved for the vector $x = (x_0, \dots, x_{d-1})$

$$a_0^{\{0\}} x_0 + a_1^{\{0\}} x_1 + \dots + a_{d-1}^{\{0\}} x_{d-1} = a^{\{0\}} \cdot x = 0 \quad (1)$$

$$a_0^{\{1\}} x_0 + a_1^{\{1\}} x_1 + \dots + a_{d-1}^{\{1\}} x_{d-1} = a^{\{1\}} \cdot x = 0 \quad (2)$$

$$\dots \quad (3)$$

$$a_0^{\{n-1\}} x_0 + a_1^{\{n-1\}} x_1 + \dots + a_{d-1}^{\{n-1\}} x_{d-1} = a^{\{n-1\}} \cdot x = 0 \quad (4)$$

is called a linear equation system

This is obviously a set of equations made from n inner products:

$$a^{\{0\}} \cdot x = 0$$

$$a^{\{1\}} \cdot x = 0$$

...

$$a^{\{n-1\}} \cdot x = 0$$

As an example, such inner products like here

$$(1, 0, 2, 3) \cdot x = 0$$

$$(-1, 1, 0, 1) \cdot x = 0$$

can be written in the sense of matrix-vector multiplication as:

$$\begin{bmatrix} 1, 0, 2, 3 \end{bmatrix} x = 0$$

$$\begin{bmatrix} -1, 1, 0, 1 \end{bmatrix} x = 0$$

This in turn is the same as:

$$\begin{bmatrix} 1, 0, 2, 3 \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \end{bmatrix} = 0$$
$$\begin{bmatrix} -1, 1, 0, 1 \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \end{bmatrix} = 0$$

Now one can fuse this into a single matrix vector multiplication:

$$\begin{bmatrix} 1 & 0 & 2 & 3 \\ -1 & 1 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \end{bmatrix} = 0$$

This conforms to the equation

$$Ax = A \begin{bmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \end{bmatrix} = 0$$

with $A = \begin{bmatrix} 1 & 0 & 2 & 3 \\ -1 & 1 & 0 & 1 \end{bmatrix}$

You can do this with more than two equations:

$$(1, 0, 2, 3) \cdot x = 0$$

$$(-1, 1, 0, 1) \cdot x = 0$$

$$(4, 3, 2, 1) \cdot x = 0$$

is equivalent to

$$Ax = A \begin{bmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \end{bmatrix} = 0$$

$$\text{with } A = \begin{bmatrix} 1 & 0 & 2 & 3 \\ -1 & 1 & 0 & 1 \\ 4 & 3 & 2 & 1 \end{bmatrix}$$

Here with n equations:

$$\begin{aligned}w^{\{0\}} \cdot x &= 0 \\w^{\{1\}} \cdot x &= 0 \\&\vdots = 0 \\w^{\{n-1\}} \cdot x &= 0\end{aligned}$$

is equivalent to

$$Ax = 0$$

with $A = \begin{bmatrix} w^{\{0\}} \\ w^{\{1\}} \\ \vdots \\ w^{\{n-1\}} \end{bmatrix}$ where each $w^{\{i\}}$ is a horizontal / row vector.

Linear equation system II

A linear equation system can be defined by

$$Ax = 0$$

$$A.shape = (n, d)$$

$$x.shape = (d, 1)$$

where one solves for the vector $x = (x_0, \dots, x_{d-1})$.

- 1 Recap
- 2 Linear equation Systems
- 3 Affine equation systems**
- 4 Solving affine systems with transformations on rows
- 5 Extending a set of k independent vectors to a basis of the whole space \mathbb{R}^d

Affine equation system

A set of n affine equations which are to be solved for the vector $x = (x_0, \dots, x_{d-1})$

$$a_0^{\{0\}} x_0 + a_1^{\{0\}} x_1 + \dots + a_{d-1}^{\{0\}} x_{d-1} = a^{\{0\}} \cdot x = b_0$$

$$a_0^{\{1\}} x_0 + a_1^{\{1\}} x_1 + \dots + a_{d-1}^{\{1\}} x_{d-1} = a^{\{1\}} \cdot x = b_1$$

...

$$a_0^{\{n-1\}} x_0 + a_1^{\{n-1\}} x_1 + \dots + a_{d-1}^{\{n-1\}} x_{d-1} = a^{\{n-1\}} \cdot x = b_{n-1}$$

is called an affine equation system. b_i are called the bias terms.

This is equivalent to

$$Ax = b$$
$$\text{with } A = \begin{bmatrix} w^{\{0\}} \\ w^{\{1\}} \\ \vdots \\ w^{\{n-1\}} \end{bmatrix}, \quad b = \begin{bmatrix} b_0 \\ b_1 \\ \vdots \\ b_{n-1} \end{bmatrix}$$

where each $w^{\{i\}}$ is a horizontal / row vector, that is of shape $(1, d)$.

① Recap

② Linear equation Systems

③ Affine equation systems

④ Solving affine systems with transformations on rows

Transformations on rows which do not change the solution

Solving affine systems using Gaussian Elimination

Main Step 2: achieve row echelon form - as general algorithm

What solutions can we get?

Identifying the solutions based on the result of the Gauss-Jordan algorithm

⑤ Extending a set of k independent vectors to a basis of the whole space \mathbb{R}^d

The first step is to find a set of operations which do not change the solution. That is, we want to perform an operation L by matrix multiplication such that

$$Ax = b$$

$$L Ax = L b$$

have the same set of solutions x .

Swapping two rows

- This just swaps equations. So it does not affect the solutions.

Which operations can we use to solve? Multiplying a row with a non-zero value

| 38

Multiplying a row with a non-zero value

- if $a \neq 0$, then

$$\begin{array}{lll} v_1x_1 + v_2x_2 + \dots + v_{n-1}x_{n-1} & = v \cdot x & = b \\ \text{and } av_1x_1 + av_2x_2 + \dots + av_{n-1}x_{n-1} & = av \cdot x & = ab \end{array}$$

have the same solutions x .

Which operations can we use to solve? Adding a multiple of a row to another row

| 39

We consider the following: Equation set S1 is defined as:

$$a^{\{i\}} \cdot x = b_i$$

$$a^{\{k\}} \cdot x = b_k$$

Equation set S2 is defined as:

$$a^{\{i\}} \cdot x = b_i$$

$$(a^{\{k\}} + \beta a^{\{i\}}) \cdot x = b_k + \beta b_i$$

The S2 was created from S1 by adding a multiple of the first equation (multiplied with β) to the second equation.

Adding a multiple of an equation to another does not change the solutions

If we take any two equations from an affine equation system, and we add a multiple of an equation, then we do not change the set of solution vectors x . The solution for the original set, and for the set with added equations are the same.

Which operations can we use to solve? Adding a multiple of a row to another row

| 40

Why is this true?

Adding a multiple of an equation to another does not change the solutions

If we take any two equations from an affine equation system, and we add a multiple of an equation, then we do not change the set of solution vectors x . The solution for the original set, and for the set with added equations are the same.

Suppose z solves $S1$, which means:

$$a^{\{i\}} \cdot z = b_i$$

$$a^{\{k\}} \cdot z = b_k$$

But then also this holds:

$$\beta a^{\{i\}} \cdot z = \beta b_i$$

... Now just add these equations

Suppose z solves S1, which means:

$$a^{\{i\}} \cdot z = b_i$$

$$a^{\{k\}} \cdot z = b_k$$

But then also this holds:

$$\beta a^{\{i\}} \cdot z = \beta b_i$$

Adding the first equation from S1 to the last equation results in S2

$$a^{\{i\}} \cdot z = b_i$$

$$(a^{\{k\}} + \beta a^{\{i\}}) \cdot z = b_k + \beta b_i$$

So we have shown, if z solves S1, then z also solves S2.

Which operations can we use to solve? Adding a multiple of a row to another row

| 42

The other way round can be also shown easily: If z solves S2, then these two equations hold:

$$a^{\{i\}} \cdot z = b_i$$

$$(a^{\{k\}} + \beta a^{\{i\}}) \cdot z = b_k + \beta b_i$$

Now add $-\beta$ times the first equation (which is satisfied by assumption that z solves the equation set S2) to the second equation. This transforms S2 into

$$a^{\{i\}} \cdot z = b_i$$

$$(a^{\{k\}} + \beta a^{\{i\}}) \cdot z - \beta a^{\{i\}} \cdot z = b_k + \beta b_i - \beta b_i$$

which is S1:

$$a^{\{i\}} \cdot z = b_i$$

$$a^{\{k\}} \cdot z = b_k$$

Conclusion:

- ⊙ We have shown that adding a multiple of a row to another row does not change the solutions of an affine equation $Ax = b$

This method is suitable for solving small systems by hand

- Step 1: rewrite the affine equation system

$$Ax = b$$

into to the augmented matrix

$$\left[A|b \right] = \left[\begin{array}{ccccc|c} a_{00} & a_{01} & a_{02} & \dots & a_{0,d-1} & b_0 \\ a_{10} & a_{11} & a_{12} & \dots & a_{1,d-1} & b_1 \\ a_{20} & a_{21} & a_{22} & \dots & a_{2,d-1} & b_2 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{n-1,0} & a_{n-1,1} & a_{n-1,2} & \dots & a_{n-1,d-1} & b_{n-1} \end{array} \right]$$

- Step 2: obtain the so-called **row echelon form**.

To understand the idea behind a row echelon form, let us consider at first a special case:

If we had a (d, d) -shaped matrix A , then, in the second step, we want to achieve at first an upper diagonal matrix. It could look like this one:

$$A = \begin{bmatrix} 1 & * & * & \dots & * & * \\ 0 & 1 & * & \dots & * & * \\ 0 & 0 & 1 & \dots & * & * \\ 0 & 0 & 0 & \dots & * & * \\ \vdots & \vdots & \vdots & \ddots & 1 & * \\ 0 & 0 & 0 & \dots & 0 & 1 \end{bmatrix}$$

then, in Step 3, we process it into a diagonal matrix

$$A = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 & 0 \\ 0 & 1 & 0 & \dots & 0 & 0 \\ 0 & 0 & 1 & \dots & 0 & 0 \\ 0 & 0 & 0 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & 1 & 0 \\ 0 & 0 & 0 & \dots & 0 & 1 \end{bmatrix}$$

Why ? If the augmented matrix looks like that,

$$A = \left[\begin{array}{cccccc|c} 1 & 0 & 0 & \dots & 0 & 0 & b_0 \\ 0 & 1 & 0 & \dots & 0 & 0 & b_1 \\ 0 & 0 & 1 & \dots & 0 & 0 & b_2 \\ 0 & 0 & 0 & \dots & 0 & 0 & b_3 \\ \vdots & \vdots & \vdots & \ddots & 1 & 0 & b_* \\ 0 & 0 & 0 & \dots & 0 & 1 & b_{n-1} \end{array} \right]$$

then the resulting solution is

$$\begin{aligned} x_0 &= b_0 \\ x_1 &= b_1 \\ &\vdots \\ x_{n-1} &= b_{n-1} \end{aligned}$$

that is, we can read off the solution right away.

The shape of matrix which we want to achieve in the step 2, if the matrix A is any matrix, is a so-called **row echelon form**:

Definition: Row echelon form

Definition: Row echelon form We look only at the part without the bias column b_0, b_1, b_2, \dots

- All rows consisting of only zeroes are at the bottom of the matrix.
- for any non-zero row: The left-most nonzero entry of this row ('leading entry of a row') is to the right of the left-most nonzero entry of every row above it.
- the leading entry in any non-zero row is 1

Examples:

$$A = \begin{bmatrix} 1 & f & c \\ 0 & 1 & g \\ 0 & 0 & 1 \end{bmatrix}, A = \begin{bmatrix} 1 & f & c \\ 0 & 1 & e \\ 0 & 0 & 0 \end{bmatrix}, A = \begin{bmatrix} 1 & b & c \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix}$$

$$A = \begin{bmatrix} 1 & a & b & c \\ 0 & 1 & d & e \\ 0 & 0 & 1 & f \\ 0 & 0 & 0 & 1 \end{bmatrix}, A = \begin{bmatrix} 1 & a & c & b \\ 0 & 1 & f & d \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}, A = \begin{bmatrix} 1 & b & c & f \\ 0 & 0 & 1 & e \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$A = \begin{bmatrix} 1 & b & c & d \\ 0 & 0 & 1 & g \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}, A = \begin{bmatrix} 1 & b & e & d \\ 0 & 1 & c & g \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}, A = \begin{bmatrix} 1 & b & e & d \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Row Echelon Form Criteria:

We look only at the part without the bias column b_0, b_1, b_2, \dots

- ⊙ All rows consisting of only zeroes are at the bottom of the matrix.
- ⊙ for any non-zero row:
The left-most nonzero entry of this row ('leading entry of a row') is to the right of the left-most nonzero entry of every row above it.
- ⊙ the leading entry in any non-zero row is 1

What would be **NOT** a row echelon form by violating the second statement ? Here a counterexample:

$$Z = \begin{bmatrix} 1 & x & e & d \\ 0 & 1 & c & g \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$Z = \begin{bmatrix} x & y & e & d \\ 1 & z & c & g \\ 0 & 1 & f & g \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

The violation occurs in which row?

Row Echelon forms are a special case of upper diagonal matrices:

Definition upper diagonal matrix

B is an upper diagonal matrix, if all entries below the main diagonal are zero. That is, if

$$i > k \Rightarrow B_{ik} = 0$$

The entries below the main diagonal have a larger row index than their column index ($B_{ik}, i > k$).

- You can use the following operations to turn $[A|b]$ into the row echelon form:
 - swapping two rows (including the b_i -term)
 - multiplying a whole row (including the b_i -term) with a non-zero number
 - Adding a multiple of one row to another row (including the b_i -term)
- We know that all these operations do not change the solution space.

- Step 3: further process the row echelon form to a reduced row echelon form:
Goal is a matrix in which every column, which contains a left-most non-zero entry, has only this entry as non-zero entry. This is a so-called **Reduced row echelon form**.

Definition: Reduced Row echelon form

Definition: Reduced Row echelon form

- It is a row echelon form
- all other entries in a column containing a left-most non-zero entry are 0

- note: Columns which have no left-most non-zero entry cannot be further processed.

Examples for the the reduced row echelon form:

$$A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}, A = \begin{bmatrix} 1 & 0 & c \\ 0 & 1 & e \\ 0 & 0 & 0 \end{bmatrix}, A = \begin{bmatrix} 1 & b & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{bmatrix}$$

$$A = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}, A = \begin{bmatrix} 1 & 0 & c & 0 \\ 0 & 1 & f & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}, A = \begin{bmatrix} 1 & b & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$A = \begin{bmatrix} 1 & b & 0 & d \\ 0 & 0 & 1 & g \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}, A = \begin{bmatrix} 1 & 0 & e & d \\ 0 & 1 & c & g \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}, A = \begin{bmatrix} 1 & 0 & e & 0 \\ 0 & 1 & c & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

We obtain the reduced row echelon form from a matrix in row echelon form by adding rows with left-most non-zero entries to columns above them:

$$A = \begin{bmatrix} 1 & a & b & c & g \\ 0 & 1 & d & e & h \\ 0 & 0 & 1 & f & i \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \rightarrow \tilde{A} = \begin{bmatrix} 1 & 0 & 0 & \tilde{c} & 0 \\ 0 & 1 & 0 & \tilde{e} & 0 \\ 0 & 0 & 1 & \tilde{f} & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

$$\left[\begin{array}{ccc|c} 3 & 3 & 3 & 6 \\ 2 & -1 & -4 & 13 \\ -3 & 2 & 2 & -20 \end{array} \right] \rightarrow (r_0^* = 1/3) \left[\begin{array}{ccc|c} 1 & 1 & 1 & 2 \\ 2 & -1 & -4 & 13 \\ -3 & 2 & 2 & -20 \end{array} \right]$$

$$\rightarrow (r_1 + = -2r_0), (r_2 + = 3r_0) \left[\begin{array}{ccc|c} 1 & 1 & 1 & 2 \\ 0 & -3 & -6 & 9 \\ 0 & 5 & 5 & -14 \end{array} \right]$$

$$\rightarrow (r_1^* = 1/-3) \left[\begin{array}{ccc|c} 1 & 1 & 1 & 2 \\ 0 & 1 & 2 & -3 \\ 0 & 5 & 5 & -14 \end{array} \right]$$

$$\rightarrow (r_2 + = -5r_1) \left[\begin{array}{ccc|c} 1 & 1 & 1 & 2 \\ 0 & 1 & 2 & -3 \\ 0 & 0 & -5 & 1 \end{array} \right]$$

$$\rightarrow (r_2^* = -1/5) \left[\begin{array}{ccc|c} 1 & 1 & 1 & 2 \\ 0 & 1 & 2 & -3 \\ 0 & 0 & 1 & -1/5 \end{array} \right]$$

Example with a row swap:

$$\left[\begin{array}{ccc|c} 2 & 4 & 6 & 5 \\ 3 & 6 & -3 & 4.5 \\ -2 & 2 & 2 & 7 \end{array} \right] \rightarrow \left[\begin{array}{ccc|c} 1 & 2 & 3 & 2.5 \\ 3 & 6 & -3 & 4.5 \\ -2 & 2 & 2 & 7 \end{array} \right]$$

$$A_{1,:} = A_{1,:} - 3A_{0,:} \rightarrow \left[\begin{array}{ccc|c} 1 & 2 & 3 & 2.5 \\ 0 & 0 & -12 & -3 \\ 0 & 6 & 8 & 12 \end{array} \right]$$

$$\text{swap here} \rightarrow \left[\begin{array}{ccc|c} 1 & 2 & 3 & 2.5 \\ 0 & 6 & 8 & 12 \\ 0 & 0 & -12 & -3 \end{array} \right]$$

$$\rightarrow \left[\begin{array}{ccc|c} 1 & 2 & 3 & 2.5 \\ 0 & 1 & 4/3 & 2 \\ 0 & 0 & -12 & -3 \end{array} \right]$$

$$\rightarrow \left[\begin{array}{ccc|c} 1 & 2 & 3 & 2.5 \\ 0 & 1 & 4/3 & 2 \\ 0 & 0 & 1 & 1/4 \end{array} \right]$$

- swap all zero rows to the bottom
- For the 0-th row:
 - If $a_{0,0}$ is zero, follow this general idea:

general idea of finding the leftmost non-zero entry

find in the current row of interest “the leftmost non-zero entry”, such that there is ... no non-zero entry which is at the same time further lower and further left of it.
You can swap rows to find this state.

Examples:

$$A = \left[\begin{array}{cccc|c} 0 & \underline{c} & d & e & b_0 \\ 0 & g & h & f & b_1 \\ \underline{r} & l & m & k & b_2 \\ \underline{s} & n & o & p & b_3 \end{array} \right] \quad \text{focus on zero-th row: not ok case !}$$
$$\rightarrow A = \left[\begin{array}{cccc|c} \underline{r} & l & m & k & b_2 \\ 0 & c & d & e & b_0 \\ 0 & g & h & f & b_1 \\ s & n & o & p & b_3 \end{array} \right] \quad \text{focus on zero-th row: okay case !}$$

This was obtained by swapping the zero-th and the second row!

Examples:

$$A = \left[\begin{array}{cccc|c} 0 & 0 & \underline{d} & e & b_0 \\ 0 & \underline{c} & h & f & b_1 \\ 0 & \underline{l} & m & k & b_2 \\ 0 & \underline{n} & o & p & b_3 \end{array} \right] \quad \text{focus on zero-th row: not ok case !}$$

$$A = \left[\begin{array}{cccc|c} 0 & \underline{c} & h & f & b_1 \\ 0 & 0 & d & e & b_0 \\ 0 & l & m & k & b_2 \\ 0 & n & o & p & b_3 \end{array} \right] \quad \text{focus on zero-th row: okay case !}$$

$$A = \left[\begin{array}{cccc|c} 0 & 0 & \underline{d} & e & b_0 \\ 0 & 0 & h & f & b_1 \\ 0 & \underline{l} & m & k & b_2 \\ 0 & \underline{n} & o & p & b_3 \end{array} \right] \quad \text{focus on zero-th row: not ok case !}$$

$$A = \left[\begin{array}{cccc|c} 0 & 0 & \underline{d} & e & b_0 \\ 0 & 0 & h & f & b_1 \\ 0 & 0 & m & k & b_2 \\ 0 & 0 & o & p & b_3 \end{array} \right] \quad \text{focus on zero-th row: okay case !}$$

Can swap rows if needed:

$$A = \left[\begin{array}{cccc|c} 0 & c & d & e & b_0 \\ 0 & g & h & f & b_1 \\ r & l & m & k & b_2 \\ s & n & o & p & b_3 \end{array} \right] \rightarrow A = \left[\begin{array}{cccc|c} r & l & m & k & b_2 \\ 0 & c & d & e & b_0 \\ 0 & g & h & f & b_1 \\ s & n & o & p & b_3 \end{array} \right]$$

- If $a_{0,0}$ is zero, swap rows from the bottom of it, until $a_{0,0} \neq 0$.
- If getting $a_{0,0} \neq 0$ by swapping rows is not possible, it means that the whole 0-th column $A_{0:end,0}$ must be zero. In that case do the same thing with one column to the right: check if $a_{0,1}$ is zero, swap rows from the bottom of it, until $a_{0,1} \neq 0$.

$$A = \left[\begin{array}{cccc|c} 0 & 0 & d & e & b_0 \\ 0 & g & h & f & b_1 \\ 0 & l & m & k & b_2 \\ 0 & n & o & p & b_3 \end{array} \right] \rightarrow A = \left[\begin{array}{cccc|c} 0 & g & h & f & b_1 \\ 0 & 0 & d & e & b_0 \\ 0 & l & m & k & b_2 \\ 0 & n & o & p & b_3 \end{array} \right]$$

- ⊙ If that is not possible, then the first column is all zero, too. Then do the same with one column to the right further : check if $a_{0,2}$ is zero ...

$$A = \left[\begin{array}{cccc|c} 0 & 0 & 0 & e & b_0 \\ 0 & 0 & h & f & b_1 \\ 0 & 0 & m & k & b_2 \\ 0 & 0 & z & p & b_3 \end{array} \right] \rightarrow A = \left[\begin{array}{cccc|c} 0 & 0 & h & f & b_1 \\ 0 & 0 & 0 & e & b_0 \\ 0 & 0 & m & k & b_2 \\ 0 & 0 & z & p & b_3 \end{array} \right]$$

Suppose you have found a leftmost $a_{0,r_0} \neq 0$.

- Suppose you have found a leftmost $a_{0,r_0} \neq 0$. r_0 can be at index 0 if we have $a_{0,0} \neq 0$.
- **next step:** multiply the 0-th row so that $a_{0,r_0} = 1$. do the same operation with the same multiplier to the bias column.

$$A = \left[\begin{array}{cccc|c} a & b & c & d & b_0 \\ e & f & g & h & b_1 \\ i & k & l & m & b_2 \\ n & o & q & q & b_3 \end{array} \right] \rightarrow \hat{A} = \left[\begin{array}{cccc|c} \underline{1} & \hat{b} & \hat{c} & \hat{d} & \hat{b}_0 \\ e & f & g & h & b_1 \\ i & k & l & m & b_2 \\ n & o & q & q & b_3 \end{array} \right]$$

- ⊙ Add multiples of the 0-th row $a_{0,:}$ to all rows, such that the r_0 -column in all other rows becomes zero. Do the same operation with the same multiplier to the bias column.

$$A = \left[\begin{array}{cccc|c} 1 & b & c & d & b_0 \\ e & f & g & h & b_1 \\ i & k & l & m & b_2 \\ n & o & p & q & b_3 \end{array} \right] \rightarrow \hat{A} = \left[\begin{array}{cccc|c} 1 & b & c & d & b_0 \\ 0 & \hat{f} & \hat{g} & \hat{h} & \hat{b}_1 \\ 0 & \hat{k} & \hat{l} & \hat{m} & \hat{b}_2 \\ 0 & \hat{o} & \hat{p} & \hat{q} & \hat{b}_3 \end{array} \right]$$

$\downarrow \dots \dots \dots \downarrow$
 add (row 0)*factor(-e) to row 1
 add (row 0)*factor(-i) to row 2
 add (row 0)*factor(-n) to row 3

or, if the 0-th column was all zeros and your $r_0 = 1$:

$$A = \left[\begin{array}{cccc|c} 0 & 1 & c & d & b_0 \\ 0 & f & g & h & b_1 \\ 0 & k & l & m & b_2 \\ 0 & o & p & q & b_3 \end{array} \right] \rightarrow \hat{A} = \left[\begin{array}{cccc|c} 0 & 1 & c & d & b_0 \\ 0 & 0 & \hat{g} & \hat{h} & \hat{b}_1 \\ 0 & 0 & \hat{l} & \hat{m} & \hat{b}_2 \\ 0 & 0 & \hat{p} & \hat{q} & \hat{b}_3 \end{array} \right]$$

Now repeat the same three steps for the row number 1:

- Find $A_{1,r_1} \neq 0$ such that there is no non-zero entry in rows below which is further left of column index r_1 . You can swap rows to achieve this
- Multiply the row with $\frac{1}{A_{1,r_1}}$
- Add multiples of the row to all rows below so that all entries in the column indexed with r_0 become zero: $A_{2,r_1} = 0, \dots, A_{n-1,r_1} = 0$

$$A = \left[\begin{array}{cccc|c} 1 & b & c & d & b_0 \\ 0 & 1 & g & h & b_1 \\ 0 & k & l & m & b_2 \\ 0 & o & p & q & b_3 \end{array} \right] \rightarrow \hat{A} = \left[\begin{array}{cccc|c} 1 & b & c & d & b_0 \\ 0 & 1 & g & h & b_1 \\ 0 & 0 & \hat{l} & \hat{m} & \hat{b}_2 \\ 0 & 0 & \hat{p} & \hat{q} & \hat{b}_3 \end{array} \right]$$

$\downarrow \dots \dots \dots \downarrow$
 add (row 1)*factor(-k) to row 2
 add (row 1)*factor(-o) to row 3

Now repeat the same three steps for all other row numbers k :

- Find $A_{k,r_k} \neq 0$ such that there is no non-zero entry in rows below which is further left of column index r_k . You can swap rows to achieve this
- Multiply the row with $\frac{1}{A_{k,r_k}}$
- Add multiples of the row to all rows below so that all entries in the column indexed with r_0 become zero: $A_{k+1,r_k} = 0, \dots, A_{n-1,r_k} = 0$

$$A = \left[\begin{array}{cccccc|c} 1 & 0 & * & * & \dots & * & * \\ 0 & 1 & * & * & \dots & * & * \\ 0 & 0 & \textcolor{teal}{1} & * & \dots & * & \textcolor{teal}{*} \\ 0 & 0 & * & * & \dots & * & * \\ 0 & 0 & * & * & \dots & * & * \\ 0 & 0 & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & * & * & \dots & * & * \end{array} \right] \rightarrow A = \left[\begin{array}{cccccc|c} 1 & * & * & * & \dots & * & * \\ 0 & 1 & * & * & \dots & * & * \\ 0 & 0 & \textcolor{teal}{1} & * & \dots & * & \textcolor{teal}{*} \\ 0 & 0 & \textcolor{red}{0} & \textcolor{blue}{*} & \dots & \textcolor{blue}{*} & \textcolor{blue}{*} \\ 0 & 0 & \textcolor{red}{0} & \textcolor{blue}{*} & \dots & \textcolor{blue}{*} & \textcolor{blue}{*} \\ 0 & 0 & \textcolor{red}{\vdots} & \textcolor{blue}{\vdots} & \ddots & \textcolor{blue}{\vdots} & \textcolor{blue}{\vdots} \\ 0 & 0 & \textcolor{red}{0} & \textcolor{blue}{*} & \dots & \textcolor{blue}{*} & \textcolor{blue}{*} \end{array} \right]$$

- ⦿ this is much simpler: you go through all columns which have a left-most-non zero entry and make all other entries in this column turn to zero.
- ⦿ you use only adding rows to others. Difference to before: you add a row to rows above of it.
- ⦿ you do less operations if you go through all columns which have a left-most-non zero entry – starting from the right. But it also works, if you would start from the left as you have done it in **main step 2**

$$\begin{aligned}
 \left[\begin{array}{ccc|c} 1 & 1 & 1 & 2 \\ 0 & 1 & 2 & -3 \\ 0 & 0 & 1 & -1/5 \end{array} \right] &\rightarrow \left[\begin{array}{ccc|c} 1 & 1 & 1 & 2 \\ 0 & 1 & 0 & -3 + 2/5 = -2.4 \\ 0 & 0 & 1 & -1/5 \end{array} \right] \begin{array}{l} \text{add row 2 to row 1} \\ \uparrow \dots \uparrow \end{array} \\
 &\rightarrow \left[\begin{array}{ccc|c} 1 & 1 & 0 & 2 + 1/5 = 2.2 \\ 0 & 1 & 0 & -2.4 \\ 0 & 0 & 1 & -1/5 \end{array} \right] \begin{array}{l} \text{add row 2 to row 0} \\ \uparrow \dots \uparrow \end{array}
 \end{aligned}$$

- ⦿ this is much simpler: you go through all columns which have a left-most-non zero entry and make all other entries in this column turn to zero.
- ⦿ you use only adding rows to others. Difference to before: you add a row to rows above of it.
- ⦿ you do less operations if you go through all columns which have a left-most-non zero entry – starting from the right. But it also works, if you would start from the left as you have done it in **main step 2**

$$\left[\begin{array}{ccc|c} 1 & 1 & 0 & 2 + 1/5 = 2.2 \\ 0 & 1 & 0 & -2.4 \\ 0 & 0 & 1 & -1/5 \end{array} \right] \rightarrow \left[\begin{array}{ccc|c} 1 & 0 & 0 & 2.2 - (-2.4) = 4.6 \\ 0 & 1 & 0 & -2.4 \\ 0 & 0 & 1 & -1/5 \end{array} \right] \begin{array}{l} \text{add row 1 to row 0} \\ \uparrow \dots \uparrow \end{array}$$

You are adding whole rows. I only put focus on what is important by **row of origin change to zero change to any value**. If there would be a fourth column, as in the following example, then it would change too:

$$\begin{bmatrix} 1 & 1 & 1 & 1 & | & 2 \\ 0 & 1 & 2 & 3 & | & -3 \\ 0 & 0 & 1 & -2 & | & -1/5 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 1 & 1 & | & 2 \\ 0 & 1 & 0 & 7 & | & -3 + 2/5 = -2.4 \\ 0 & 0 & 1 & -2 & | & -1/5 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 1 & 1 & 1 & | & 2 \\ 0 & 1 & 0 & 7 & | & -2.4 \\ 0 & 0 & 1 & -2 & | & -1/5 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 1 & 0 & 3 & | & 2 + 1/5 = 2.2 \\ 0 & 1 & 0 & 7 & | & -2.4 \\ 0 & 0 & 1 & -2 & | & -1/5 \end{bmatrix}$$

Observe: Anything in columns left of the left-most-non zero entry will not be changed

What solutions can we get? One possibility in case of full-rank $d \times d$ matrices

69

When A is of shape $d \times d$ and has full rank, then you will obtain an equation which looks like this:

$$A = \left[\begin{array}{ccccc|c} 1 & 0 & 0 & \dots & 0 & b_0 \\ 0 & 1 & 0 & \dots & 0 & b_1 \\ 0 & 0 & 1 & \dots & 0 & b_2 \\ 0 & 0 & 0 & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & 1 & b_{d-1} \end{array} \right]$$

This corresponds to a directly read-off solution

$$x_0 = b_0$$

$$x_1 = b_1$$

$$x_2 = b_2$$

$$\vdots$$

$$x_{d-1} = b_{d-1}$$

If you obtain a row, where all entries are zeros, but the bias is not zero, $b \neq 0$,

$$A = \left[\begin{array}{cccccc|c} 1 & 0 & * & * & \dots & * & * \\ 0 & 1 & * & * & \dots & * & * \\ 0 & 0 & 0 & 0 & \dots & 0 & b \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \end{array} \right]$$

then you know that this has no solution: This corresponds to trying to solve

$$0x_0 + 0x_1 + \dots + 0x_{d-1} = b \neq 0$$

This is common to be observed when $n > d$. That is when you have more equations n than dimensions d in your solution space.

In this case: a row with all zeros in the matrix part, and a non-zero bias $b \neq 0$, no solution exists.

If you obtain such a condition with a zero bias $b = 0$

$$A = \left[\begin{array}{cccccc|c} 1 & 0 & * & * & \dots & * & b_0 \\ 0 & 1 & * & * & \dots & * & b_1 \\ 0 & 0 & 0 & 0 & \dots & 0 & 0 \end{array} \right],$$

it is a row stating

$$0x_0 + 0x_1 + \dots + 0x_{d-1} = 0$$

It poses no constraint on the solution. This is satisfied by any values for x_0, x_1, \dots, x_{d-1} . Therefore you can remove or ignore such a row.

Zero rows in A can also happen when $n \leq d$, and some of the rows of A are a linear combination of other rows of A . Then Gauss-Jordan transforms these rows into zero-rows.

Example for a case with $n \leq d$ and zero rows in A after transformation:

$$A = \begin{bmatrix} 1 & 2 & 1 \\ -2 & -4 & -2 \\ 1 & 0 & 1 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 2 & 1 \\ 0 & 0 & 0 \\ 1 & 0 & 1 \end{bmatrix}$$

What solutions can we get? Rows with zero bias and conditions with many solutions

| 72

If you obtain, after removing conditions with zero bias, a structure like that:

$$A = \left[\begin{array}{cccccc|c} 1 & 0 & * & * & \dots & * & b_0 \\ 0 & 1 & * & * & \dots & * & b_1 \end{array} \right],$$

then you have as solution an affine space with as many dimensions, as columns with stars *.
How to solve it ?

- ⦿ If you just need any solution, then the simplest way is to plug in a zero ($= 0$) for all the variables with a *.

In the above example:

$$x_0 = b_0$$

$$x_1 = b_1$$

$$x_2 = x_3 = \dots = x_{d-1} = 0$$

What solutions can we get? Rows with zero bias and conditions with many solutions

| 73

- For other solutions: You can plug in any values for the variables with a *, and move the resulting numbers to the existing bias term. Then solve for the terms having a 1 coefficient.

Example here:

$$A = \left[\begin{array}{cccc|c} 1 & 0 & -2 & 3.5 & 5 \\ 0 & 1 & 3 & -1 & 7 \end{array} \right]$$

You could plug in $x_2 = 4$. This results in $-2 * 4$ and $3 * 4$ for the column corresponding to x_2 . You could plug in $x_3 = 2$. This results in $3.5 * 2$ and $-1 * 2$ for the column corresponding to x_3 . Bring this to the side of the bias:

$$x_2 = 4, x_3 = 2 \text{ and}$$

$$A = \left[\begin{array}{cccc|c} 1 & 0 & 0 & 0 & 5 - (-2 * 4) - 3.5 * 2 \\ 0 & 1 & 0 & 0 & 7 - (3 * 4) - (-1 * 2) \end{array} \right] \Rightarrow x_0 = 6, x_1 = -3$$

The solution to the affine equation

One can see this by looking at the reduced row echelon form after processing $Ax = b$

- it can have no solution, a single solution, or an affine space of solutions.
- it has no solution, if the Gauss-Jordan Algorithm results (at any step) in a matrix where one row has only zeros in its matrix part, but the corresponding bias term of that row is non-zero
- special case: it has exactly one solution if A is of shape (d, d) and the Gauss-Jordan Algorithm in Main step 2 results in an upper diagonal matrix with only non-zero entries on the diagonal
- it has a vector space of solutions if the result of the Gauss-Jordan Algorithm shows that
 - all rows which have only zeros in their matrix part have also their bias term being zero.
 - there exists a column which has no left-most non-zero entry
- it has exactly one solution if the result of the Gauss-Jordan Algorithm shows that
 - all rows which have only zeros in their matrix part have also their bias term being zero.
 - all columns have a left-most non-zero entry

These cases have also certain namings for the equation systems

- ⦿ no solution: **inconsistent** equation system
- ⦿ a single solution: **consistent and determinate** equation system
- ⦿ an affine space of solutions: **consistent and indeterminate** equation system

An affine space is a vector space with an added offset vector, that is:

$$A = \{a + v, v \in V\}$$

where V is a vector space.

An affine space can be described in terms of a basis v_0, \dots, v_{d-1} of a vector space V the set of

$$A = \{a + \sum_{k=0}^{d-1} b_k v_k, b_k \in \mathbb{R}\}$$

Why the solutions of this define an affine space?

$$A = \left[\begin{array}{cccc|c} 1 & 0 & -2 & 3.5 & 5 \\ 0 & 1 & 3 & -1 & 7 \end{array} \right],$$

Lets write them out:

$$x_0 - 2x_2 + 3.5x_3 = 5$$

$$x_1 + 3x_2 - 1x_3 = 7$$

This transforms into:

$$x_0 = 5 + 2x_2 - 3.5x_3$$

$$x_1 = 7 - 3x_2 + 1x_3$$

Now write this in terms of a vector

$$\begin{bmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \end{bmatrix}$$

$$\begin{aligned}x_0 &= 5 + 2x_2 - 3.5x_3 \\x_1 &= 7 - 3x_2 + 1x_3 \\ \Rightarrow \begin{bmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \end{bmatrix} &= \begin{bmatrix} 5 \\ 7 \\ 0 \\ 0 \end{bmatrix} + x_2 \begin{bmatrix} 2 \\ -3 \\ 1 \\ 0 \end{bmatrix} + x_3 \begin{bmatrix} -3.5 \\ 1 \\ 0 \\ 1 \end{bmatrix}\end{aligned}$$

This conforms to the definition of an affine space

$$A = \{x = a + b_0 v_0 + b_1 v_1, b_0 \in \mathbb{R}, b_1 \in \mathbb{R}\} \text{ with}$$

$$a = \begin{bmatrix} 5 \\ 7 \\ 0 \\ 0 \end{bmatrix}, b_0 = x_2, v_0 = \begin{bmatrix} 2 \\ -3 \\ 1 \\ 0 \end{bmatrix}, b_1 = x_3, v_1 = \begin{bmatrix} -3.5 \\ 1 \\ 0 \\ 1 \end{bmatrix}$$

- ① Recap
- ② Linear equation Systems
- ③ Affine equation systems
- ④ Solving affine systems with transformations on rows
- ⑤ Extending a set of k independent vectors to a basis of the whole space \mathbb{R}^d

The next sections are about understanding what the solutions are for linear and affine equation systems. At first we need a helping construction when we are given a set of vectors $w^{\{0\}}, \dots, w^{\{k-1\}}$.

What is our goal here ?

- We start with a given set of k vectors $w^{\{0\}}, \dots, w^{\{k-1\}}$. The space has dimensionality d .
- We want to add $d - k$ vectors $w^{\{k\}}, w^{\{k+1\}}, \dots, w^{\{d-1\}}$ such that

(Prop 1) the union of all vectors $w^{\{0\}}, \dots, w^{\{k-1\}}, w^{\{k\}}, w^{\{k+1\}}, \dots, w^{\{d-1\}}$ are a basis of \mathbb{R}^d
Rmb: Basis means independent + every vector from \mathbb{R}^d can be expressed as a linear combination of the basis vectors

(Prop 2) any vector from $w^{\{0\}}, \dots, w^{\{k-1\}}$ (the first part) is orthogonal to any vector from $w^{\{k\}}, w^{\{k+1\}}, \dots, w^{\{d-1\}}$ (the second part). Mathematically one can write this condition as:

$$r \leq k - 1, s \geq k \Rightarrow w^{\{r\}} \cdot w^{\{s\}} = 0$$

$$r \leq k - 1, s \geq k \Rightarrow w^{\{r\}} \cdot w^{\{s\}} = 0$$

insight

Any vector, which is a linear combination from the first part, is orthogonal to any vector, which is a linear combination from the second part.

One can also say that the vectors from the first part create a vector space which is orthogonal to the vector space created by the vectors from the second part.

Out of exams: Why is this true?

$$u = \sum_{r=1}^{k-1} a_r w^{\{r\}} \text{ see: summing of } r \text{ runs only over the first part}$$

$$v = \sum_{s=k}^{d-1} a_s w^{\{s\}} \text{ see: summing of } r \text{ runs only over the 2nd part}$$

$$u \cdot v = \left(\sum_{r=1}^{k-1} a_r w^{\{r\}} \right) \cdot \left(\sum_{s=k}^{d-1} a_s w^{\{s\}} \right)$$

$$= \sum_{r=1}^{k-1} a_r \left(w^{\{r\}} \cdot \left(\sum_{s=k}^{d-1} a_s w^{\{s\}} \right) \right) \text{ use linearity of inner product in its components!}$$

$$= \sum_{r=1}^{k-1} a_r \sum_{s=k}^{d-1} a_s \underbrace{(w^{\{r\}} \cdot w^{\{s\}})}_{=0} = 0 \text{ due to } r \leq k-1, s \geq k$$

Examples for such extension of a set of vectors:

- We start with the vector $(1, 0, 0)$ (spans the space of all $(a, 0, 0)$). We can extend $(1, 0, 0)$ by $(0, 1, 0)$ and $(0, 0, 1)$. These three vectors allow to represent any vector $(a_0, a_1, a_2) \in \mathbb{R}^3$
- We can extend $(1, 0, 0)$ also by $(0, 1, -1)$ and $(0, 1, 1)$
- We can extend $(1, 1, 0)$ by $(-2, 2, 0)$ and $(0, 0, 1)$

How to get a set with Prop1 and Prop2?

Gram-Schmid will be the tool to achieve that extension of an independent set!

Remember that Gram-Schmid results in an orthogonal set of vectors $\tilde{v}^{\{0\}}, \dots, \tilde{v}^{\{d-1\}}$, if the input is an independent set, that is $\tilde{v}^{\{r\}} \cdot \tilde{v}^{\{s\}} = 0$ whenever $r \neq s$

- we know that the d -dimensional space of real valued vectors has this set of vectors as basis:

$$e^{\{0\}} = \underbrace{(1, 0, \dots, 0)}_{d \text{ of them}}, e^{\{1\}} = (0, 1, 0, \dots, 0), e^{\{2\}} = (0, 0, 1, \dots, 0), e^{\{d-1\}} = (0, 0, 0, \dots, 1)$$

These are one hot-vectors. They are zero, except at a single position, where they are one.

- start with this initialization:

$$\begin{aligned} v^{\{0\}} &= w^{\{0\}}, v^{\{1\}} = w^{\{1\}}, \dots, v^{\{k-1\}} = w^{\{k-1\}}, \text{ } k \text{ vecs} \\ v^{\{k\}} &= e^{\{0\}}, v^{\{k+1\}} = e^{\{1\}}, \dots, v^{\{k+d-1\}} = e^{\{d-1\}} \text{ } d \text{ vecs} \end{aligned}$$

that is we create a sequence which starts with $w^{\{0\}}, \dots, w^{\{k-1\}}$, then we add the other basis vectors to the sequence behind these k vectors.

- Next: we run Gram-Schmidt-Orthogonalization on this set.

Next: we run Gram-Schmidt-Orthogonalization on this set.

- ⊙ **Important 1:** the vector space spanned by $v^{\{0\}}, \dots, v^{\{k-1\}}$ will be the same as the space spanned by the outputs $\tilde{v}^{\{0\}}, \dots, \tilde{v}^{\{k-1\}}$. See the last lecture for a proof.¹
- ⊙ **Important 2:** when we run Gram-Schmidt past index $k - 1$, then it will remove the components of all vectors $v^{\{0\}}, \dots, v^{\{k-1\}}$ (up to index $k - 1$) from all the following vectors starting at index k , that is from $v^{\{k\}}, \dots, v^{\{k+d-1\}}$.
- ⊙ So the resulting $\tilde{v}^{\{k\}}, \dots, \tilde{v}^{\{k+d-1\}}$ will be orthogonal to all vectors $v^{\{0\}}, \dots, v^{\{k-1\}}$ – as we wanted to find!

¹I did not prove that Gram-Schmidt does not increase the space of representable vectors.

- ◉ We know that the space is d -dimensional - that is every basis set has exactly d vectors, but we have given $k + d$ vectors as input for the algorithm.
- ◉ From the first k vectors, none of the will become zero (because we assumed that they are linearly independent). Therefore k of the last d basis vectors $e^{\{0\}}, \dots, e^{\{d-1\}}$ must turn into zero vectors after orthogonalization.

Solution: We will use $\tilde{v}^{\{0\}}, \dots, \tilde{v}^{\{k-1\}}$ for the first part. We will use those $d - k$ vectors from the set $\tilde{v}^{\{k\}}, \dots, \tilde{v}^{\{k+d-1\}}$ for the second part, which did not become zero vectors .

Together they are d orthogonal vectors, and therefore they must be a basis of the vector space \mathbb{R}^d as desired.

Next step: We do the same, but for the case, when the input set of k vectors

$$w^{\{0\}}, \dots, w^{\{k-1\}}$$

is possibly not independent.

We do this very similarly to the previous application of Gram-Schmidt.

- start with this initialization:

$$v^{\{0\}} = w^{\{0\}}, v^{\{1\}} = w^{\{1\}}, \dots, v^{\{k-1\}} = w^{\{k-1\}}, \text{ that are } k \text{ vecs}$$
$$v^{\{k\}} = e^{\{0\}}, v^{\{k+1\}} = e^{\{1\}}, \dots, v^{\{k+d-1\}} = e^{\{d-1\}}, \text{ that are } d \text{ vecs}$$

that is we create a sequence which starts with $w^{\{0\}}, \dots, w^{\{k-1\}}$, then we add the other basis vectors to the sequence behind these k vectors.

- Next: we run Gram-Schmidt on this set.

- There is one difference:

Since the first k vectors are not independent, there can be zeros among the output of the first k vectors. If this happens, then the last d vectors would have a few zero vectors less. So we have to look at the whole output (of length $k + d$) $\tilde{v}^{\{0\}}, \dots, \tilde{v}^{\{k+d-1\}}$ to remove uninteresting zero vectors.

Solution:

- ⊙ We will use those vectors from $\tilde{v}^{\{0\}}, \dots, \tilde{v}^{\{k-1\}}$, which did not become zero vectors, for the first part.
- ⊙ We will use those vectors from $\tilde{v}^{\{k\}}, \dots, \tilde{v}^{\{k+d-1\}}$, which did not become zero vectors, for the second part.

They are d orthogonal vectors, and therefore they must be a basis of the vector space \mathbb{R}^d as desired.