

# Mathematics

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# Outline

Algebra

Number Theory

Combinatorics

Geometry

## Sum of Powers

$$\sum_{k=1}^n k^2 = \frac{1}{6}n(n+1)(2n+1)$$

$$\sum k^3 = \left(\sum k\right)^2 = \left(\frac{1}{2}n(n+1)\right)^2$$

- ▶ Pretty useful in many random situations
- ▶ Memorize above!

## Fast Exponentiation

- Recursive computation of  $a^n$ :

$$a^n = \begin{cases} 1 & n = 0 \\ a & n = 1 \\ (a^{n/2})^2 & n \text{ is even} \\ a(a^{(n-1)/2})^2 & n \text{ is odd} \end{cases}$$

## Implementation (recursive)

```
double pow(double a, int n) {  
    if(n == 0) return 1;  
    if(n == 1) return a;  
    double t = pow(a, n/2);  
    return t * t * pow(a, n%2);  
}
```

- ▶ Running time:  $O(\log n)$

## Implementation (non-recursive)

```
double pow(double a, int n) {  
    double ret = 1;  
    while(n) {  
        if(n%2 == 1) ret *= a;  
        a *= a; n /= 2;  
    }  
    return ret;  
}
```

- You should understand how it works

# Linear Algebra

- ▶ Solve a system of linear equations
- ▶ Invert a matrix
- ▶ Find the rank of a matrix
- ▶ Compute the determinant of a matrix
- ▶ All of the above can be done with Gaussian elimination

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## Greatest Common Divisor (GCD)

- ▶  $\gcd(a, b)$ : greatest integer divides both  $a$  and  $b$
- ▶ Used very frequently in number theoretical problems
- ▶ Some facts:
  - $\gcd(a, b) = \gcd(a, b - a)$
  - $\gcd(a, 0) = a$
  - $\gcd(a, b)$  is the smallest positive number in  $\{ax + by \mid x, y \in \mathbf{Z}\}$

## Euclidean Algorithm

- ▶ Repeated use of  $\gcd(a, b) = \gcd(a, b - a)$
- ▶ Example:

$$\begin{aligned}\gcd(1989, 867) &= \gcd(1989 - 2 \times 867, 867) \\ &= \gcd(255, 867) \\ &= \gcd(255, 867 - 3 \times 255) \\ &= \gcd(255, 102) \\ &= \gcd(255 - 2 \times 102, 102) \\ &= \gcd(51, 102) \\ &= \gcd(51, 102 - 2 \times 51) \\ &= \gcd(51, 0) \\ &= 51\end{aligned}$$

## Implementation

```
int gcd(int a, int b) {  
    while(b){int r = a % b; a = b; b = r;}  
    return a;  
}
```

- ▶ Running time:  $O(\log(a + b))$
- ▶ Be careful:  $a \% b$  follows the sign of  $a$ 
  - $5 \% 3 == 2$
  - $-5 \% 3 == -2$

## Congruence & Modulo Operation

- ▶  $x \equiv y \pmod{n}$  means  $x$  and  $y$  have the same remainder when divided by  $n$
- ▶ Multiplicative inverse
  - $x^{-1}$  is the inverse of  $x$  modulo  $n$  if  $xx^{-1} \equiv 1 \pmod{n}$
  - $5^{-1} \equiv 3 \pmod{7}$  because  $5 \cdot 3 \equiv 15 \equiv 1 \pmod{7}$
  - May not exist (e.g., inverse of 2 mod 4)
  - Exists if and only if  $\gcd(x, n) = 1$

## Multiplicative Inverse

- ▶ All intermediate numbers computed by Euclidean algorithm are integer combinations of  $a$  and  $b$ 
  - Therefore,  $\gcd(a, b) = ax + by$  for some integers  $x, y$
  - If  $\gcd(a, n) = 1$ , then  $ax + ny = 1$  for some  $x, y$
  - Taking modulo  $n$  gives  $ax \equiv 1 \pmod{n}$
- ▶ We will be done if we can find such  $x$  and  $y$

## Extended Euclidean Algorithm

- ▶ Main idea: keep the original algorithm, but write all intermediate numbers as integer combinations of  $a$  and  $b$
- ▶ Exercise: implementation!

## Chinese Remainder Theorem

- ▶ Given  $a, b, m, n$  with  $\gcd(m, n) = 1$
- ▶ Find  $x$  with  $x \equiv a \pmod{m}$  and  $x \equiv b \pmod{n}$
- ▶ Solution:
  - Let  $n^{-1}$  be the inverse of  $n$  modulo  $m$
  - Let  $m^{-1}$  be the inverse of  $m$  modulo  $n$
  - Set  $x = ann^{-1} + bmm^{-1}$  (check this yourself)
- ▶ Extension: solving for more simultaneous equations

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## Binomial Coefficients

- ▶  $\binom{n}{k}$  is the number of ways to choose  $k$  objects out of  $n$  distinguishable objects
- ▶ same as the coefficient of  $x^k y^{n-k}$  in the expansion of  $(x + y)^n$ 
  - Hence the name “binomial coefficients”
- ▶ Appears everywhere in combinatorics

## Computing Binomial Coefficients

- ▶ Solution 1: Compute using the following formula:

$$\binom{n}{k} = \frac{n(n-1) \cdots (n-k+1)}{k!}$$

- ▶ Solution 2: Use Pascal's triangle
- ▶ Can use either if both  $n$  and  $k$  are small
- ▶ Use Solution 1 carefully if  $n$  is big, but  $k$  or  $n - k$  is small

# Fibonacci Sequence

- ▶ Definition:
  - $F_0 = 0, F_1 = 1$
  - $F_n = F_{n-1} + F_{n-2}$ , where  $n \geq 2$
- ▶ Appears in many different contexts

## Closed Form

- ▶  $F_n = (1/\sqrt{5})(\varphi^n - \bar{\varphi}^n)$ 
  - $\varphi = (1 + \sqrt{5})/2$
  - $\bar{\varphi} = (1 - \sqrt{5})/2$
- ▶ Bad because  $\varphi$  and  $\sqrt{5}$  are irrational
- ▶ Cannot compute the exact value of  $F_n$  for large  $n$
- ▶ There is a more stable way to compute  $F_n$ 
  - ... and any other recurrence of a similar form

## Better “Closed” Form

$$\begin{bmatrix} F_{n+1} \\ F_n \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} F_n \\ F_{n-1} \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^n \begin{bmatrix} F_1 \\ F_0 \end{bmatrix}$$

- ▶ Use fast exponentiation to compute the matrix power
- ▶ Can be extended to support any linear recurrence with constant coefficients

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# Geometry

- ▶ In theory: not that hard
- ▶ In programming contests: more difficult than it looks
- ▶ Will cover basic stuff today
  - Computational geometry in week 9

## When Solving Geometry Problems

- ▶ Precision, precision, precision!
  - If possible, don't use floating-point numbers
  - If you have to, always use `double` and never use `float`
  - Avoid division whenever possible
  - Introduce small constant  $\epsilon$  in (in)equality tests
    - ▶ e.g., Instead of `if(x == 0)`, write `if(abs(x) < EPS)`
- ▶ No hacks!
  - In most cases, randomization, probabilistic methods, small perturbations won't help



## 2D Vector Operations

- ▶ Have a vector  $(x, y)$
- ▶ Norm (distance from the origin):  $\sqrt{x^2 + y^2}$
- ▶ Counterclockwise rotation by  $\theta$ :

$$\begin{bmatrix} \cos \theta & -\sin \theta \\ \sin \theta & \cos \theta \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix}$$

- Make sure to use correct units (degrees, radians)
- ▶ Normal vectors:  $(y, -x)$  and  $(-y, x)$
- ▶ Memorize all of them!

## Line-Line Intersection

- ▶ Have two lines:  $ax + by + c = 0$  and  $dx + ey + f = 0$
- ▶ Write in matrix form:

$$\begin{bmatrix} a & b \\ d & e \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = - \begin{bmatrix} c \\ f \end{bmatrix}$$

- ▶ Left-multiply by matrix inverse

$$\begin{bmatrix} a & b \\ d & e \end{bmatrix}^{-1} = \frac{1}{ae - bd} \begin{bmatrix} e & -b \\ -d & a \end{bmatrix}$$

- Memorize this!
- ▶ Edge case:  $ae = bd$ 
  - The lines coincide or are parallel

## Circumcircle of a Triangle

- ▶ Have three points  $A, B, C$
- ▶ Want to compute  $P$  that is equidistance from  $A, B, C$
- ▶ Don't try to solve the system of quadratic equations!
- ▶ Instead, do the following:
  - Find the (equations of the) bisectors of  $AB$  and  $BC$
  - Compute their intersection

## Area of a Triangle

- ▶ Have three points  $A, B, C$
- ▶ Want to compute the area  $S$  of triangle  $ABC$
- ▶ Use cross product:  $2S = |(B - A) \times (C - A)|$
- ▶ Cross product:

$$(x_1, y_1) \times (x_2, y_2) = \begin{vmatrix} x_1 & x_2 \\ y_1 & y_2 \end{vmatrix} = x_1 y_2 - x_2 y_1$$

- Very important in computational geometry. Memorize!

## Area of a Simple Polygon

- ▶ Given vertices  $P_1, P_2, \dots, P_n$  of polygon  $P$
- ▶ Want to compute the area  $S$  of  $P$
- ▶ If  $P$  is convex, we can decompose  $P$  into triangles:

$$2S = \left| \sum_{i=2}^{n-1} (P_{i+1} - P_1) \times (P_i - P_1) \right|$$

- ▶ It turns out that the formula above works for non-convex polygons too
  - Area is the absolute value of the sum of “signed area”
- ▶ Alternative formula (with  $x_{n+1} = x_1, y_{n+1} = y_1$ ):

$$2S = \left| \sum_{i=1}^n (x_i y_{i+1} - x_{i+1} y_i) \right|$$

## Conclusion

- ▶ No need to look for one-line closed form solutions
- ▶ Knowing “how to compute” (algorithms) is good enough
- ▶ Have fun with the exercise problems
  - ... and come to the practice contest if you can!