

# 2026 MC Tennis Parity Puzzle

Charles Tan 2028'

## Setup and Observations

- Let  $p \in (0, 1)$  be the probability that Player A wins a *service* game (when A serves)
- Let  $q \in (0, 1)$  be the probability that Player A wins a *return* game (when B serves)
- The same player serves first in *every* set; denote this first server by  $S \in \{A, B\}$
- A set ends at score  $(a, b)$  with  $\max(a, b) = 6$  and  $|a - b| \geq 2$  (no 7+ games occur)
- Empirical constraints (the “parity observations”):
  - (a) If A wins the set, then  $a + b$  is *odd*
  - (b) If B wins the set, then  $a + b$  is *even*
- A won the match (best of  $2n+1$  sets)

## Single-Set DP Likelihood

I will model each set as a Markov chain on states  $(a, b, t)$ , where  $a, b \in \{0, \dots, 6\}$  are games won by A and B, and  $t \in \{0, 1\}$  is the next server ( $t=0 \Rightarrow$  A serves,  $t=1 \Rightarrow$  B serves) The initial server is determined by  $S$ : start at  $(0, 0, t=0)$  if  $S=A$ , or  $(0, 0, t=1)$  if  $S=B$

Transitions:

$$\begin{aligned} t=0 : (a, b, 0) &\rightarrow (a+1, b, 1) \text{ w.p. } p, & (a, b+1, 1) \text{ w.p. } 1-p, \\ t=1 : (a, b, 1) &\rightarrow (a+1, b, 0) \text{ w.p. } q, & (a, b+1, 0) \text{ w.p. } 1-q \end{aligned}$$

A terminal set occurs at any  $(a, b)$  with  $\max(a, b) = 6$  and  $|a - b| \geq 2$ . Let  $\mathbb{P}((a, b) \mid S)$  be the total absorbed mass at terminal  $(a, b)$  given  $S$

**Filter by parity observations, we define the following indicators:**

$$\begin{aligned} I_{(a,b)}^A &= \mathbf{1}[\max(a, b) = 6, |a - b| \geq 2, a > b, (a+b) \text{ odd}], \\ I_{(a,b)}^B &= \mathbf{1}[\max(a, b) = 6, |a - b| \geq 2, b > a, (a+b) \text{ even}] \end{aligned}$$

Then the likelihood that a single set is consistent with the observations and is won by A (resp. B) under first server  $S$  is

$$\pi_{\text{set},A}^{(S)} = \sum_{a,b} \mathbb{P}((a, b) \mid S) I_{(a,b)}^A, \quad \pi_{\text{set},B}^{(S)} = \sum_{a,b} \mathbb{P}((a, b) \mid S) I_{(a,b)}^B$$

## Match Likelihood and Posterior

For a best-of  $m = 2n+1$  sets with independent sets conditional on  $S$ , the probability that A wins the match *and* every set satisfies the observations is

$$L_S = \sum_{k=(m+1)/2}^m \binom{m}{k} (\pi_{\text{set,A}}^{(S)})^k (\pi_{\text{set,B}}^{(S)})^{m-k}$$

With prior  $\mathbb{P}(S=A) = \mathbb{P}(S=B) = \frac{1}{2}$ , the posterior is

$$\mathbb{P}(S=A \mid \mathcal{O}) = \frac{L_A}{L_A + L_B}$$

**Proposition 1** (Symmetry). *For any  $p, q \in (0, 1)$  and any integer  $n \geq 0$ ,  $L_A = L_B$ . Consequently,*

$$\mathbb{P}(S=A \mid \mathcal{O}) = \frac{1}{2}$$

*Sketch.* If  $S = A$ , admissible A-won sets (odd length) end on A's serve; B-won sets (even length) end on B's serve. If  $S = B$ , the same terminal scorelines occur but the serving labels are swapped at each game. The parity filter depends only on (winner, total-games parity), and there is a measure-preserving bijection between paths under  $S = A$  and  $S = B$ . Thus the filtered mass is identical, giving  $L_A = L_B$ .  $\square$

## Algorithms

Please find full code at my github: <https://github.com/CharlesTanShi/Tennis-MC.git>

### Exact DP for One Set

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#### Algorithm 1 SETLIKELIHOODS( $p, q, S$ )

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1: Initialize DP tensor  $P[a, b, t] \leftarrow 0$  for  $a, b \in \{0, \dots, 6\}$ ,  $t \in \{0, 1\}$ .
2:  $t_0 \leftarrow 0$  if  $S=A$  else 1;  $P[0, 0, t_0] \leftarrow 1$ .
3: function TERMINAL( $a, b$ ) return  $(\max(a, b) = 6) \wedge (|a - b| \geq 2)$ 
4: for  $a = 0..6$  do
5:   for  $b = 0..6$  do
6:     for  $t \in \{0, 1\}$  do
7:        $w \leftarrow P[a, b, t]$ ; if  $w = 0$  or TERMINAL( $a, b$ ) then continue
8:       if  $t = 0$  then ▷ A serves
9:          $P[a+1, b, 1] += w \cdot p$ ;  $P[a, b+1, 1] += w \cdot (1 - p)$ 
10:      else ▷ B serves
11:         $P[a+1, b, 0] += w \cdot q$ ;  $P[a, b+1, 0] += w \cdot (1 - q)$ 
12:  $\pi_A, \pi_B \leftarrow 0, 0$ 
13: for  $a = 0..6$  do
14:   for  $b = 0..6$  do
15:     if TERMINAL( $a, b$ ) and  $\max(a, b) = 6$  then
16:        $w \leftarrow P[a, b, 0] + P[a, b, 1]$ 
17:       if  $a > b$  and  $(a+b)$  odd then  $\pi_A += w$ 
18:       if  $b > a$  and  $(a+b)$  even then  $\pi_B += w$ 
19: return  $(\pi_A, \pi_B)$ 

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