1. (1) one set with 2 blocks, affect $0 \sim 6 \Rightarrow 5$ bits word offset $2^5 = 32$ words in a block lndex 5 bits \Rightarrow 32 sets totally set 2 block block $32 \sim 100$ word. $32 \sim 100$ set 2 block block $32 \sim 100$ set 2 block block $32 \sim 100$ set $20 \sim 100$ bits.

cache $size = 32 \times 2 \times (2 + 20 + 32 \times 32) = 66944$ bits

(2)		, word off	index	tag	miss/Hit
Ο,	0	0	Ö	O	M
4	(00		0	0	
20	0 00	101	Ò	B	H
136	000 000	(0	1	O	M
231_	1/1101000	11010)	O	H
164	[6100100	0 (00/	1	0	H
1024	(acc) 000 0cm	0	1000	0	M
30	11110	00111	0	0	
140	1000 1100	00011	1	0	H
ઢો૦૦	110000011100	00	11000	0	\mathcal{M}
176	10110000	0//00	J	0	H
2180	1000 0000)00	0000	1000	\circ	M

- 2. (1) Pr clock time: 1.18 ns = clock rate = $\frac{1}{1.18 \times 10^9}$ = 847457627.1 HzPz clock time: 2.22 ns => clock rate = $\frac{1}{2.22 \text{ ns}}$ = 450450450.5 Hz.
 - (2) AMAT $P_1 = 1.18 + 4.3\%. \cdot 70 = 4.19 \text{ ns}$ AMAT $P_2 = 2.22 + 2.7\%. \cdot 70 = 4.11 \text{ ns}$
 - (3) $CPIP_1 = |+ 36\% \cdot 4.3\% \cdot \frac{70}{1.18} = |.9|8$ $CPIP_2 = |+ 36\% \cdot 2.7\% \cdot \frac{70}{2.22} = |.306|$

Time Pi: $(.918 \times 1.18 = 2.26 \text{ ns})$, Time $P_2 = 1.306 \times 2.22 = 2.42 \text{ ns}$. .: Pi is better

3(1) 3 blocks per set, 2 words per block 24 words => 4 sets.

27 tag, 2 set intex, 1 worf, 2 byte off
index bits, tag bits, offset bits, H/M

M3: 00 10 0 1 0 011 180; 10. \mathcal{M} 101 100 43: 0 M(011 3: 00 \mathcal{D} $O \mid l$ H (0) | 11 M 191: 11 8J: 11 10 001 W 190: [0] | 10 H M 14: 0 10 b2. [8]; 0 0 H 44: () H 100 H 186: 11 0 0 [6] 252: 1 [1] 60 M b3.

Index Tag WO WI Tag WO WI Tag WO WI.

00 27'60 M[0] M[1]

01 (26'60)1 M[10] M[11] 27'60 M[2] M[3]

10 (24'60)0/ M(44] M(45)

1 1 (24'60)101 M(46) M(47] (25'60)11 M(22] M(23) (24'60)17 M(62) M(63)

I word per block, & blocks per set. I set. index bits, tag bits, H/M \mathcal{M} bi V 3 - \bigcirc 10/10/ \mathcal{N} b2 ∨ 180 : b3. \mathcal{M} 100 43? H 3 : \bigcirc 191: 101/11 M64 M89: 010 b5 196: 0|||0 \mathcal{M} 14: [[6 -181: 10/10/ \mathcal{M} 44: [0] \bigvee 186: 0/110 PS 252: 111111 replace bs. $M[16] \rightarrow M[63]$ mis rate= Tag Tag W lag W Tag Index 0 W \mathcal{W} MEO] 30b0 (24/60)14101 M45] (24'b0)||111 Mb3] (24'b0)|0||11 M[47] Tag Tag W Tag W Tag \bigvee \mathbb{W}

(28'bo) 11 M[3] (26'b0) 1011 M[11] (24'b0)101110 M[46]

(25'bo) | ollo

M[22]

(3)	one set,	4 blocks w	ith 2 wo	rds.	
	tag	word at.	H/M	(ZRU)	
3;	0	0	M	Ы	
180%	10/10	1	1/2	bz	
43:		0	\mathcal{M}	13	
3:	O	D	H	bi v.	
191:	101/1		\mathcal{M}	M	
89:	[0]	D	\mathcal{M}	b2.	
190:	0	7	1-1	b4 v	
14:		1	\mathcal{M}	b3	
18]:	10110	l	\bigwedge	bı	
44;	0		//	b2	
186:	0	Ò	H	b4 v	
252:		1	\mathcal{M}	63	niss rate = $\frac{9}{12}$ = 0.75
	tag	word dt.	H/M	(MRU).	
3;	G	0		ы	
180%	10/10	l	\bigwedge	2	
43:	0	Ó	\mathcal{M}	b3 /	
3:	O	O	H	Þı V	
191:	10111		\mathcal{M}	b4 \	
89:	[0]	D	\mathcal{M}	b4	
190:	0	7	/ /\		
14:		1	M	b24	
8] :	10 10		H	b2_	
,	(0)		-	63	
186:	0	0	\mathcal{M}	la3 _	C ₁
252: 14) 	Miles parte	b3	$y_{\text{MSS}} \text{ rate} = \frac{9}{12} = 0.75$

If choosing LRU or MRU, Miss rate will be 0.75, if combine LRU and MRU, we can have at most 5 hits, so miss rate will be $\frac{7}{12} = 0.583$.