

COMP3821

Maximal Flow Practice

Charlie Bradford z5114682

June 21, 2018

1. You are running a dating agency and have m guys and f girls as customers. Each guy and each girl have reviewed all profiles of the candidates of opposite sex and have sent you their corresponding lists of people whose profiles they liked. Your task is to organise a largest possible number of first dates so that everyone meets at most one person of opposite sex and so that both parties liked each others profile (i.e., have included the other party on their list of people whose profile they liked).
 - Connect a super-source to all the women
 - Set all outgoing edges of the super-source to have capacity one
 - Connect a super-sink to all the men
 - Set all incoing edges of the super-sink to have capacity one
 - Create edges of capacity one between all the women and men that like eachother
 - Use the Edmonds-Karp algorithm to find the maximal flow of the graph
 - All links that have flow are first dates between their respected nodes
2. Assume that you are the administrator of a network of computers; each computer is connected by unidirectional fiberoptic cables to a few other computers on the same network (so the network can be modeled by a directed graph). You noticed that computers P_1, P_2, \dots, P_n are mounting an attack on computers Q_1, Q_2, \dots, Q_m . The total number of computers on the network is $N > m + n$. Since it is a real emergency, you must disconnect some of the optical cables of the network so that none of computers P_1, P_2, \dots, P_n can send packets to any of Q_1, Q_2, \dots, Q_m . Since you must send crews to disconnect some of the fiberoptic cables, for each cable c_{ij} for traffic from a computer X_i to a computer X_j there is an associated cost c_{ij} for disconnecting it. Your task is to design an algorithm for determining which cables to disconnect to isolate computers Q_1, Q_2, \dots, Q_m from all of the computers P_1, \dots, P_n so that the total cost incurred is minimal.
 - Create a maximal flow graph, with each node being a computer and each edge being being a link with capacity equal to its bandwidth
 - Connect a super-source to all computers P
 - Set the capacity of the edges between the super-source and each node to be equal to the total outgoing capacity of the node
 - Connect a super-sink to all computers Q
 - Set the capacity of the edges between each node and the super-sink to be equal to the total incoming capacity of the node
 - Use the Edmonds-Karp algorithm to find the maximal flow
 - All nodes accesible via augmenting paths from the super sink form one side of the minimal cut
 - Disconnect all outgoing cables from those nodes that have flow
 - As this is the minimal cut, you have disconnected the cables with least total bandwidth, but still cut computers P of from computers Q

3. You work for a new private university which wants to keep the sizes of classes small. Each class is assigned its maximal capacity - the largest number of students which can enrol in it. Students pay the same tuition fee for each class they get enrolled in. Students can apply to be enrolled in as many classes as they wish, but each of them will eventually be enrolled to at most 5 classes at any given semester. You are given the wish lists of all students, containing for each student the list of all classes they would like to enrol this particular semester and you have to choose from the classes they have put on their wish lists in which classes you will enrol them, without exceeding the maximal enrollment of any of the classes and without enrolling any student into more than 5 classes. Your goal is, surprisingly, to maximise the income from the tuition fees for your university. Design an efficient algorithm for such a task.

- Create a source that connects to all students
- Set capacity of links between source and each student to 5
- Create sink that connects to all classes
- Set capacity of links between sink and classes equal to the max number of students in that class
- Connect students to all classes on their wish list with links of capacity 1
- Use Edmonds-Karp algorithm to find maximal flow
- If there is flow on a link between a student and a class then the student is in that class

4. Assume each student can borrow at most 10 books from the library, and the library has three copies of each title in its inventory. Each student submits a list of books he wishes to borrow. You have to assign books to students, so that a maximal number of volumes is checked out.

- Source connects to all student
- Capacity 10 between source and student
- Students connect to all books on wishlist
- Books all connect to a sink, with links capacity 3
- Find max flow
- Flow between book and student means student borrows book

5. The emergency services are responding to a major earthquake that has hit a wide region, and left n people injured who need to be sent to a hospital. Let P be the set of n people and H be the set of k hospitals. Several hospitals are available to treat these people, but there are some constraints:

- (a) Each injured person needs to be sent to a hospital no further than one hour drive away. Let H_p be the set of hospitals that are within range for person p .
- (b) Each hospital h has a capacity c_h , the maximum number of people that the hospital can receive.

Design an efficient algorithm that determines whether it is possible to assign each person to a hospital in a way that satisfies these constraints, and returns such an assignment if so.

- Connect source to all people
- Capacity between source and each person is 1
- Connect each person p to all hospitals H_p
- Connect each hospital h to a sink, link capacity c_h
- Find max flow
- If max flow is n , then task is possible
- Flow between a person and a hospital means that person goes to that hospital