CSC 212: Data Structures and Abstractions 06: Dynamic Arrays

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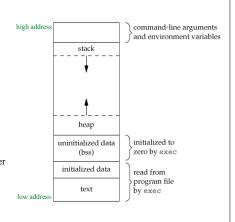
Memory model

- What is the C/C++ memory model?
 - a formal specification that defines how programs interact with memory at a high level, ensuring safe and predictable behavior
 - implementation details are delegated to the compiler and CPU architecture
 - the memory model establishes a "contract" between programmer
 and compiler
- Memory Layout
 - memory is divided into multiple segments
 - each segment serves a specific purpose and has different properties

C/C++ memory model

Memory layout

- · Text Segment (code)
 - contains instructions generated by the compiler
 - marked as <u>read-only</u> to prevent accidental modification
- · Data Segment (global/static variables)
 - contains multiple subsections (e.g. initialized data, uninitialized data, constant data)
 - size determined at compilation, addresses resolved during linking
- Heap
 - ✓ dedicated to dynamic memory allocation
 - requires explicit management by the programmer
- · Stack (function calls, local variables)
 - implements last-in-first-out (LIFO) for function calls and local variables
 - √ no explicit deallocation required



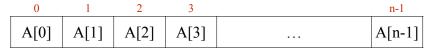
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```
#include <iostream>
                                                                          ./prog | sort -k 4
float pi = 3.1416;
const int min = 100:
// uninitialized global variable
                                                                     address of foo 0x105499fc0
int sum;
                                                                     address of main 0x10549a0f0
                                                                     address of min 0x10549af48
void foo(int arg) {
                                                                     address of pi 0x10549c000
                                                                     address of sum 0x10549c004
   int i = 1;
   std::cout << "address of arg\t" << &arg << std::endl;</pre>
                                                                     value of A
                                                                                         0x7fd29f705e90
   std::cout << "address of i\t" << &i << std::endl;</pre>
                                                                     address of i
                                                                                         0x7ff7baa66438
                                                                     address of arg 0x7ff7baa6643c
                                                                     address of A 0x7ff7baa66460
int main() {
                                                                    NOTE: leading zeros are ignored (64-bit addresses)
   int *A = new int[10];
   std::cout << "address of pi\t" << &pi << std::endl;</pre>
   std::cout << "address of min\t" << &min << std::endl;</pre>
   std::cout << "address of sum\t" << &sum << std::endl;</pre>
                                                                            Can you tell what are the
    std::cout << "value of A\t" << A << std::endl;</pre>
   std::cout << "address of A\t" << &A << std::endl;</pre>
                                                                            memory locations grouped
   std::cout << "address of main\t" << (void*) &main << std::endl;</pre>
                                                                               by different colors?
   std::cout << "address of foo\t" << (void*) &foo << std::endl;</pre>
   foo(5):
                                                                            What happens if you run
   delete [] A;
                                                                           the program multiple times?
```

Dynamic arrays

C-style arrays

- Contiguous sequence of elements of identical type
 - random access: base_address + index * sizeof(type)



array name: A

array length: n

- Statically allocated arrays
 - ✓ allocated in the stack (fixed-length), size known at compile time
- Dynamic allocated arrays
 - ✓ allocated in the heap (fixed-length), size may be determined at runtime

Dynamic arrays

- Limitations of C-style arrays
 - size must be known at compile time or use dynamic memory allocation
 once created the array size does not change
 - \checkmark provide $\Theta(1)$ read/write cost, but inflexible
- Dynamic arrays
 - √ can grow or shrink in size during runtime
 - essential for many applications, for example, a server keeping track of a queue of requests
 - combine the flexibility of dynamic memory allocation with the efficiency of fixed-length arrays
 - ✓ e.g. std::vector in C++, ArrayList in Java, List in Python, Array in JavaScript, List in C#, Vec in Rust, etc.

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Dynamic array class in C++

```
class DynamicArray {
    private:
        int *arr;
                                              // pointer to the (internal) array
                                              // total number of elements that can be stored
        int capacity;
        int size:
                                              // number of elements currently stored
    public:
        DynamicArray();
        ~DynamicArray();
                                             // destructor
        void push_back(int val);
                                             // add an element to the end
        void pop_back();
                                              // remove the last element
        const int& operator[](int idx) const; // read-only access at a specific index
        int& operator[](int idx);
                                             // access at a specific index (can modify)
        void insert(int val, int idx);
                                             // insert an element at a specific index
        void erase(int idx):
                                              // remove an element at a specific index
        void resize(int len);
                                              // change the capacity of the array
        int size():
                                              // return the number of elements
                                             // return the capacity
        int capacity();
        bool empty();
                                              // check if the array is empty
        void clear();
                                              // remove all elements, maintaining the capacity
        // additional methods can be added here
};
```

A class definition specifies the **data members** and **member functions** of the class. The data members are the attributes of the class, and the member functions are the operations that can be performed on the data members. The class definition is a blueprint for creating objects of the class.

Grow by one

- When array is full, new capacity: current + 1
 - ✓ starting from an empty array, <u>count number of array accesses</u> (**reads and writes**) for adding *n* elements (ignore cost of allocating/deallocating memory)

n	сору	append	
1	2 x 0	1	
2	2 x 1	1	
3	2 x 2	1	
4			
5			
6			
n-1			
n			

$$T(n) = n + \sum_{i=0}^{n-1} 2i$$

$$= n + 2\left(\frac{n(n-1)}{2}\right)$$

$$= \Theta(n^2) \xrightarrow{\text{cost of addin} \\ \text{n elements}}$$

The amortized cost of inserting an element is $\Theta(n)$, meaning that any sequence of n insertions takes at most $\Theta(n^2)$ time in total.

Resizing dynamic arrays

Grow

when the array is full (Size == capacity), allocate a new array with increased capacity, copy elements from old to new array, deallocate old array

• Shrink

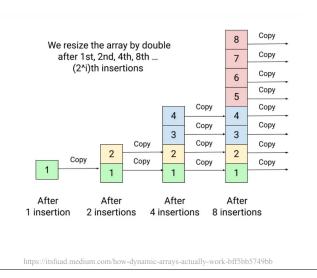
optional optimization, used when the number of elements is "significantly" less than the capacity, allocate a new array with decreased capacity, copy the elements from old to new array, and deallocate the old array



4 3 -5 0 9 -2 7 1

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Repeated doubling



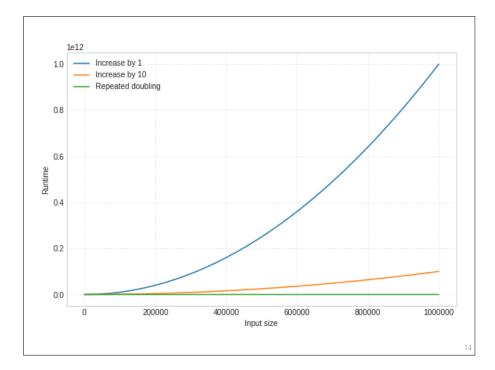
Grow by factor

- → When array is full, new capacity: current * factor
 - repeated doubling when factor == 2
 - ✓ starting from an empty array, <u>count number of array accesses</u> (**reads and writes**) for adding *n* elements assume n is a power of 2 (ignore cost of allocating/deallocating memory)

n	copy	append	
1	0	1	
2	2 * 1	1	
3	2 * 2	1	
4	0	1	
5	2 * 4	1	
6	0	1	
7	0	1	
8	0	1	
9	2 * 8	1	
10	0	1	
n-1			
n			

$T(n) = n + 2\sum_{i=0}^{\log n - 1} 2^{i}$
$= n + 2\left(2^{\log n} - 1\right)$
= n + 2n - 2
$=\Theta(n)$ \leftarrow cost of adding n elements

The amortized cost of inserting an element is $\Theta(1)$, meaning that any sequence of n insertions takes at most $\Theta(n)$ time in total.



Shrinking the array

- May half the capacity when array is **one-half** full
 - worst-case when the array is full and we <u>alternate between adding and</u> removing elements
 - each alternating operation would require resizing the array
- · More efficient resizing
 - √ half the capacity when the array is one-quarter full
- In practice ...
 - most standard implementations do not automatically shrink capacity
 - avoids performance penalties from frequent resizing
 - instead, they provide explicit operations like shrink_to_fit() that allow the programmer to request size reduction when deemed necessary

Live coding ...

• Complete the implementation of the DynamicArray class:

```
class DynamicArray {
   private:
        int *arr:
                                                 // pointer to the (internal) array
        size_t capacity;
                                                 // total number of elements that can be stored
        size t size:
        float growth_rate;
        DynamicArray(float growth_rate);
        ~DynamicArray();
                                                 // destructor
        void push_back(int val);
                                                 // add an element to the end
        void pop_back();
                                                 // remove the last element
        const int& operator[](size t idx) const; // read-only access at a specific index
        int& operator[](size_t idx);
                                                 // access at a specific index (can modify)
        void insert(int val, size_t idx);
                                                 // insert an element at a specific index
       void erase(int uidx);
                                                 // remove an element at a specific index
        void resize(int ulen);
                                                 // change the capacity of the array
        size t get size();
                                                 // return the number of elements
       size_t get_capacity();
                                                 // return the capacity
       bool empty();
                                                 // check if the array is empty
       void clear();
                                                 // remove all elements, maintaining the capacity
        // additional methods can be added here
};
```

References in C++ non-const References int i = 2;int& ri = i; // reference to i ri and i refer to the same object / memory location: cout << i <<'\n'; // 2 cout << ri <<'\n'; // 2 cout << i <<'\n'; // 5 cout << ri <<'\n'; // 5 ri = 88; cout << i <<'\n'; // 88 cout << ri <<'\n'; // 88 · references cannot be "null", i.e., they must always refer to an object · a reference must always refer to the same memory location · reference type must agree with the type of the referenced object int& ri = i; // reference to i // assigns value of k to i (target of ri) // * COMPILER ERROR: reference must be initialized double& r3 = i; // x COMPILER ERROR: types must agree

Growth factors by language

- . C++ (std::vector)
 - ✓ grow by 1.5 times the current capacity
 - ✓ shrink when the array is one-quarter full
- Java (ArrayList)
 - grow by 1.5 of the current capacity
 - ✓ shrink when the array is one-half full
- Python (list)
 - √ grow by 1.125 times the current capacity
 - ✓ shrink when the array is one-quarter full
- Rust (std::vec::Vec)
 - grow by 2 times the current capacity
 - ✓ shrink when the array is one-half full

Information taken from claude.ai (to be confirmed)

bottom line: growth factors range from ~1.2 to ~2 depending on language used

Practice

 \cdot Complete the following table with rates of growth using Θ notation

https://hackingcpp.com/cpp/lang/references.html

✓ assume we implement a dynamic array with repeated doubling and no shrinking

Best case	Average case	Worst case
	Best case	Best case Average case

Dynamic arrays in the STL (C++)

std::vector

```
Defined in header <vector>
template<
    class T
                                                                             (1)
    class Allocator = std::allocator<T>
> class vector;
namespace pmr
    template< class T >
                                                                              (2) (since C++17)
    using vector = std::vector<T, std::pmr::polymorphic_allocator<T>>;
```

- 1) std::vector is a sequence container that encapsulates dynamic size arrays.
- 2) std::pmr::vector is an alias template that uses a polymorphic allocator.

The elements are stored contiguously, which means that elements can be accessed not only through iterators, but also using offsets to regular pointers to elements. This means that a pointer to an element of a vector may be passed to any function that expects a pointer to an element of an array.

The storage of the vector is handled automatically, being expanded as needed. Vectors usually occupy more space than static arrays, because more memory is allocated to handle future growth. This way a vector does not need to reallocate each time an element is inserted, but only when the additional memory is exhausted. The total amount of allocated memory can be queried using capacity() function. Extra memory can be returned to the system via a call to

Reallocations are usually costly operations in terms of performance. The reserve() function can be used to eliminate reallocations if the number of elements is known beforehand.

The complexity (efficiency) of common operations on vectors is as follows:

- Random access constant O(1).
- Insertion or removal of elements at the end amortized constant $\mathcal{O}(1)$.
- Insertion or removal of elements linear in the distance to the end of the vector $\mathcal{O}(n)$.

std::vector (for T other than bool) meets the requirements of Container, Allocator Aware Container (since C++11), SequenceContainer, ContiguousContainer(since C++17) and ReversibleContainer.

https://en.cppreference.com/w/cpp/container/vector

std::vector

```
#include <iostream>
#include <vector>
int main()
    // create a vector containing integers
    std::vector<int> v = \{8, 4, 5, 9\};
    // add two more integers to vector
    v.push back(6);
    v.push_back(9);
    // overwrite element at position 2
    v[2] = -1;
    // print out the vector
    for (int n : v)
        std::cout << n << ' ';
    std::cout << '\n';</pre>
```

Member functions constructs the vector (constructor)

destructs the vector (destructor) assigns values to the container operator= assigns values to the container assign_range (C++23) assigns a range of values to the container returns the associated allocator get allocator

cend (C++11)

rbegin crbegin (C++11)

crend (C++11)

Element access access specified element with bounds checking access specified element operator[] access the first element front access the last element direct access to the underlying contiguous storage Iterators begin cbegin (C++11) returns an iterator to the beginning returns an iterator to the end

returns a reverse iterator to the beginning

returns a reverse iterator to the end

Capacity

checks whether the container is empty returns the number of elements returns the maximum possible number of elements max size reserves storage reserve public member function) eturns the number of elements that can be held in currently allocated storage reduces memory usage by freeing unused memory shrink to fit(DR*) Modifiers inserts elements insert

inserts a range of elements insert_range(C++23) constructs element in-place emplace (C++11) adds an element to the end push back constructs an element in-place at the end emplace back(C++11) append_range (C++23) adds a range of elements to the end

removes the last element pop_back changes the number of elements stored resize swaps the contents

https://en.cppreference.com/w/cpp/container/vector