NanoWasm Specification

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NanoWasm is a small language with simple types and instructions.

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ONE

ABSTRACT SYNTAX

The *abstract syntax* of types is as follows:

```
mut ::= mut
valtype ::= i32 \mid i64 \mid f32 \mid f64
functype ::= valtype^* \rightarrow valtype^*
globaltype ::= mut^? valtype
```

Instructions take the following form:

```
\begin{array}{llll} const & ::= & 0 \mid 1 \mid 2 \mid \dots \\ instr & ::= & \mathsf{nop} \\ & \mid & \mathsf{drop} \\ & \mid & \mathsf{select} \\ & \mid & \mathit{valtype}.\mathsf{const} \ const \\ & \mid & \mathsf{local}.\mathsf{get} \ localidx \\ & \mid & \mathsf{local}.\mathsf{set} \ localidx \\ & \mid & \mathsf{global}.\mathsf{get} \ globalidx \\ & \mid & \mathsf{global}.\mathsf{set} \ globalidx \end{array}
```

The instruction nop does nothing, drop removes an operand from the stack, select picks one of two operands depending on a condition value. The instruction t const c pushed the constant c to the stack. The remaining instructions access local and global variables.

CHAPTER

TWO

VALIDATION

NanoWasm instructions are type-checked under a context that assigns types to indices:

$$context ::= \{globals \ globaltype^*, locals \ valtype^*\}$$

2.1 nop

The nop nop is valid with the $\epsilon \to \epsilon \epsilon \to \epsilon$.

$$\overline{C \vdash \mathsf{nop} : \epsilon \to \epsilon}$$

2.2 drop

The drop drop is valid with the $t \to \epsilon \, t \to \epsilon$.

$$\overline{C \vdash \mathsf{drop} : t \to \epsilon}$$

2.3 select

The select select is valid with the $t\ t$ i32 $\ \rightarrow\ t\ t\ t$ i32 $\ \rightarrow\ t$.

$$\overline{C \vdash \mathsf{select} : t \ t \ \mathsf{i32} \to t}$$

2.4 const

The $(t.\mathsf{const}\ c)\ (t.\mathsf{const}\ c)$ is valid with the $\epsilon\ \to\ t\ \epsilon\ \to\ t$.

$$\overline{C \vdash t.\mathsf{const}\ c : \epsilon \to t}$$

2.5 local.get

The (local.get x) (local.get x) is valid with the $\epsilon \to t \, \epsilon \to t$ if:

- The C.locals[x] C.locals[x] exists.
- The C.locals[x] C.locals[x] is of the form t.

$$\frac{C.\mathsf{locals}[x] = t}{C \vdash \mathsf{local.get} \ x : \epsilon \to t}$$

2.6 local.set

The (local.set x) (local.set x) is valid with the $t \to \epsilon t \to \epsilon$ if:

- The C.locals[x] C.locals[x] exists.
- The $C.\mathsf{locals}[x]$ $C.\mathsf{locals}[x]$ is of the form t.

$$\frac{C.\mathsf{locals}[x] = t}{C \vdash \mathsf{local set} \ x : t \to \epsilon}$$

2.7 global.get

The (global.get x) (global.get x) is valid with the $\epsilon \to t \, \epsilon \to t$ if:

- The C.globals[x] C.globals[x] exists.
- The C.globals[x] C.globals[x] is of the form (mut[?] t).

$$\frac{C.\mathsf{globals}[x] = \mathsf{mut}^?\ t}{C \vdash \mathsf{global.get}\ x : \epsilon \to t}$$

2.8 global.set

The (global.get x) (global.get x) is valid with the $t \to \epsilon t \to \epsilon$ if:

- The C.globals[x] C.globals[x] exists.
- The C.globals[x] C.globals[x] is of the form (mut t).

$$\frac{C.\mathsf{globals}[x] = \mathsf{mut}\ t}{C \vdash \mathsf{global.get}\ x : t \to \epsilon}$$

THREE

EXECUTION

NanoWasm execution requires a suitable definition of state and configuration:

```
addr
                     0 | 1 | 2 | ...
module inst
                     {globals addr^*}
                     {\tt const}\ valtype\ const
         val
                     \{globals \ val^*\}
       store
              ::=
                     \{locals \ val^*, module \ module inst \}
     frame
       state
                     store; frame
     config
              ::=
                     state; instr^*
```

We define the following auxiliary functions for accessing and updating the state:

```
\begin{array}{lll} \operatorname{local}((s;f),x) & = & f.\operatorname{locals}[x] \\ \operatorname{global}((s;f),x) & = & s.\operatorname{globals}[f.\operatorname{module.globals}[x]] \\ \operatorname{update}_{local}((s;f),x,v) & = & s;f[.\operatorname{locals}[x]=v] \\ \operatorname{update}_{qlobal}((s;f),x,v) & = & s[.\operatorname{globals}[f.\operatorname{module.globals}[x]]=v];f \end{array}
```

With that, execution is defined as follows:

3.1 nop

1. Do nothing.

$$\mathsf{nop} \;\hookrightarrow\; \epsilon$$

3.2 drop

- 1. Assert: Due to validation, a value is on the top of the stack.
- 2. Pop the value val from the stack.

$$val \; \mathsf{drop} \;\; \hookrightarrow \;\; \epsilon$$

3.3 select

- 1. Assert: Due to validation, value type is on the top of the stack.
- 2. Pop the value (i32.const c) from the stack.
- 3. Assert: Due to validation, a value is on the top of the stack.
- 4. Pop the value val_2 from the stack.
- 5. Assert: Due to validation, a value is on the top of the stack.
- 6. Pop the value val_1 from the stack.

- 7. If $c \neq 0$, then:
 - a. Push the value val_1 to the stack.
- 8. Else:
 - a. Push the value val_2 to the stack.

```
val_1 \ val_2 \ (\mbox{i32.const} \ c) \ \mbox{select} \quad \hookrightarrow \quad val_1 \quad \mbox{if} \ c \neq 0 val_1 \ val_2 \ (\mbox{i32.const} \ c) \ \mbox{select} \quad \hookrightarrow \quad val_2 \quad \mbox{otherwise}
```

3.4 local.get x

- 1. Let z be the current state.
- 2. Let val be local(z, x).
- 3. Push the value *val* to the stack.

$$z; (\mathsf{local}.\mathsf{get}\ x) \quad \hookrightarrow \quad z; val \quad \text{if}\ val = \mathrm{local}(z,x)$$

3.5 local.set x

- 1. Assert: Due to validation, a value is on the top of the stack.
- 2. Pop the value val from the stack.

$$z; val \text{ (local.set } x) \hookrightarrow z'; \epsilon \text{ if } z' = \text{update}_{local}(z, x, val)$$

3.6 global.get x

- 1. Let z be the current state.
- 2. Let val be global(z, x).
- 3. Push the value *val* to the stack.

$$z; (\mathsf{global}.\mathsf{get}\; x) \;\; \hookrightarrow \;\; z; val \quad \mathsf{if}\; val = \mathsf{global}(z,x)$$

3.7 global.set x

- 1. Assert: Due to validation, a value is on the top of the stack.
- 2. Pop the value val from the stack.

```
z; val \text{ (global.set } x) \hookrightarrow z'; \epsilon \text{ if } z' = \text{update}_{qlobal}(z, x, val)
```

FOUR

BINARY FORMAT

The following grammars define the binary representation of NanoWasm programs.

First, constants are represented in LEB format:

Types are encoded as follows:

```
valtype ::=
                     0x7F
                                                                     i32
                     0x7E
                                                                     i64
                     0x7D
                                                                     f32
                     0x7C
                                                                     f64
         mut ::= 0x00
                     0x01
                                                                     mut
\verb|globaltype| ::= t: \verb|valtype| mut: \verb|mut|
                                                                     mut\ t
resulttype ::= n:u32 (t:valtype)^n
  functype := 0x60 t_1^*:resulttype t_2^*:resulttype \Rightarrow
                                                                     t_1^* \rightarrow t_2^*
```

Finally, instruction opcodes:

```
globalidx ::= x:u32
                                                           \boldsymbol{x}
 localidx ::= x:u32
                                                           \boldsymbol{x}
      instr := 0x01
                                                           nop
                       0x1A
                                                    \Rightarrow drop
                        0x1B
                                                    \Rightarrow select
                        0x20 x:localidx \Rightarrow
                                                          local.get x
                        \texttt{0x21} \ x \texttt{:localidx} \quad \Rightarrow \quad \mathsf{local.set} \ x
                        \texttt{0x23} \ x : \texttt{globalidx} \ \Rightarrow \ \ \texttt{global.get} \ x
                        0x24 x:globalidx \Rightarrow global.set x
                        0x41 n:u32 \Rightarrow i32.const n
                        0x42 n:u64
                                                 \Rightarrow i64.const n
                                                \Rightarrow f32.const p
                        0x43 p:f32
                                                   \Rightarrow f64.const p
                        0x44 p:f64
```