

# CS323 Project1

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## Lexical Analysis & Syntax Analysis

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## I. Overview

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In this project, we are required to implement lexical analysis and syntax analysis with lexer, syntaxer and other useful files in order to output a parser tree or possible lexical and syntax error information for given SPL(Sustech Programming Language) code. Our files can be run successfully with GCC version 7.4.0, GNU Flex version 2.6.4 and GNU Bison version 3.0.4 .

## II. Design and Implementation

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### A. Lexer

In lexer part, we define a new class named `sp1_node` which record the information of matched token for syntax analysis and final output.

```
enum class Node_TYPE{
    INT,
    FLOAT,
    CHAR,
    NAME,    //IF ELSE ASSIGN etc
    ID,      //ALL IDENTIFIERS
    DTYPE,   //DATATYPE, INCLUDING INT,FLOAT & CHAR
    LINE
};

class Node{
private:
    string name;
    //string id_name;// == string_value
    Node_TYPE TYPE;
    union{
        int lineno;
        int int_value;
        float float_value;
        char char_value;
    };
    vector <Node*> child;

public:
    Node();

    void print(int depth);
```

```

Node(int val);
Node(float val);
Node(char val);

Node(string name);
Node(string name, Node_TYPE type);
Node(string name, int line_no);
Node(string name, int line_no, vector<Node*>& child);

void set_child(vector<Node*>&);

void show(int depth);

};

```

We define the variable `has_err` to record whether there exists possible lexical and syntax error for final output. We also support single line comment symbol `//` besides given matching rules. We declare the illegal and error situations in detail so that we can handle all possible errors. All lexical error will be reported here with line number.

```

40  %%
41  "/" { char c = yyinput();while(c!='\n'){c=yyinput();}}
42  "\n" {yycolumn = 1;}
43  "int" { yylval.value = new Node("INT",Node_TYPE::DTYPE);return TYPE; }
44  "float" { yylval.value = new Node("FLOAT",Node_TYPE::DTYPE);return TYPE; }
45  "char" { yylval.value = new Node("CHAR",Node_TYPE::DTYPE);return TYPE; }
46  "struct" { yylval.value = new Node("STRUCT");return STRUCT; }
47  "if" { yylval.value = new Node("IF");return IF; }
48  "else" { yylval.value = new Node("ELSE");return ELSE; }
49  "while" { yylval.value = new Node("WHILE");return WHILE; }
50  "return" { yylval.value = new Node("RETURN");return RETURN; }
51  "." { yylval.value = new Node("DOT");return DOT; }

```

```

42  INT_10_ERR_OUTRANGE [1-9]{num}+
43  INT_10_ERR_0STA 0{num}+
44
45  CHAR_ERR_TOOLONG '.{2,}'
46
47  INT_16_ERR_CHARILL 0[xX](({num_16}*[g-zG-Z])+{num_16}*)
48  INT_16_ERR_MORE0 0[xX](0+{num_16}+)
49  INT_16_ERR_TOOLONG 0[xX]([1-9A-Fa-f]{num_16}{8,})
50
51  CHAR_16_ERR_TOOLONG '\\x.{3,}'
52  CHAR_16_ERR_CHARILL '\\x{num_16}*[G-Zg-z]+{num_16}*'
53  CHAR_16_ERR_MORE0 '\\x0+{num_16}'
54
55
56  ID_ERR_TOOLONG {letter_}({letter_}|{num_16}){32,}
57  ID_ERR_NUMSTA [0-9]({letter_}|{num_16}){1,31}

```

```

145     {CHAR_ERR_TOOLONG} {
146         has_err = 1;
147         fprintf(LEX_ERR_OP,"Error type A at Line %d: too many chars\n",yylineno);
148         return ERR_TOKEN;
149     }
150
151
152     {INT_16_ERR_TOOLONG} {
153         has_err = 1;
154         fprintf(LEX_ERR_OP,"Error type A at Line %d: int number overflow \n",yylineno);
155         return ERR_TOKEN;
156     }
157     {INT_16_ERR_CHARILL} {
158         has_err = 1;
159         fprintf(LEX_ERR_OP,"Error type A at Line %d: hexadecimal int with illegal char\n",yylineno);
160         return ERR_TOKEN;
161     }

```

Figure.1 Some matching rules we defined in `lex.1`

## B. Syntaxer

As mentioned in Appendix B, we construct our syntaxer which will accept tokens and make actions or report errors(type B). In this part we will take use of class `Node` not only for tokens' information but also for level differentiation. Finally the nodes will form up a tree to record those tokens' information when the program has no lexical and syntax error.

In syntaxer, we construct `vector<Node*> vec` to record the child nodes of the current node. If necessary, we can traverse the tree from root and output the whole parser tree as required.

```

/* expression */
Exp:
    Exp ASSIGN Exp { vector<Node*> vec = {$1, $2, $3}; $$ = new Node("Exp", @$$.first_line, vec); }
    | Exp AND Exp { vector<Node*> vec = {$1, $2, $3}; $$ = new Node("Exp", @$$.first_line, vec); }
    | Exp OR Exp { vector<Node*> vec = {$1, $2, $3}; $$ = new Node("Exp", @$$.first_line, vec); }
    | Exp LT Exp { vector<Node*> vec = {$1, $2, $3}; $$ = new Node("Exp", @$$.first_line, vec); }
    | Exp LE Exp { vector<Node*> vec = {$1, $2, $3}; $$ = new Node("Exp", @$$.first_line, vec); }
    | Exp GT Exp { vector<Node*> vec = {$1, $2, $3}; $$ = new Node("Exp", @$$.first_line, vec); }
    | Exp GE Exp { vector<Node*> vec = {$1, $2, $3}; $$ = new Node("Exp", @$$.first_line, vec); }
    | Exp NE Exp { vector<Node*> vec = {$1, $2, $3}; $$ = new Node("Exp", @$$.first_line, vec); }
    | Exp EQ Exp { vector<Node*> vec = {$1, $2, $3}; $$ = new Node("Exp", @$$.first_line, vec); }

```

Figure.2 Syntax Design

## C. Other useful files

`sp1_node.cpp` and `sp1_node.hpp` are used to define the class `Node` and declare the fields and functions about it.

`main.cpp` are main function to start parsing and output parse tree.

## III. Test Cases

For evaluation purpose, our test cases contain 1 correct code, 2 type A errors and 2 type B errors. Including missing right curly braces and illegal identifier.

## IV. Instructions

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Using `make splc` to create `splc` in `./bin` root for spl codes' parsing.