

Graph Theory and Representation

Some Applications of Graph Theory

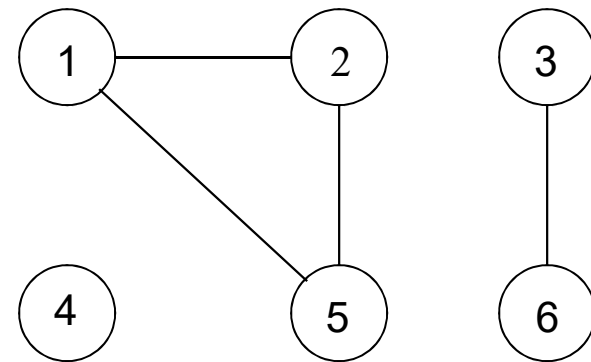
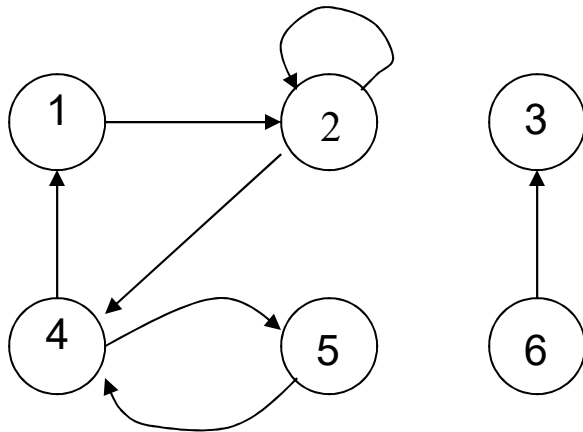
- Some high level problems in networks, eg
- 1 Topology planning
- 2 Dimensioning
- 3 Routing
- 4 Traffic engineering

Graphs ↔ Networks

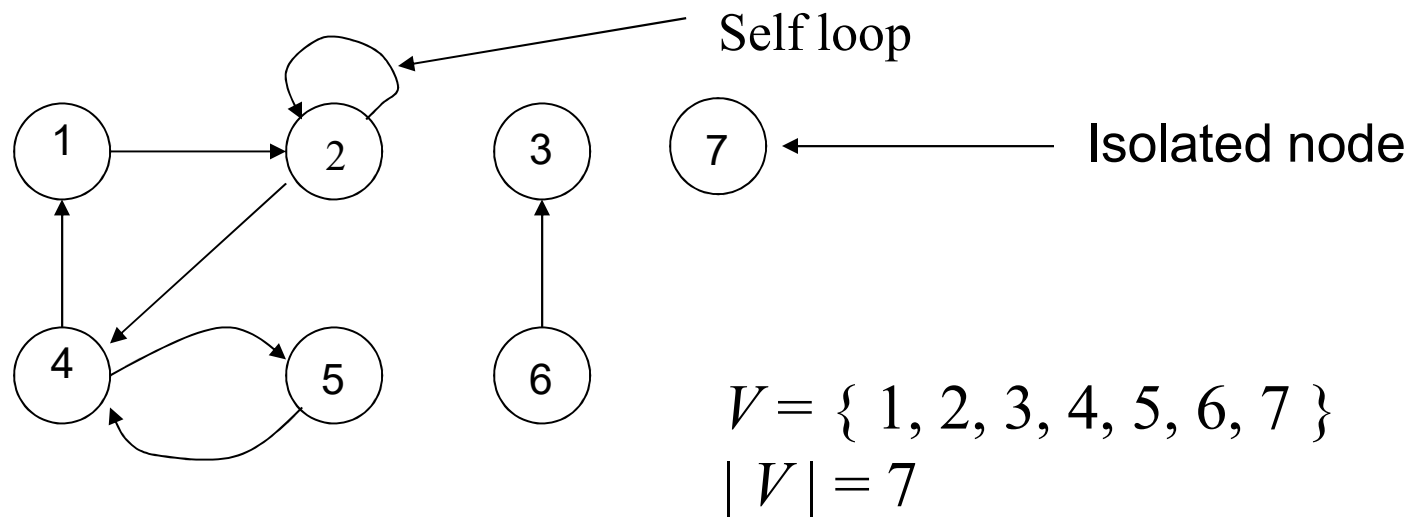
| Graph (Network) | Vertexes (Nodes) | Edges (Arcs) | Flow |
|--------------------|------------------------------------------------|-------------------------------------------|-------------------------------------|
| Communications | Telephones exchanges, computers, satellites | Cables, fiber optics, microwave relays | Voice, video, packets |
| Circuits | Gates, registers, processors | Wires | Current |
| Mechanical | Joints | Rods, beams, springs | Heat, energy |
| Hydraulic | Reservoirs, pumping stations, lakes | Pipelines | Fluid, oil |
| Financial | Stocks, currency | Transactions | Money |
| Transportation | Airports, rail yards, street intersections | Highways, railbeds, airway routes | Freight, vehicles, passengers |

What is a Graph?

- Informally a *graph* is a set of nodes joined by a set of lines or arrows.



A **directed graph**, also called a **digraph** G is a pair (V, E) , where the set V is a finite set and E is a binary relation on V . The set V is called the **vertex set** of G and the elements are called vertices. The set E is called the **edge set** of G and the elements are *edges* (also called *arcs*). A edge from node a to node b is denoted by the ordered pair (a, b) .

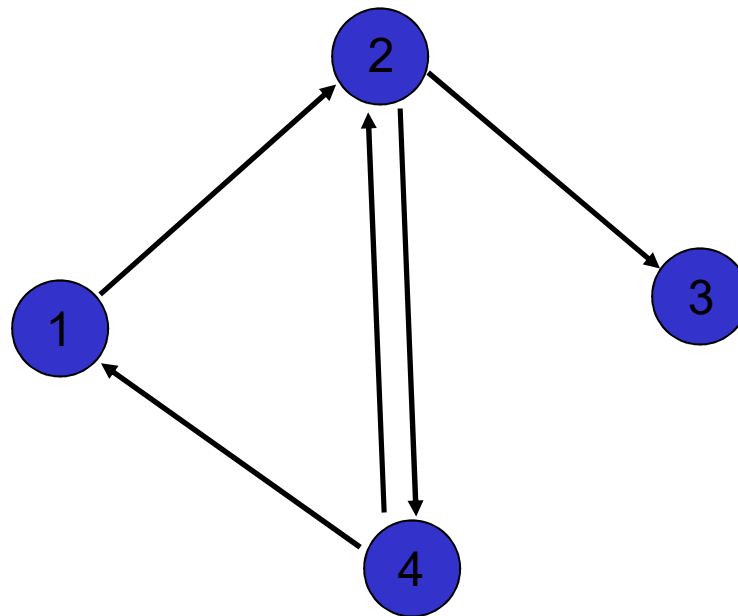


$$E = \{ (1,2), (2,2), (2,4), (4,5), (4,1), (5,4), (6,3) \}$$

$$|E| = 7$$

Directed Graph

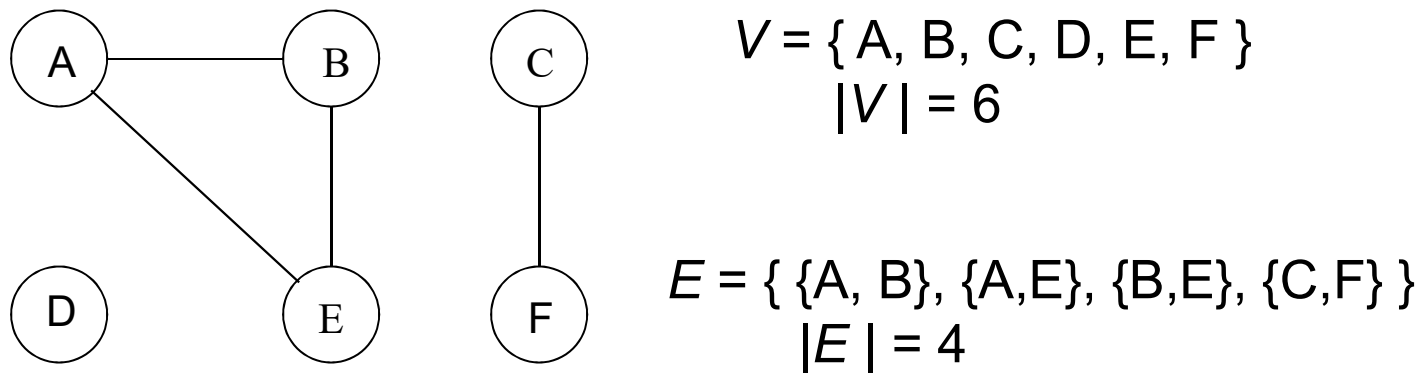
An edge $e \in E$ of a directed graph is represented as an ordered pair (u,v) , where $u, v \in V$. Here u is the initial vertex and v is the terminal vertex. Also assume here that $u \neq v$



$$V = \{ 1, 2, 3, 4 \}, |V| = 4$$

$$E = \{(1,2), (2,3), (2,4), (4,1), (4,2)\}, |E| = 5$$

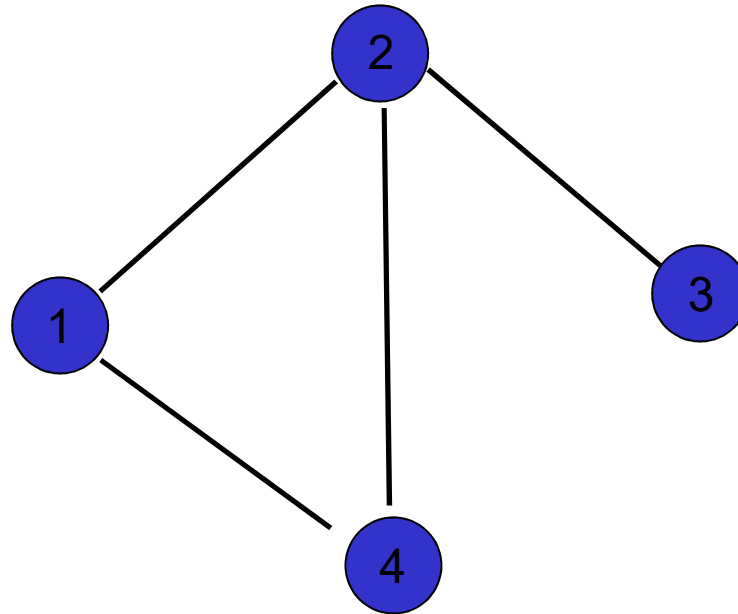
An **undirected graph** $G = (V, E)$, but unlike a digraph the edge set E consist of unordered pairs. We use the notation (a, b) to refer to a directed edge, and $\{a, b\}$ for an undirected edge.



Some texts use (a, b) also for undirected edges.
So (a, b) and (b, a) refers to the same edge.

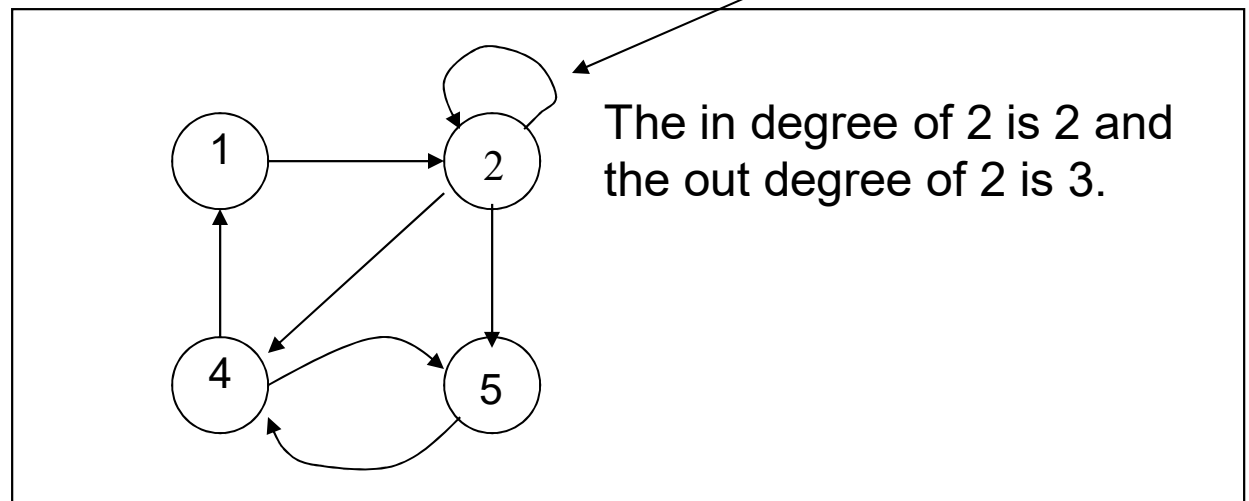
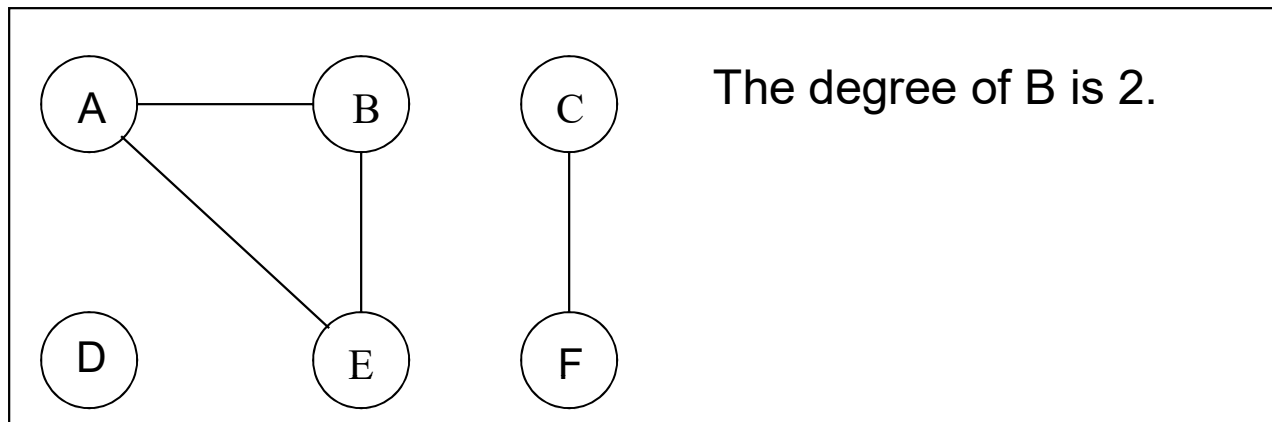
Undirected Graph

An edge $e \in E$ of an undirected graph is represented as an unordered pair $(u,v)=(v,u)$, where $u, v \in V$. Also assume that $u \neq v$

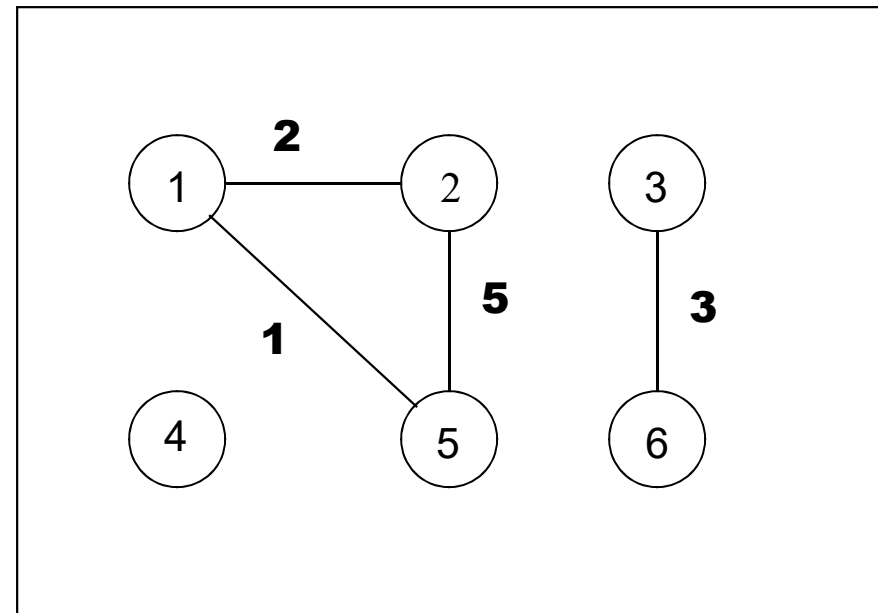
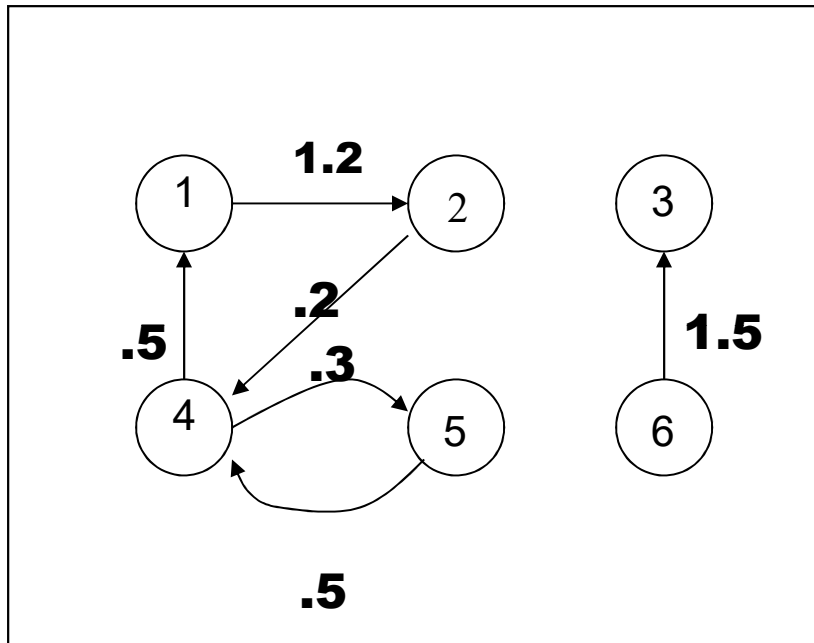


$$V = \{ 1, 2, 3, 4 \}, |V| = 4$$
$$E = \{(1,2), (2,3), (2,4), (4,1)\}, |E| = 4$$

Degree of a Vertex in an undirected graph is the number of edges incident on it. In a directed graph, the **out degree** of a vertex is the number of edges leaving it and the **in degree** is the number of edges entering it.

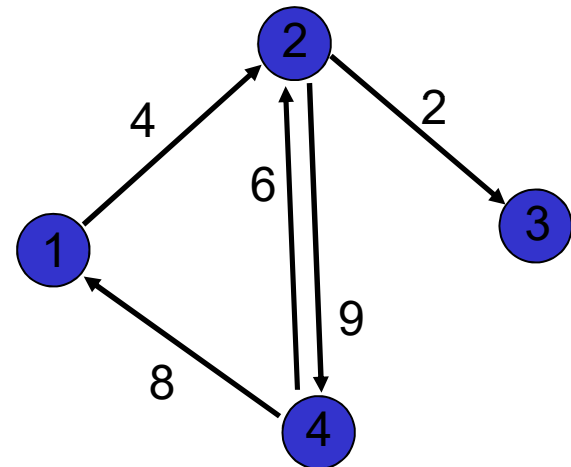
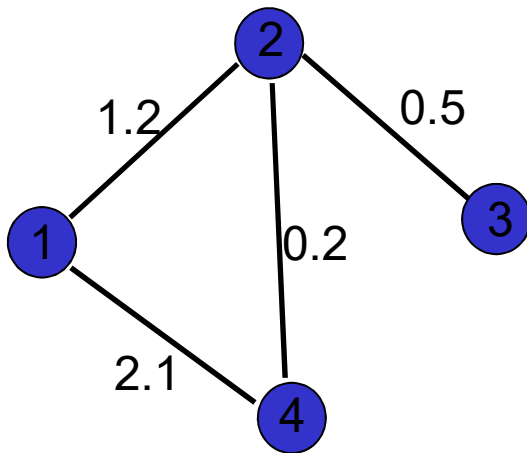


A **weighted graph** is a graph for which each edge has an associated **weight**, usually given by a **weight function** $w: E \rightarrow \mathbf{R}$.

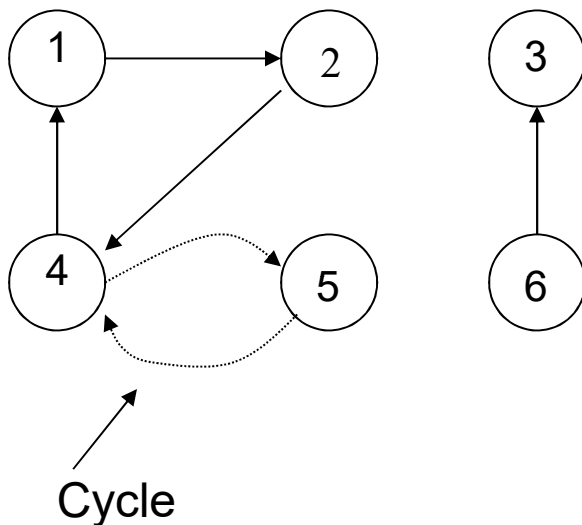


Weighted Graph

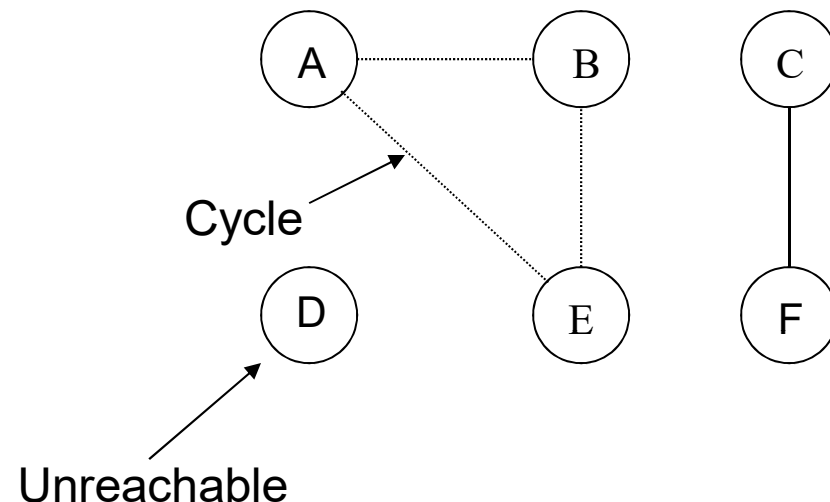
A *weighted graph* is a graph for which each edge has an associated *weight*, usually given by a *weight function* $w: E \rightarrow \mathbb{R}$



A **path** is a sequence of vertices such that there is an edge from each vertex to its successor. A path from a vertex to itself is called a **cycle**. A graph is called **cyclic** if it contains a cycle; otherwise it is called **acyclic**. A path is **simple** if each vertex is distinct.

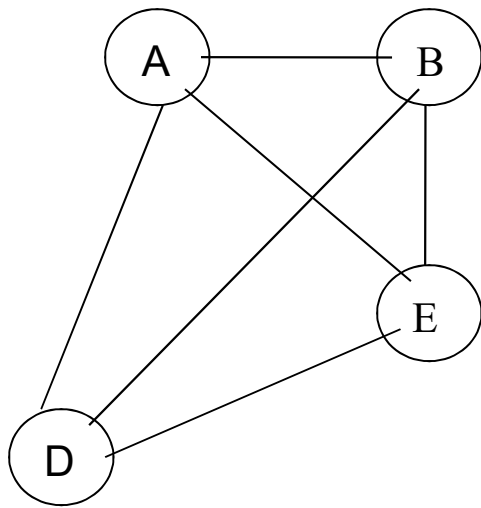


Simple path from 1 to 5
= (1, 2, 4, 5)
 or as in our text
 ((1, 2), (2, 4), (4,5))



If there is path p from u to v then we say v is **reachable** from u via p .

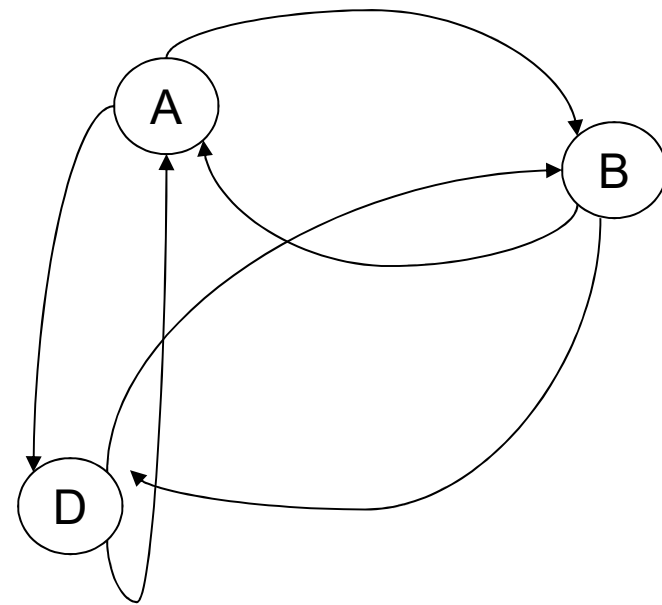
A **Complete graph** is an undirected/directed graph in which every pair of vertices is adjacent. If (u, v) is an edge in a graph G , we say that vertex v is **adjacent** to vertex u .



4 nodes and $(4*3)/2$ edges

V nodes and $V*(V-1)/2$ edges

Note: if self loops are allowed $V(V-1)/2 + V$ edges



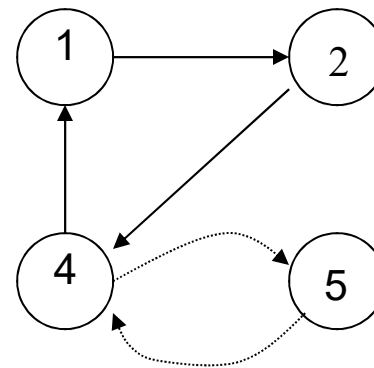
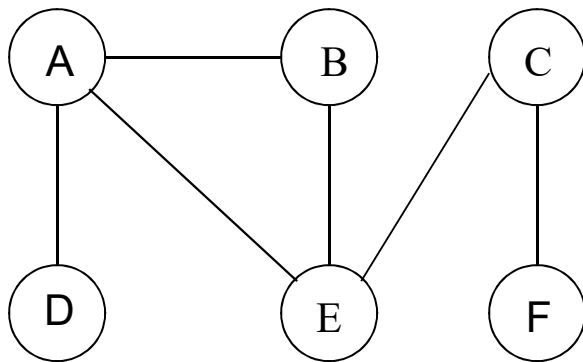
3 nodes and $3*2$ edges

V nodes and $V*(V-1)$ edges

Note: if self loops are allowed V^2 edges

An undirected graph is **connected** if you can get from any node to any other by following a sequence of edges OR any two nodes are connected by a path.

A directed graph is **strongly connected** if there is a directed path from any node to any other node.



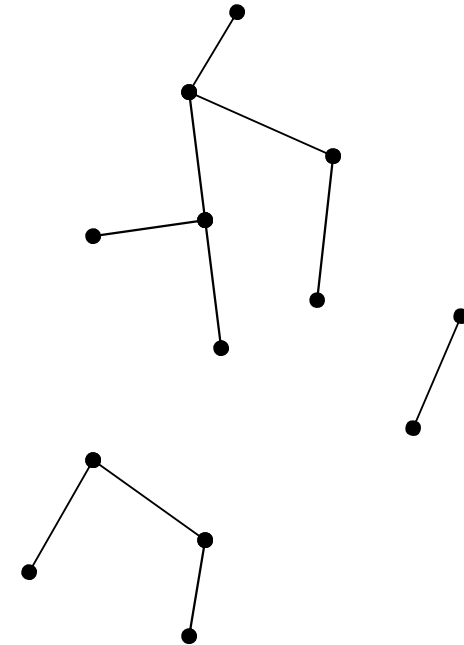
- A graph is **sparse** if $|E| \approx |V|$
- A graph is **dense** if $|E| \approx |V|^2$.

A **free tree** is an acyclic, connected, undirected graph. A **forest** is an acyclic undirected graph. A **rooted tree** is a tree with one distinguished node, **root**.

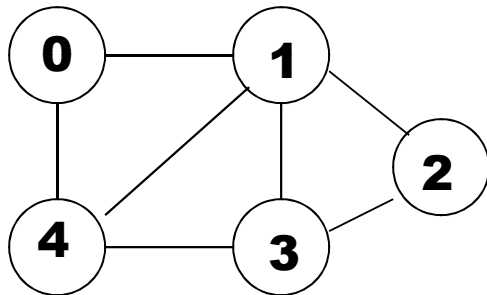
Let $G = (V, E)$ be an undirected graph.

The following statements are equivalent.

1. G is a tree
2. Any two vertices in G are connected by unique simple path.
3. G is connected, but if any edge is removed from E , the resulting graph is disconnected.
4. G is connected, and $|E| = |V| - 1$
5. G is acyclic, and $|E| = |V| - 1$
6. G is acyclic, but if any edge is added to E , the resulting graph contains a cycle.

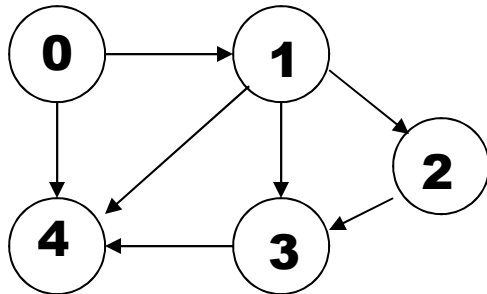


Adjacency-matrix-representation of a graph $G = (V, E)$ is a $|V| \times |V|$ matrix $A = (a_{ij})$ such that
 $a_{ij} = 1$ (or some Object) if $(i, j) \in E$ and
 0 (or null) otherwise.



| | 0 | 1 | 2 | 3 | 4 |
|---|---|---|---|---|---|
| 0 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 | 1 |
| 2 | 0 | 1 | 0 | 1 | 0 |
| 3 | 0 | 1 | 1 | 0 | 1 |
| 4 | 1 | 1 | 0 | 1 | 0 |

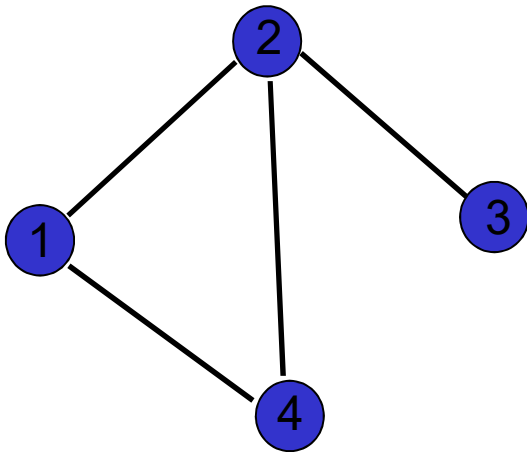
Adjacency Matrix Representation for a Directed Graph



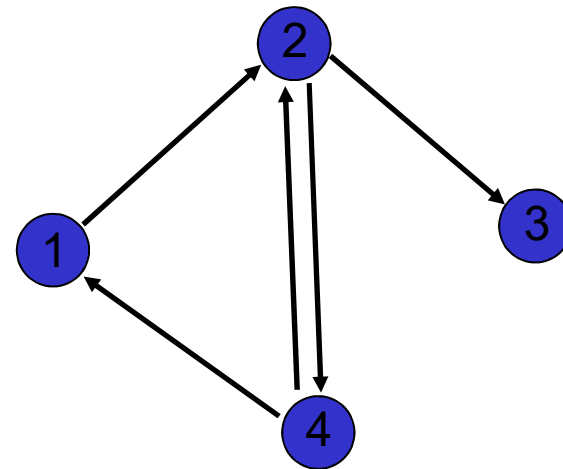
| | 0 | 1 | 2 | 3 | 4 |
|---|---|---|---|---|---|
| 0 | 0 | 1 | 0 | 0 | 1 |
| 1 | 0 | 0 | 1 | 1 | 1 |
| 2 | 0 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 | 1 |
| 4 | 0 | 0 | 0 | 0 | 0 |

Degree of a Vertex

Degree of a vertex in an undirected graph is the number of edges incident on it. In a directed graph, the *out degree* of a vertex is the number of edges leaving it and the *in degree* is the number of edges entering it



The *degree* of vertex 2 is 3



The *in degree* of vertex 2 is 2 and the *in degree* of vertex 4 is 1