

Today!

Choose one:

1) Snakes + Mayer-Vietoris

2) Persistence

Last time:

A chain complex (of vector spaces) is a sequence of vec. spaces + linear maps

$$\dots \leftarrow C_{n-1} \xrightarrow{d_n} C_n \xrightarrow{d_{n+1}} C_{n+1} \leftarrow \dots$$

such that (for all $n \dots$)

$$d_n^2 = 0 \quad (d_n \circ d_{n+1} = 0)$$

equivalently,

$$\text{im } d_{n+1} \subseteq \ker d_n$$

[for now, finite dimensional vector spaces]

A chain complex is called exact if $\text{im } d_n = \ker d_n$ for all n .

The homology of a chain complex is

$$H_n = \frac{\ker d_n}{\text{im } d_{n+1}}$$

for $n = 0, 1, \dots$

A short exact sequence is an exact chain complex w/ 5 elements, the first and last being 0.

Observation: if C_\bullet is exact, then $H_\bullet = 0$

IF $d_n = 0$ for all n , $H_\bullet = C_\bullet$.

$$0 \xrightarrow{s} A \xrightarrow{f} B \xrightarrow{g} C \xrightarrow{z} 0$$

1) exactness means that $0 = \text{im } s = \ker f$
so f is injective

2) But, $\text{im } (g) = \ker z = C$
so g is surjective

3) Because f is injective,
 $\ker g = \text{im } f \cong A$

By 1st Isomorphism Thm,
 $\text{im } g \cong B / \ker g$

$$C \cong B / A$$

For vector spaces, this implies

$$\dim(A) + \dim(C) = \dim(B)$$

$$\text{i.e. } A \oplus C \cong B$$

In groups, it only tells you $B \cong A \times C$

[now, infinite is fair game!]

Reminder: A singular n -simplex in X is a map $\Delta^n \rightarrow X$

$$C_n^{\text{sing}}(X) = \text{vector space whose basis is singular } n\text{-simp. at } X$$

This is ghastly. Typically, X is uncountable, so $\{\Delta^n \rightarrow X\}$ is uncountable, so $C_n^{\text{sing}}(X)$ is uncountably ∞ -dimensional.

The boundary faces of a simplex



$$\{(t_0, t_1, t_2) \mid \begin{matrix} t_i \geq 0 \\ \sum t_i = 1 \end{matrix}\}$$

0th boundary face of Δ^2 is

$$\{(0, t_1, t_2) \mid \begin{matrix} t_1 + t_2 = 1 \\ 0 \leq t_i \end{matrix}\}$$

1st boundary face:

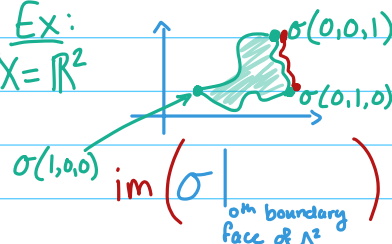
$$(t_0, 0, t_2)$$

$$2\text{nd: } (t_0, t_1, 0)$$

Note: each boundary face "is" Δ^1

So, let $\sigma: \Delta^2 \rightarrow X$

Ex:
 $X = \mathbb{R}^2$



$$\sigma: \Delta^2 \rightarrow X$$

$$\sigma|_A: A \rightarrow X$$

$$\text{notice: } \Delta^1 = \{(s_0, s_1) \mid \begin{matrix} s_i \geq 0 \\ s_0 + s_1 = 1 \end{matrix}\}$$

$$\text{define } f_0(s_0, s_1) = (0, s_0, s_1)$$

$$f_0: \Delta^1 \rightarrow \Delta^2 \quad \text{im}(f_0) = 0^{\text{th}} \text{ boundary face}$$

So, we can really say:

$$\sigma|_{0^{\text{th}} \text{ boundary face}} \circ f_0: \Delta^1 \rightarrow X$$

In other words, $\sigma|_{0^{\text{th}} \text{ boundary face}}$ can be thought of as a singular 1-simplex.
(explicitly, $\sigma \circ f_0$ is a singular 1-simplex)

Punchline:

$$\text{define: } f_i: \Delta^{n-1} \rightarrow \Delta^n$$

$(s_0, \dots, s_{n-1}) \rightarrow (s_0, \dots, \underset{\substack{\downarrow \text{ } i^{\text{th}} \text{ coordinate}}}{0}, \dots, s_n)$

This is the i^{th} boundary inclusion of Δ^n .

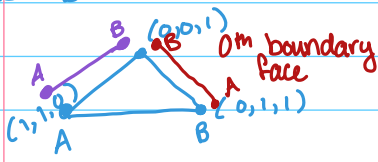
If $\sigma: \Delta^n \rightarrow X$ is a singular n -simplex, then

$\sigma \circ f_i: \Delta^{n-1} \rightarrow X$ is a singular $(n-1)$ -simplex.

$$d[\sigma] = \sum_{i=0}^n (-1)^i [\sigma \circ f_i] \quad \text{notation: } [\sigma] \text{ is the vector in } C_n \text{ corresponding to the function } \sigma.$$

$B = 0^{\text{th}} \text{ boundary face of } \Delta^2$

$$d(\Delta^1) = B - A$$



$$d([\Delta^2]) = \sum_{i=0}^2 (-1)^i \sigma \circ f_i$$

$$= \sum_{i=0}^2 (-1)^i \begin{matrix} \text{red line } A \rightarrow B \\ \text{purple line } A \rightarrow C \\ \text{blue line } B \rightarrow C \end{matrix}$$

$$= \text{red line} - \text{purple line} + \text{blue line}$$

$$d(\text{red line}) = \Delta - \Delta$$

$$d(\text{purple line}) = \Delta - \Delta$$

$$d(\text{blue line}) = \Delta - \Delta$$

$$d(d(\Delta^2)) = d(\text{red line} - \text{purple line} + \text{blue line})$$

$$= \Delta - \Delta - (\Delta - \Delta) + \Delta - \Delta = 0$$

Lemma:

d from topology really is a chain map,
meaning $\text{im } d \subseteq \ker d$

PF: $\ddot{!}$ exercise.

Let's prove some Theorems!

$$\begin{array}{ccc} \text{Top} & X & \rightarrow Y \\ & \Downarrow & \end{array}$$

$$\text{Vec} \quad H_k^{\text{sing}}(X) \rightarrow H_k^{\text{sing}}(Y) \quad \forall k$$

Question: $X \xrightarrow{f} Y$

idea: $\Delta^n \rightarrow X \rightarrow Y$ gives us $\Delta^n \rightarrow Y$
so $C_n(X) \rightarrow C_n(Y)$

$$\begin{array}{ccc} \dots \rightarrow C_n(X) & \rightarrow & C_{n-1}(X) \rightarrow \dots \\ \downarrow f^\# & & \downarrow f^\# \\ \dots \rightarrow C_n(Y) & \xrightarrow{d_n} & C_{n-1}(Y) \rightarrow \dots \end{array}$$

Does this commute?

meaning: $f^\# \circ d_n \stackrel{?}{=} d_n \circ f^\#$

Answer: Yes.

What we really need to verify, though:

I claim that this gets us $f^\# : H_n(X) \rightarrow H_n(Y)$

so we should check...

$$\frac{\ker d_n^{(X)}}{\text{im } d_{n+1}^{(X)}} \stackrel{?}{\rightarrow} \frac{\ker d_n^{(Y)}}{\text{im } d_{n+1}^{(Y)}}$$

Answer: Yes

PF: Exercise