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# A Streaming Algorithm for Graph Clustering

## Supplementary Material

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### Abstract

We here provide the proofs of the main results of the paper.

## 1 Proof of Lemma 1

This lemma gives an expression of  $Q_{t+1}$  in function of  $Q_t$  when a new edge  $(i, j)$  arrives.

**Lemma 1.** *If  $e_{t+1} = (i, j)$  and if  $P_{t+1} = P_t$ ,  $Q_{t+1}$  can be expressed in function of  $Q_t$  as follows*

$$Q_{t+1} = Q_t + 2 \left[ \delta(i, j) - \frac{\text{Vol}_t(C(i)) + \text{Vol}_t(C(j)) + 1 + \delta(i, j)}{w} \right]$$

where  $C(v)$  denotes the community of  $v$  in  $P_t$ , and  $\delta(i, j) = 1$  if  $i$  and  $j$  belongs to the same community of  $P_t$  and 0 otherwise.

*Proof.* Given a new edge  $e_{t+1} = (i, j)$ , we have the following relation between quantities  $\text{Int}(C)$  and  $\text{Vol}(C)$  at times  $t$  and  $t + 1$ .

$$\text{Int}_{t+1}(C) = \text{Int}_t(C) + 1_{i \in C} 1_{j \in C}$$

and

$$\text{Vol}_{t+1}(C) = \text{Vol}_t(C) + 1_{i \in C} + 1_{j \in C}.$$

This gives us the following equation for  $(\text{Vol}_{t+1}(C))^2$

$$(\text{Vol}_{t+1}(C))^2 = (\text{Vol}_t(C))^2 + \begin{cases} 0 & \text{if } C \neq C(i) \text{ and } C \neq C(j) \\ 2\text{Vol}_t(C(i)) + 1 & \text{if } C = C(i) \\ 2\text{Vol}_t(C(j)) + 1 & \text{if } C = C(j) \end{cases}$$

in the case  $C(i) \neq C(j)$ , and

$$(\text{Vol}_{t+1}(C))^2 = (\text{Vol}_t(C))^2 + \begin{cases} 0 & \text{if } C \neq C(i) \text{ and } C \neq C(j) \\ 4\text{Vol}_t(C(i)) + 4 & \text{if } C = C(i) = C(j) \end{cases}$$

in the case  $C(i) = C(j)$ .

Finally, the definition of  $Q_{t+1}$

$$Q_{t+1} = \sum_{C \in P_{t+1}} \left[ 2\text{Int}_{t+1}(C) - \frac{(\text{Vol}_{t+1}(C))^2}{w} \right]$$

gives us the wanted result. □

## 2 Proof of Lemma 2

Lemma 2 gives an expression for the variation of  $Q_t$  when node  $i$  joins community  $C(j)$ .

**Lemma 2.**

$$\Delta Q_t = Q_t^{(a)} - Q_t^{(c)} = 2 \left[ L_t(i, C(j)) - L_t(i, C(i)) - \frac{(w_t(i))^2}{w} \right]$$

where

$$L_t(i, C) = \sum_{(i', j') \in S_t} \left[ 1_{i' \in C} \left( 1_{j'=i} - \frac{w_t(i)}{w} \right) + 1_{j' \in C} \left( 1_{i'=i} - \frac{w_t(i)}{w} \right) \right].$$

*Proof.*  $Q_t$  is defined as a sum over all communities of partition  $P_t$ . Only terms depending on  $C(i)$  and  $C(j)$  are modified by action (a). Thus, we have:

$$\begin{aligned} \Delta Q_t &= 2 [\text{Int}_t(C(i) \setminus \{i\}) + \text{Int}_t(C(j) \cup \{i\}) - \text{Int}_t(C(i)) - \text{Int}_t(C(j))] \\ &\quad - \frac{(\text{Vol}_t(C(i)) - w_t(i))^2 + (\text{Vol}_t(C(j)) + w_t(i))^2 - (\text{Vol}_t(C(i)))^2 - (\text{Vol}_t(C(j)))^2}{w}. \end{aligned}$$

This leads to:

$$\begin{aligned} \Delta Q_t &= 2 \sum_{(i', j') \in S_t} [1_{j'=i}(1_{i' \in C(j)} - 1_{i' \in C(i)}) + 1_{i'=i}(1_{j' \in C(j)} - 1_{j' \in C(i)})] \\ &\quad - 2 \frac{w_t(i) \text{Vol}_t(C(j)) - w_t(i) \text{Vol}_t(C(i)) + (w_t(i))^2}{w}. \end{aligned}$$

Using the definition of  $\text{Vol}_t$ , we obtain the wanted expression for  $\Delta Q_t$ . □

## 3 Proof of Theorem 1

Theorem 1 gives a sufficient condition in order to have a positive variation  $\Delta Q_{t+1}$  of the modularity when  $i$  joins  $C(j)$ .

**Theorem 1.** If  $\text{Vol}_t(C(i)) \leq \text{Vol}_t(C(j))$ , then:

$$\text{Vol}_t(C(j)) \leq v_t(i, j) \implies \Delta Q_{t+1} \geq 0$$

where

$$v_t(i, j) = \frac{1 - (w_t(i) + 1)^2/w}{l_t(i, C(i)) - l_t(i, C(j))}.$$

*Proof.* From Lemma 1, we obtain

$$\begin{aligned} \Delta Q_{t+1} &= Q_t^{(a)} + 2 \left[ 1 - \frac{(\text{Vol}_t(C(j)) + w_i(t)) + (\text{Vol}_t(C(j)) + w_i(t)) + 2}{w} \right] \\ &\quad - Q_t^{(b)} - 2 \left[ 0 - \frac{(\text{Vol}_t(C(i)) - w_i(t)) + (\text{Vol}_t(C(j)) + w_i(t)) + 1}{w} \right], \end{aligned}$$

which gives us:

$$\Delta Q_{t+1} = \Delta Q_t + 2 \left[ 1 - \frac{\text{Vol}_t(C(j)) - \text{Vol}_t(C(i)) + 2w_i(t) + 1}{w} \right] \quad (1)$$

Then, Equation (1) and Lemma 2 gives us the following expression for  $\Delta Q_{t+1}$

$$\begin{aligned} \Delta Q_{t+1} &= 2 \left[ 1 + \left( l_t(i, C(j)) - \frac{1}{w} \right) \text{Vol}_t(C(j)) - \left( l_t(i, C(i)) - \frac{1}{w} \right) \text{Vol}_t(C(i)) \right. \\ &\quad \left. - \frac{(w_t(i) + 1)^2}{w} \right]. \end{aligned}$$

Thus,  $\Delta Q_{t+1} \geq 0$  is equivalent to

$$\left(l_t(i, C(i)) - \frac{1}{w}\right) Vol_t(C(i)) - \left(l_t(i, C(j)) - \frac{1}{w}\right) Vol_t(C(j)) \leq 1 - \frac{(w_t(i) + 1)^2}{w}. \quad (2)$$

We use  $u_t(i, j)$  to denote the left-hand side of this inequality. If  $Vol_t(C(i)) \leq Vol_t(C(j))$ , then we have

$$u_t(i, j) \leq [l_t(i, C(i)) - l_t(i, C(j))] Vol_t(C(j))$$

Thus, the following inequality

$$[l_t(i, C(i)) - l_t(i, C(j))] Vol_t(C(j)) \leq 1 - \frac{(w_t(i) + 1)^2}{w}$$

implies inequality (2), which proves the theorem.  $\square$