



OPTIMIZING CLIENT ASSOCIATION IN 60 GHz WIRELESS ACCESS NETWORKS

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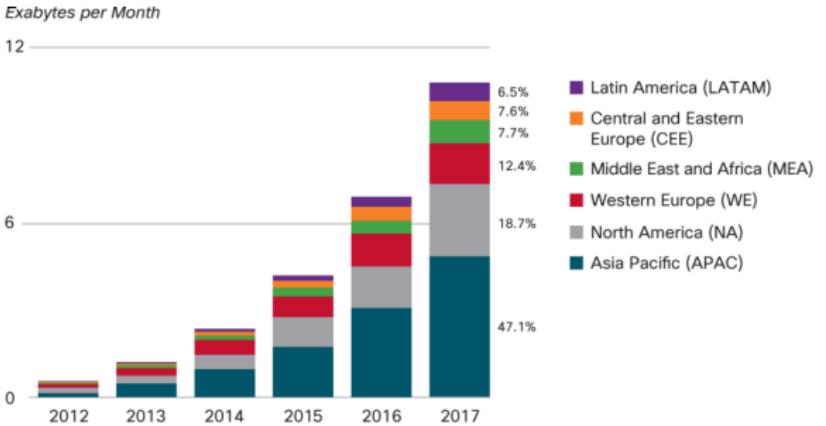
Outline

- Past, Present and Future in wireless communications
- 60 GHz millimeterWave wireless technology
- Optimizing resource allocation
- Distributed client association (DAA)
- Numerical analysis of DAA
- Conclusions and open research topics

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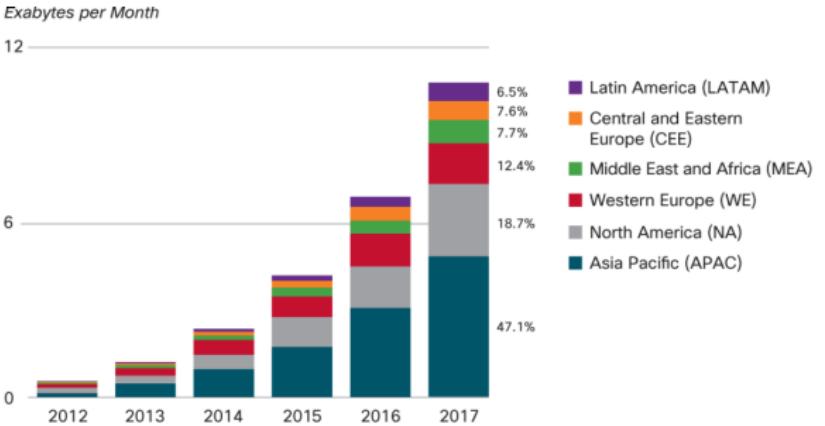
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Growth of Mobile Traffic



Source: Cisco VNI Mobile Forecast, 2013

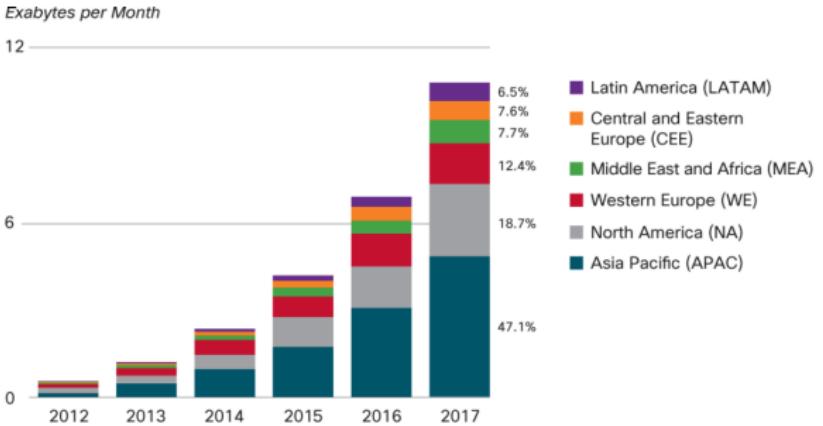
Growth of Mobile Traffic



Source: Cisco VNI Mobile Forecast, 2013

2012: 70% mobile traffic growth compared to 2011

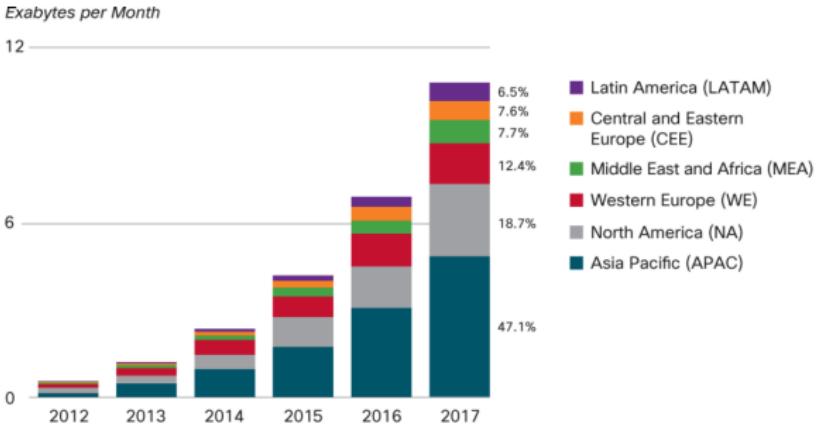
Growth of Mobile Traffic



Source: Cisco VNI Mobile Forecast, 2013

2012: 12 times the entire global internet traffic in 2000

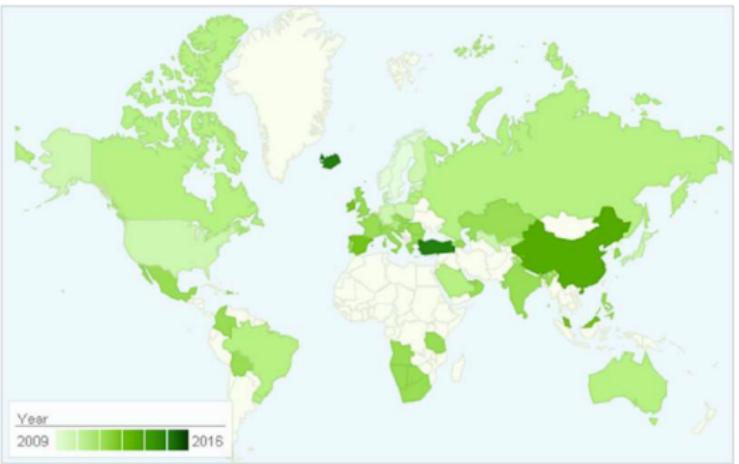
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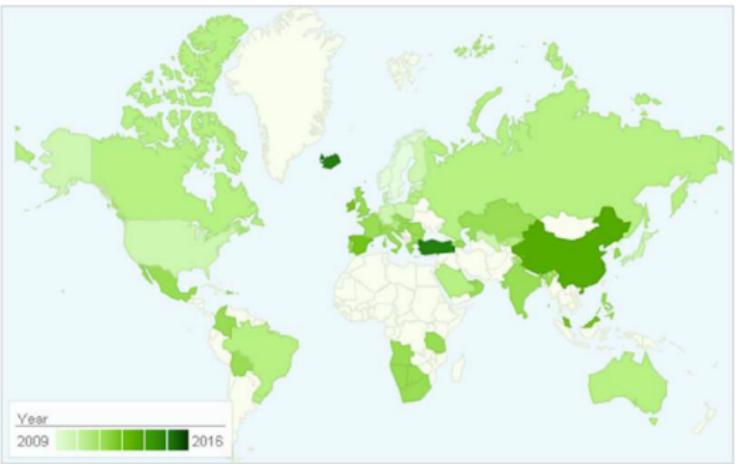
2017: 13 times higher compared to 2012

LTE Deployment



Source: Cisco, 2013

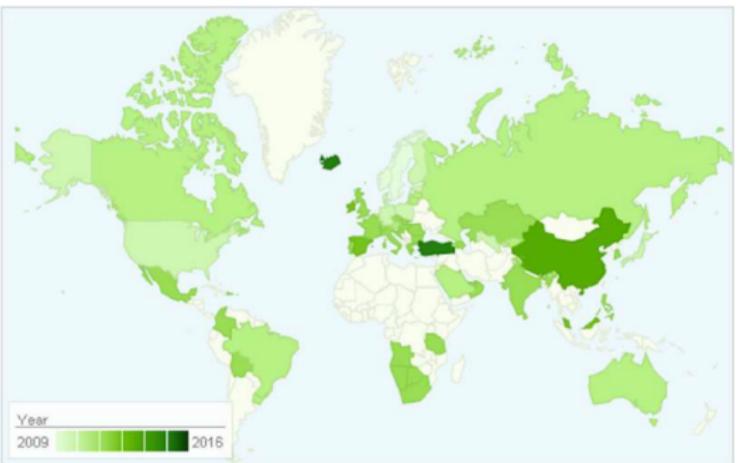
LTE Deployment



Source: Cisco, 2013

End of 2013: # of mobile connected devices > earth population

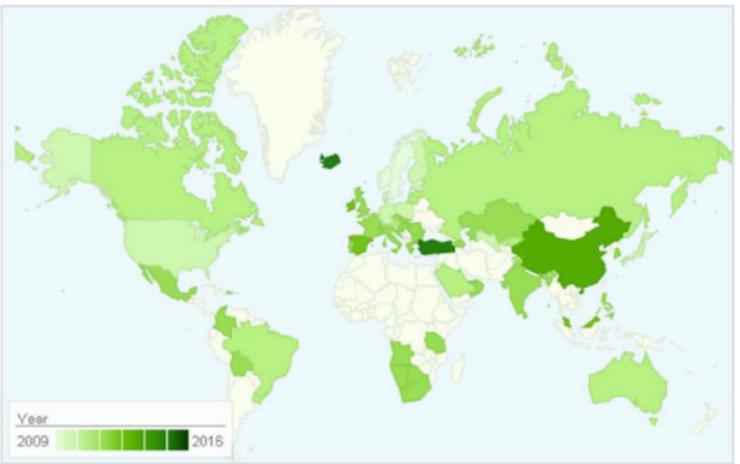
LTE Deployment



Source: Cisco, 2013

2017: 4G will represent 45% of mobile traffic

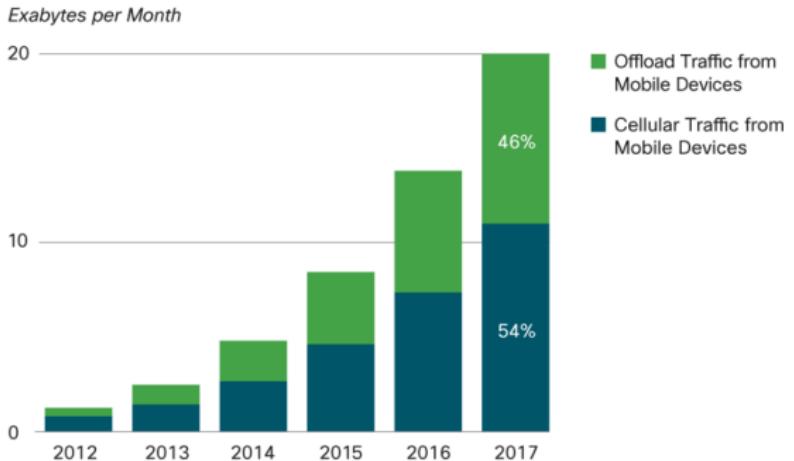
LTE Deployment



Source: Cisco, 2013

2017: 2.7 GB mobile traffic per smartphone per month

Mobile Traffic Offloading



Source: Cisco VNI Mobile Forecast, 2013

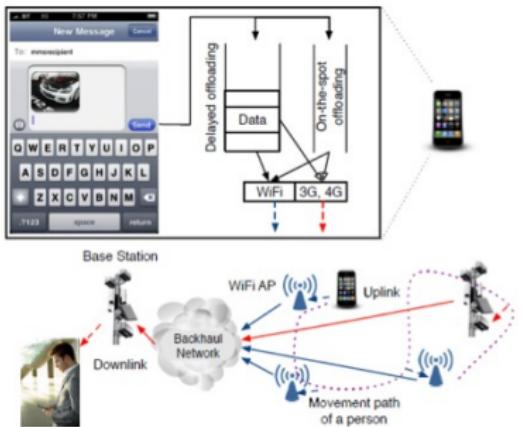
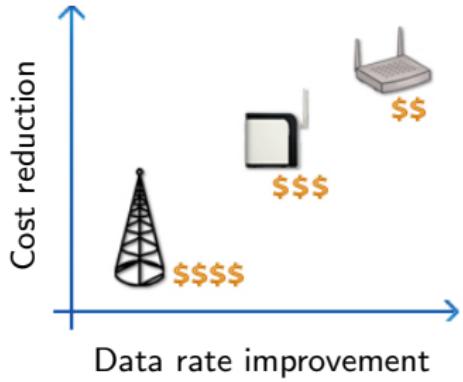
The amount of traffic offloaded from smartphones will be 46%, and the amount of traffic offloaded from tablets will be 71% in 2017 (Cisco, 2013)

90% of all cellular base stations will be small cells by 2016 (IEEE Spectrum Magazine, Informa Telecoms & Media, 2013)

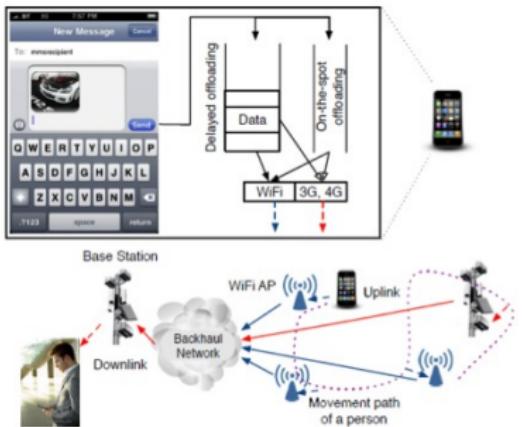
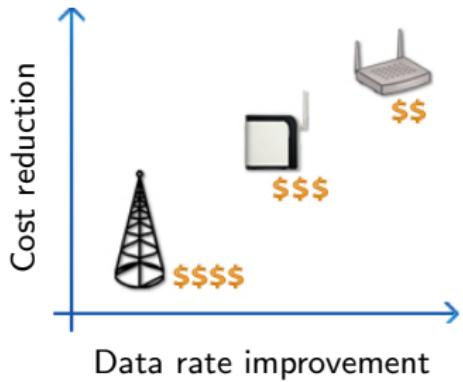
Priorities

- High bandwidth
- High coverage
- Green
- Cheap

Mobile Data Offloading



Mobile Data Offloading



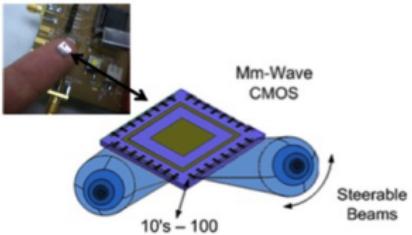
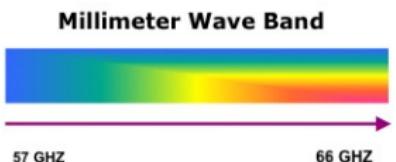
Are there better solutions for mobile data offloading?

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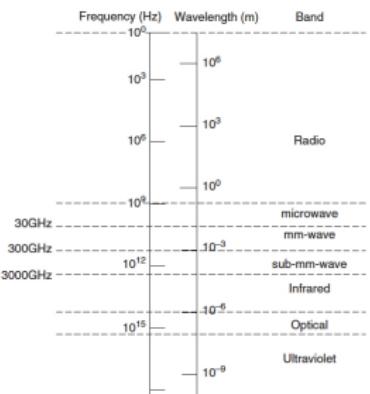
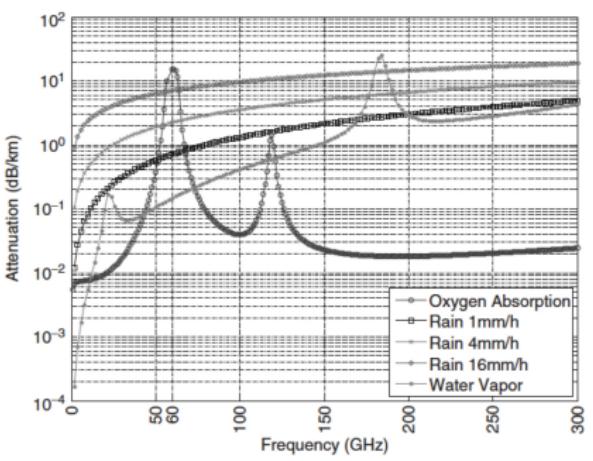
60 GHz Wireless Access Networks

- Unlicensed short range transmissions in the 60 GHz millimeter wave (mmW) band
- Achieve Gbps communication
- Reduced interference
- Low-cost mmW transceivers



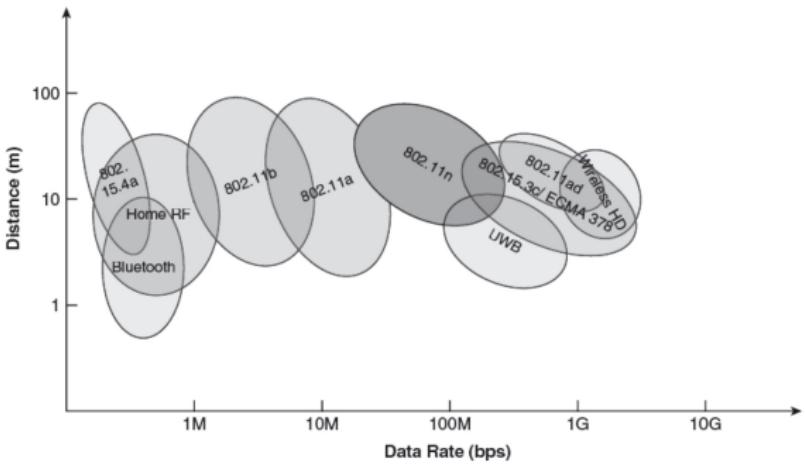
MillimeterWave Band

- History (J.C. Bose, 1897)
- High path loss
- High oxygen absorption

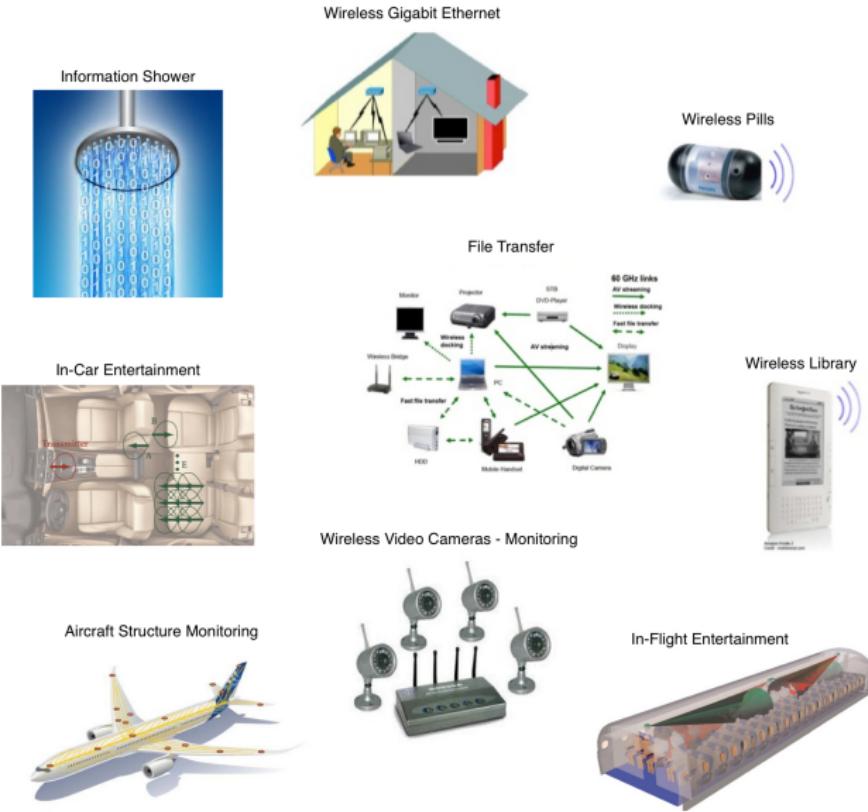


60 GHz Wireless Standards

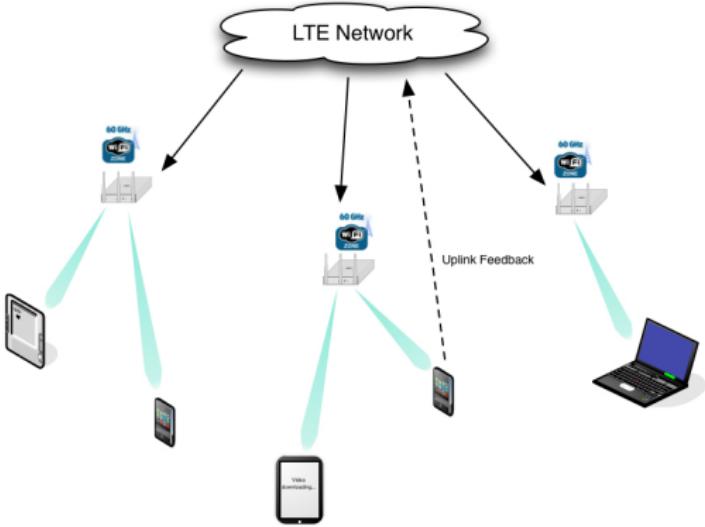
- IEEE 802.11ad
- WiGig
- IEEE 802.15.3c
- WirelessHD
- ECMA-387



Applications



60 GHz to the Mobile

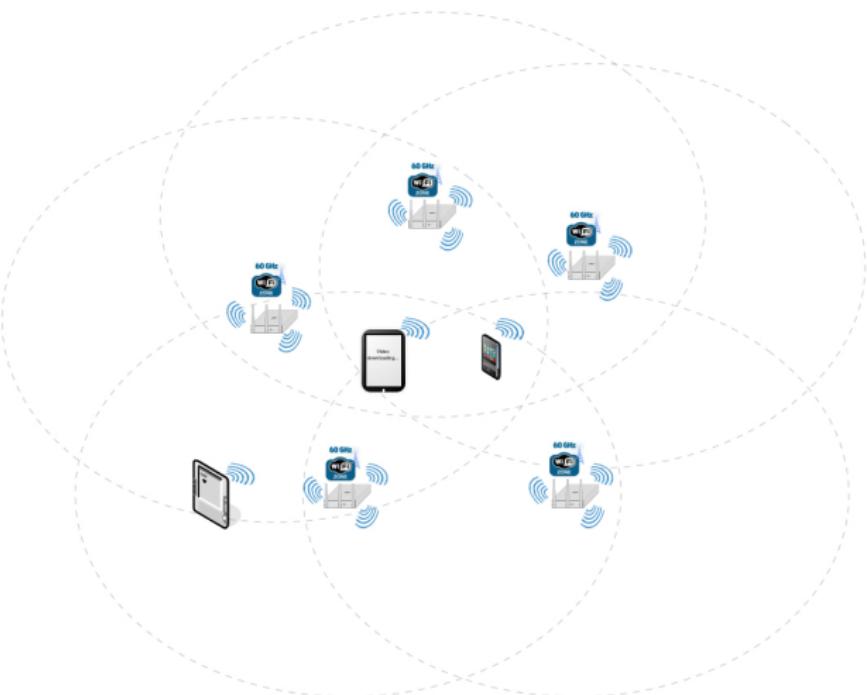


- 60 GHz small base stations (eg. on lamp posts)
- Downlink offload traffic in 60 GHz band with uplink LTE feedback
- 60 GHz radio on mobile device in receive-only mode

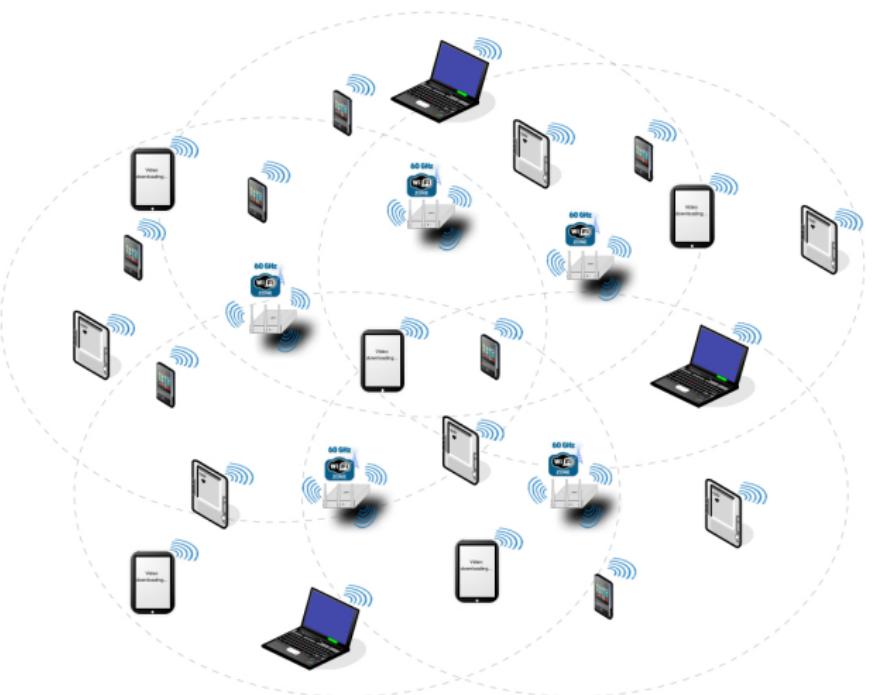
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Optimizing Client Association in 60 GHz Wireless Access Networks



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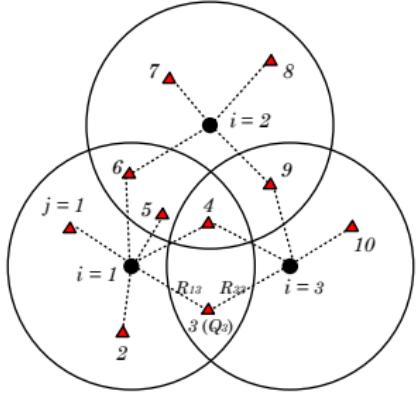


Optimizing Client Association in 60 GHz Wireless Access Networks

- **Goal:**
 - Distributed client association and fair load distribution
 - find the client association that minimizes the maximum AP utilization
- **Solution method:** Based on Lagrangian duality theory
- **Results:** Theoretical and numerical analysis

G. Athanasiou, P. C. Weeraddana, C. Fischione and L. Tassiulas, "Optimizing Client Association in 60 GHz Wireless Access Networks", *arXiv:1301.2723, Cornell University Library*, 2013 [Online]. Available: <http://arxiv.org/abs/1301.2723>

System Model



W	System bandwidth
P_{ij}	Transmission power of AP i to client j
G_{ij}	Power gain from AP i to client j
N_0	Power spectral density of the noise
I_j	Interference spectral density at client j
\mathcal{M}_i	Set of clients that can be associated to AP i
\mathcal{N}_j	Set of APs that client j could be associated with

- $\mathcal{N} = \{1, \dots, N\}$ APs and $\mathcal{M} = \{1, \dots, M\}$ clients
- Achievable rate from AP i to client $j \in \mathcal{M}_i$ is

$$R_{ij} = W \log_2 \left(1 + \frac{P_{ij} G_{ij}}{(N_0 + I_j) W} \right)$$

- Q_j is the demanded data rate of client j

System Model

- Channel utilization between AP i and client j is

$$\beta_{ij} = \frac{Q_j}{R_{ij}}$$

- Utilization of AP i is

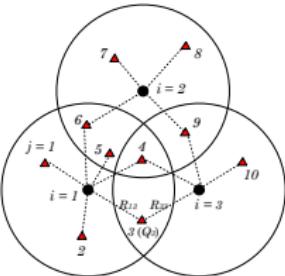
$$\sum_{j \in \mathcal{M}_i} \beta_{ij} x_{ij}$$

- $(x_{ij})_{j \in \mathcal{M}_i}$ are binary decision variables, which indicate the client association

$$x_{ij} = \begin{cases} 1 & \text{if client } j \text{ is associated to AP } i \\ 0 & \text{otherwise} \end{cases}$$

Client Association Problem Formulation

$$\begin{aligned} & \text{minimize} && \max_{i \in \mathcal{N}} \sum_{j \in \mathcal{M}_i} \beta_{ij} x_{ij} \\ & \text{subject to} && \sum_{i \in \mathcal{N}_j} x_{ij} = 1, \quad j \in \mathcal{M} \\ & && x_{ij} \in \{0, 1\}, \quad j \in \mathcal{M}, i \in \mathcal{N}_j \end{aligned}$$



- Variable: $(x_{ij})_{i \in \mathcal{N}, j \in \mathcal{M}_i}$
- Main problem parameters: $(\beta_{ij})_{i \in \mathcal{N}, j \in \mathcal{M}_i}$, $(Q_j)_{j \in \mathcal{M}}$, $(R_{ij})_{i \in \mathcal{N}, j \in \mathcal{M}_i}$
- Constraints: a) Client j can only be assigned to one AP, b) The decision variables are binary

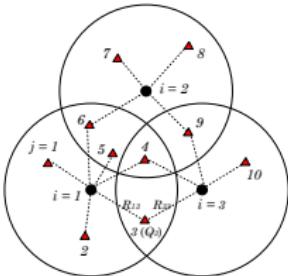
Equivalent Epigraph Form

minimize t

subject to $\sum_{j \in \mathcal{M}_i} \beta_{ij} x_{ij} \leq t, i \in \mathcal{N}$

$\sum_{i \in \mathcal{N}_j} x_{ij} = 1, j \in \mathcal{M}$

$x_{ij} \in \{0, 1\}, j \in \mathcal{M}, i \in \mathcal{N}_j$



- Variable: $(x_{ij})_{i \in \mathcal{N}, j \in \mathcal{M}_i}$ and t
- Main problem parameters: $(\beta_{ij})_{i \in \mathcal{N}, j \in \mathcal{M}_i}$
- Mixed integer linear program (MILP)

Solution Method Challenges

- Existing **MILP** solvers are **centralized**
- Typically based on global branch and bound algorithms \Rightarrow the **worst-case complexity grows exponentially** with the problem size
- **Even small** problems, with a few tens of variables, can take a **very long time**

Solution Method Challenges

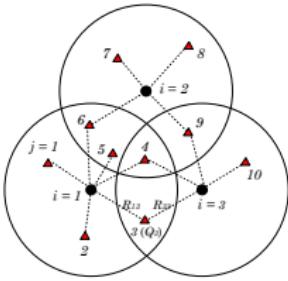
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- **Even small** problems, with a few tens of variables, can take a **very long time**
- **Our approach: efficient, optimality ?**

Lagrangian Duality

Partial Lagrangian

$$\begin{aligned}
 L(t, \mathbf{x}, \boldsymbol{\lambda}) &= t + \sum_{i \in \mathcal{N}} \lambda_i \left(\sum_{j \in \mathcal{M}_i} \beta_{ij} x_{ij} - t \right) \\
 &= t \left(1 - \sum_{i \in \mathcal{N}} \lambda_i \right) + \sum_{j \in \mathcal{M}} \sum_{i \in \mathcal{N}_j} \beta_{ij} \lambda_i x_{ij}
 \end{aligned}$$

$\boldsymbol{\lambda} = (\lambda_i)_{i \in \mathcal{N}}$: multipliers for the first set of inequality constraints



Lagrangian Duality

Partial Lagrangian

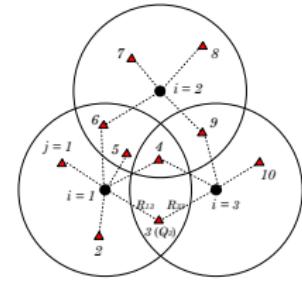
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Dual function

$$g(\boldsymbol{\lambda}) = \inf_{\substack{t \in \mathbb{R} \\ \mathbf{x} \in \mathcal{X}}} L(t, \mathbf{x}, \boldsymbol{\lambda})$$

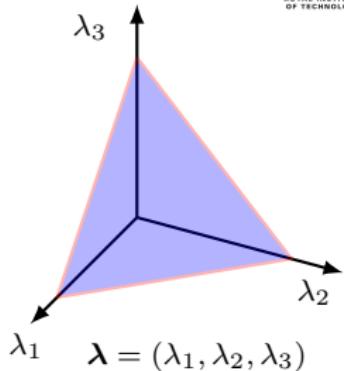
$$\mathcal{X} = \left\{ \mathbf{x} = (x_{ij})_{j \in \mathcal{M}, i \in \mathcal{N}_j} \mid \sum_{i \in \mathcal{N}_j} x_{ij} = 1, \quad x_{ij} \in \{0, 1\}, \quad j \in \mathcal{M}, i \in \mathcal{N}_j \right\}.$$



Lagrangian Duality

Dual problem

$$\begin{aligned} & \text{maximize} && g(\boldsymbol{\lambda}) = \sum_{j \in \mathcal{M}} g_j(\boldsymbol{\lambda}) \\ & \text{subject to} && \sum_{i \in \mathcal{N}} \lambda_i = 1 \\ & && \lambda_i \geq 0, \quad i \in \mathcal{N} \end{aligned}$$



- variables: $\boldsymbol{\lambda} = (\lambda_i)_{i \in \mathcal{N}}$
- $g_j(\boldsymbol{\lambda})$ is the optimal value of the **subproblem**

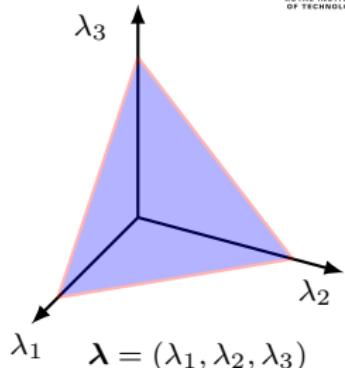
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with the variable \mathbf{x}_j . Solved @ client j

Lagrangian Duality

Dual problem

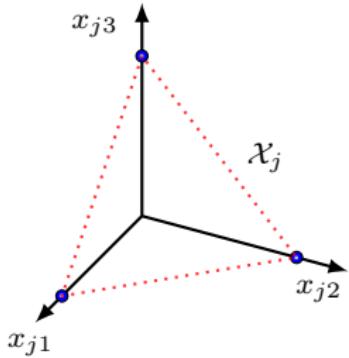
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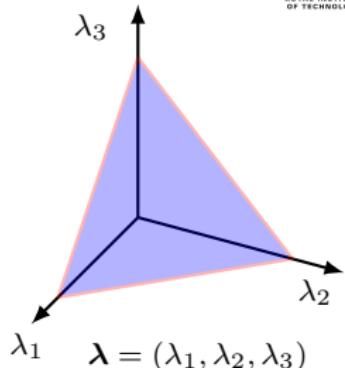


$$\mathbf{x}_j = (x_{j1}, x_{j2}, x_{j3})$$

Lagrangian Duality

Dual problem

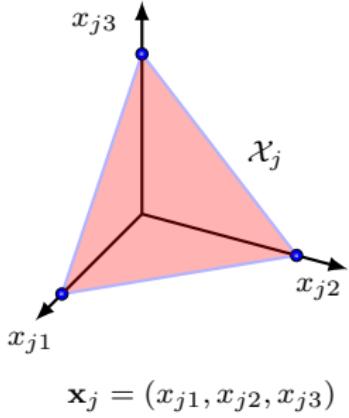
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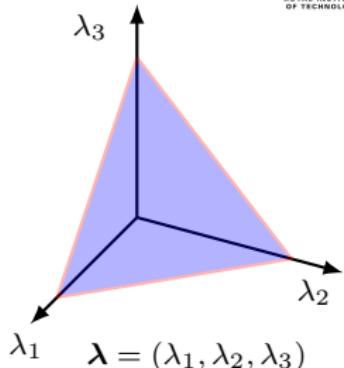
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Lagrangian Duality

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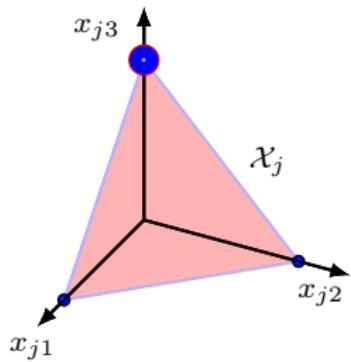
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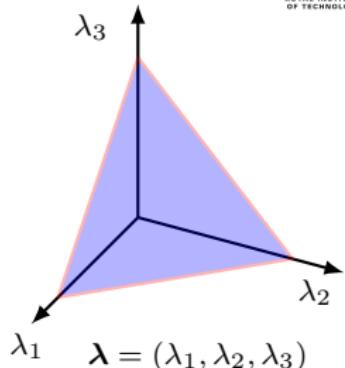
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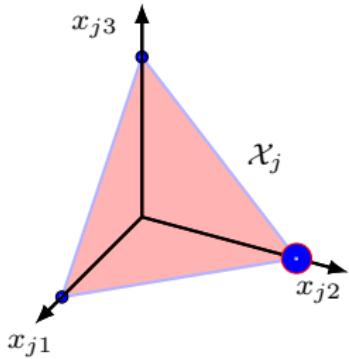
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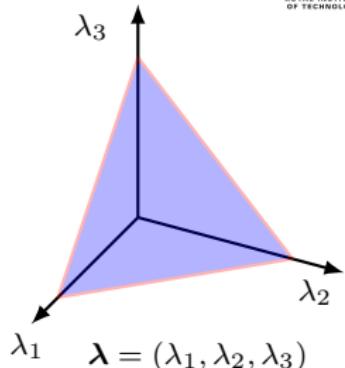
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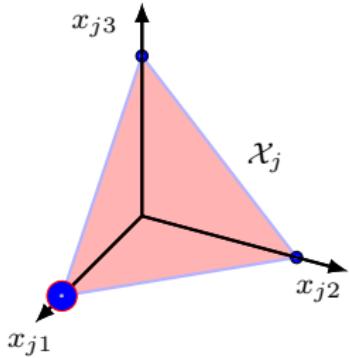
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Subgradient Method

Proj. subgradient method to the dual problem

$$\boldsymbol{\lambda}^{(k+1)} = P(\boldsymbol{\lambda}^{(k)} - \alpha_k \mathbf{u}^{(k)})$$

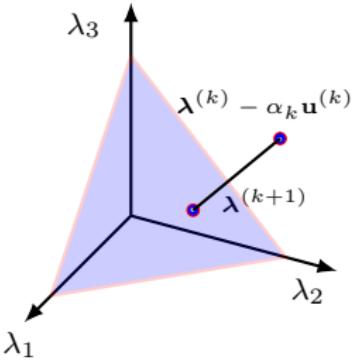
- P : Euclidean projection onto the unit simplex

$$\Pi = \{\boldsymbol{\lambda} \mid \sum_{i \in \mathcal{N}} \lambda_i = 1, \lambda_i \geq 0\}$$

- $\alpha_k > 0$ is the k th step size
- $\mathbf{u}^{(k)} = \left(u_i^{(k)}\right)_{i \in \mathcal{N}}$: a subgradient of $-g$ at $\boldsymbol{\lambda}^{(k)}$, where

$$u_i^{(k)} = -\sum_{j \in \mathcal{M}_i} \beta_{ij} x_{ij}^*,$$

and $(x_{ij}^*)_{j \in \mathcal{M}_i}$ is the solution of the i th subproblems with $\boldsymbol{\lambda} = \boldsymbol{\lambda}^{(k)}$



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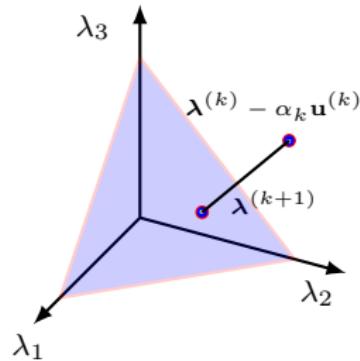
→ price update

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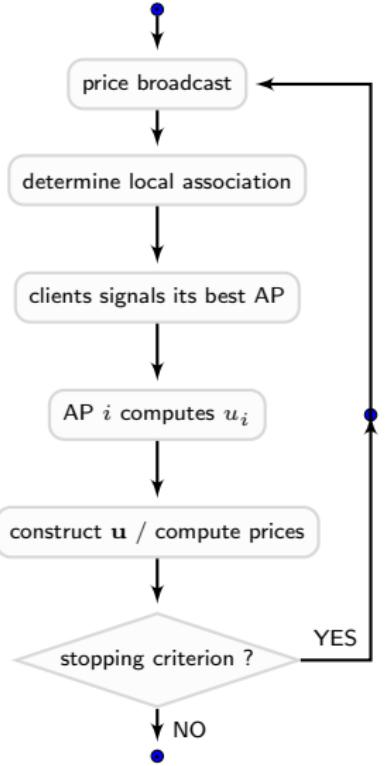
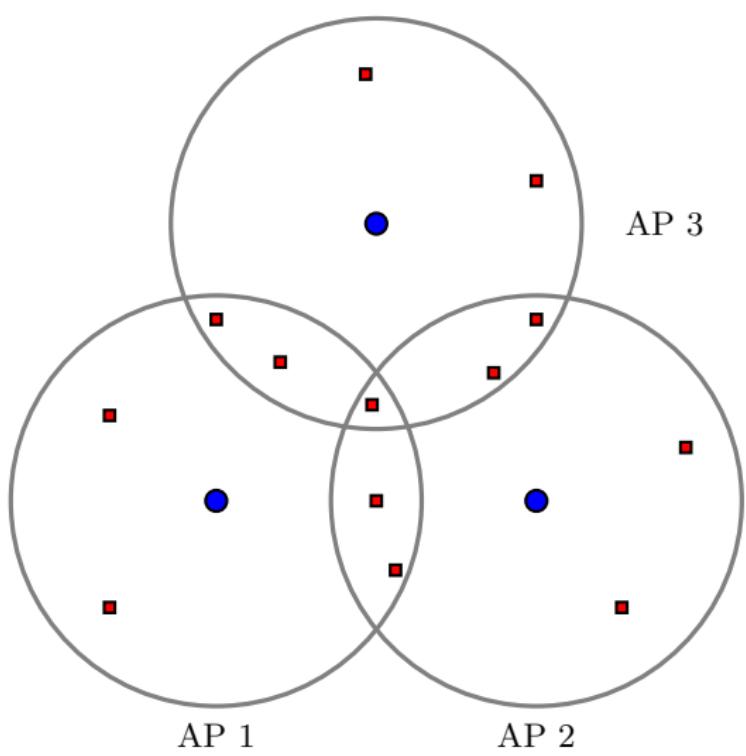


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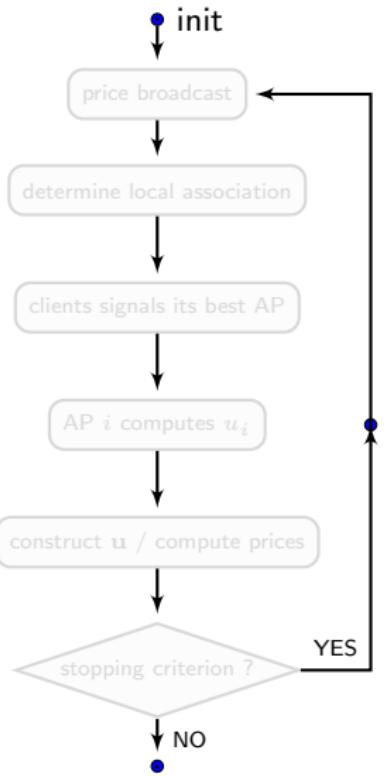
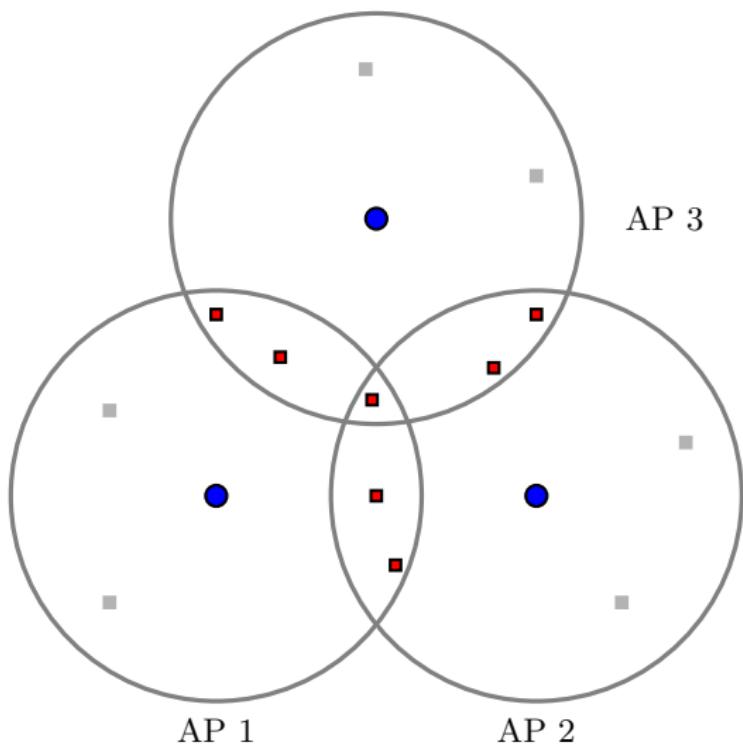
Outline

- Past, Present and Future in wireless communications
- 60 GHz millimeterWave wireless technology
- Optimizing resource allocation
- **Distributed client association (DAA)**
- Numerical analysis of DAA
- Conclusions and open research topics

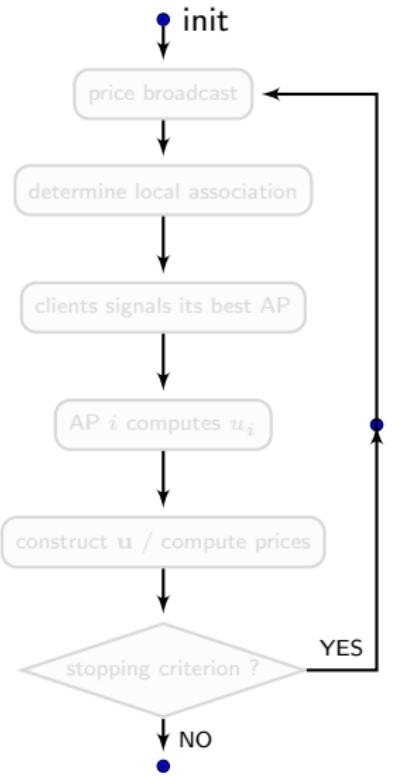
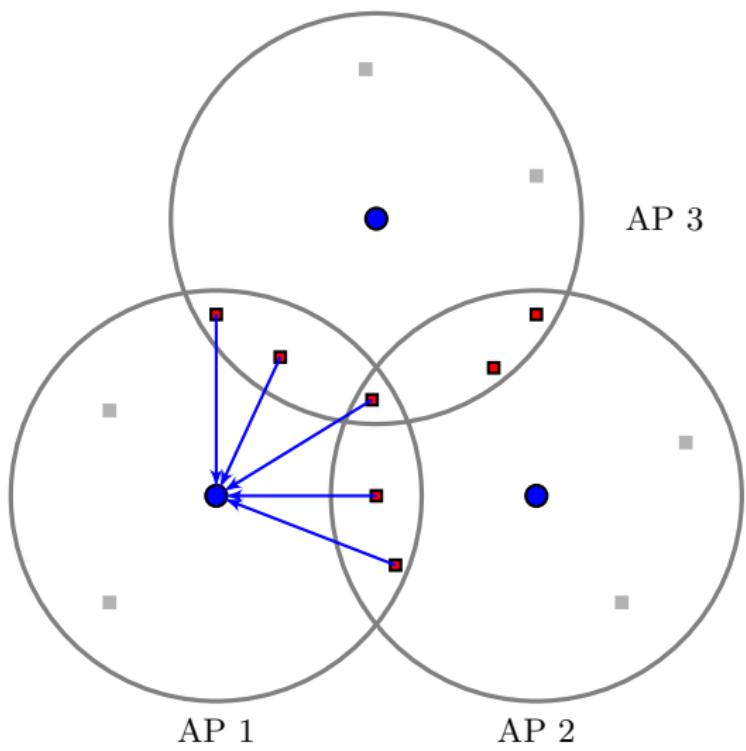
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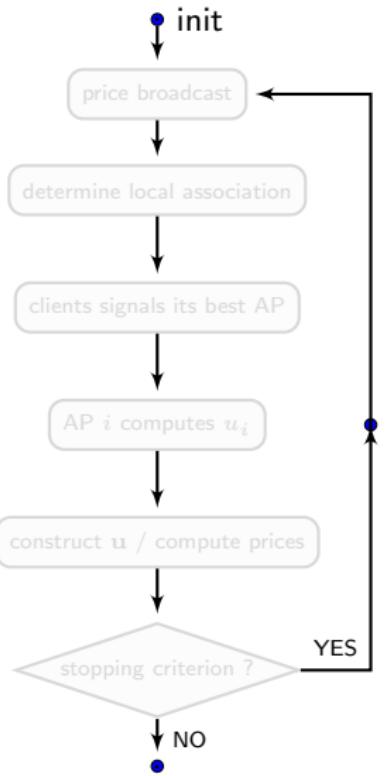
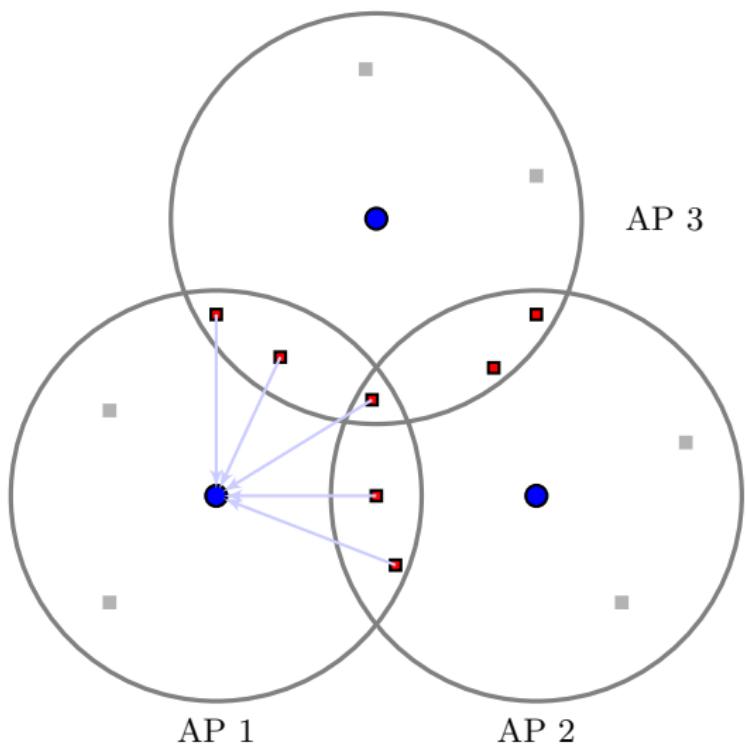
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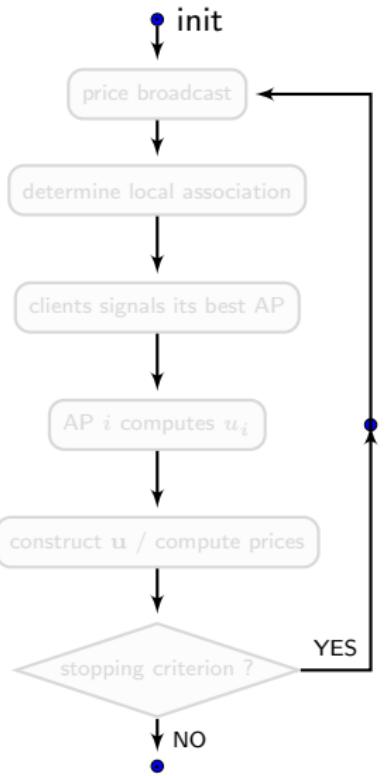
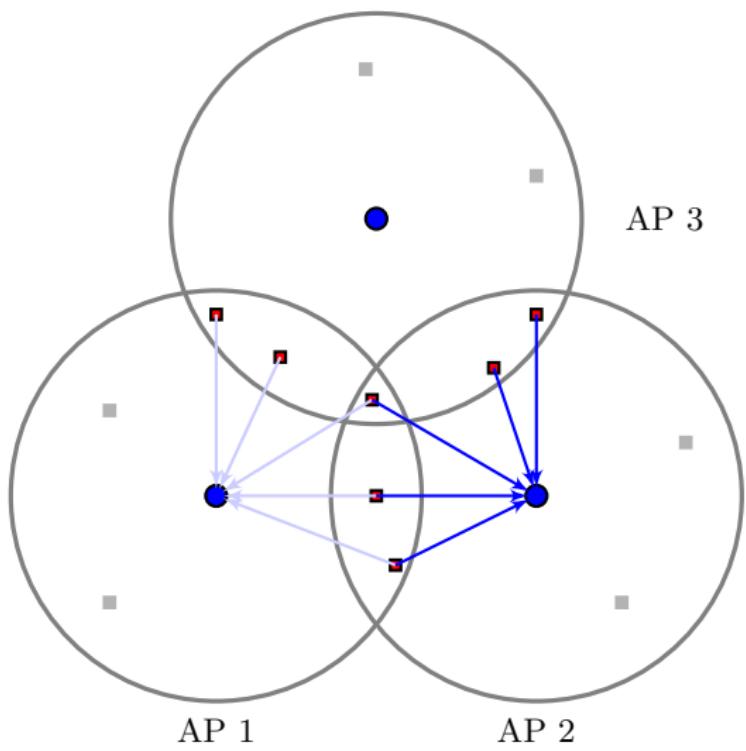
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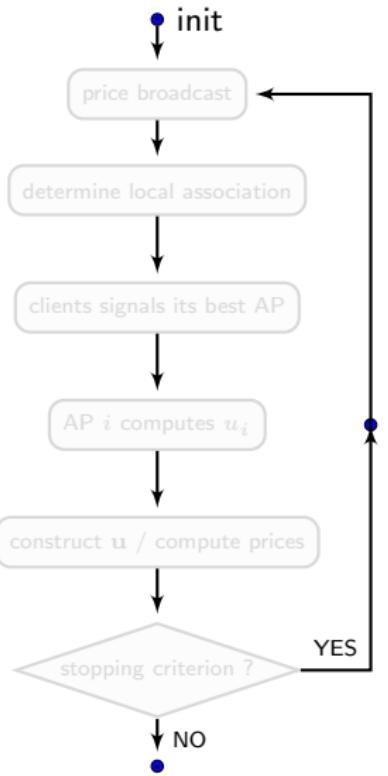
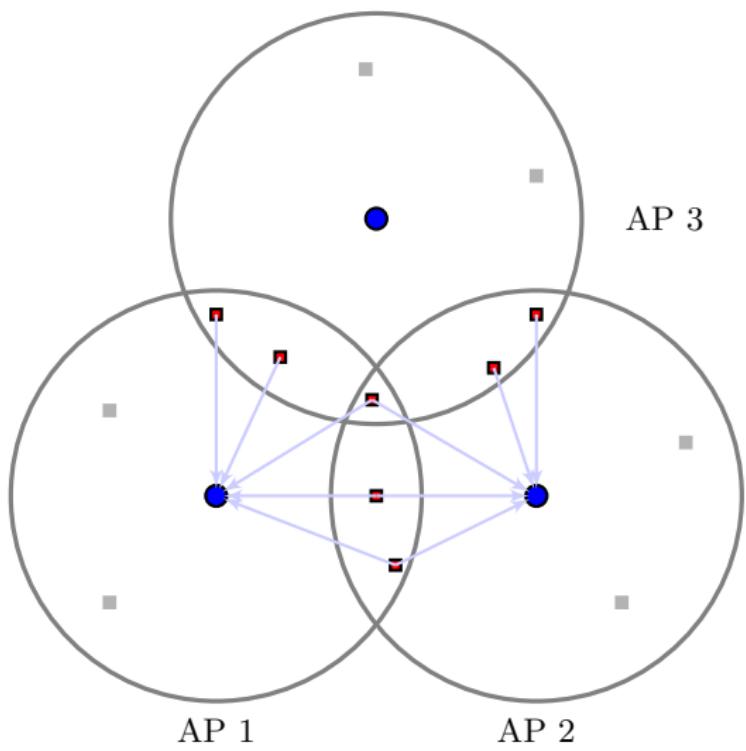
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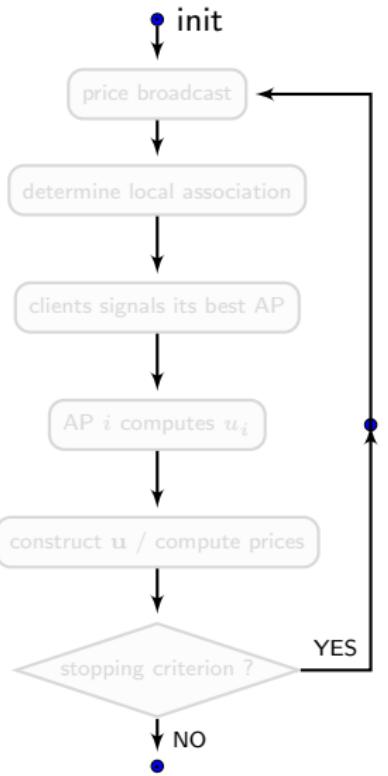
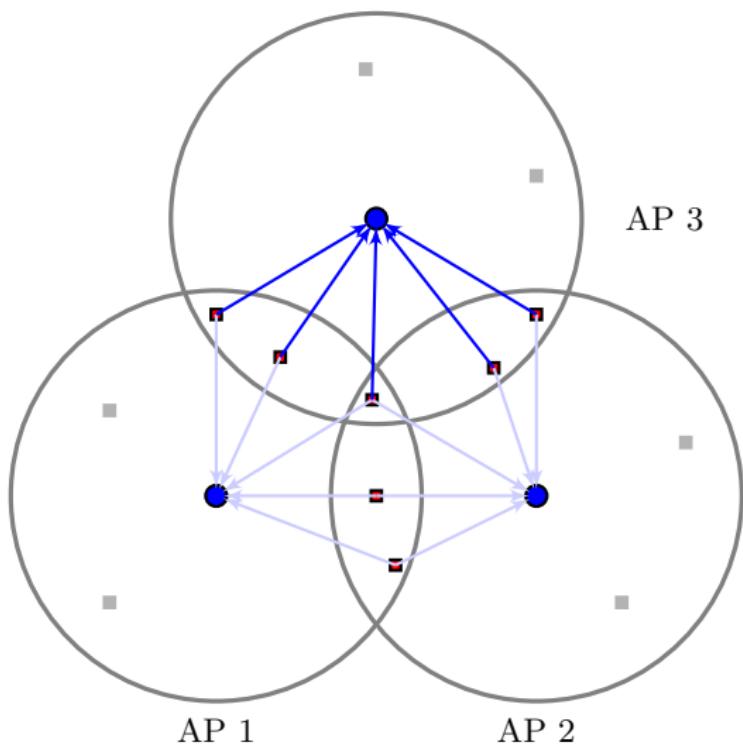
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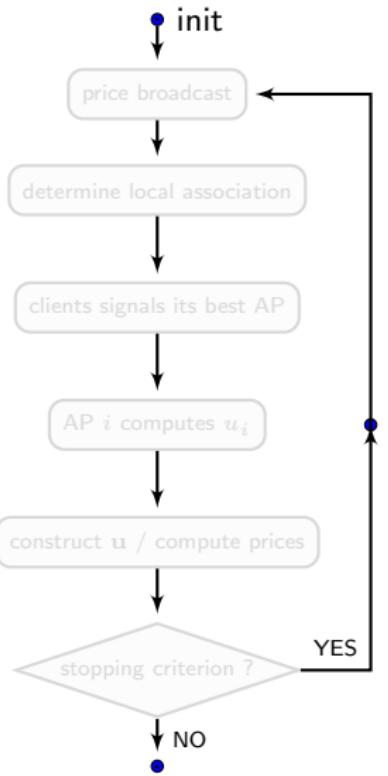
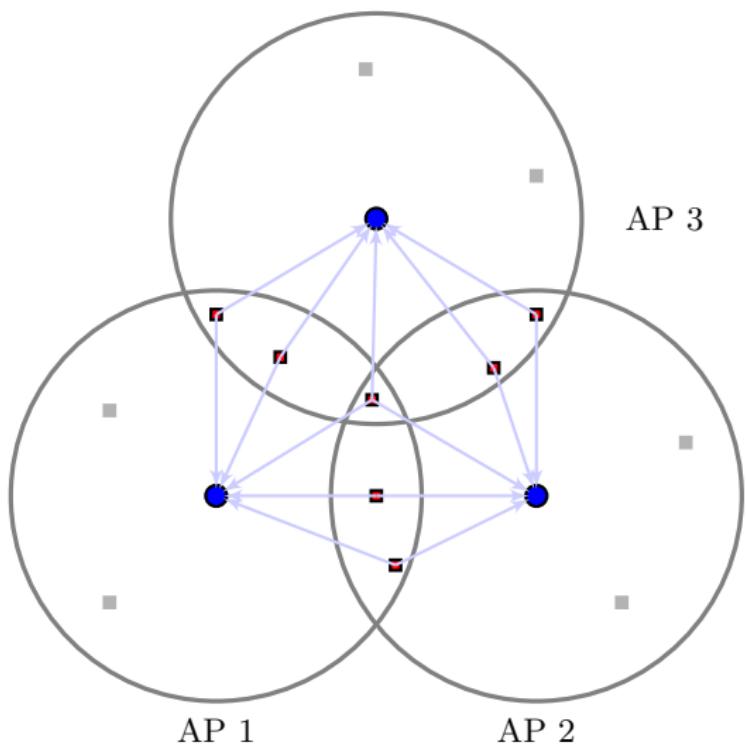
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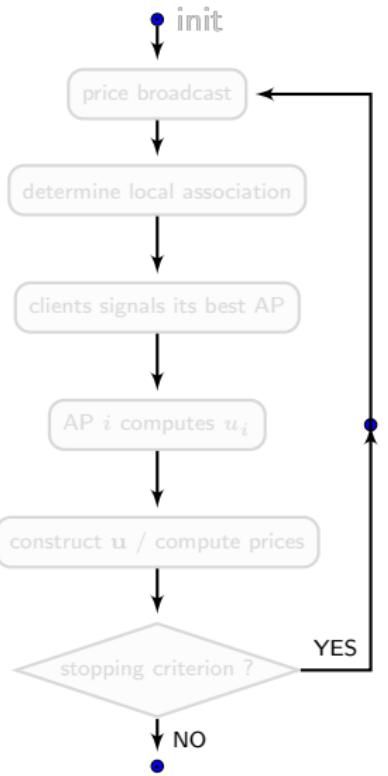
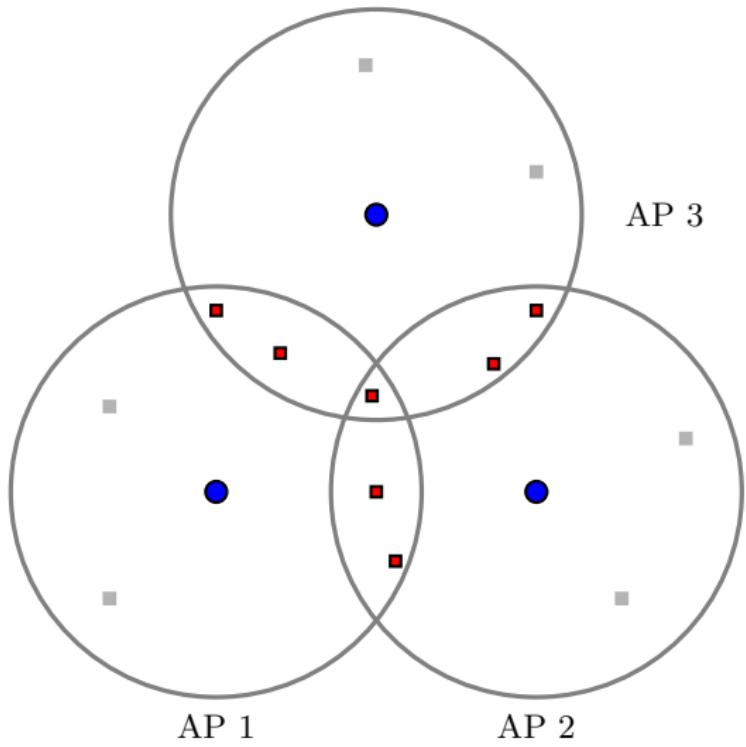
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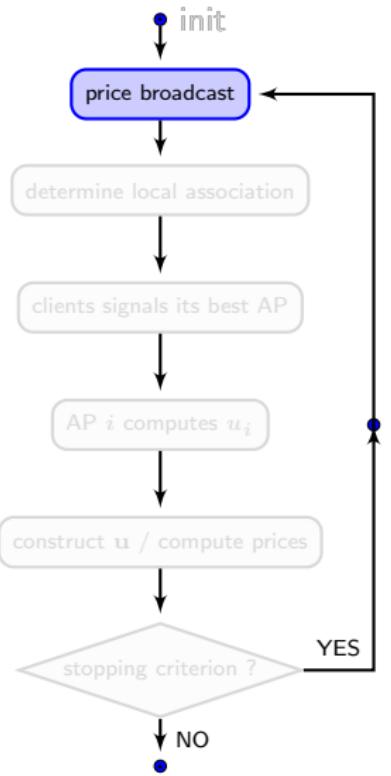
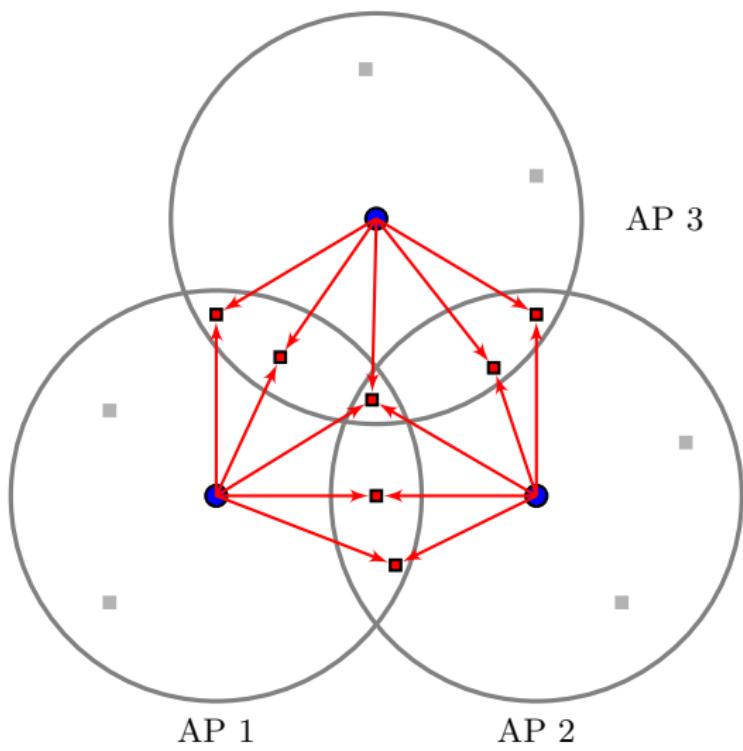
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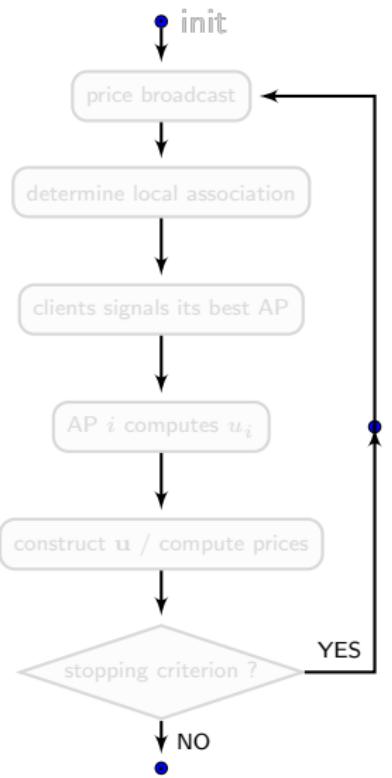
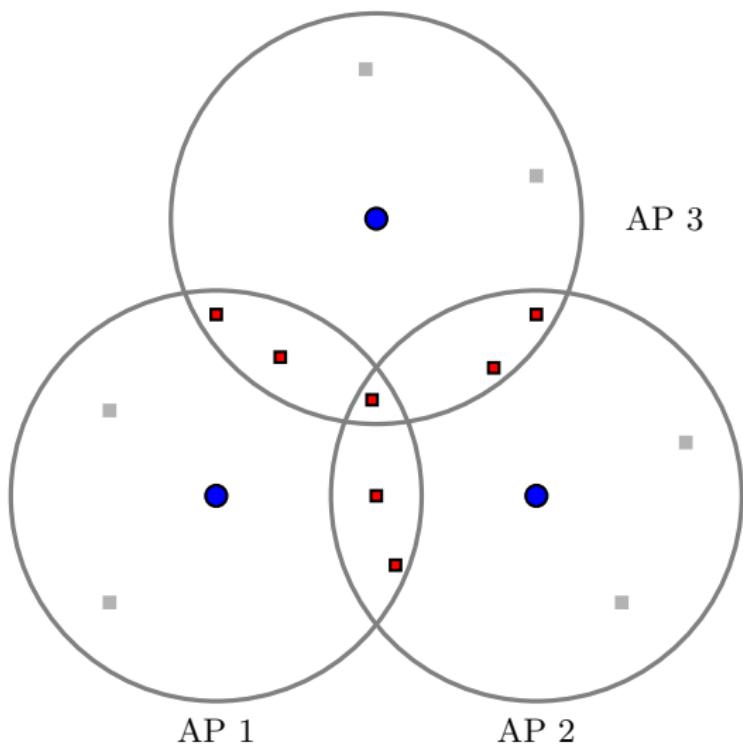
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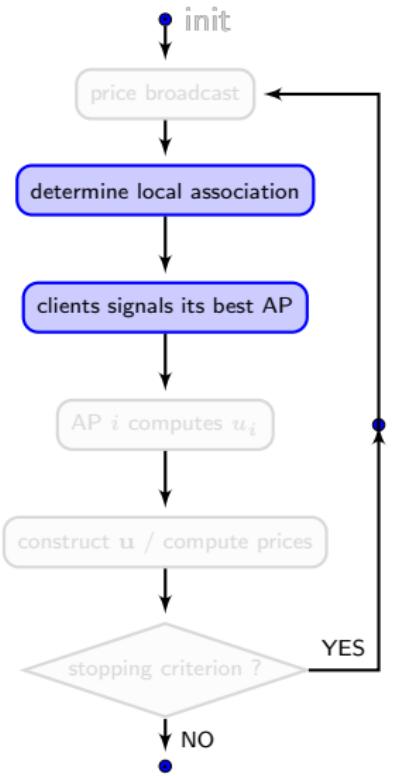
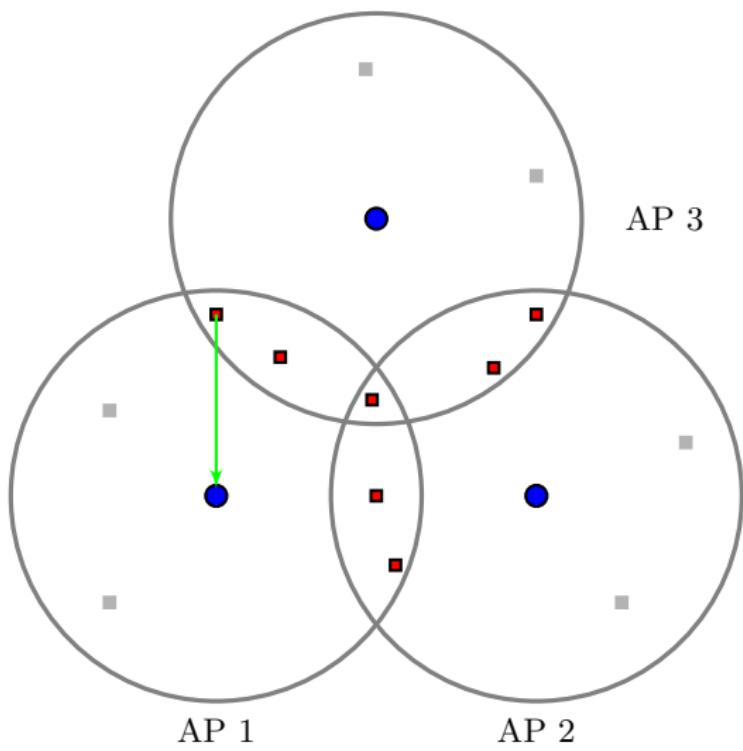
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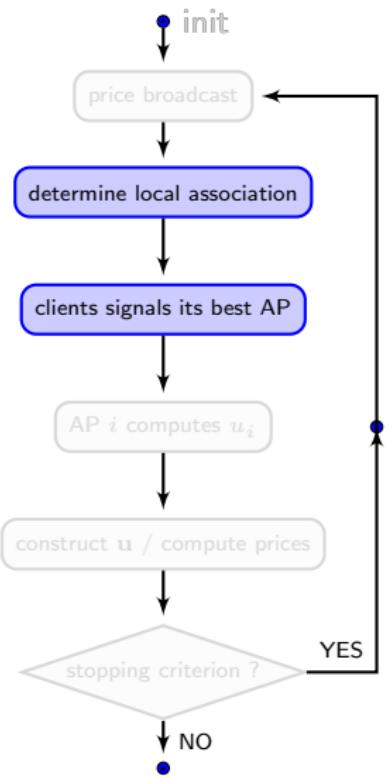
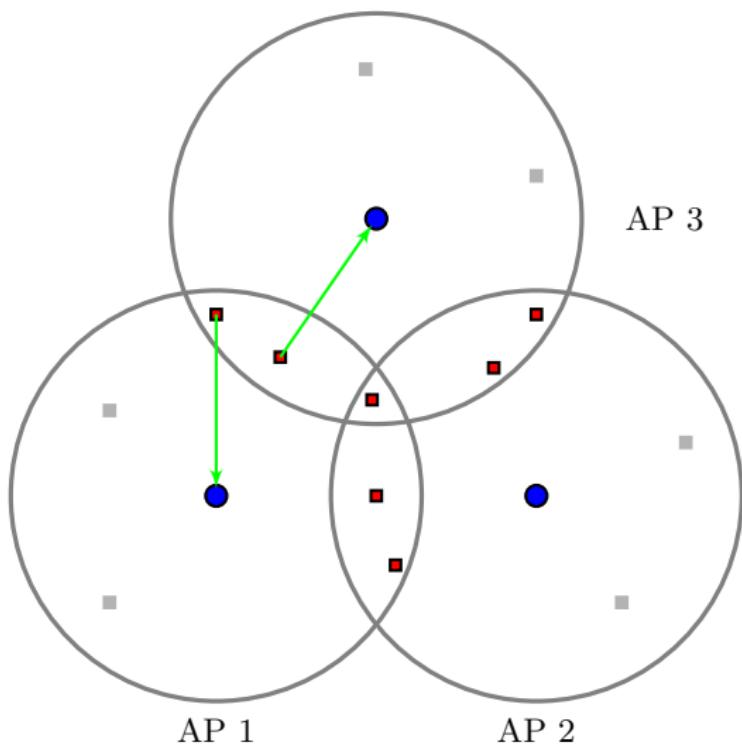
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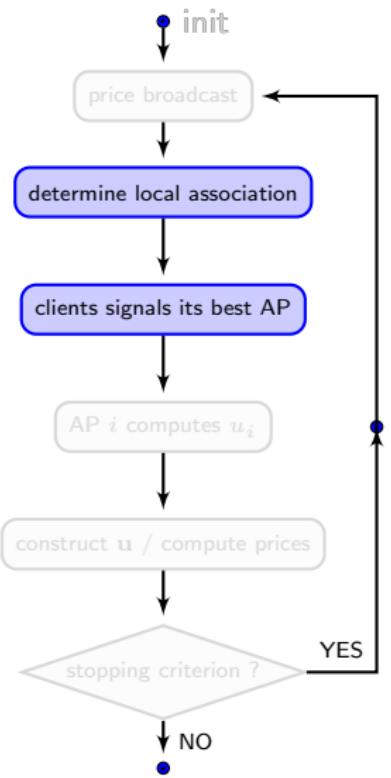
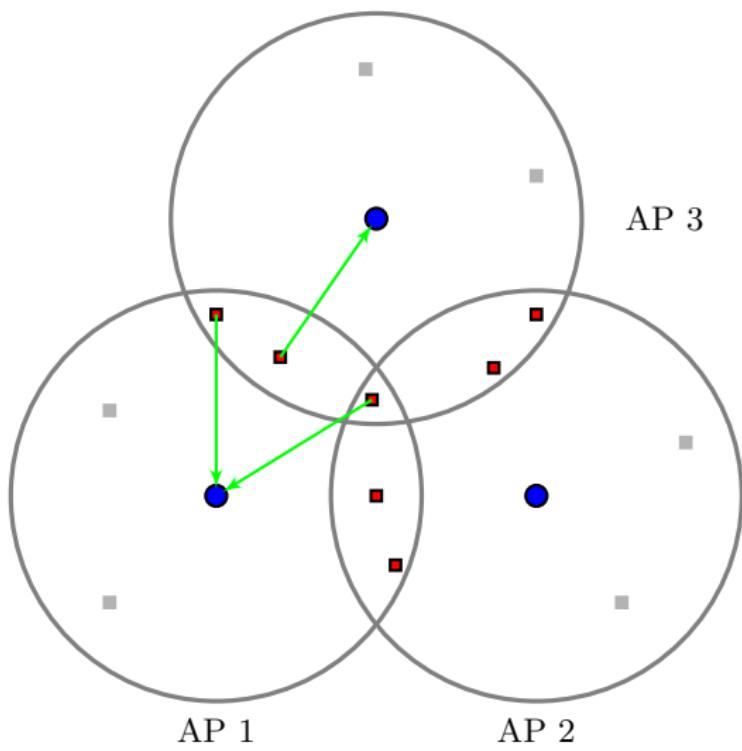
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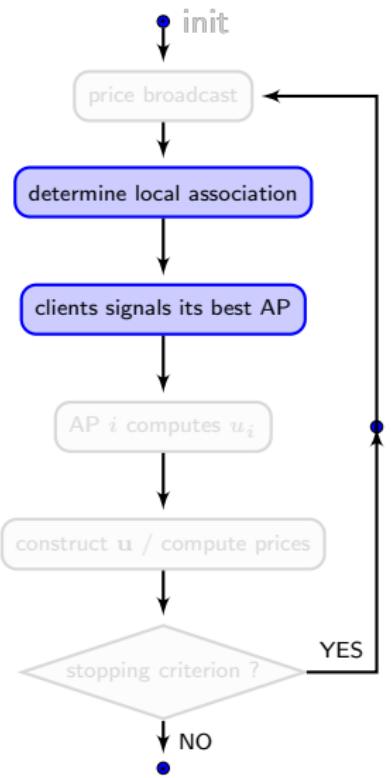
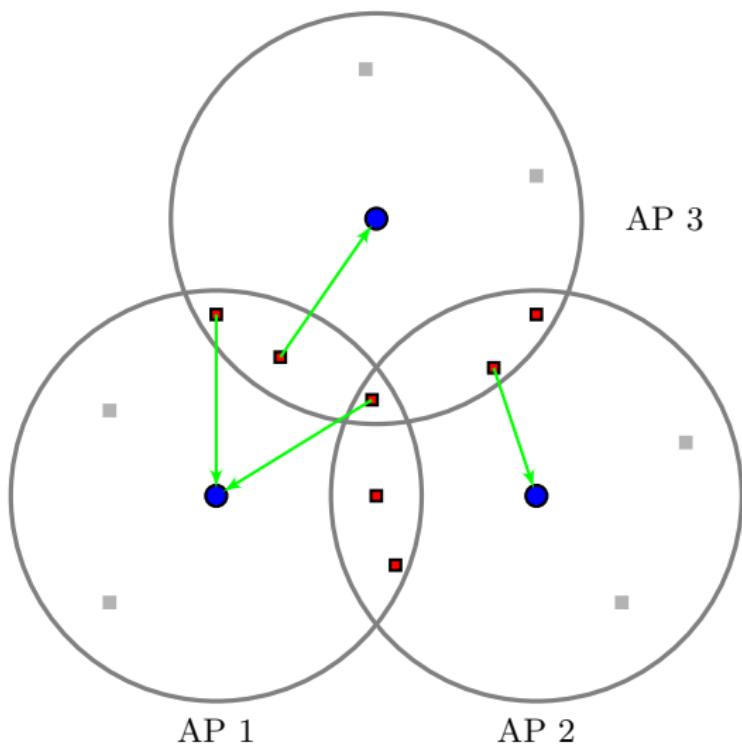
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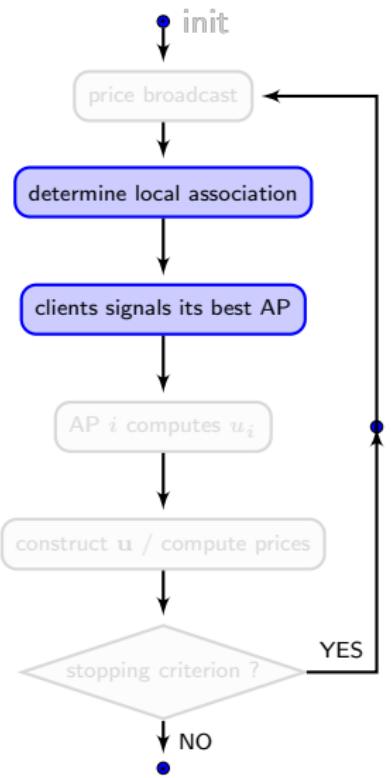
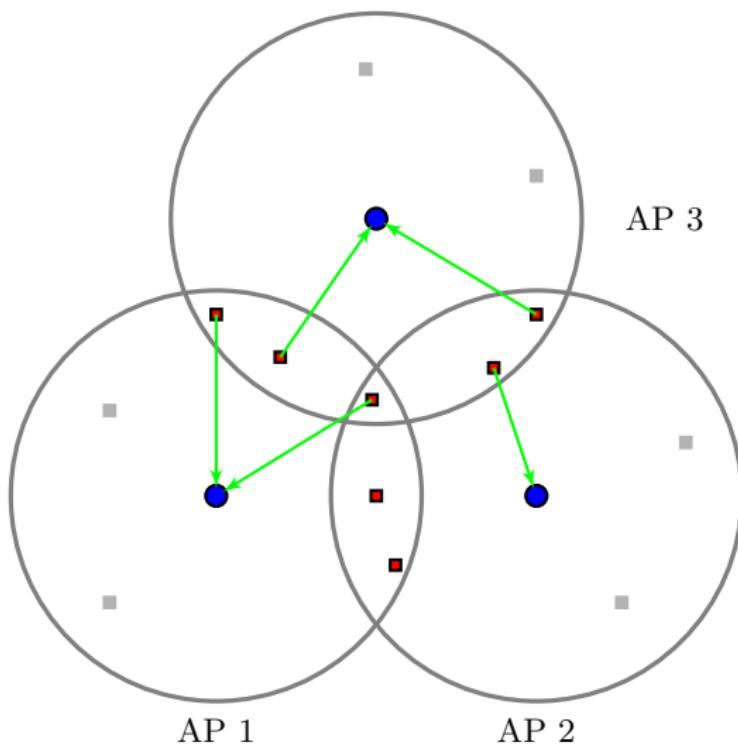
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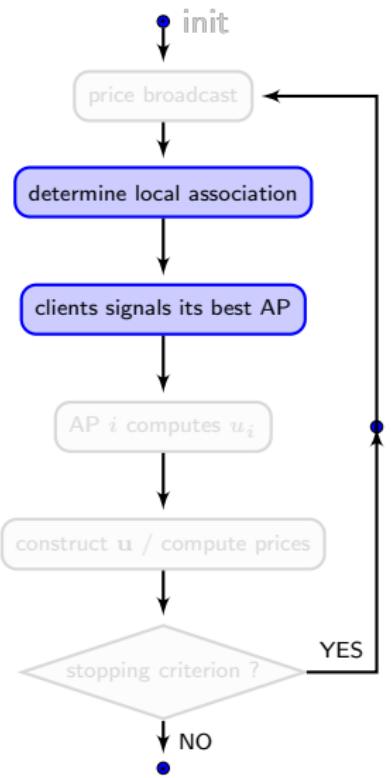
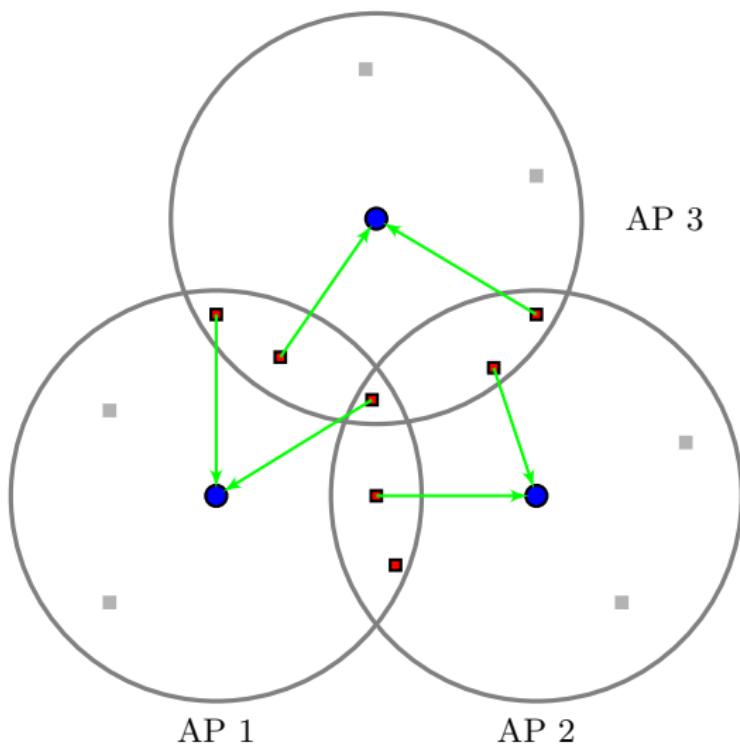
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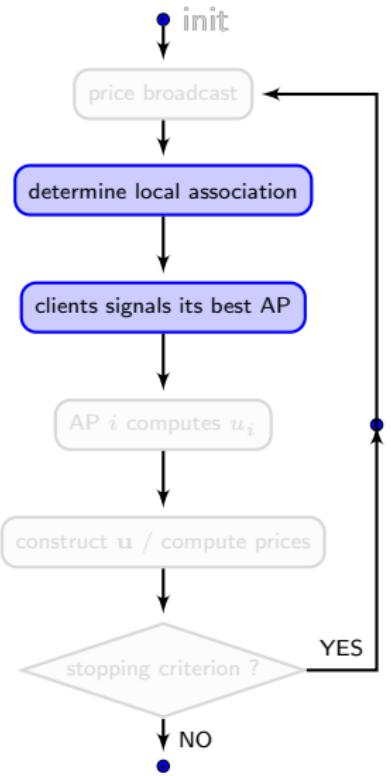
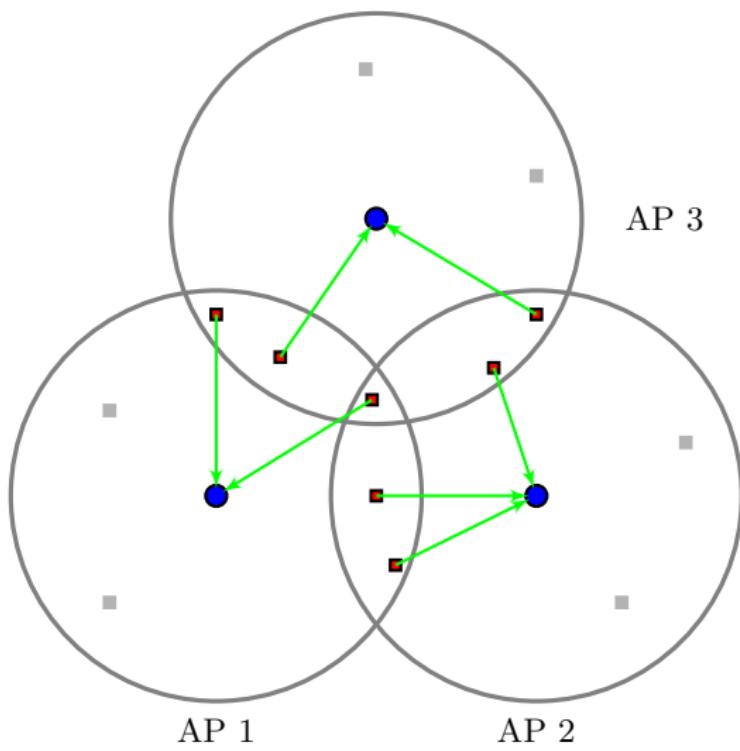
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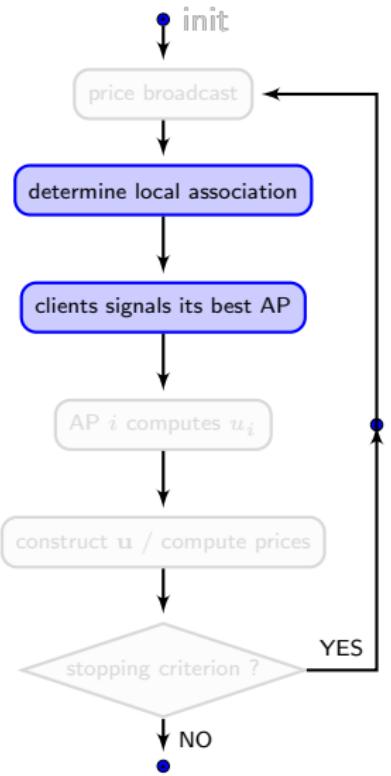
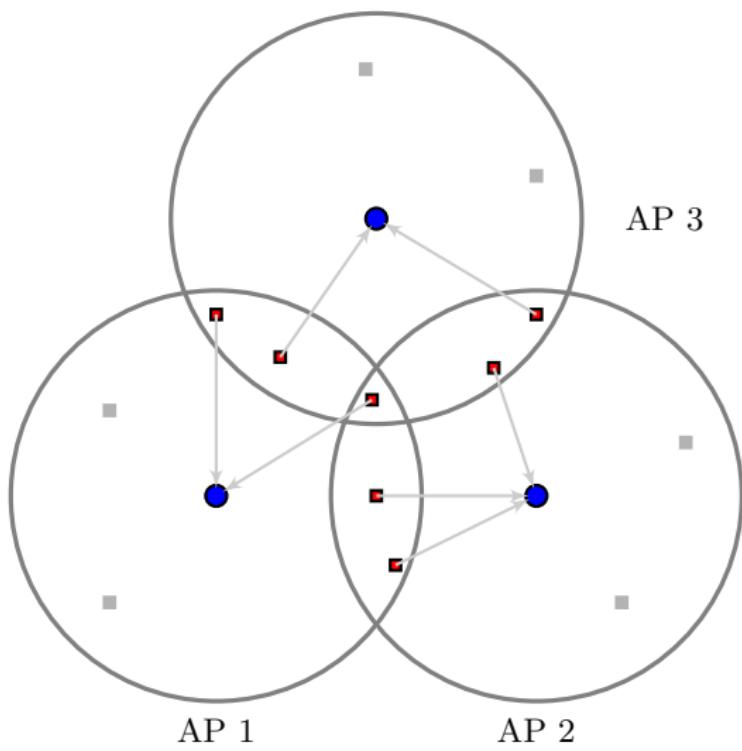
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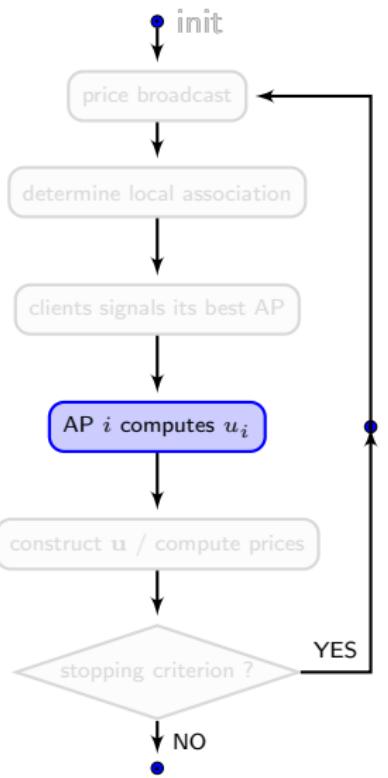
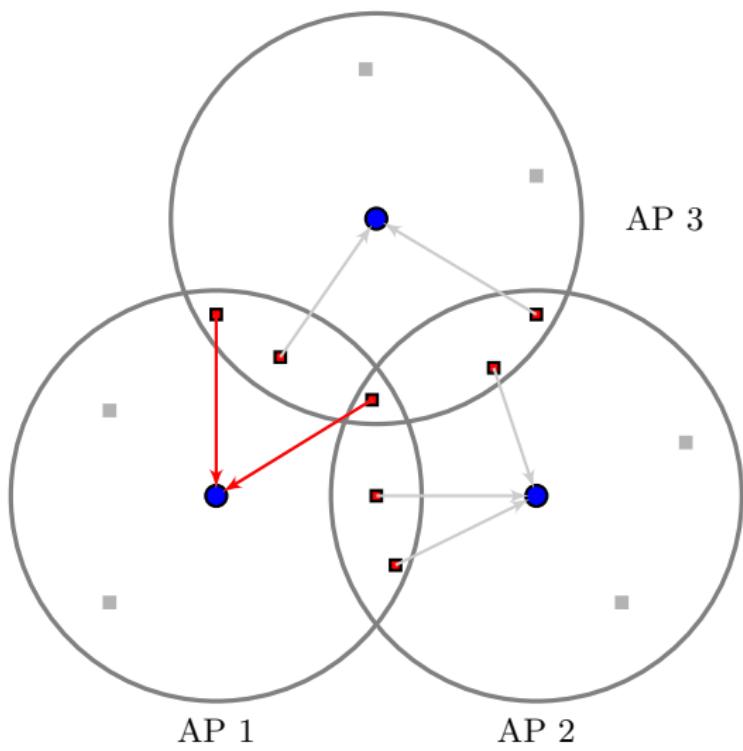
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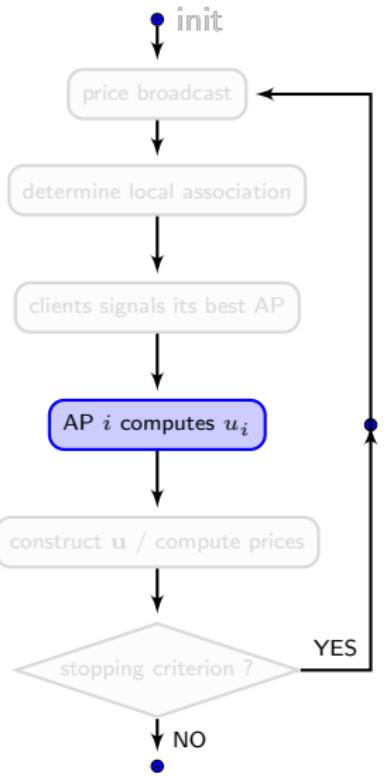
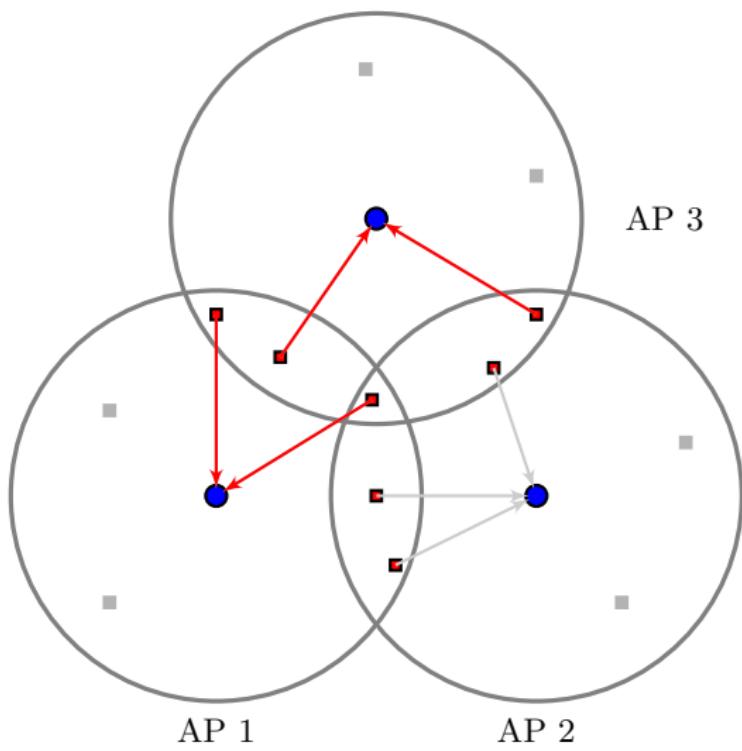
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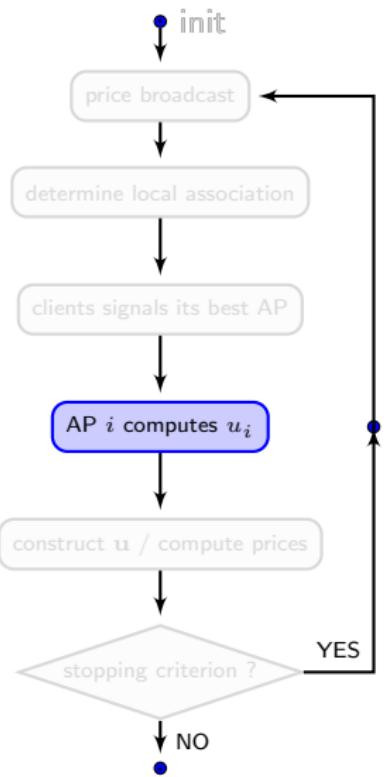
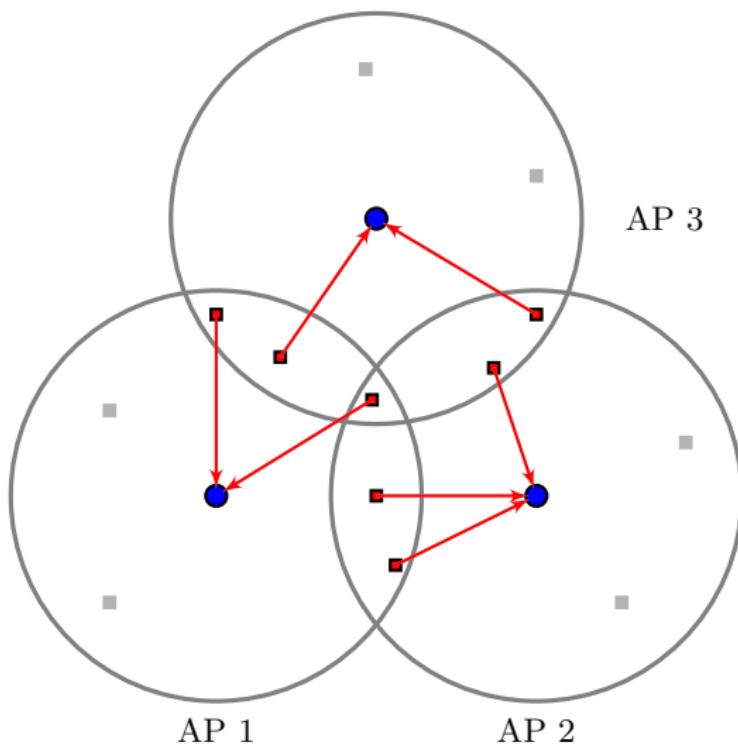
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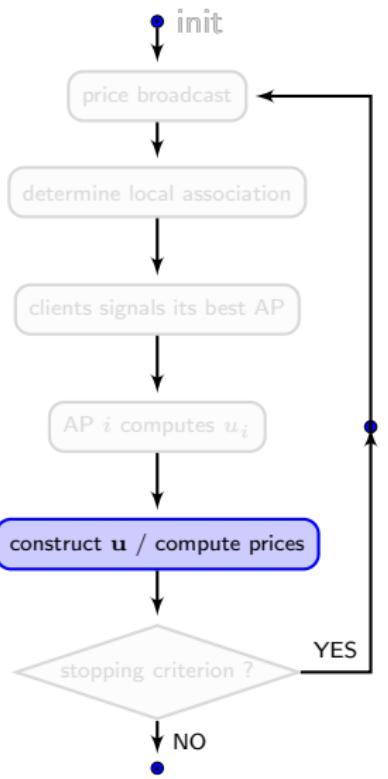
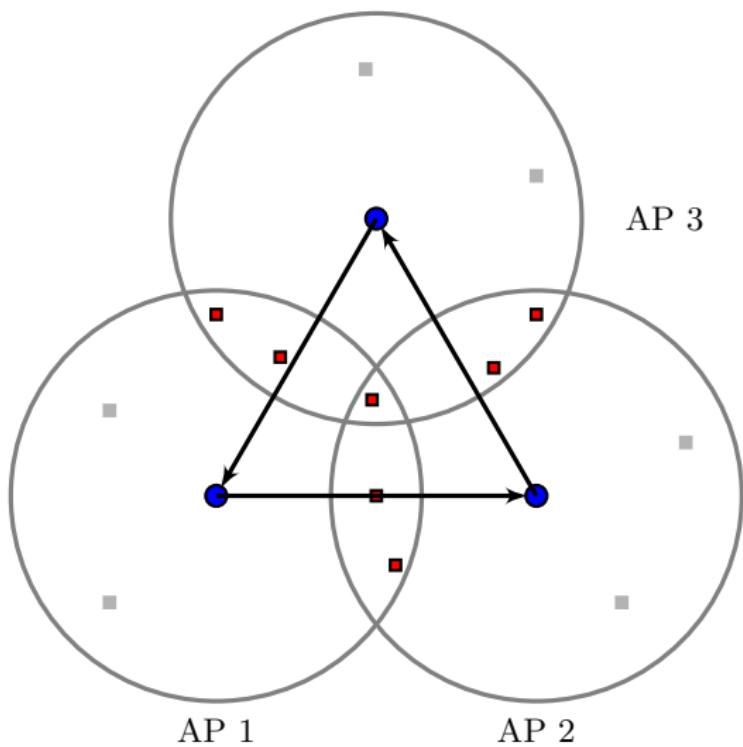
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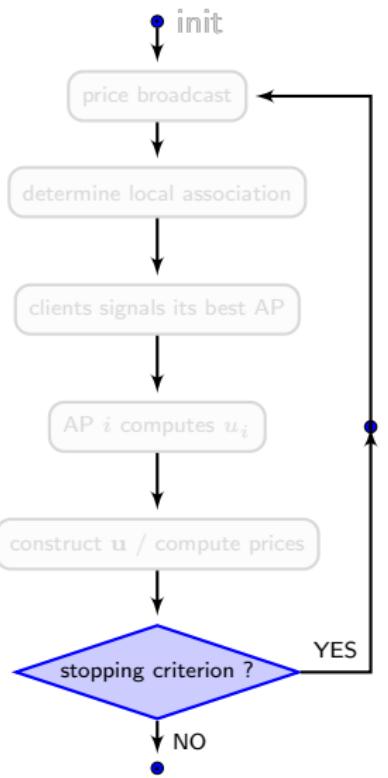
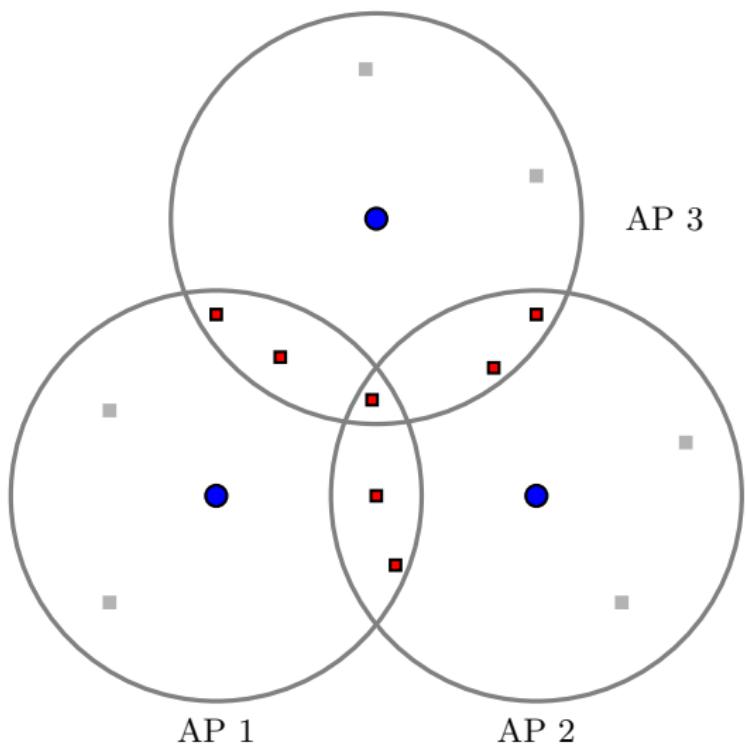
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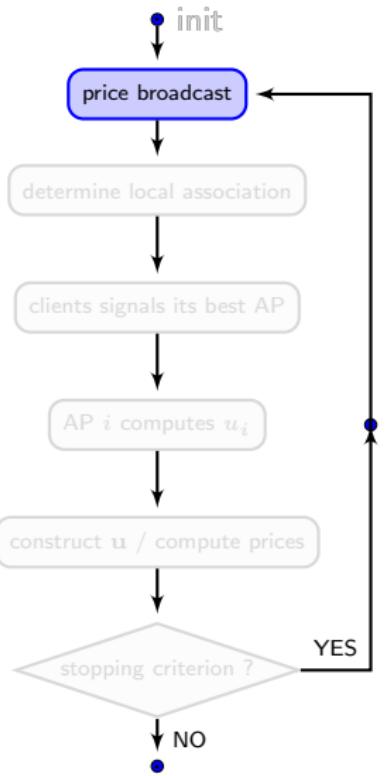
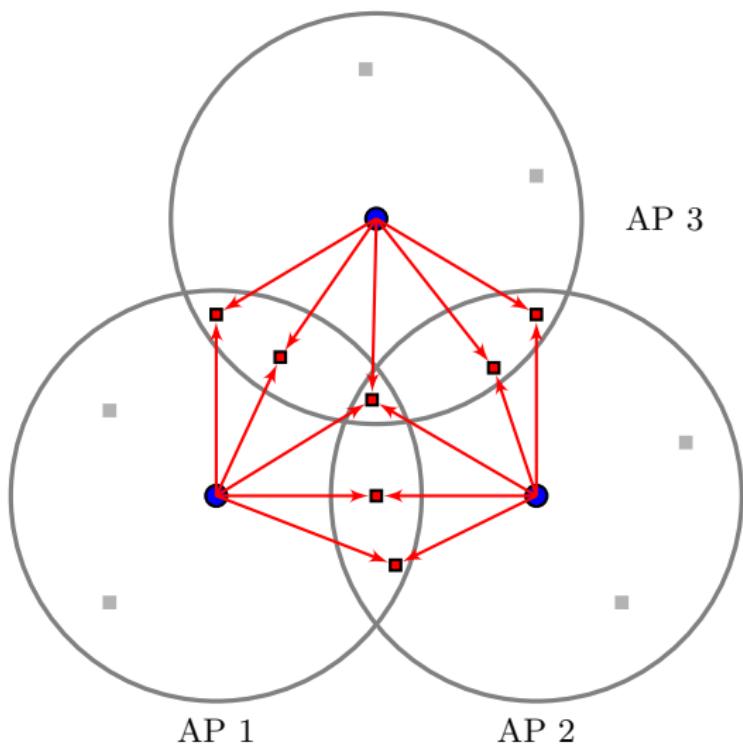
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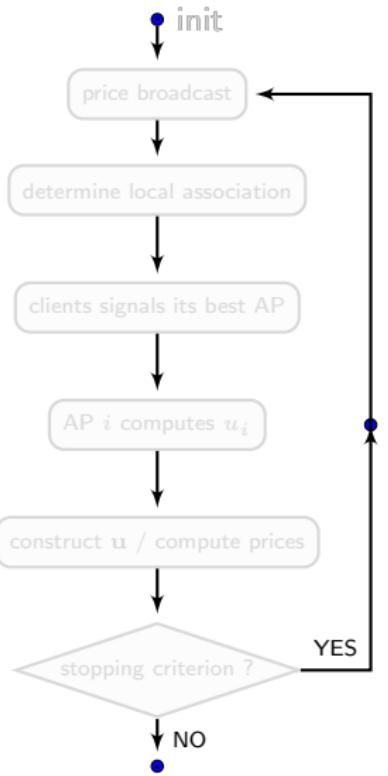
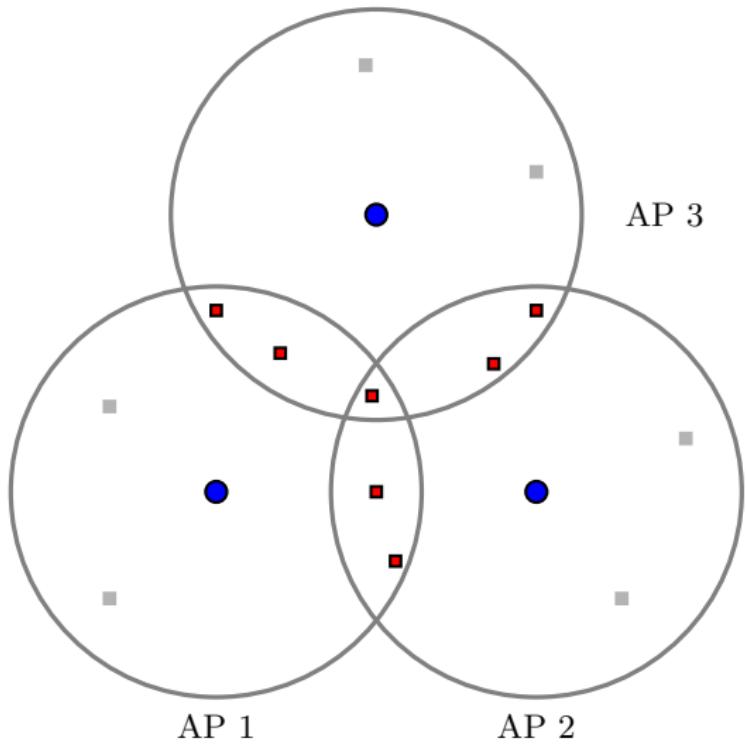
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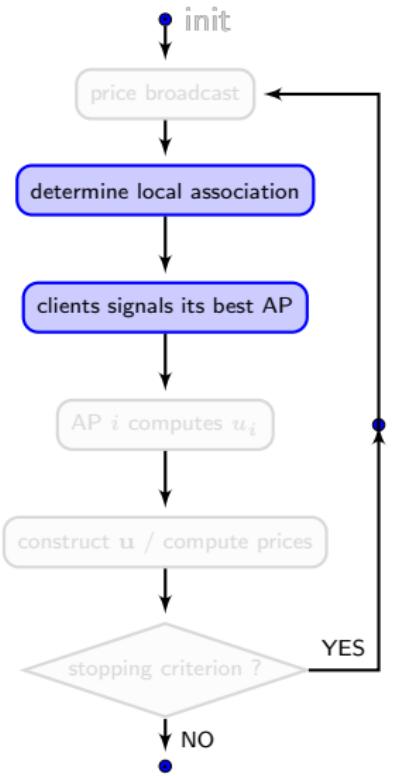
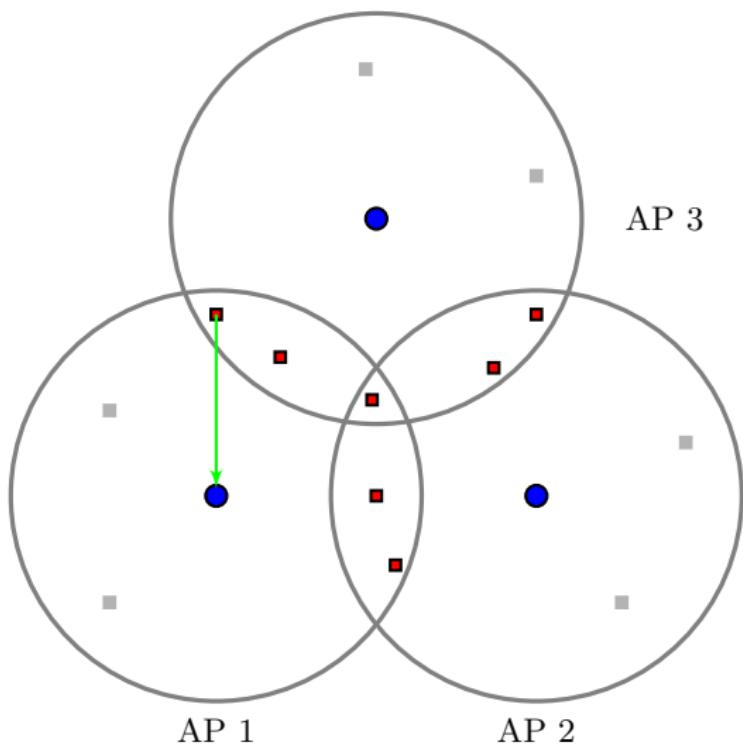
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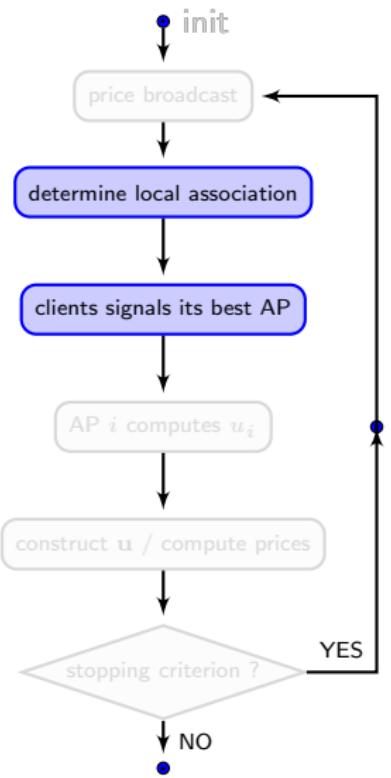
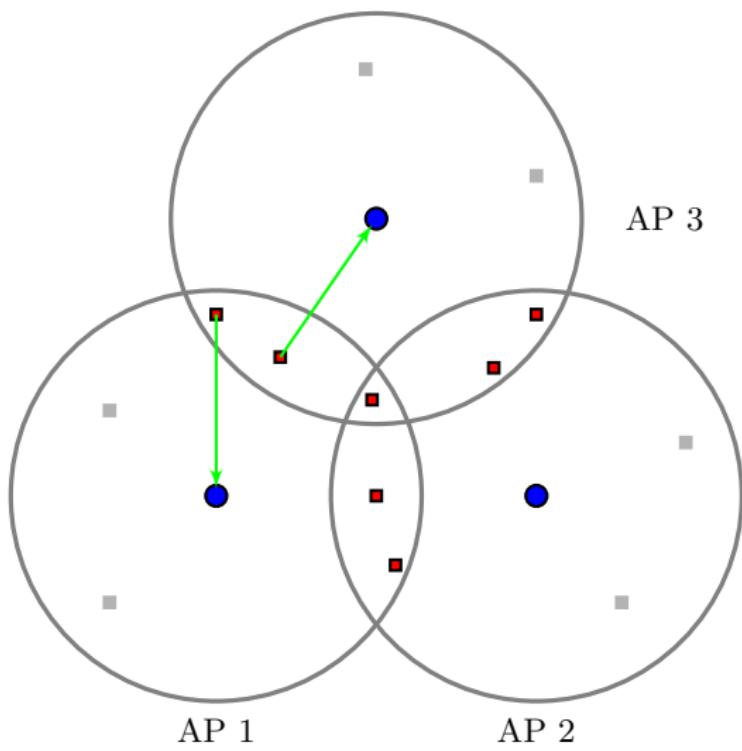
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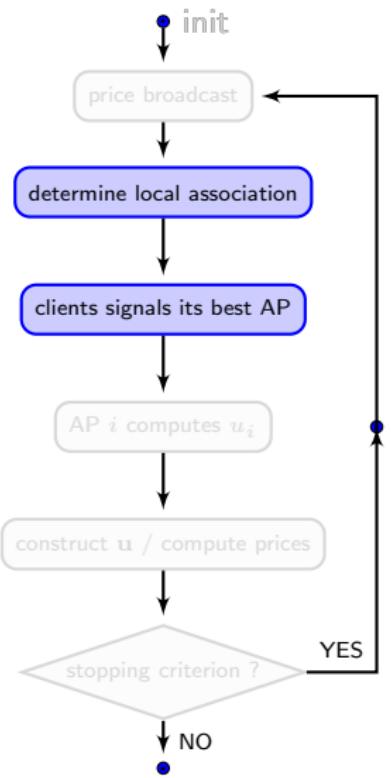
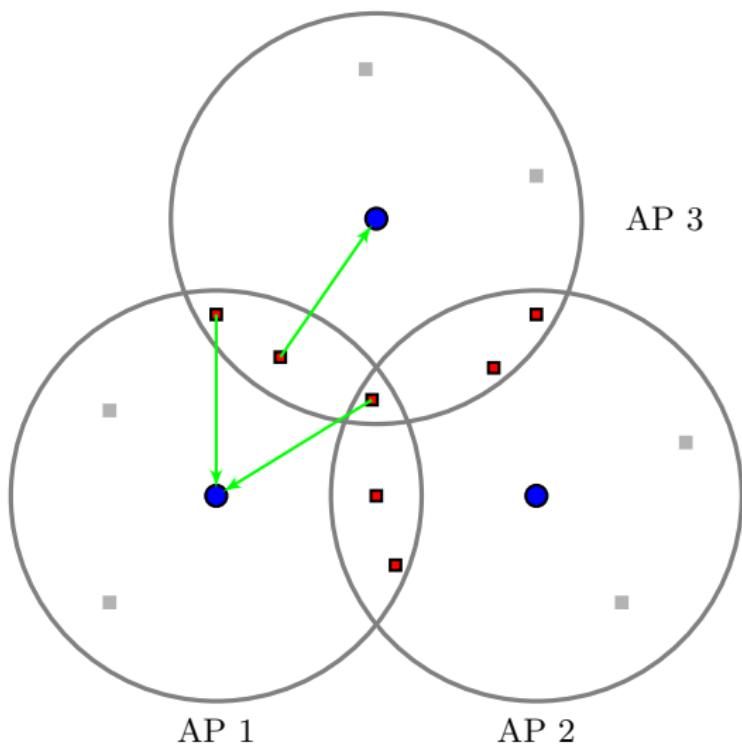
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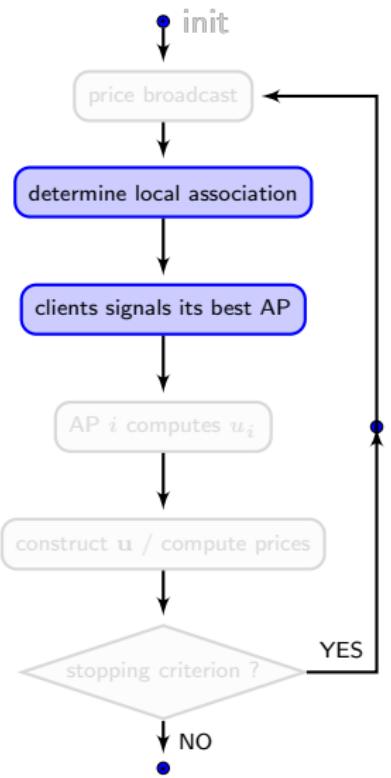
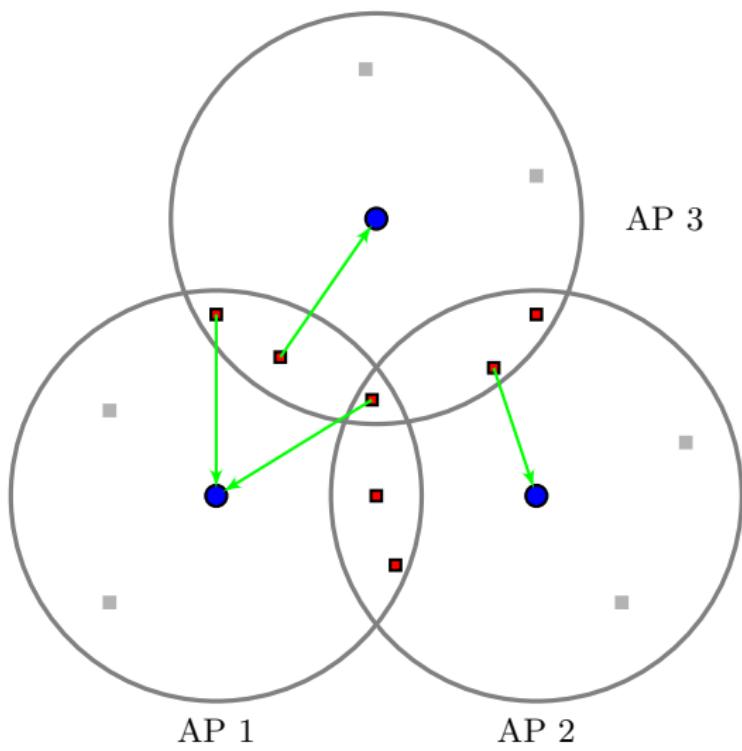
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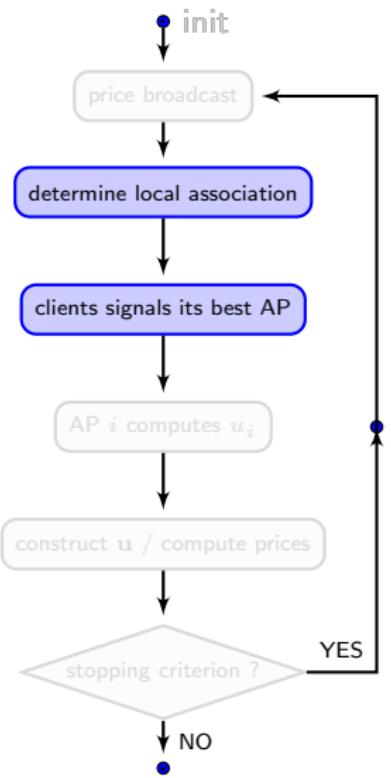
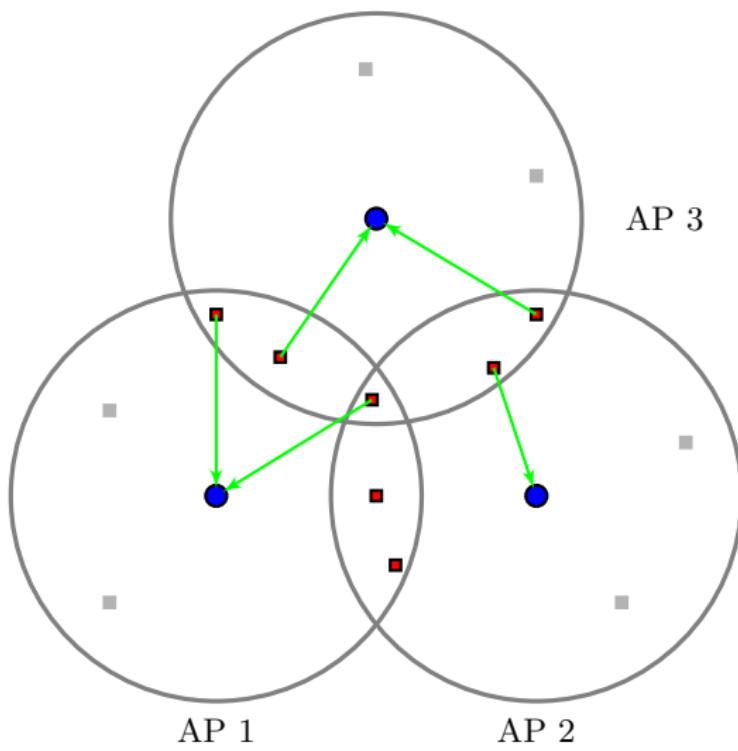
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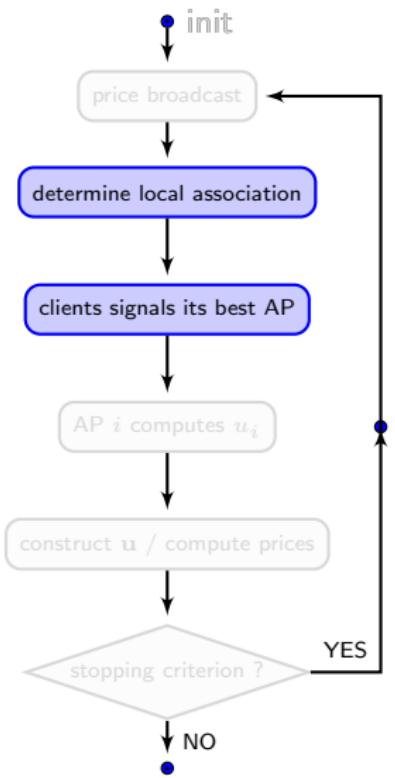
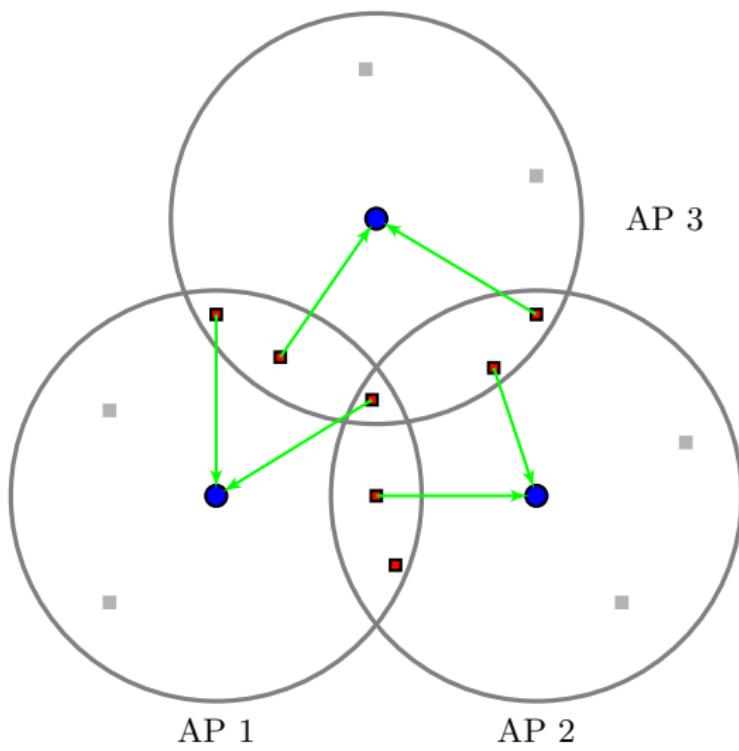
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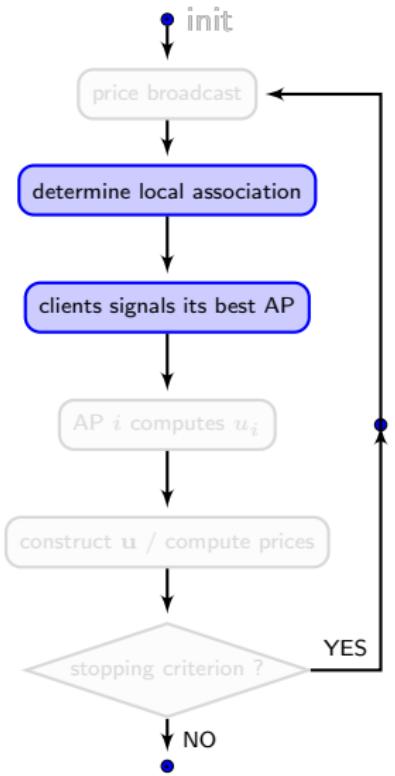
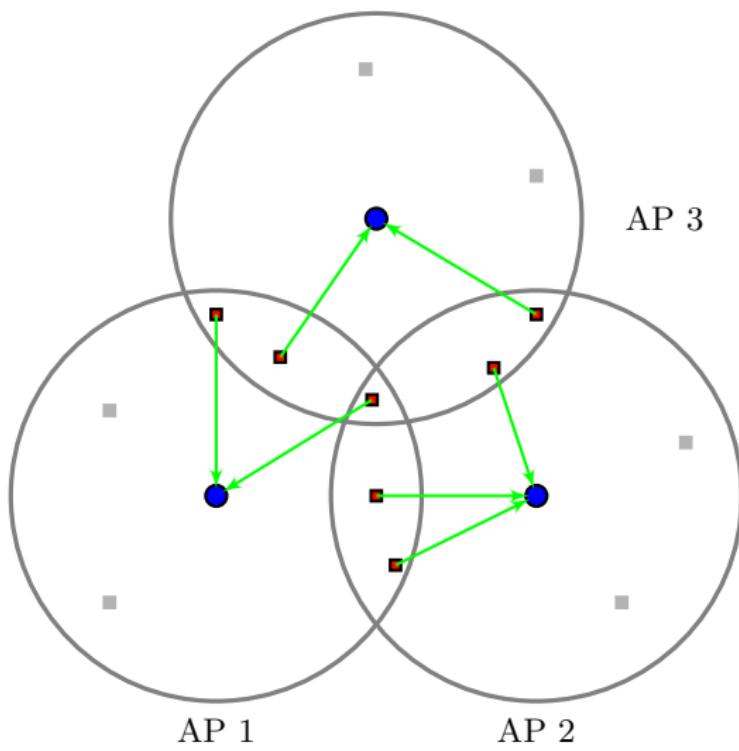
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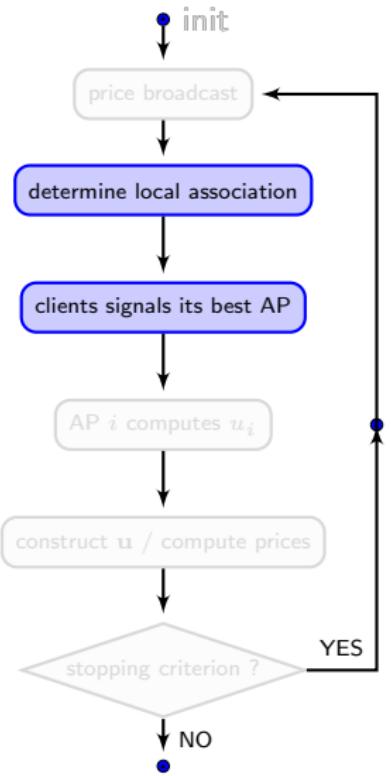
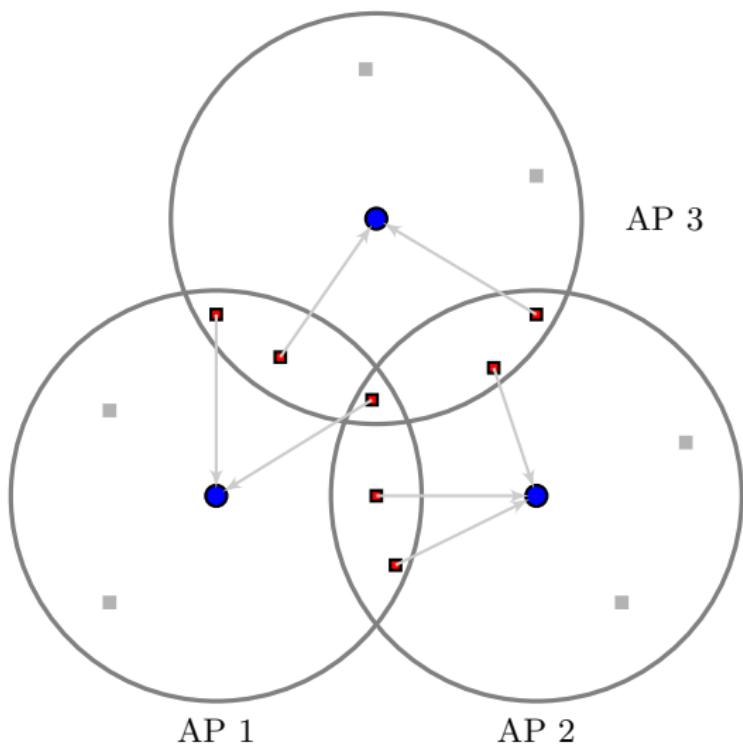
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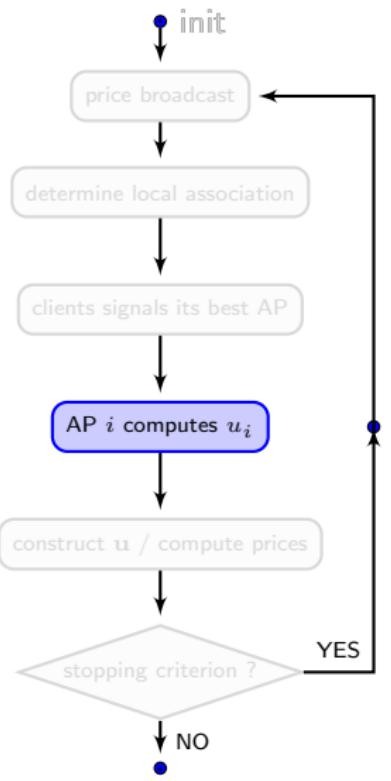
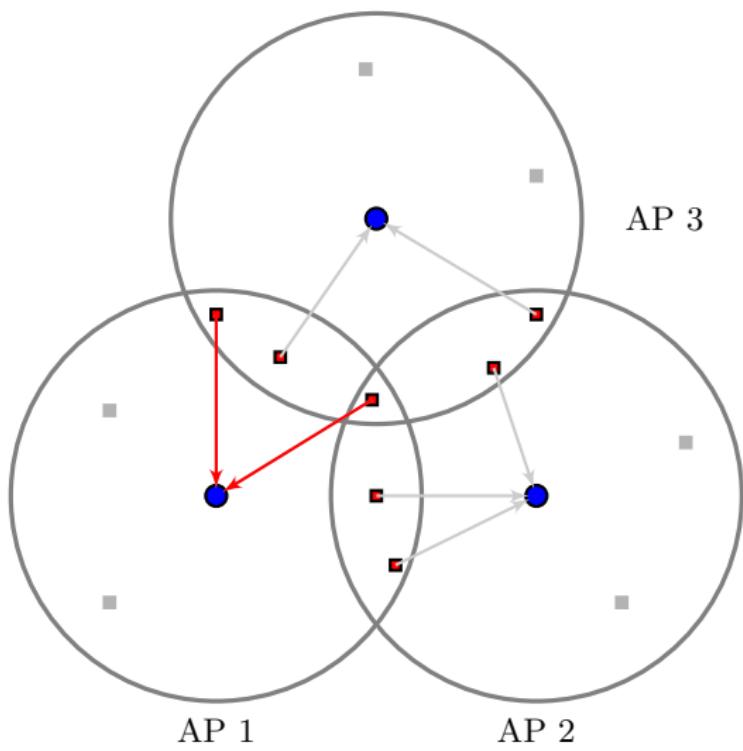
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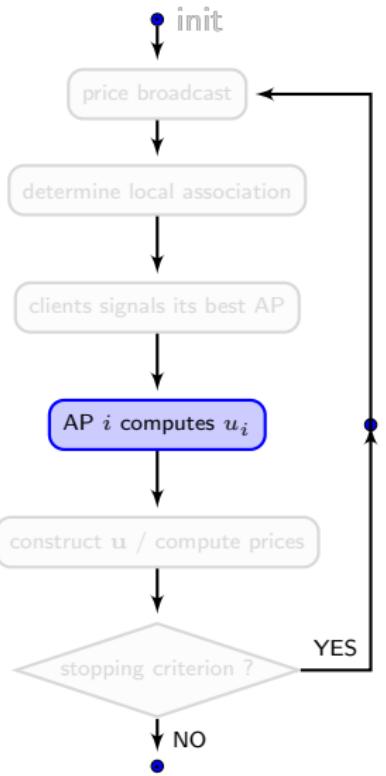
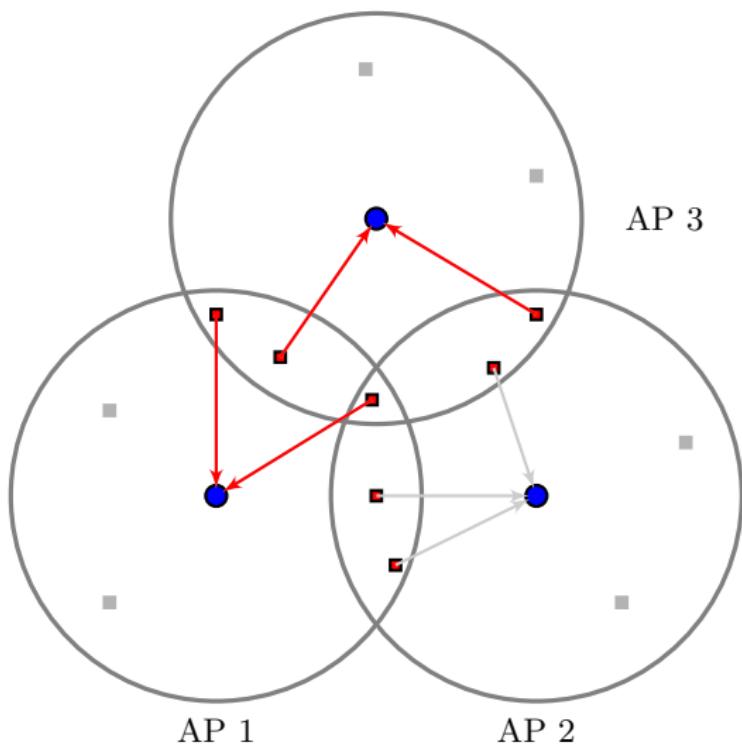
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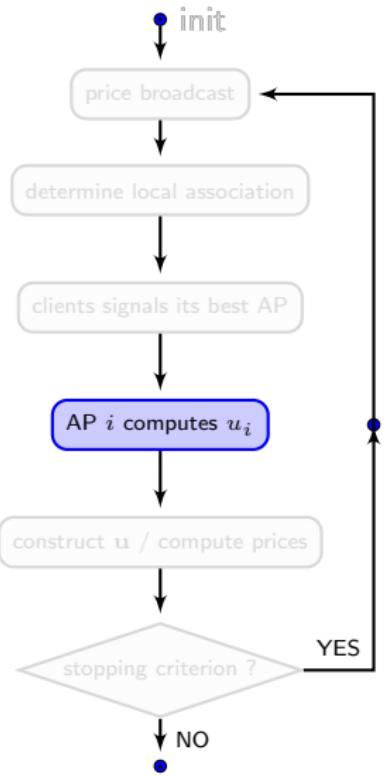
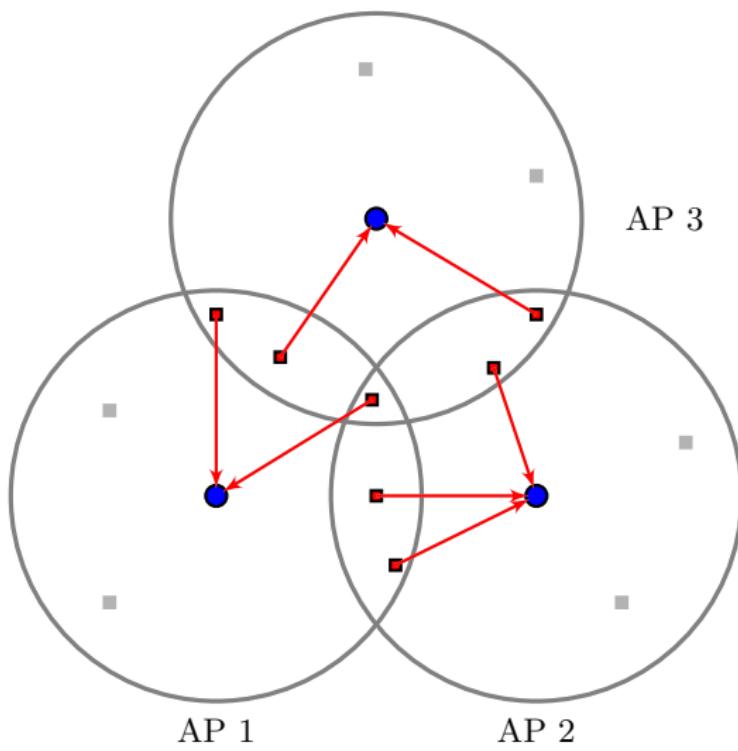
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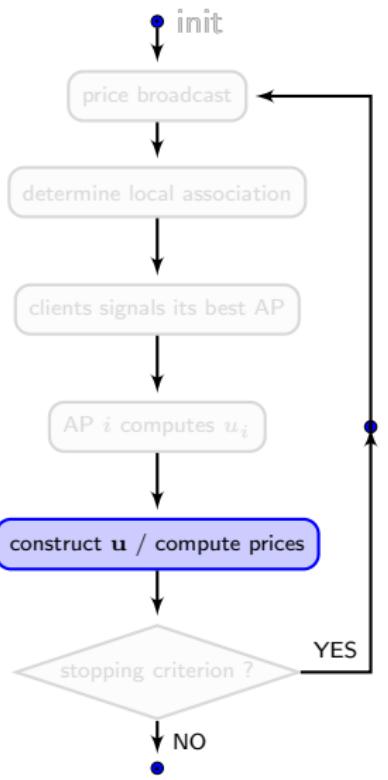
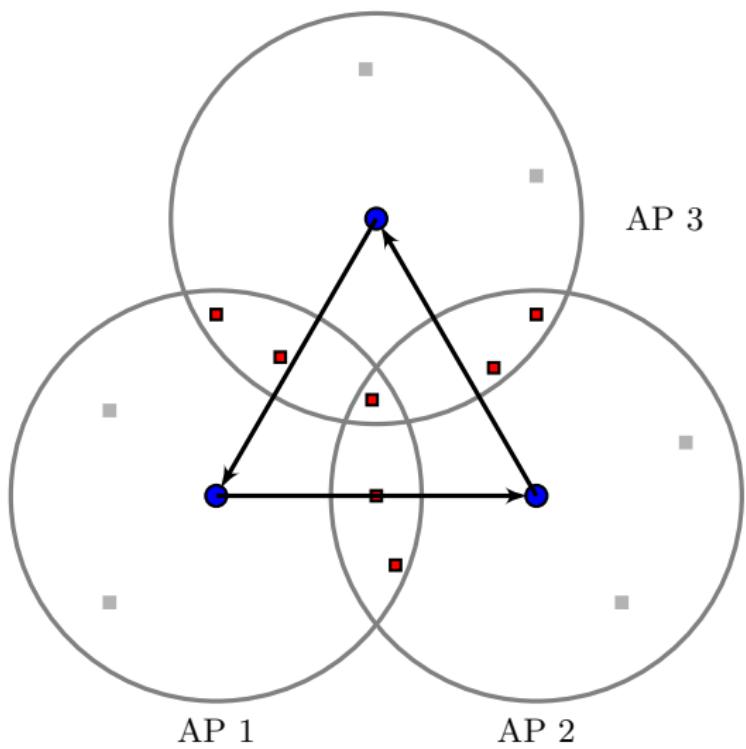
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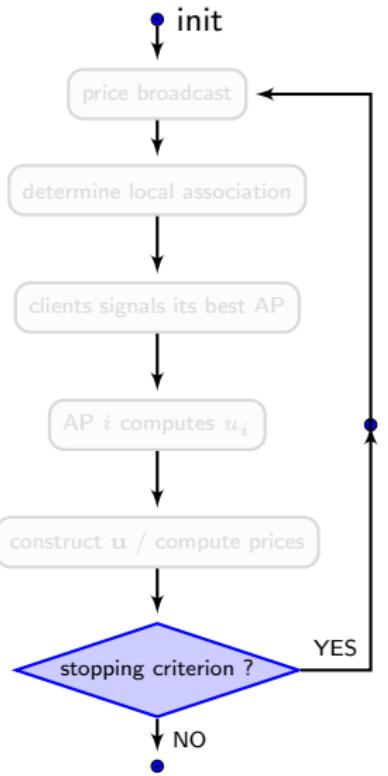
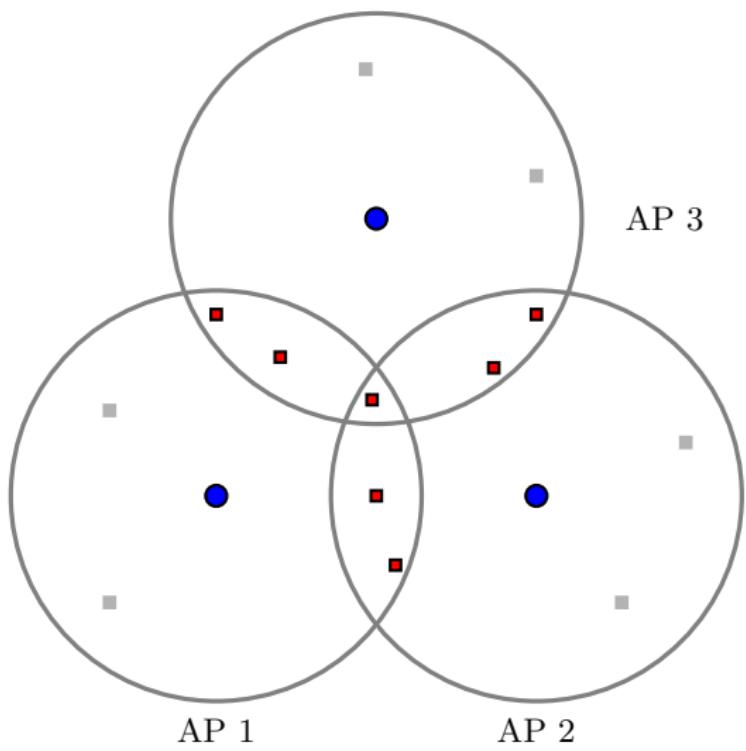
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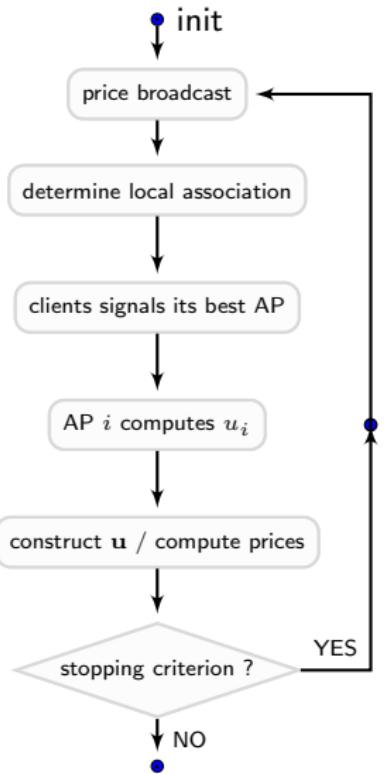
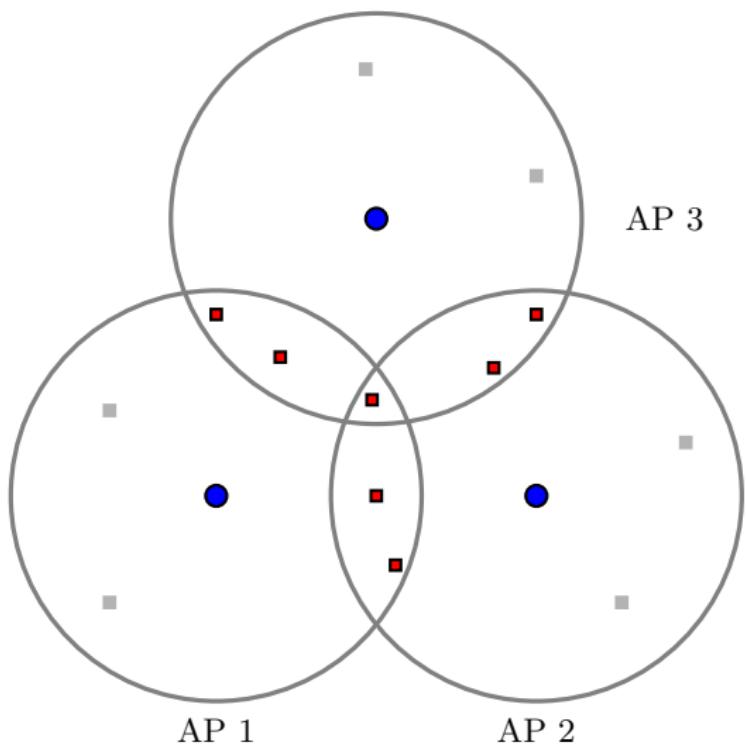
DAA Illustration



DAA Illustration



DAA Illustration



DAA Properties

Proposition (Convergence)

Let $g_{\text{best}}^{(k)}$ denote the best dual objective value found after k subgradient iterations, i.e., $g_{\text{best}}^{(k)} = \max\{g(\boldsymbol{\lambda}^{(1)}), \dots, g(\boldsymbol{\lambda}^{(k)})\}$. Then, $\forall \epsilon > 0 \exists n \geq 1$ such that $\forall k \geq n \Rightarrow (d^* - g_{\text{best}}^{(k)}) < \epsilon$.

Proof.

Standard convergence proof for the projected subgradient method [Boy07, § 3.2][Ber99] □

DAA Properties

Theorem (Constant Duality Gap and Relative Duality Gap)

The optimal duality gap of the mixed integer linear program is bounded as follows:

$$p^* - d^* \leq (N + 1)(\varrho + \max_{j \in \mathcal{M}} \varrho_j) ,$$

where $\varrho = \max_{i \in \mathcal{N}, j \in \mathcal{M}_i} \beta_{ij}$ and $\varrho_j = \min_{i \in \mathcal{N}_j} \beta_{ij}$. Moreover, the relative duality gap $(p^* - d^*)/p^*$ diminishes to 0 as $M \rightarrow \infty$.

Proof.

Here we capitalized on

- the equivalent problem reformulations
- the separability of the problem
- problem's geometry

The essential guide lines are provided in [Ber98, § 5.6.1]



DAA Properties

minimize $\sum_{j \in \mathcal{M}} f_j(\mathbf{y}_j)$
subject to $\mathbf{y}_j \in \mathcal{Y}_j, j \in \mathcal{M}$
 $\sum_{j \in \mathcal{M}} \mathbf{h}_j(\mathbf{y}_j) - \mathbf{b} \leq \mathbf{0},$

variables: \mathbf{y}_j .

DAA Properties

$$\sum_{j \in \mathcal{M}} f_j(\mathbf{y}_j)$$

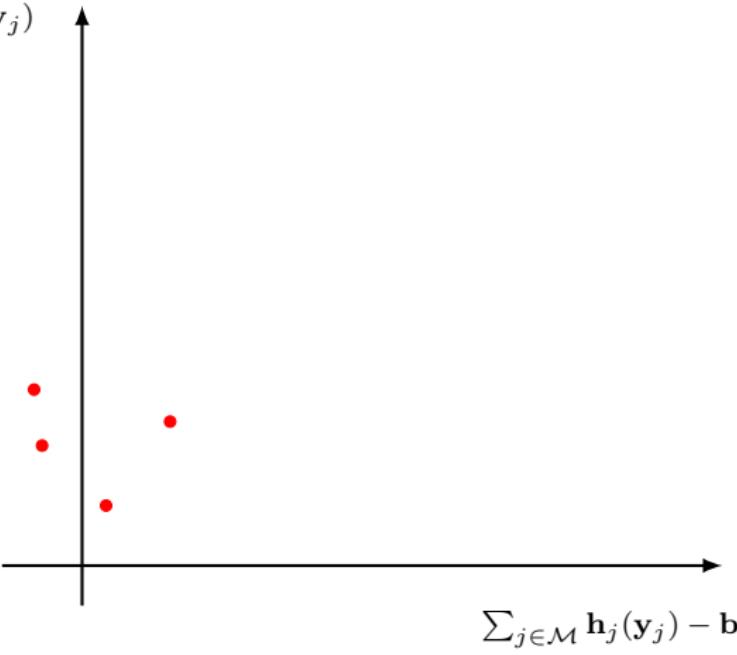
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e.g. 1: $\mathcal{Y}_j = \{0, 1\}, \mathcal{M} = \{1, 2\}$



DAA Properties

$$\sum_{j \in \mathcal{M}} f_j(\mathbf{y}_j)$$

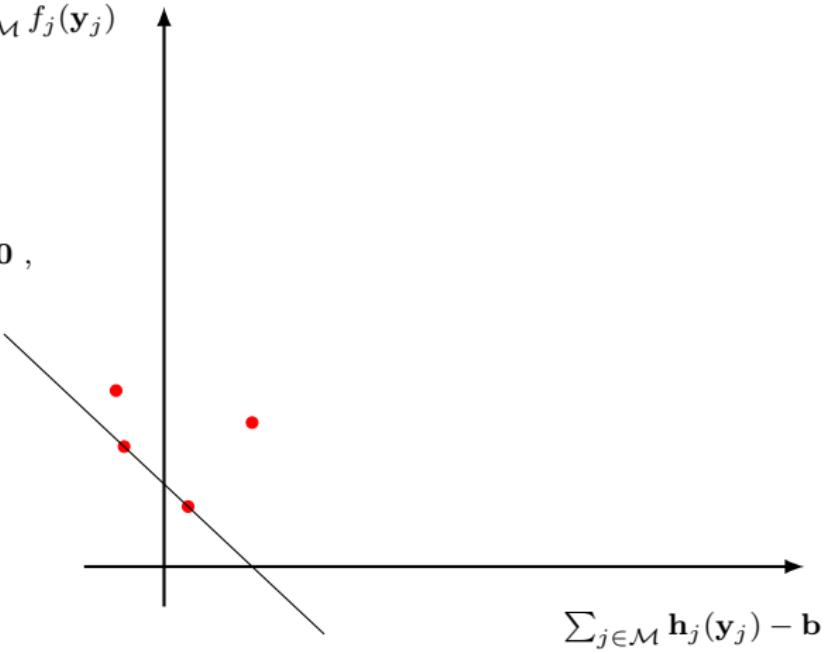
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DAA Properties

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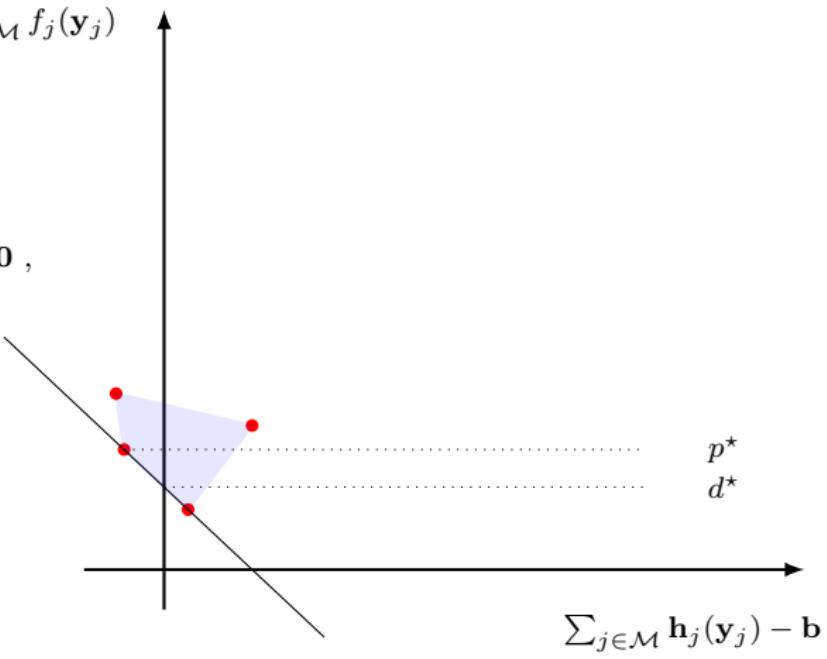
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DAA Properties

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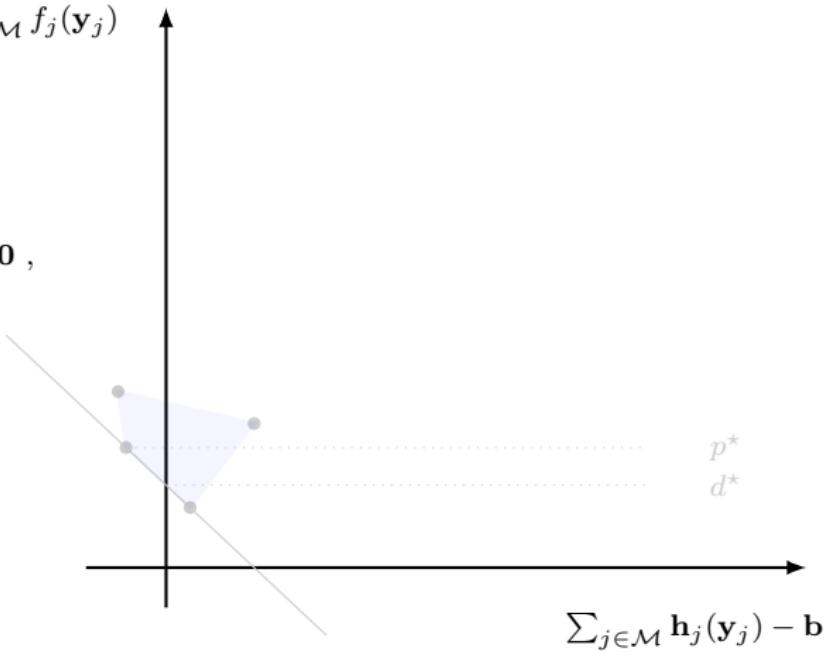
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subject to $\mathbf{y}_j \in \mathcal{Y}_j, j \in \mathcal{M}$

$$\sum_{j \in \mathcal{M}} \mathbf{h}_j(\mathbf{y}_j) - \mathbf{b} \leq \mathbf{0},$$

variables: \mathbf{y}_j .

e.g. 1: $\mathcal{Y}_j = \{0, 1\}, \mathcal{M} = \{1, 2\}$



DAA Properties

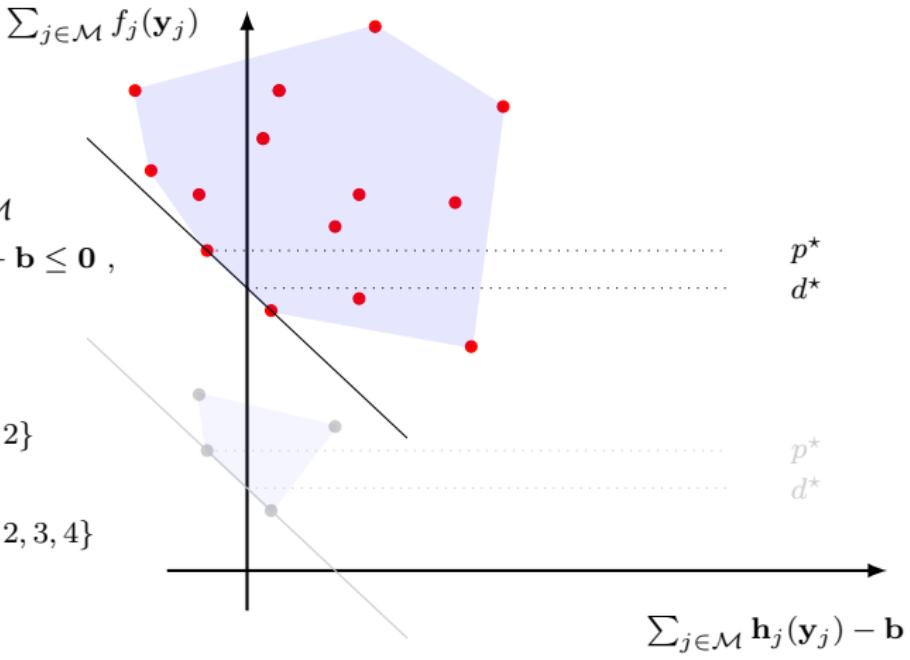
minimize $\sum_{j \in \mathcal{M}} f_j(\mathbf{y}_j)$
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e.g. 2: $\mathcal{Y}_j = \{0, 1\}, \mathcal{M} = \{1, 2, 3, 4\}$



DAA Properties

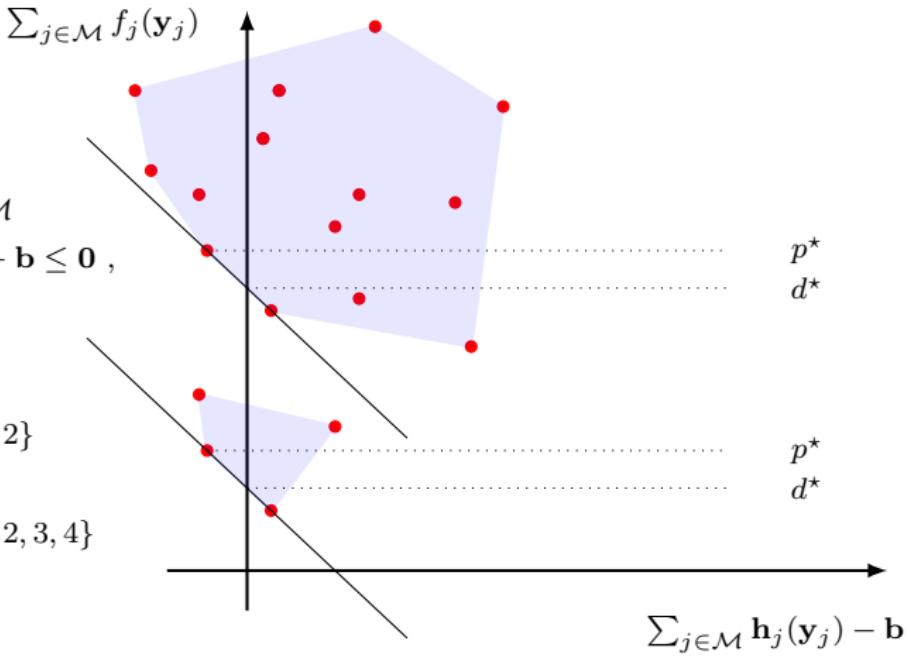
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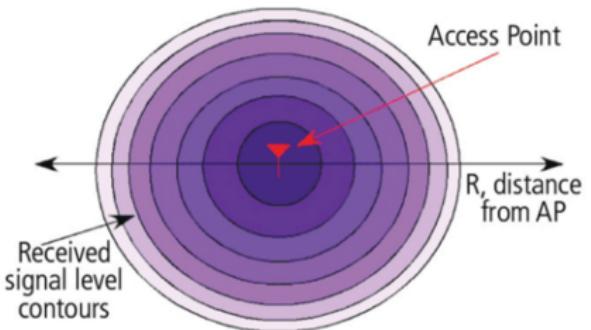
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e.g. 2: $\mathcal{Y}_j = \{0, 1\}, \mathcal{M} = \{1, 2, 3, 4\}$

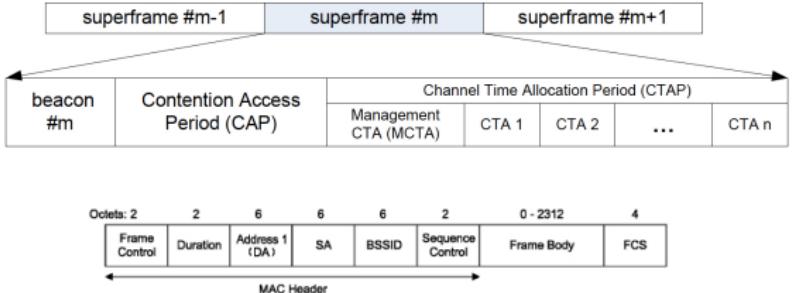


Implementation Over Existing Standards



- Initially the clients follow the RSSI-based association policy that IEEE 802.11ad and IEEE 802.15.3c define
- DAA** is periodically executed to correct possible suboptimal client associations by reallocating the available resources

Implementation Over Existing Standards



- APs trigger the initialization of **DAA** by setting a special bit into the beacon frame
- Information exchange is performed through the control frames or piggy-backing the data frames

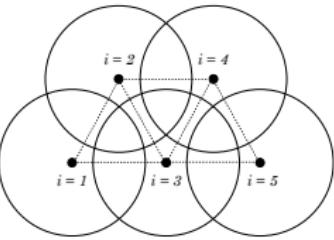
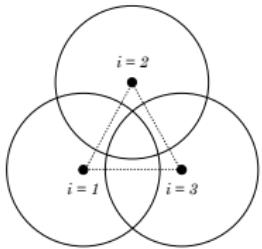
Outline

- Past, Present and Future in wireless communications
- 60 GHz millimeterWave wireless technology
- Optimizing resource allocation
- Distributed client association (DAA)
- **Numerical analysis of DAA**
- Conclusions and open research topics

Numerical Analysis

- Consider a multi-user multi-cell environment
- Compare DAA to
 - Random association
 - RSSI-based association (IEEE 802.11)
 - Optimal association (IBM CPLEX)
- Measure
 - Convergence
 - Scalability
 - Efficiency (time)
 - Fairness

Topologies



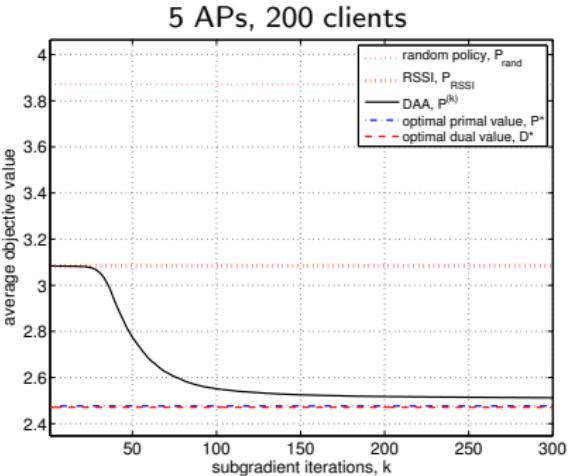
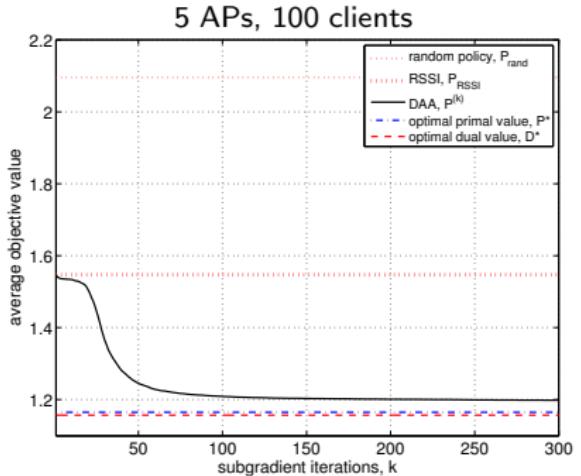
λ	Wavelength
d_0	Far field reference distance
η	Path loss exponent

- SNR operating point at a distance d from any AP

$$\text{SNR}(d) = \begin{cases} P_0 \lambda^2 / (16\pi^2 N_0 W) & d \leq d_0 \\ P_0 \lambda^2 / (16\pi^2 N_0 W) \cdot (d/d_0)^{-\eta} & \text{otherwise} \end{cases}$$

- Radius of each cell r is chosen such that $\text{SNR}(r) = 10$ dB
- Clients are uniformly distributed at random, among the circular cells

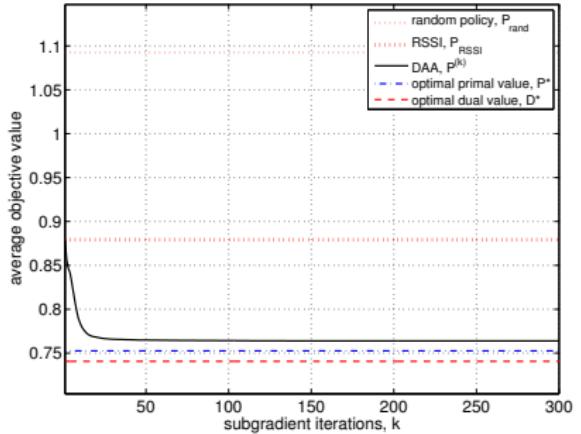
Convergence of DAA



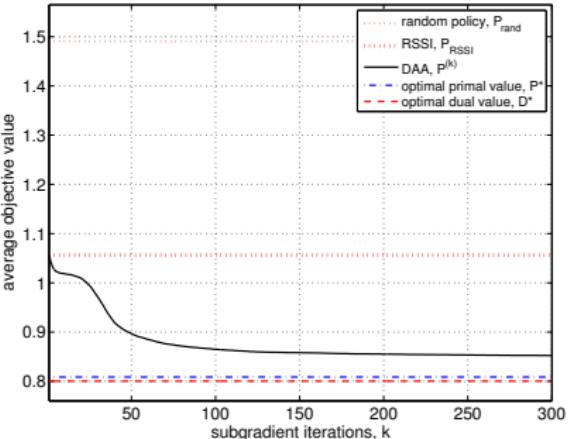
- Average primal objective value from **DAA** after k subgradient iterations: $P^{(k)} = (1/\bar{T}) \sum_{T=1}^{\bar{T}} p_{\text{best}}^{(k)}(T)$
- Average dual optimal value by **DAA**: $D^* = (1/\bar{T}) \sum_{T=1}^{\bar{T}} d^*(T)$

Convergence of DAA

3 APs, 30 clients

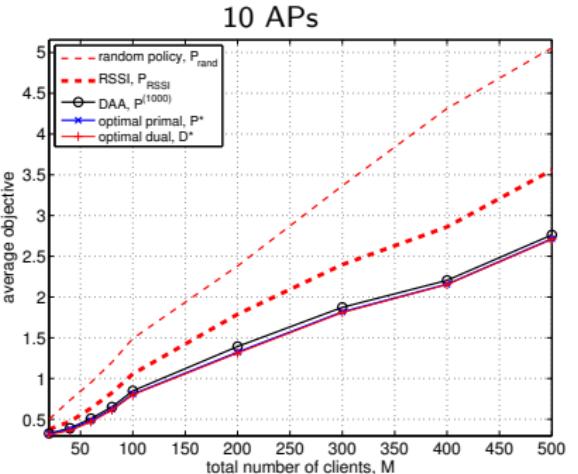
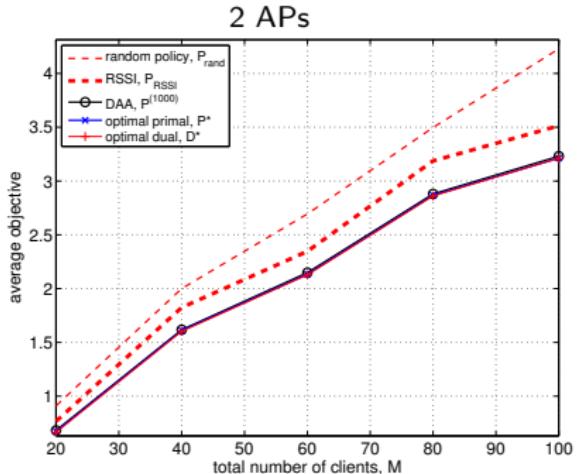


10 APs, 100 clients



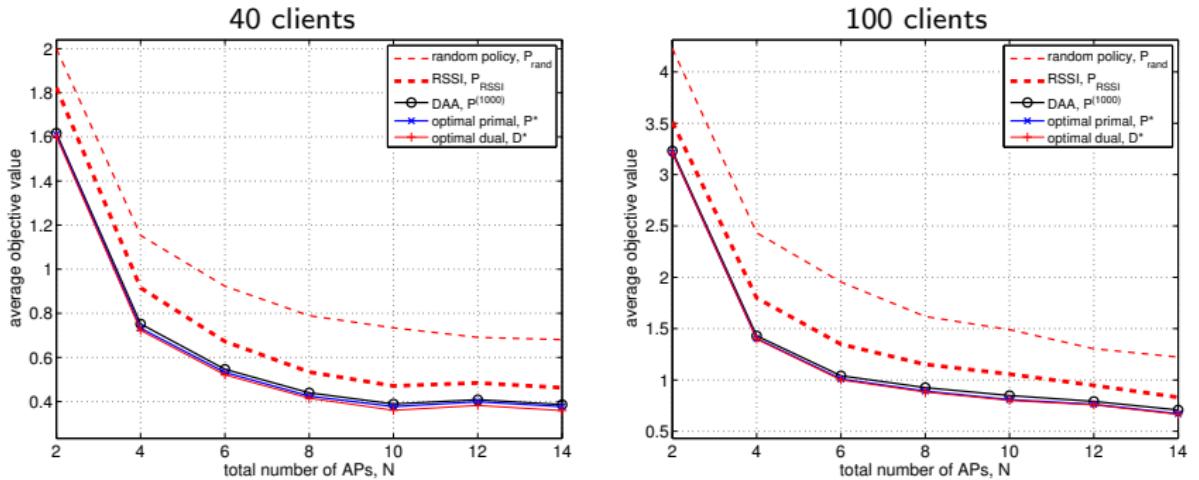
- Convergence time is affected by the number of APs and clients: The smaller the network, the faster **DAA** converges

Scalability of DAA



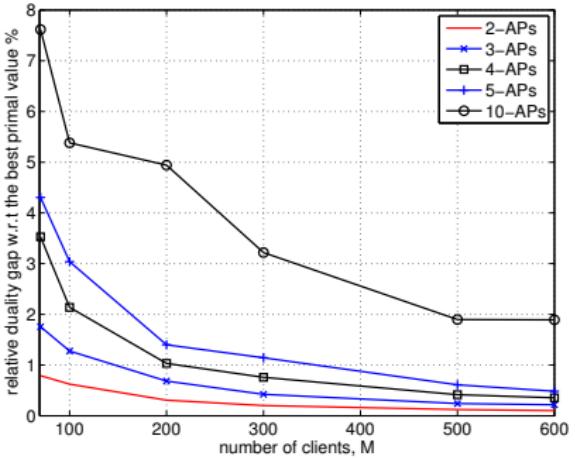
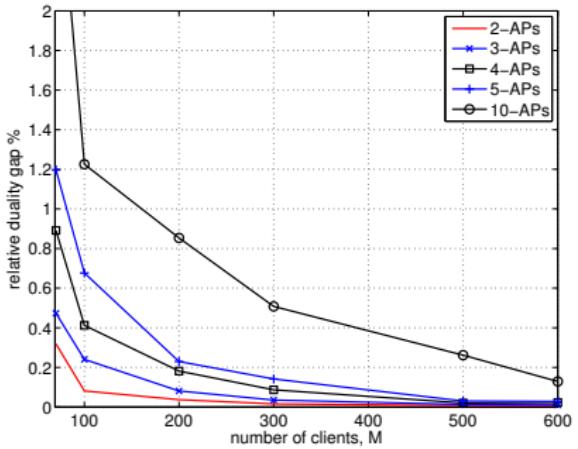
- **DAA** performs close to optimal

Scalability of DAA



- Considering constant load, the average objective value decreases while the number of APs increases
- **DAA** performs close to optimal

Optimality of DAA

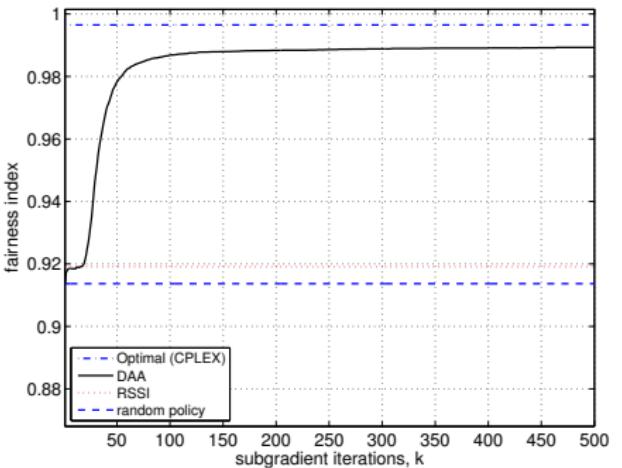


- Average relative duality gap:

$$\text{Ave-RDG} = (1/\bar{T}) \sum_{T=1}^{\bar{T}} (p^*(T) - d^*(T)) / p^*(T)$$

- Average relative duality gap taking into account the best primal feasible objective value from **DAA** after K iterations at time slot T :
- $$\text{Ave-RDG-best-achieved} = (1/\bar{T}) \sum_{T=1}^{\bar{T}} (p_{\text{best}}^{(K)}(T) - d^*(T)) / p_{\text{best}}^{(K)}(T)$$

Fairness Achieved by DAA

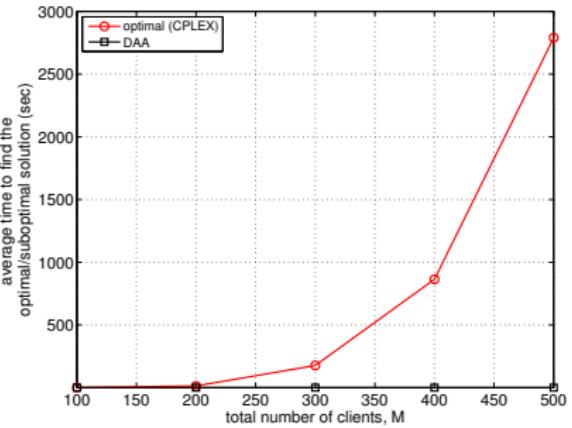
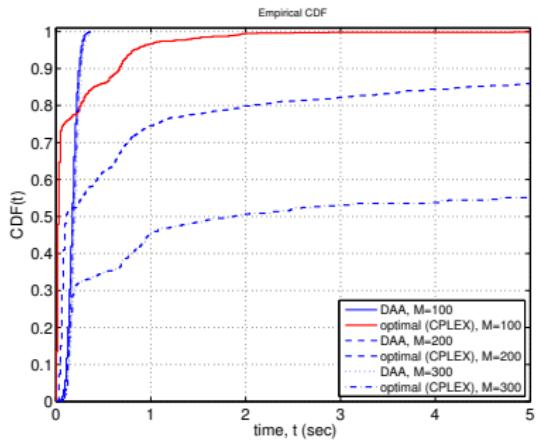


- Jain's fairness index:

$$J^{(k)}(T) = \left(\sum_{i \in \mathcal{N}} Y_i^{(k)}(T) \right)^2 / \left(N \sum_{i \in \mathcal{N}} Y_i^{(k)}(T)^2 \right),$$

$Y_i^{(k)}(T) = \sum_{j \in \mathcal{M}_i} \beta_{ij} x_{ij}^{(k)}(T)$ and $x_{ij}^{(k)}(T)$ is the solution (best feasible) resulted from **DAA** at time slot T after k iterations

Speed and Resources Used by DAA



- Empirical CDF plots of the number of iterations for $M = 100, 200, 300$ clients, with $N = 10$ APs
- Trade-off between optimality and complexity

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Conclusions

- 60 GHz wireless technology: characteristics, benefits, challenges, applications
- Distributed association algorithm (DAA) for optimizing resource allocation in 60 GHz wireless access networks
- Performance evaluation of **DAA**: Asymptotically optimal, convergence, time efficiency and fairness
- Integration of **DAA** into current standards

Open Research Topics

- 60 GHz channel modeling
- Medium Access Control (MAC)
- Connectivity maintenance, blockage and directivity
- Coexistence and cooperation with existing wireless technologies
- Multi-hop communications
- ...

Thank you



OPTIMIZING CLIENT ASSOCIATION IN 60 GHz WIRELESS ACCESS NETWORKS

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