

# Network Security - Fall 2018

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Lecture 3 - Term Project

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## Make up Classes

- ▶ Will have 3 missed classes in September (AFRL Reporting Talks)
- ▶ First one Sep 4, 2018 (Make up class: Friday Sep 7, 2018)
- ▶ Second one Sep 18, 2018 (Makeup class: Students decide)
- ▶ Third one Sep 20, 2018 (Makeup class: Students decide)

# Term Project Discussion

- ▶ The Linux VM Image will be posted **TODAY** after class
- ▶ Instructor will share a DropBox link for download
- ▶ Distribution is based on **Ubuntu**
- ▶ Distribution has a lot of vulnerabilities
- ▶ Goal: To find the vulnerabilities (easy part) and fix them (hard part)
- ▶ Hack-a-thon at the end of semester **BEFORE** Thanksgiving

# Term Project Discussion

- ▶ Root Access: msfadmin\msfadmin
- ▶ Use a **Virtual Machine** (Virtualbox or VMware preferred)
- ▶ **Do not get owned by unknown actors!** Please!
- ▶ **Networking:** Supports *only*:
  1. NAT
  2. Host Only
- ▶ Vulnerabilities: Many are documented
- ▶ However some vulnerabilities are not documented
- ▶ That is the fun part

# Term Project Discussion: Attack System

- ▶ Use a virtual machine for this as well
- ▶ Best options: Kali Linux (Debian) or Backbox Linux
- ▶ Many other options: Parrot Security OS etc.
- ▶ Many are hard to install on Windows 10 Systems (UEFI)
- ▶ Should you use any one of these systems?
- ▶ Why or Why Not?
- ▶ HIDE, IMPERSONATE, FOOL TARGET

# Term Project Discussion: Target System

- ▶ Active Security (take active measures)
  - ▶ Cover your back (look for basic things)
  - ▶ Have policies that **actively** protect system
  - ▶ AUDIT AUDIT AUDIT
  - ▶ Have a a designated sysadmin (to blame)
  - ▶ Port scans: Should you allow them?
  - ▶ Should Kali or Backbox distros raise a flag?
  - ▶ Protect root
- ▶ Make sure that you wear both hats as a team
- ▶ Good Luck! ENJOY!

# Diffie-Hellman Key Exchange

- ▶ Client and server agree on a large prime (shared prime  $p$ )
- ▶ Both of them agree on an encryption generation method (AES)
- ▶ Client comes up prime number  $p_1 \neq p$
- ▶ Server comes up with prime number  $p_2 \neq p$
- ▶  $p_1$  and  $p_2$  are not shared and used as private keys
- ▶ Private key ( $p_1$  or  $p_2$ ) and  $p$  used with AES
  - ▶ Client: Generate public key  $p_{k_1}$
  - ▶ Server: Generate public key  $p_{k_2}$
  - ▶ Exchange generated public keys
- ▶ Server uses  $p_{k_1}$ ,  $p$  and  $p_2$  to generate encryption key
- ▶ Client does the same but with  $p_{k_2}$ ,  $p$  and  $p_1$

## What is the Magic

We will study in detail going forward

# Diffie-Hellman Key Exchange

- ▶ Client and server agree on a large prime (shared prime  $p$ )
- ▶ Good to ensure  $\frac{p-1}{2}$  is also prime
- ▶ Client and server *also* agree on another shared key  $g$
- ▶ Client comes up prime number  $p_1 \neq p$
- ▶ Client computes  $p_{k_1} = g^{p_1} \bmod p$
- ▶ Server comes up with prime number  $p_2 \neq p$
- ▶ Server computes  $p_{k_2} = g^{p_2} \bmod p$
- ▶ Exchange  $p_{k_1}$  and  $p_{k_2}$
- ▶  $p_1$  and  $p_2$  are not shared and used as private keys
- ▶ Private key ( $p_1$  or  $p_2$ ) and  $p$  used with AES
- ▶ Client shared key:  $p_{k_1}^{p_2} \bmod p = g^{p_1 p_2} \bmod p$
- ▶ Server shared key:  $p_{k_2}^{p_1} \bmod p = g^{p_1 p_2} \bmod p$
- ▶ Choose  $g$  such that  $g^{(p-1)} = 1 \bmod p$



# Diffie-Hellman Key Exchange

## Question

Is it possible to launch a man-in-the-middle attack on the DH Key Exchange