Fundamentos de los Sistemas Operativos (FSO)

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Part 2: Process Management

Seminar 3
UNIX system calls for processes





Goals:

- To know the services provided by UNIX to create processes
- To analyse C language examples with process related system calls
- To know the signal concept
- Bibliografía:
 - "UNIX System Programming", Kay A. Robbins, Steven
 Robbins, Prentice Hall, ISBN 0-13-042411-0

- Identification
- Creation
- Waiting
- Ending
- Signals

getppid

UNIX process related system calls Processes fork Create a child process exit End process wait Wait for a process ending exec Execute a program getpid Get process ID attribute

Get parent ID attribute

	Signals
kill	Send signals
alarm	Generate an alarm (clock signal)
sigemptyset	Init a mask with no signals set
sigfillset	Init a mask with signals set
sigaddset	Append a signal to a signal set
sigdelset	Delete a signals in a signal set
sigismember	Check if a signals belongs to a signal set
sigprocmask	Check/Modify/Set a signal mask
sigaction	Capture/Manage a signal
sigsuspend	Wait signals capture

Identification

- Creation
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- Every process must have an ID
- The creator process is the parent while the created process is the child. To know them we use:
 - Process ID (PID)
 - Parent process ID (PPID)

```
pid_t getpid(void);
```

pid_t getppid(void);

```
/*** ej1 getpid.c *****/
                         #include <stdio.h>
                         int main(void)
                            printf("\nProcess ID: %ld\n", (long)getpid());
                            printf("Parent process ID: %ld\n", (long)getppid());
$ gcc -o ej1 ej1_getpid.c
 ./ej1 &
                            while(1);
[1] 2959
                            return 0;
Process ID: 2959
Parent process ID: 1060
$ ps -1
                                           WCHAN S ADDR
 UID PID
           PPID F
                     CPU PRI NI SZ
                                       RSS
                                                                   TTY
                                                                           TIME
                                                                                  CMD
                               2435548 1088 -
                                                 S ffffff80136e3d50 ttys000 0:00.06 bash
     1060 1059 4006 0
     2959
           1060 4006 0
                                                 R ffffff80140d8300 ttys000 0:04.65 ./a.out
                            0 2434832
```

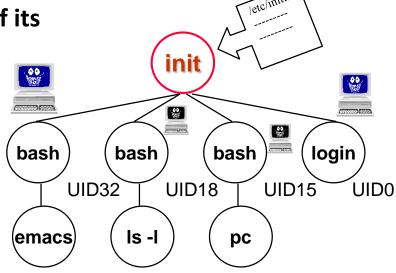
- Identification
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 Unix uses a copying mechanism for process creation pid_t fork(void);

The child process is an exact replica of its parent process

The child process inherits most of the attributes from its parent:

- Memory image
- UID, GID
- Current directory
- Opened file descriptors
- Unix assigns an identifier to every process
 called PID at the time of its creation
- Every process knows the identifier of its parent process, PPID
- The child execution is concurrent with the parent and independent
- Processes are organized into a hierarchical tree



fork: process creation

```
#include <sys/types.h>
#include <unistd.h>
pid_t fork(void)
```

Description

- It creates a child process which is a "clon" from the parent: the child process inherits the most part of the parent attributes.
- NON Inherited attributes: PID, PPID, pendant signals, file locks and time/accounting.

Returning value

- 0 to the child
- Child PID to the parent
- -1 to the parent if an error happens

— Errors

Not enough resources to create the process

```
Parent and child resume execution with the instruction following fork()
    /**ej2 fork.c **/
     #include <stdio.h>
     #include <sys/types.h>
     int main(void)
     f printf("Process %ld creates another process\n", (long)getpid());
       fork();
       printf("Process %ld with parent %ld\n",(long)getpid(),(long)getppid());
       sleep(5);
       return 0;
                                                      bash
        $ ps
                                                                       pid=1060
          PID TTY
                             TIME CMD
                                                    /*ej2*/
         1060 ttys000 0:00.07 -bash
        $ gcc -o ej2 ej2 fork.c
                                                                       ppid=1060
                                                                parent
                                                                       pid= 3242
         • ./ej2
         rocess 3242 creates another process
                                                    /*ej2*/
                                                                       ppid=3242
         Process 3242 with parent 1060
                                                                child
                                                                       pid= 3243
       rocess 3243 with parant 3242
                                                    fork()
                                                                             Pág. 10
```

Parent and child execute diferent code

```
// ej3_fork.c
#include <sys/types.h>
#include <unistd.h>
int main() {
   pid_t val;
   int var = 0;
   printf("PID before fork(): %d\n", (long) getpid());
  val = fork();
   if (val > 0) {
    printf("Parent PID: %d\n", (long) getpid());
    var++;
   } else {
    printf("Child PID: %d\n", (long) getpid());
   printf("Process [%d]-> var=%d\n", (long) getpid(), var);
   return 0;
```

How many processes print this message?

What value(s) of "var" are shown?

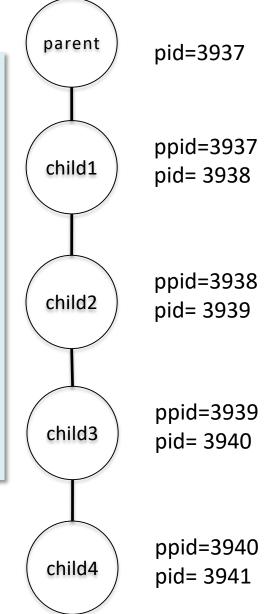
Example

```
// ej4 fork.c
#include <stdio.h>
#include <sys/types.h>
#include <unistd.h>
int main(void)
    pid_t pid=fork();
    switch (pid) {
    case -1:
         printf("Could not create child process \n");
         break;
    case 0:
         printf("I am the child with PID %ld and my parent is %ld \n",
                   (long)getpid(), (long)getppid());
         break;
    default:
         printf("I am the parent with PID %ld and my child is %ld \n",
                   (long)getpid(), pid);
                       $gcc-o ej3 ej3 fork.c
    sleep(5);
                       $ ./ej3
    return 0;
                       I am the parent with PID 3702 and my child is 3704
                       I am the child with PID 3704 and my parent is 3702
```

Creating processes in chain

```
// ej5 proc chain.c
#include <stdio.h>
                                            Variants
#include <sys/types.h>
                                                 if (pid > 0)
#define NPROCESSES 4
                                                  if (pid==0)
                                                 if (pid<0)
int main(void)
     pid_t pid;
     int i;
     for (i=0; i<NPROCESSES; i++) {</pre>
          pid=fork();
          if (pid!=0)
               break;
          printf("I am the child with PID %ld
                   my parent is %ld\n",
                   (long)getpid(), (long)getppid());
     sleep(5);
     return 0;
```

```
$ ./a.out
I am the child with PID 3938 my parent is 3937
I am the child with PID 3939 my parent is 3938
I am the child with PID 3940 my parent is 3939
I am the child with PID 3941 my parent is 3940
```



exec()

- The call to fork () creates a child who is a copy of the calling process
- What if we want to run a different program? Use exec() call
- There are different versions exec depending on the parameters specified:

- Variant I: arguments are provided separately
- Variant v: arguments are provided through a pointer to a vector
- Variant p: the location of file is searched on the PATH
- Variant e: the environment is provided to the child through envp, it is not inherited from the parent

exec () features

- It changes the process memory image by the one defined in the new executable file.
- The executable file in specified by its name or its absolute path.
- Some process attributes remain:
 - Signal handlers, exept the ones captured and treated by the default action.
 - PID and PPID
 - Time accounting
 - File descriptors
 - Working directory, root directory and file creation mode mask
- If the SETUID bit of the executable file is set then EXEC sets as efective UID the UID of the executable file owner
 - The same happens with the SETGID bit
- Errors
 - Executable file doesn't exist or it is not recognized as executable
 - Permissions
 - Incorrect arguments
 - Not enough meory or resources
- Returning value
 - If EXEC returns to the calling program it is because an error happened, the returning value is -1

Exec() example: the child process lists the current directory

```
// ej6_exec.c
#include <stdio.h>
                                    Another way:
#include <sys/types.h>
                                     execl("/bin/ls", "ls","-l",NULL)
int main(void)
     int status;
    pid_t pid=fork();
    char* arguments [] = { "ls", "-l", 0 };
     switch (pid) {
     case -1:
         printf("The child process could not be created \n");
          break;
     case 0:
         printf("I am the child with PID %ld, the current directory content is: \n",
                  (long)getpid());
          if (execvp("ls", arguments) == -1) {
              printf("Error in exec\n");
              exit(0);
          break;
    default:
         printf("I am the parent process with PID %ld and my child is %d.\n",
                  (long)getpid(), pid)
    return 0;
```

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The parent has to wait for its child to finish:

```
#include <sys/types.h>
#include <sys/wait.h>
pid_t wait(int *status);
pid_t waitpid(pid_t pid, int *status, int options);
```

 wait and waitpid stop the calling process execution until a child ends, or until the calling process receives a signal wait() call

```
pid_t wait(int *status);
```

- It suspends the execution of the calling process until any of its children finishes
 - If there is any zombie child then wait returns immediately
- status: is not a NULL pointer, it contains information about child termination:
 - Child finished with exit or child finished by a signal



- Returning value:
 - -1 if error or no children
 - Otherwise the finished child PID

Wait() example:

```
// ej7 wait.c
#include <stdio.h>
#include <sys/types.h>
#include <errno.h>
int main(void)
     int status;
     pid_t pid=fork();
     switch (pid) {
     case -1:
          printf(" The child process could not be created \n");
          break;
     case 0:
          printf("I am the child process with PID %ld and my parent is %ld\n",
                  (long)getpid(), (long)getppid());
          sleep(20);
          printf("I have finished \n");
          break;
     default:
          printf("I am the parent process with PID %ld and my child is %d.
                  Waiting ...\n", (long)getpid(), pid);
          if (wait(&status) != -1)
               printf("My child has ended ok \n");
     return 0;
```

waitpid() call

```
pid_t waitpid(pid_t pid, int *status, int options);
```

- It waits for a particular child
- Parameters:
 - **pid**: PID of the child to wait for. If its value is -1 it waits until the first child ends (like **wait**)
 - status: like wait()
 - options:
 - Value 0: blocking (common value)
 - WNOHANG: non blocking
 - Returning value:
 - 0 : no child ended (non-blocking version),
 - -1 : error
 - > 0 : pid of the returning child

waitpid() example: fan process creation

```
// ej8_waitpid.c
#include <stdio.h>
#include <sys/types.h>
                                                                 parent
#define NPROCESSES 4
int main(void)
    pid_t pid[NPROCESSES];
    int i, status;
                                                     child1
                                                             child2
                                                                      child3
                                                                              child4
    for (i=0; i<NPROCESSES; i++) {</pre>
        pid[i]=fork();
         if (pid[i]==0){
             printf("I am the child %ld my parent is %ld\n",
                        (long)getpid(), (long)getppid());
             sleep(10);
             exit(0);
                                                   $gcc -o ej8 ej8 waitpid.c
    // Now wait for the third child
    if (waitpid(pid[2],&status,0)==pid[2])
        printf(" My third child has finished \n");
    return 0;
```

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- A process ends completely when:
 - The process itself ends (normally or abnormally) AND
 - Its parent process does a wait call
- The normal termination is done by calling to exit

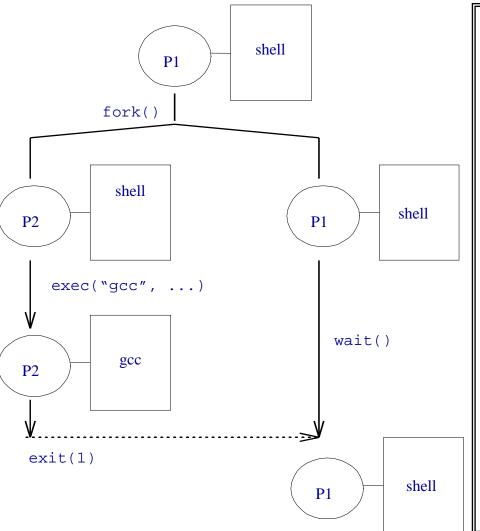
```
void exit (int status)
```

- The status parameter is used to communicate to the parent how the child process ends
- By convention, this value is typically 0 if the process ends normally and any other value means abnormal termination
- The parent process can obtain this value through the system call wait

Abnormal ending:

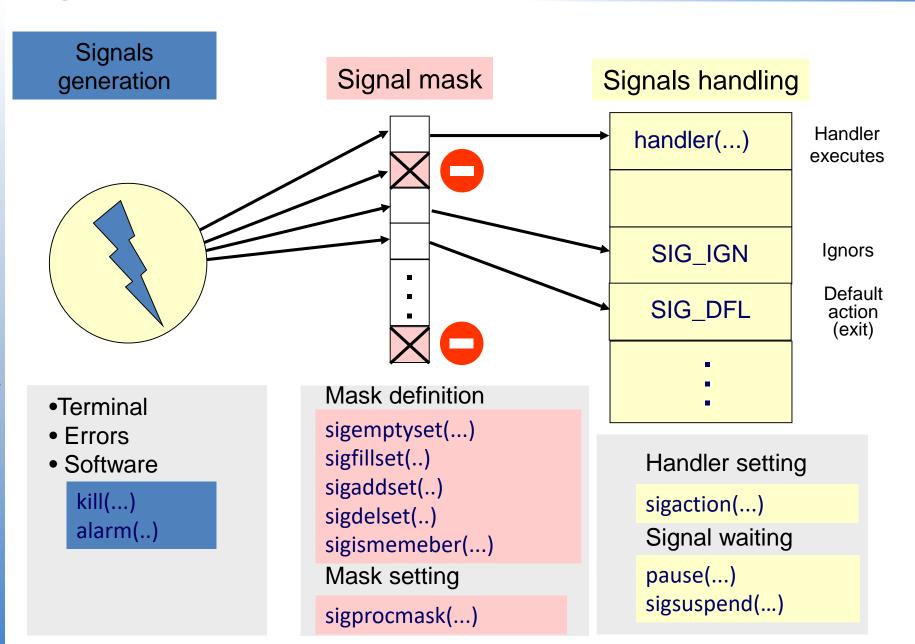
- The process is ended by the operating system because of an error condition (memory boundary violation, arithmetic errors) or by some other process initiative
 - Signals
- Zombie: If the process ends before his parent calls wait():
- **Orphan**: If the parent process terminates before the child:
 - An orphan process is adopted by process root init

Simplified structure of UNIX shell



```
while(TRUE) {
 print_prompt();
 read_command(command, parameters);
 p=fork();
                  /* creates a child */
 if (p != 0) { /* parent code */
  waitpid(-1, &status, 0); /* waiting for
                                the child */
   else {
             /* child code */
  exec(command, params, 0); /* memory image
                               change */
  error("Command cannot be executed");
  exit(1);
```

- Identification
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- A signal is the mechanism used by the OS to inform processes about certain events
- They all follow the same pattern:
 - It is generated due to the occurrence of an event
 - It is supplied to the process
 - It should be treated by the default handling routine or by the one defined by a specific process
- A signal can:
 - Be handled: a handling routine must be installed
 - Be masked: its handling is delayed
 - Be ignored: its occurrence if not informed