

SURNAME		NAME		Group
ID		Signature		

- Keep the exam sheets stapled.
- Write your answer inside the reserved space.
- Use clear and understandable writing. Answer briefly and precisely.
- The exam has 7 questions, everyone has its score specified.

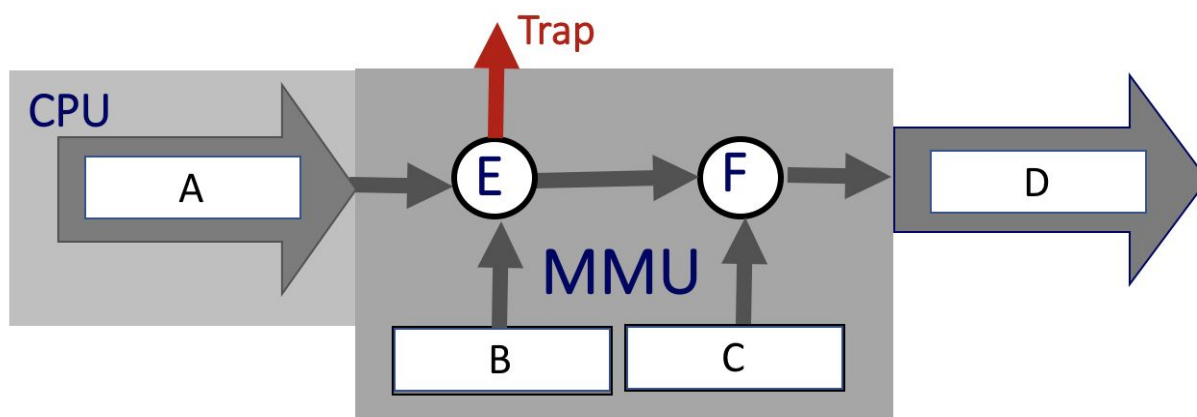
1. Given the following statements about the memory map of a 32-bit Linux process indicate which are true and which are false. (**Note.** An error voids a correct answer)

(1,0 points)

1	Statement	T/F
	The region for initialized data can be identified by checking that it has rw-p permissions and that it has support on the executable file.	
	When creating a local variable, the size of the heap region increases in order to allocate it in memory.	
	When linking a library using static linking, the library information appears in a different specific region on the process memory map.	
	If we show the pointer value of a variable declared inside the main function, it will be located in the region of uninitialized data .	
	Using memory-mapped files improves access time to file information.	

2. The following figure represents the basic operation of an MMU (Memory Management Unit), where the name of some elements have been removed, and they have been replaced by letters from A to F. On the following table, fill column NAME with the name of the elements. Considering that the value of element A is 1000, put in column VALUES the following values: 200, 1200 and 1500, as they correspond in boxes from B to D, assuming that the issued address is correct. .

(1,0 points)



2			
	ELEMENTS	NAME	VALUES
	A		
	B		
	C		
	D		
	E		
	F		

3. A memory management system, based on two level paging, allows up to 16 first level descriptors and second level tables with 256 descriptors each. The page size is 4 Kbytes and every page descriptor, both 1st and 2nd level, is 16 bytes. Addressable main memory is 16 Mbytes. Give your explained answer to the following points:

(1.5 points = 0.5 + 0.25 + 0.75)

3	<p>a) Logical and physical address formats, with name and number of bits for every field or element</p> <p><i>Logical address (name and number of bits for every element)</i></p> <div data-bbox="274 1144 1283 1209" style="border: 1px solid black; height: 29px; width: 632px;"></div> <p><i>Physical address (name and number of bits for every element)</i></p> <div data-bbox="274 1272 1283 1337" style="border: 1px solid black; height: 29px; width: 632px;"></div>
	<p>b) Obtain the hexadecimal range for logical addresses associated with the first 32 pages of a process in the system. Use the following format for your answer [First address, Last address]</p>
	<p>c) Explain how many 1st and 2nd level descriptors requires a process with 2K pages. Also, obtain the total number of pages occupied by the process' 2nd level tables.</p> <p>Nº of 1st level descriptors:</p> <p>Nº of 2nd level descriptors:</p> <p>Nº of pages occupied by 2nd level tables:</p>

4. In a system with virtual memory and demand paging there are currently running two processes: A and B. For these processes, the system has a total of 5 frames. The following table describes the evolution of frame allocation to pages demanded by both processes from time $t = 120$ to $t = 130$. Each cell details the process involved, the accessed page id (in hexadecimal) and the value of the reference bit. In the first column the starting state includes in parentheses the times for the last reference to the page.

	t= 120		t= 121		t= 122		t= 123		t= 124		t= 125	
↓frame / page →	B:d34		B:d34		B:c08		B:c08		A:002		A:555	
c0	A:36a (t=25)	1	A:36a	1	A:36a	1	A:36a	1	A:36a	1	A:36a	1
c1	A:450 (t=50)	0	A:450	0	A:450	0	A:450	0	A:450	0	A:555	1
c2	B:d34	1	B:d34	1	B:d34	1	B:d34	1	B:d34	1	B:d34	1
c3					B:c08	1	B:c08	1	B:c08	1	B:c08	1
c4									A:002	1	A:002	1

	t= 126		t= 127		t= 128		t= 129		t= 130	
↓frame / page →	A:555		B:3af		B:4b0		B:d00		A:0fe	
c0	A:36a	1	A:36a	1	A:36a	1	A:36a	1	A:36a	0
c1	A:555	1	A:555	1	A:555	1	A:555	1	A:555	0
c2	B:d34	1	B:3af	1	B:3af	1	B:d00	1	B:d00	1
c3	B:c08	1	B:c08	0	B:4b0	1	B:4b0	0	B:4b0	0
c4	A:002	1	A:002	1	A:002	1	A:002	1	A:0fe	1

Answer the following items relying on the former information:

(2 points = 0.5 + 0.5 + 0.5 + 0.5)

4	<p>a) Indicate whether the replacement algorithm used is 2nd chance or LRU and if its scope is local or global. Just in case, indicate the first event that discards one of the algorithms. Indicate also the first event that discriminates between local or global scope.</p>																												
	<p>b) If pages accessed at instants $t = 131$ and $t = 132$ are A:36a and A:070, respectively, complete this elements as a continuation of the former table:</p> <table border="1" style="margin-left: auto; margin-right: auto;"> <tr> <th colspan="2">t= 131</th><th colspan="2">t= 132</th></tr> <tr> <th colspan="2">A:36a</th><th colspan="2">A:070</th></tr> <tr><td></td><td></td><td></td><td></td></tr> <tr><td></td><td></td><td></td><td></td></tr> <tr><td></td><td></td><td></td><td></td></tr> <tr><td></td><td></td><td></td><td></td></tr> <tr><td></td><td></td><td></td><td></td></tr> </table>	t= 131		t= 132		A:36a		A:070																					
	t= 131		t= 132																										
	A:36a		A:070																										
<p>c) Assuming that a working set policy is applied, with a window size of four, determine the working sets for processes A and B after access completion at time $t = 130$.</p>																													
<p>d) Explain if at time $t = 130$ thrashing happens for processes A and B, relying on the working set criteria and assuming global frame management.</p>																													

5. Given the process inheritance mechanisms in Unix and the POSIX system calls, answer the following items: :

(1.5 points = 0.5 + 1.0)

5	<p>a) Consider that the command : <code>\$cat < proc.c grep "for" > loop.txt</code> has to run and end correctly. Draw the file descriptor tables for every process involved, describing the content of non empty descriptors. Note. Name the pipe as "fdpipe"</p>
	<p>b) Complete the following C program with the sentences and system calls required in order to perform the same operation as command: <code>\$ /bin/ls /bin/grep ".c" > my_programs</code> Note. Consider that system calls are error free.</p> <pre> /** Prog.c: \$/bin/ls /bin/grep ".c" > my_programs */ #include <unistd.h> #define newfile (O_WRONLY O_CREAT O_TRUNC) #define mode644 (S_IRUSR S_IWUSR S_IRGRP S_IROTH) int main() { int fdpipe[2]; int fd; if (fork()) { execl("/bin/ls", "ls", NULL); } else { fd = open("my_programs", newfile, mode644); } return 0; } </pre>

6. Given the following listing of the content of directory **XXXX/** (with unknown name) on a POSIX system:

```

i-node  permissions links user  group  size  date  time  name
13504322 drwxrwxr-x  3  aperez  disca  4096  dic  21  19:33  .
12200756 drwx-----  5  aperez  disca  4096  dic  21  19:20  ..
13504405 drwxrwxr-x  2  aperez  disca  4096  dic  21  19:29  back
13504425 -rwxr-xr-x  1  aperez  disca 151024  dic  21  19:33  cp
13504425 -rwsr-xr--  1  aperez  disca 151024  dic  21  19:33  cp2
13504417 -rw----r--  1  aperez  disca  6364  dic  21  19:22  dat1
13504421 -rw-rw-r--  3  ana     fso    1134  dic  21  19:25  dat2
13504421 -r--r--r--  3  juan    aic    1134  dic  21  19:25  dat3

```

```
./back:
13504405 drwxrwxr-x 2 aperez disca 4096 dic 21 19:29 .
13504322 drwxrwxr-x 3 aperez disca 4096 dic 21 19:33 ..
13504420 lrwxrwxrwx 1 manolo disca 7 dic 21 19:29 old1 -> ../dat1
13504421 -rw-rw-r-- 3 maria fso 1134 dic 21 19:25 old2
```

(1.5 points = 1 + 0,5)

6

a) cp and cp2 programs are copies of command cp that copies the contents of the file received as the first argument into another file which name is indicated as the second argument. That is, "cp a b" copies the contents of the file a to file b, if b does not exist then b is created and the copy is done. Fill the table indicating whether the corresponding command works, the EUID and EGID (effective UID and GID for the process that executes the command) and in case of error, what is the missing permission(s).

(UID,GID)	COMMAND (from directory XXXX/)	Does it work?	(EUID,EGID)	If it fails tell the missing permission(s)
(ana,fso)	cp dat1 dat4			
(ana, disca)	cp2 dat1 back/old3			
(ana,disca)	cp back/old1 dat2			
(maria,fso)	cp dat3 ..			

b) After running command: `ls -li ..`
we get:

```
i-node  permissions links user group size  date   time  name
12200756 drwx----- 5 aperez aperez 4096 dic 26 11:52 .
10223617 drwxr-xr-x 149 aperez aperez 12288 dic 25 01:10 ..
13764791 drwxrwxr-x 2 aperez aperez 4096 dic 26 11:51 alumnos
13504322 drwxrwxr-x 3 aperez aperez 4096 dic 26 11:51 organizacion
13636461 drwxrwxr-x 2 aperez aperez 4096 dic 26 11:50 practicas
```

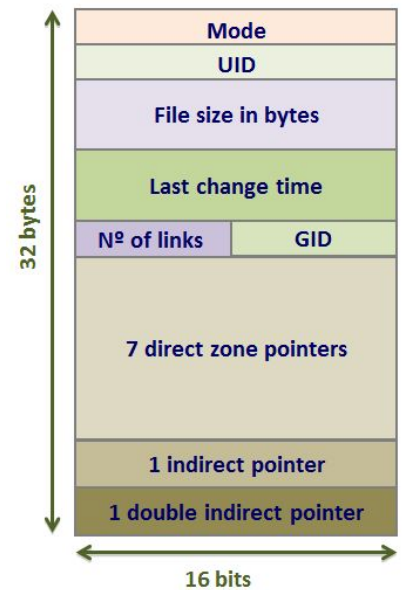
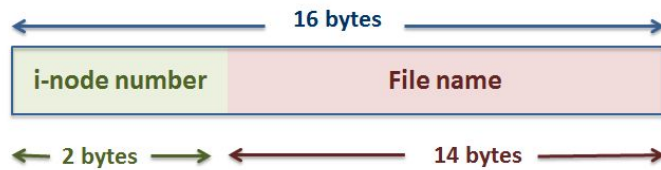
What is the name of directory XXXX/?

7. The following figures refer to the sizes and structures of the elements used to format a partition with a MINIX file system, with the following configuration :

- N° of i-nodes : 3424
- Block size : 1 KByte
- 1 zone = 1 block
- Partition size : 10 MBytes

Notice that all the i-node fields are 16-bit, except “N° of links” and GID that are 8-bit.

Directory entries have the following format:



(1,5 points = 0,5 + 0,2 + 0,8)

a) Number of blocks that occupy the following header elements: i-nodes bitmap, data zones bitmap and i-nodes .

b) Block index for the first block in the data area .

c) Just after formatting the partition a 64 KByte file named "message.txt", is created. Knowing that the first block assigned to the file is 113, and that blocks are allocated in order, obtain in what block is located byte 62459 from file “message.txt”