

CS 5/7320

Artificial Intelligence

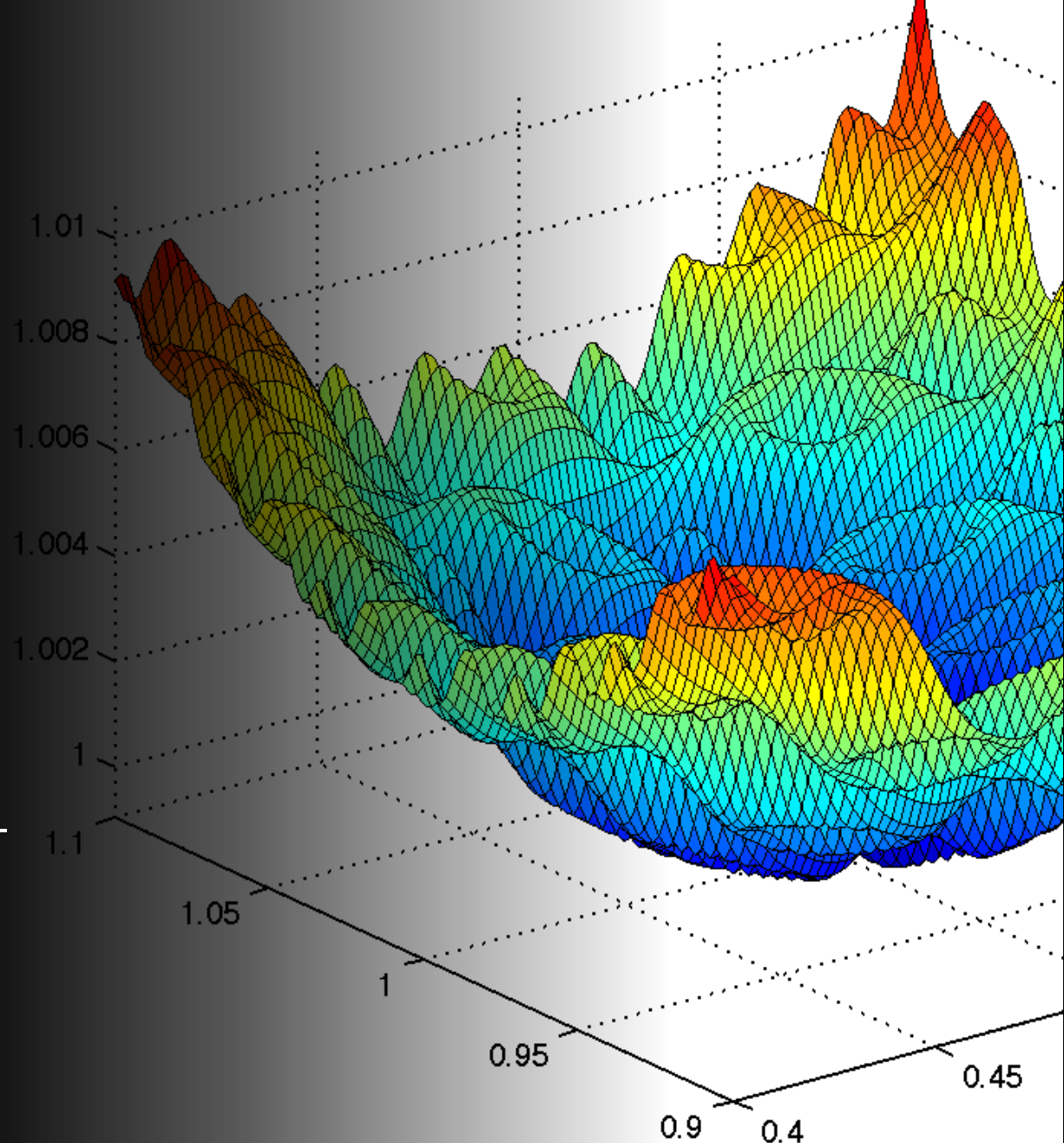
Local Search

AIMA Chapters 4.1 & 4.2

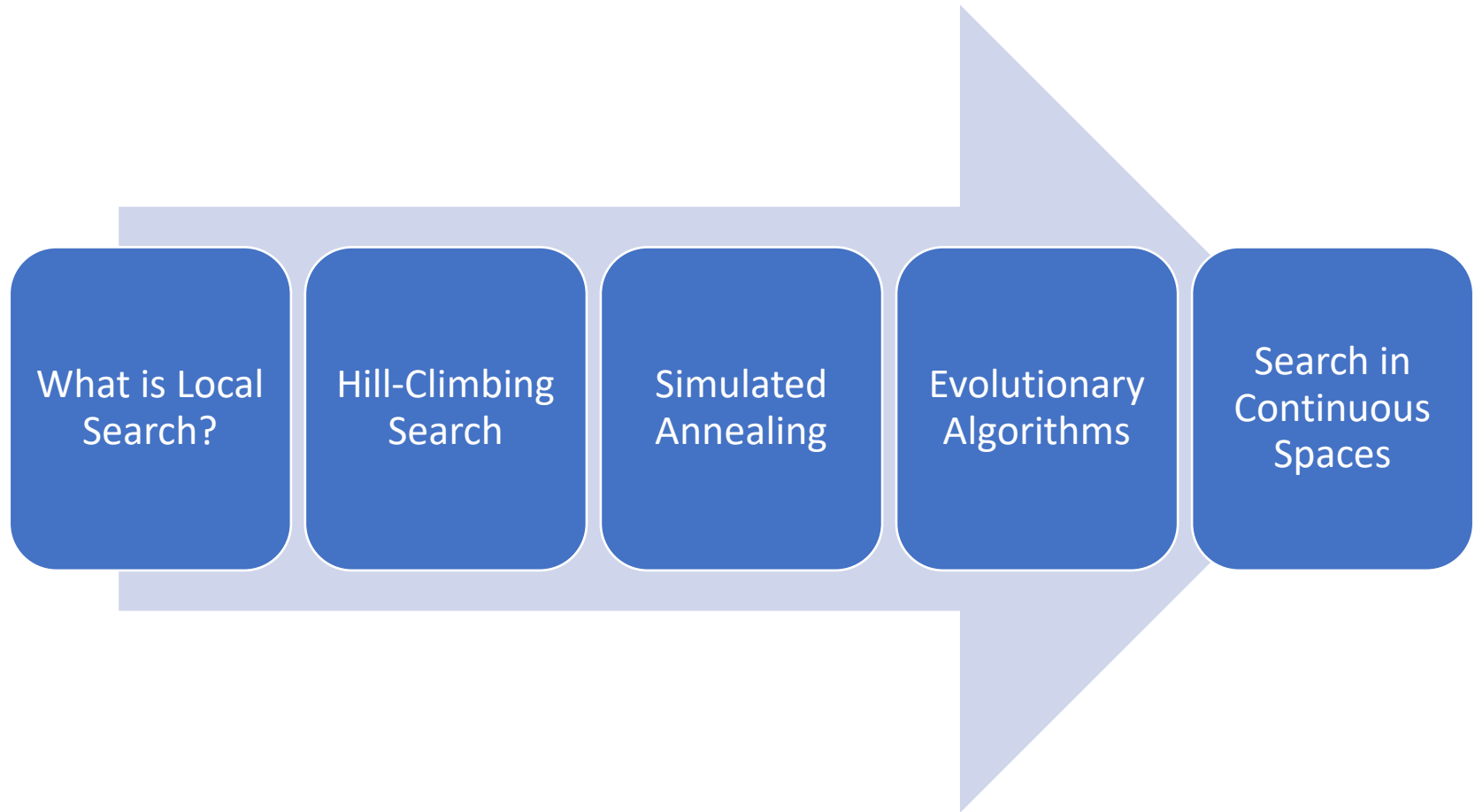
Slides by Michael Hahsler
based on slides by Svetlana Lazepnik
with figures from the AIMA textbook.



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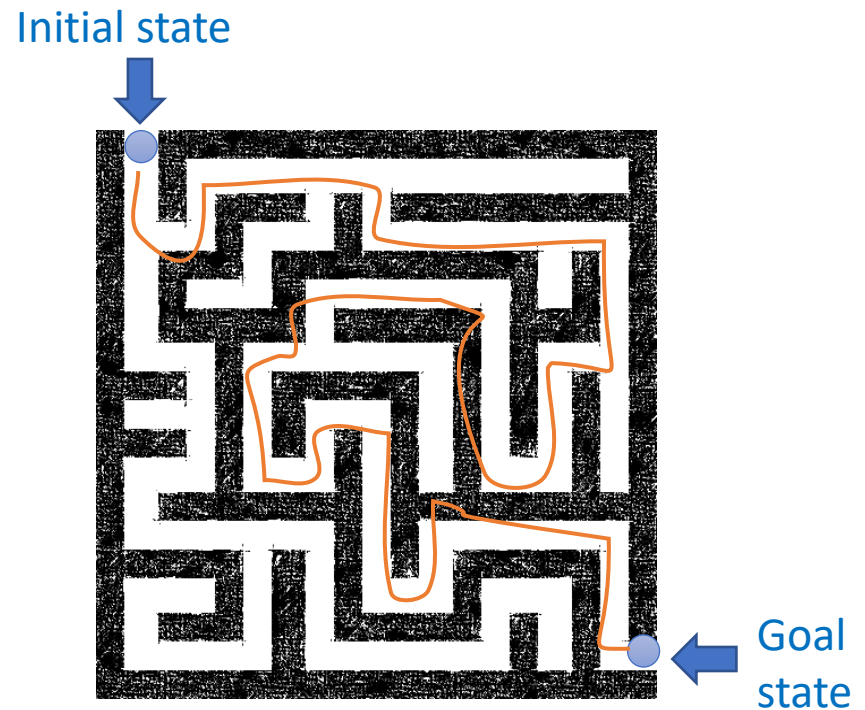
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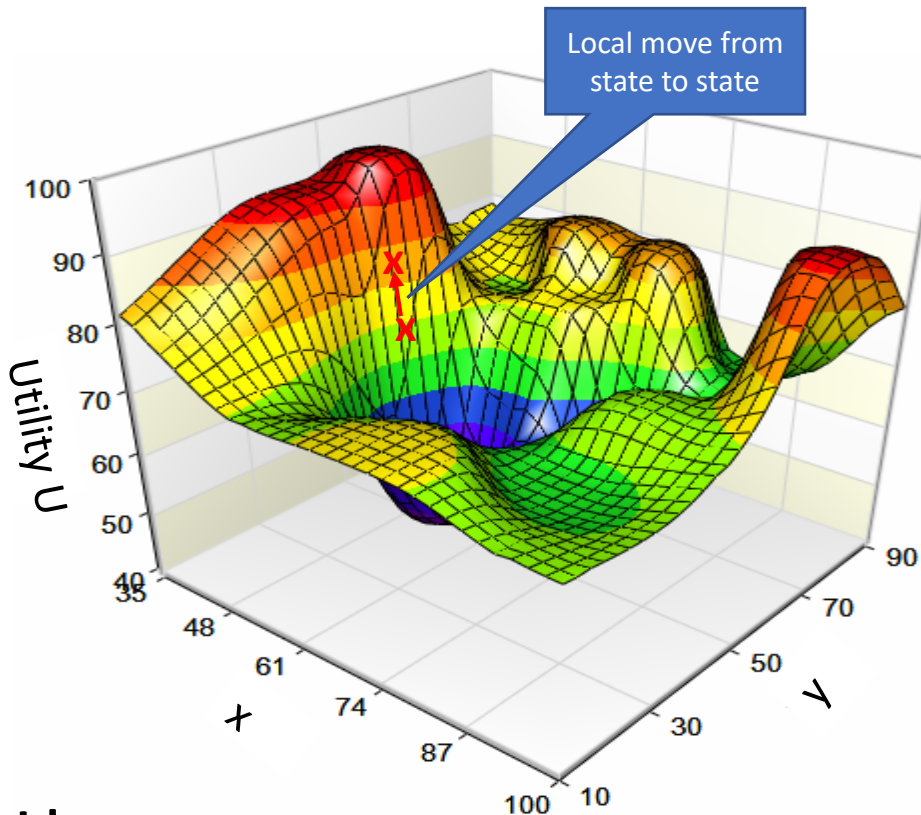
Recap: Uninformed and Informed Search

Tries to **plan** the
best path
from a
given initial state
to a
given goal state.

- Often comes with optimality guarantees (BFS, A* Search, IDS).
- Typically searches a large portion of the search space (needs time and memory).



Local Search Algorithms



- What if we do not know the goal state, but the utility of different states is given by a utility function $U = u(s)$?

- We use a factored state description. Here $s = (x, y)$
- We could try to identify the best or at least a “good” state?

- This is the **optimization problem**:
$$s^* = \operatorname{argmax}_{s \in S} u(s)$$

- We need a fast and memory-efficient way to find the best/a good state.

Idea:

Start with a current solution (a state) and improve the solution by moving from the current state to a “neighboring” better state (a.k.a. performing a series of **local moves**).

Local Search Algorithms

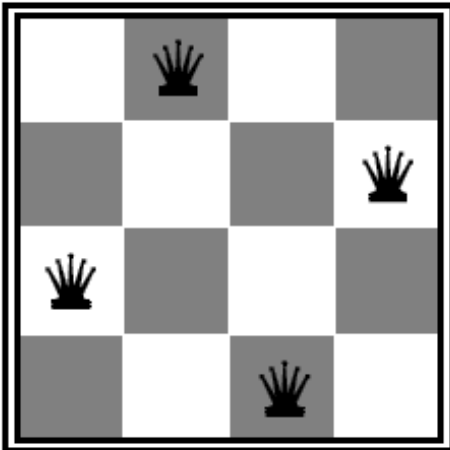
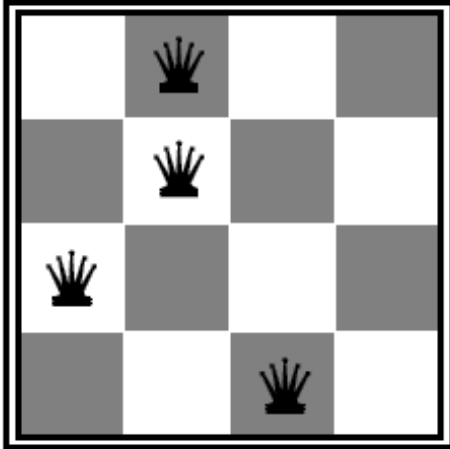
Difference to search from the previous chapter:

- a) **Goal state is unknown**, but we know or can calculate the utility for each state. We want to identify the state with the highest utility.
- b) Often no explicit initial state + **path to goal and path cost are not important**.
- c) **No search tree**. Just stores the current state and move to a “better” state if possible.

Use in AI

- **Goal-based agent**: Identify a good goal state with a good utility before planning a path to that state.
- **Utility-based agent**: Always move to a neighboring state with higher utility. A simple greedy method used for
 - complicated/large state spaces or
 - online search.
- **General optimization**: $u(s)$ can be replaced by a general objective function. Local search is an effective heuristic to find good solutions in large or continuous search spaces. E.g., stochastic gradient descent to train neural networks learns to approximate a function by using the prediction error as the objective function.

states



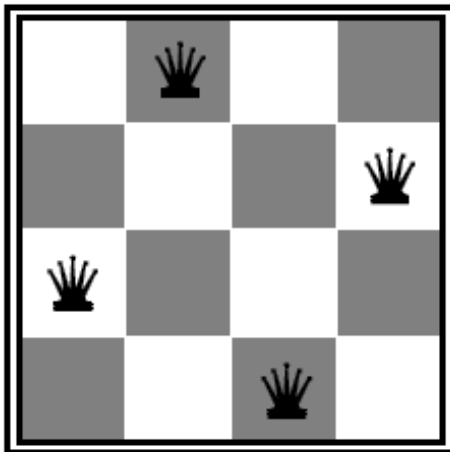
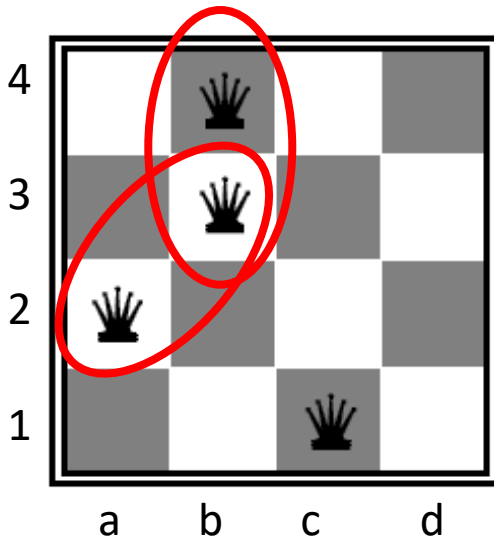
Example: *n*-Queens Problem

Goal: Put n queens on an $n \times n$ board with no two queens on the same row, column, or diagonal.

Defining the search problem:

- **State space:** All possible n -queen configurations. How many are there?
- **State representation:** How do we define a factored representation?
- **Objective function:** What is a possible utility function given the state representation?
- **Local neighborhood:** What states are close to each other?

2 conflicts = utility of -2



0 conflicts = utility of 0

Example: *n*-Queens Problem

Defining the search problem:

- **State space:** All possible *n*-queen configurations. How many are there?
4-queens problem: $\binom{16}{4} = 1820$
- **State representation:** How do we define a factored representation?
E.g. $(a2, b3, b4, c1)$
- **Objective function:** What is a possible utility function given the state representation?
Maximizing utility means minimize the number of pairwise conflicts based on the state representation.
- **Local neighborhood:** What states are close to each other?
Move a single queen.


Has its optimum at the goal state. Similar to a heuristic in A* search.

Defines a transition function.



- **Goal:** Find the shortest tour connecting a given set of cities
- **State space:** all possible tours (states are not individual cities!)
- **State representation:** Order of cities in the tour.
- **Objective function:** minimize the length of the tour.
- **Local neighborhood:** Change the order of visiting a few cities.

Note: We have solved a **different** problem with uninformed/informed search! Each city was defined as a state and the path was the solution.



Hill-Climbing Search aka Greedy Local Search

Idea: keep a single “current” state and try to find better neighboring states.

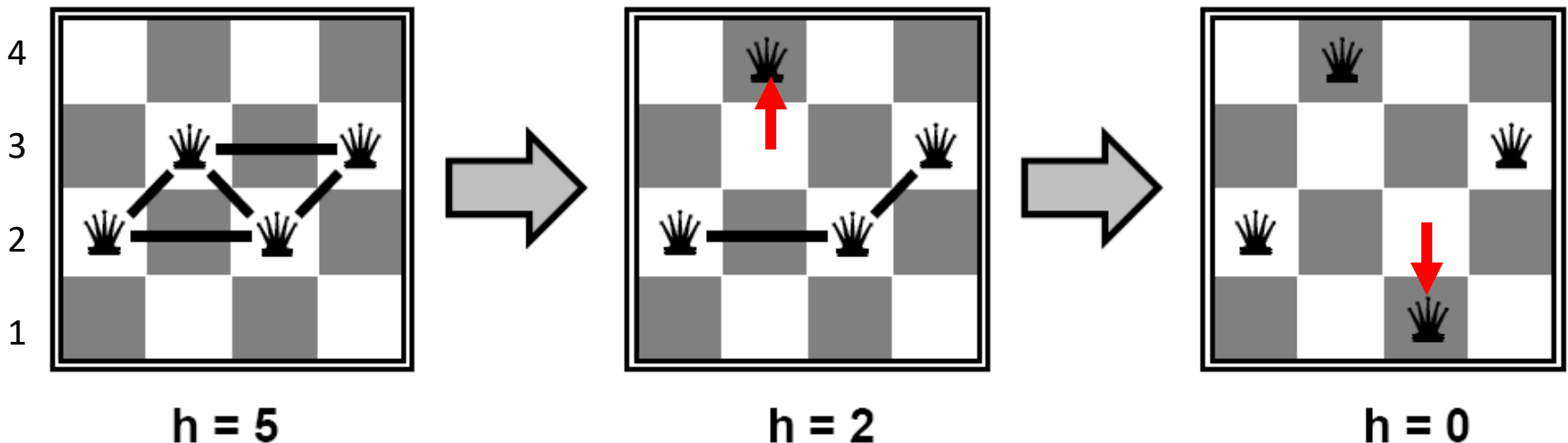
Example: n -Queens Problem

- **Goal:** Put n queens on an $n \times n$ board with no two queens on the same row, column, or diagonal.
- **State space:** all possible n -queen configurations. We can restrict the state space: Only one queen per column.
- **State representation:** row position of each queen in its column (e.g., 2, 3, 2, 3)
- **Objective function:** minimize the number of pairwise conflicts.
- **Local neighborhood:** Move one queen anywhere in its column.

State space is reduced from 1820 to $4^4 = 256$

Improvement strategy

- Find a local neighboring state (move one queen within its column) to reduce conflicts



Example: n -Queens Problem

To find the best local move, we must evaluate all local neighbors (moving a single queen in its column while leaving the others in place) and calculate the objective function.

Objective value after moving the queen in column 1 to this square

18	12	14	13	13	12	14	14
14	16	13	15	12	14	12	16
14	12	18	13	15	12	14	14
15	14	14	♙	13	16	13	16
♙	14	17	15	♙	14	16	16
17	♙	16	18	15	♙	15	♙
18	14	♙	15	15	14	♙	16
14	14	13	17	12	14	12	18

Current objective value: $h = 17$

Best local improvement has $h = 12$

Notes:

- There are many options with $h = 12$. We must choose one!
- Calculating all the objective values may be expensive!

Example: n -Queens Problem

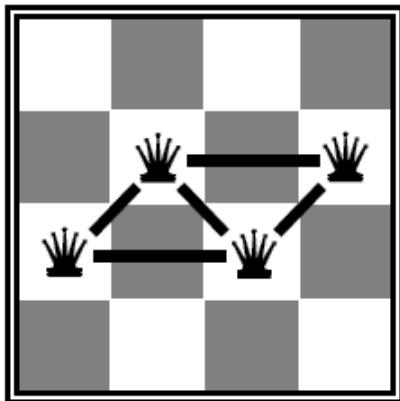
Formulation as an optimization problem:

Find the best state s^* representing an arrangement of queens.

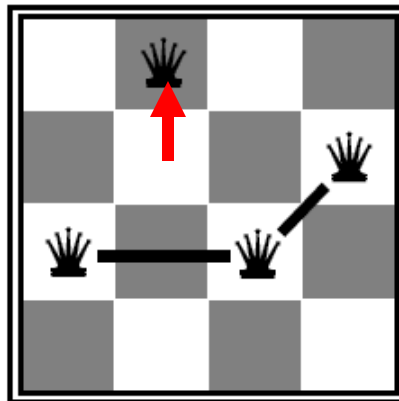
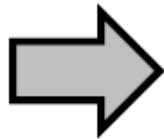
$$s^* = \operatorname{argmin}_{s \in \mathcal{S}} \operatorname{conflicts}(s)$$

subject to: s has one queen per column

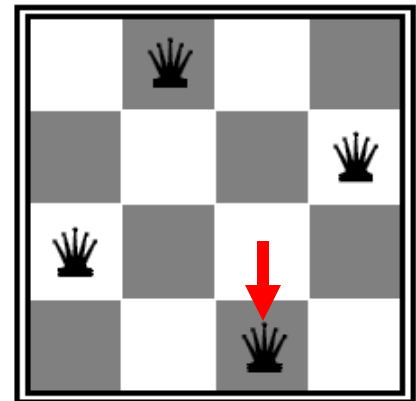
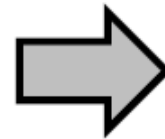
Remember: This makes the problem a lot easier.



$h = 5$



$h = 2$



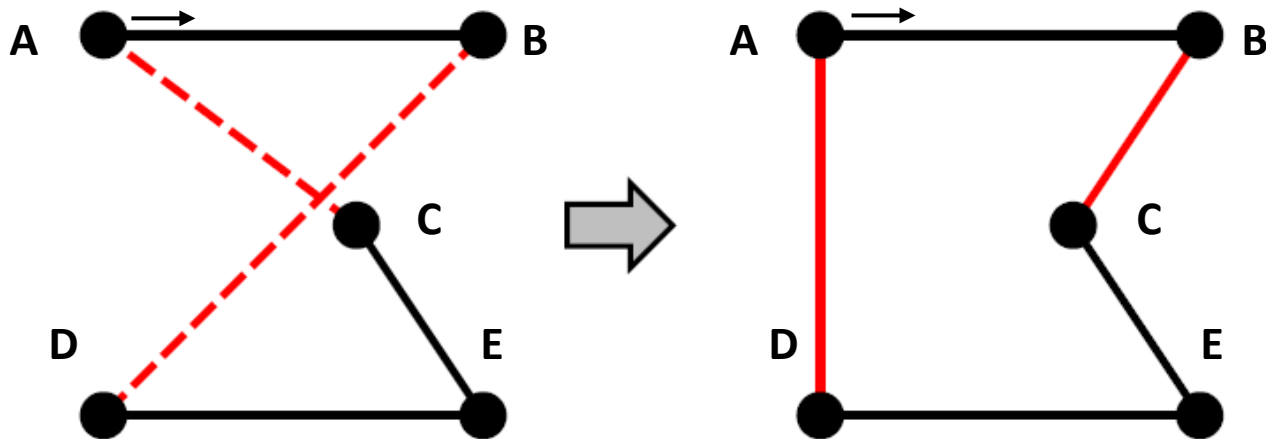
$h = 0$

Example: Traveling Salesman Problem



- **Goal:** Find the shortest tour connecting n cities
- **State space:** all possible tours
- **State representation:** tour (order in which to visit the cities) = a permutation. There are $n!$ Many permutations.
- **Objective function:** length of tour
- **Local neighborhood:** reverse the order of visiting a few cities

Local move to reverse the order of cities C, E and D:



State representation
(permutation):

ABDEC

ABCED

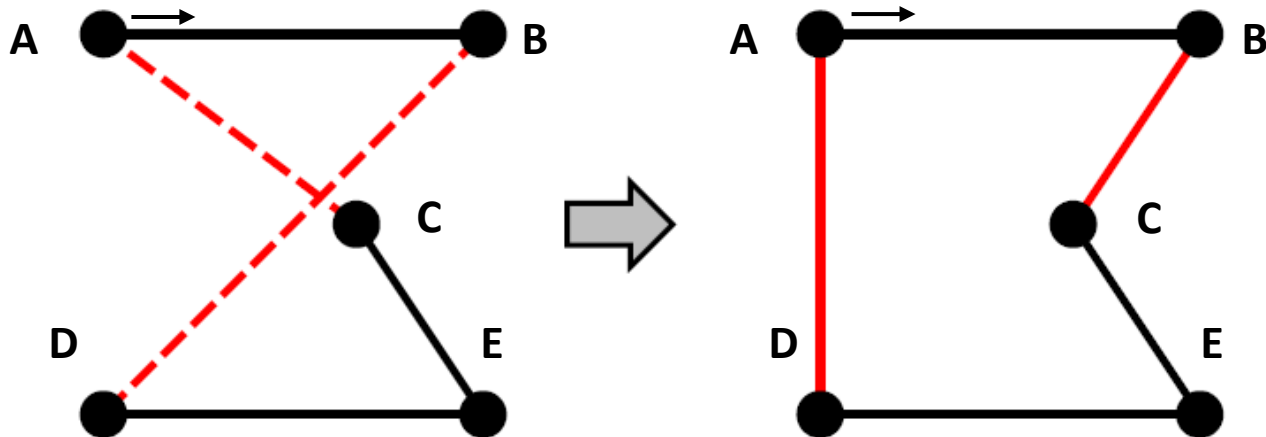
Example: Traveling Salesman Problem

Formulation as an optimization problem:
Find the best tour π

$$\pi^* = \operatorname{argmin}_{\pi} \operatorname{tourLength}(\pi)$$

s.t. π is a valid permutation (i.e., sub-tour elimination)

Local move to reverse the order of cities C, E and D:



State representation: **ABDEC**

ABCED



Hill-Climbing Search (Greedy Local Search)

```
function HILL-CLIMBING(problem) returns a state that is a local maximum
  current  $\leftarrow$  problem.INITIAL
  while true do
    neighbor  $\leftarrow$  a highest-valued successor state of current
    if VALUE(neighbor)  $\leq$  VALUE(current) then return current
    current  $\leftarrow$  neighbor
```

Typically, we start with a random state

Variants:

Steepest-ascend hill climbing

- Check all possible successors and choose the highest-valued successors.

Stochastic hill climbing

- Choose randomly among all uphill moves, or
- generate randomly one new successor at a time and only move to better ones = first-choice hill climbing – the most popular variant, this is what people often mean when they say “stochastic hill climbing”

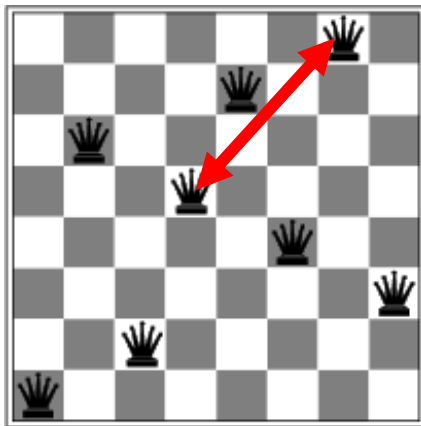
18	12	14	13	13	12	14	14
14	16	13	15	12	14	12	16
14	12	18	13	15	12	14	14
15	14	14	♔	13	16	13	16
♔	14	17	15	♔	14	16	16
17	♔	16	18	15	♔	15	♔
18	14	♔	15	15	14	♔	16
14	14	13	17	12	14	12	18

Local Optima

Hill-climbing search is like greedy best-first search with the objective function as a (maybe not admissible) heuristic and no frontier (just stops in a dead end).

Is it complete/optimal?

- No – can get stuck in local optima



$$h = 1$$

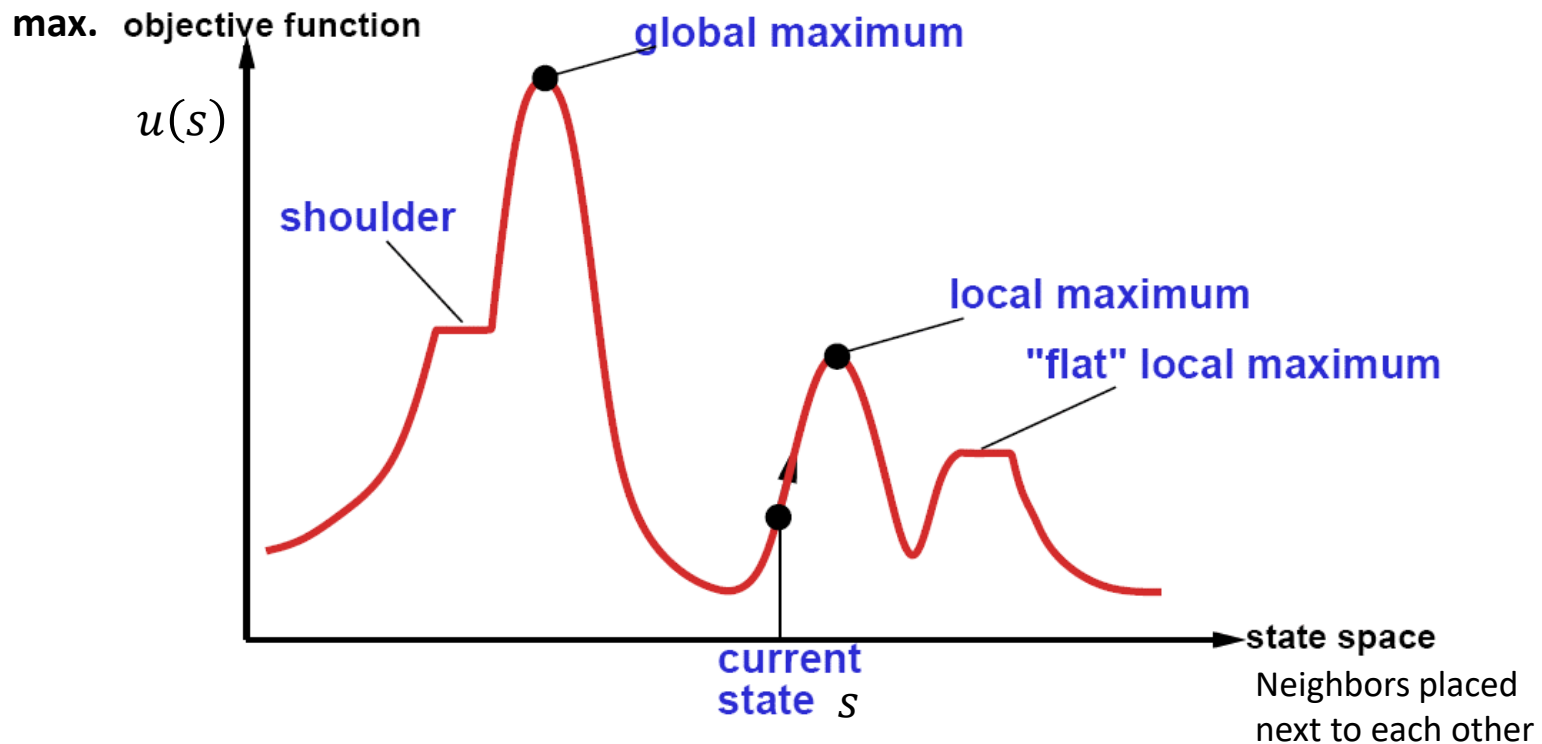
Example: local optimum for the 8-queens problem. No single queen can be moved within its column to improve the objective function.

Simple approach that can help with local optima:

Random-restart hill climbing: Restart hill-climbing many times with random initial states and return the best solution.

The State Space “Landscape”

We can get the utility (objective function value) from the state description using $u(s)$.



How to escape local maxima?

→ Random restart hill-climbing can help.

What about “shoulders” (called “ridges” in higher dimensional space)?

→ Hill-climbing that allows sideways moves and uses momentum.

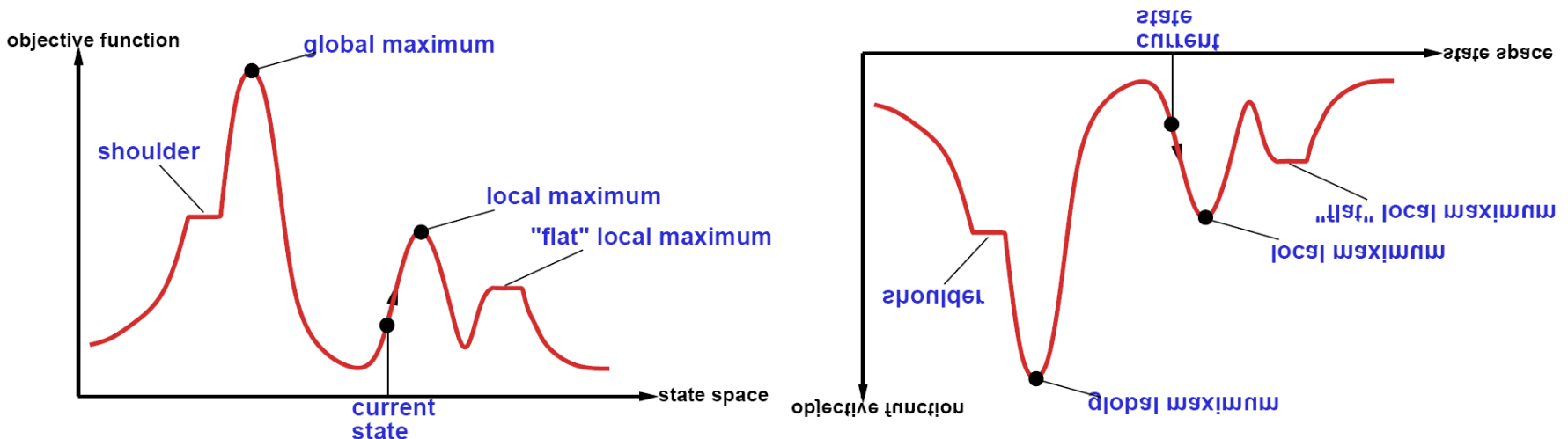
Minimization vs. Maximization

- The name **hill climbing** implies **maximizing a function**.
- Optimizers like to state problems as **minimization problems** and call hill climbing **gradient descent** instead.
- Both types of problems are equivalent:

$$\max(f(x))$$

\Leftrightarrow

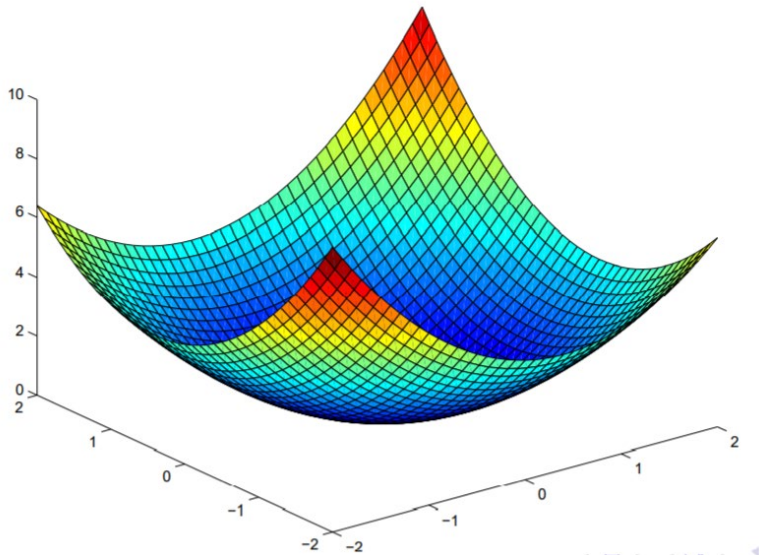
$$\min(-f(x))$$



Convex vs. Non-Convex Optimization Problems

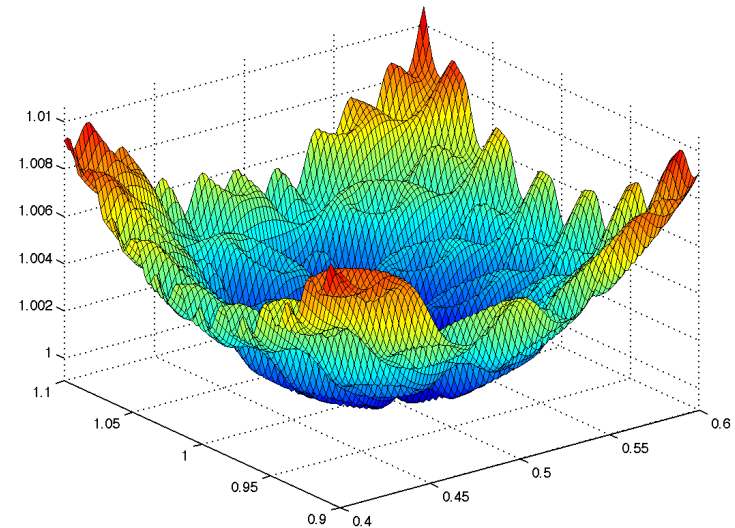
Minimization problems

Convex Problem



One global optimum +
smooth function → calculus
makes it easy

Non-convex Problem



Many local optima → hard

Many discrete optimization
problems are like this.

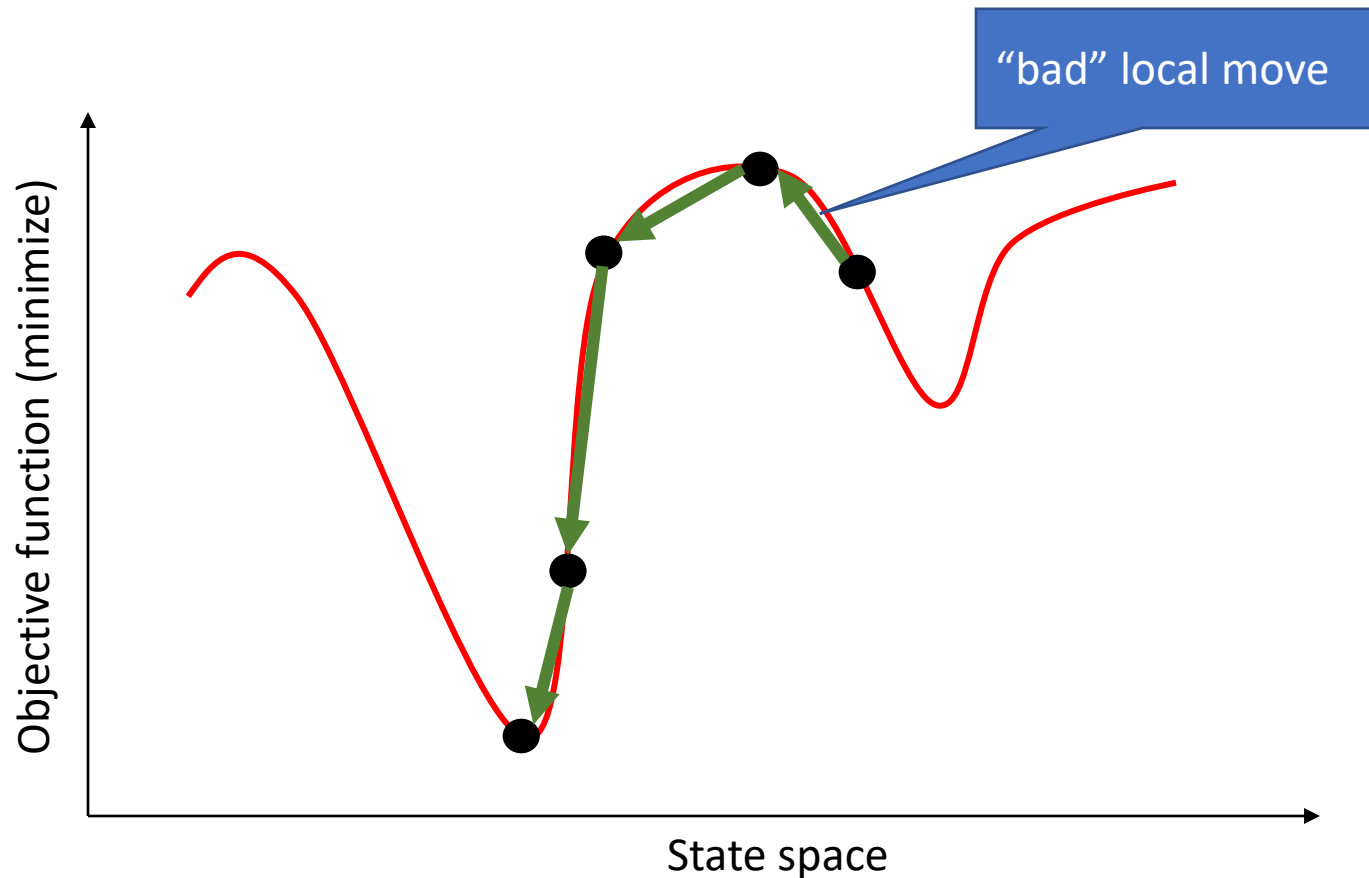
A glowing yellow-orange metal bar is being processed in a dark industrial machine. The bar is long and rectangular, with a bright, intense glow along its length. It is positioned horizontally, and the machine's components, including rollers and structural frames, are visible in the background. The lighting is dramatic, with the bright glow of the metal contrasting sharply with the dark, shadowed machinery.

Simulated Annealing

Use heat to escape local optima...

Simulated Annealing

- **Idea:** First-choice stochastic hill climbing + escape local minima by **allowing some “bad” moves** but gradually decrease their frequency.
- Inspired by the process of controlled cooling of glass or metals by decreasing the temperature (here chance of accepting bad moves) gradually.



Simulated Annealing

- **Idea:** First-choice stochastic hill climbing + escape local minima by allowing some “bad” moves but gradually decreasing their frequency as we get closer to the solution.
- Annealing tries to reach a low energy state so a negative ΔE means the solution gets better.
- The probability of accepting “bad” moves follows the **annealing schedule** that reduces the temperature T over time t .

Maximization

function SIMULATED-ANNEALING(*problem, schedule*) **returns** a solution state

current \leftarrow *problem*.INITIAL

Typically, we start with a random state

for $t = 1$ **to** ∞ **do**

$T \leftarrow$ *schedule*(t)

if $T = 0$ **then return** *current*

next \leftarrow a randomly selected successor of *current*

$\Delta E \leftarrow$ VALUE(*current*) – VALUE(*next*)

if $\Delta E < 0$ **then** *current* \leftarrow *next*

else *current* \leftarrow *next* only with probability $e^{-\Delta E/T}$

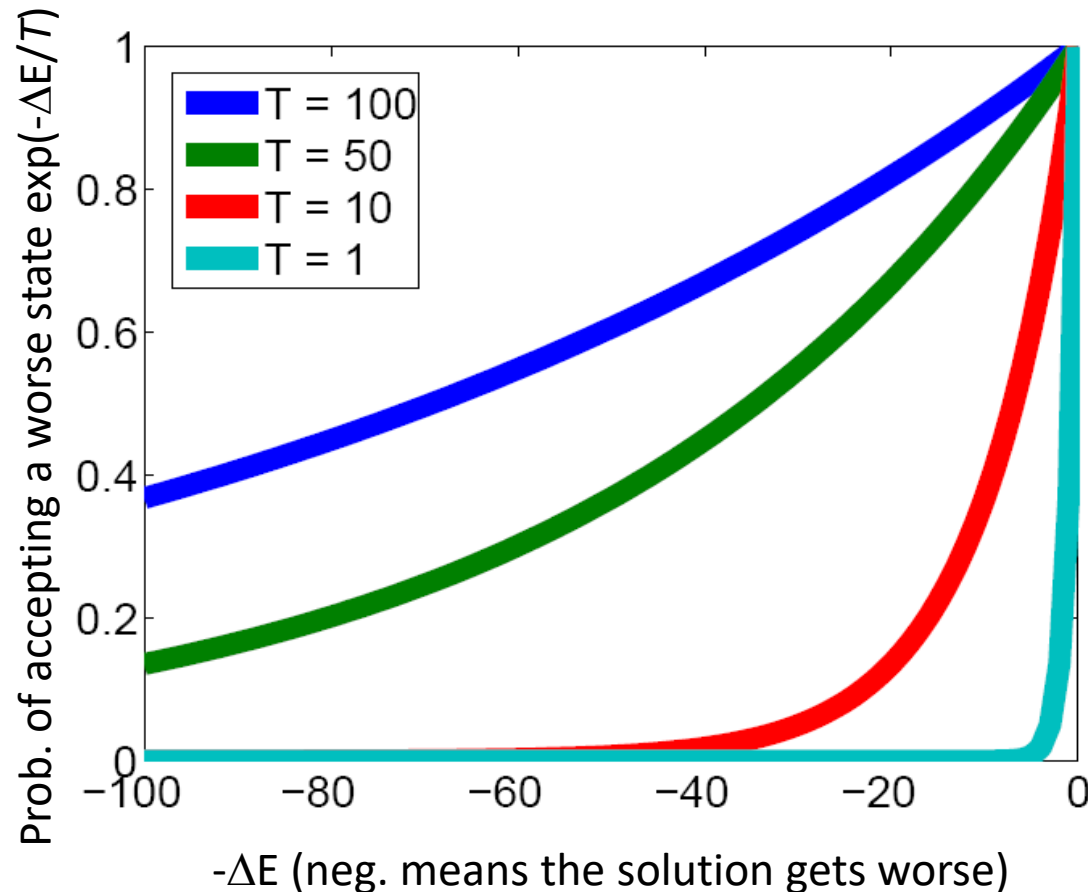
Always accept good moves that reduce the energy.

Accept “bad” moves with a probability inspired by the acceptance criterion in the Metropolis–Hastings MCMC algorithm.

Note: Use VALUE(*next*) – VALUE(*current*) for minimization

The Effect of Temperature

Convert the changes due to “bad” moves into an acceptance probability depending on the temperature. The criterion uses the negative part of the exponential function.



Cooling Schedule

The cooling schedule is very important.

Popular schedules for the temperature at time t :

- **Classic simulated annealing:** $T_t = T_0 \frac{1}{\log(1+t)}$
- **Exponential cooling** (Kirkpatrick, Gelatt and Vecchi; 1983)

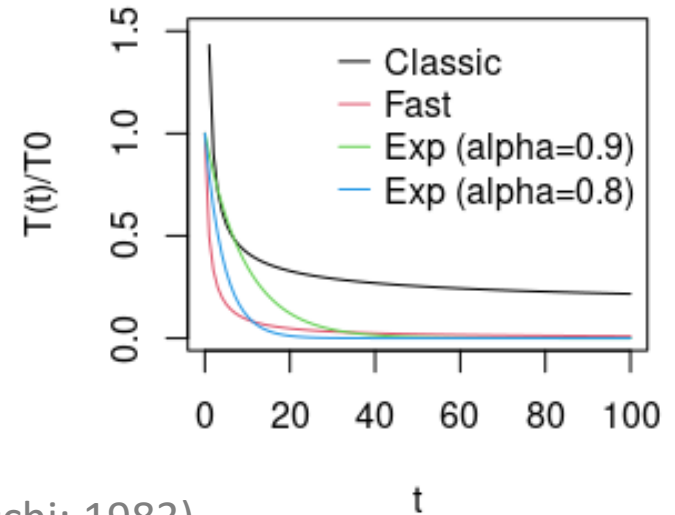
$$T_t = T_0 \alpha^t \quad \text{for } 0.8 < \alpha < 1$$

- **Fast simulated annealing** (Szy and Hartley; 1987)

$$T_t = T_0 \frac{1}{1+t}$$

Notes:

- Choose T_0 to provide a high probability $p_0 = e^{-\frac{\Delta E}{T_0}}$ that any move will be accepted at time $t = 0$. ΔE is determined by the worst possible move.
- T_t will not become 0 but very small. Stop when $T < \epsilon$ (ϵ is a very small constant).
- The best schedule (cooling rate) is typically determined by trial-and-error. The goal is to have a low chance of getting stuck in a local optima.

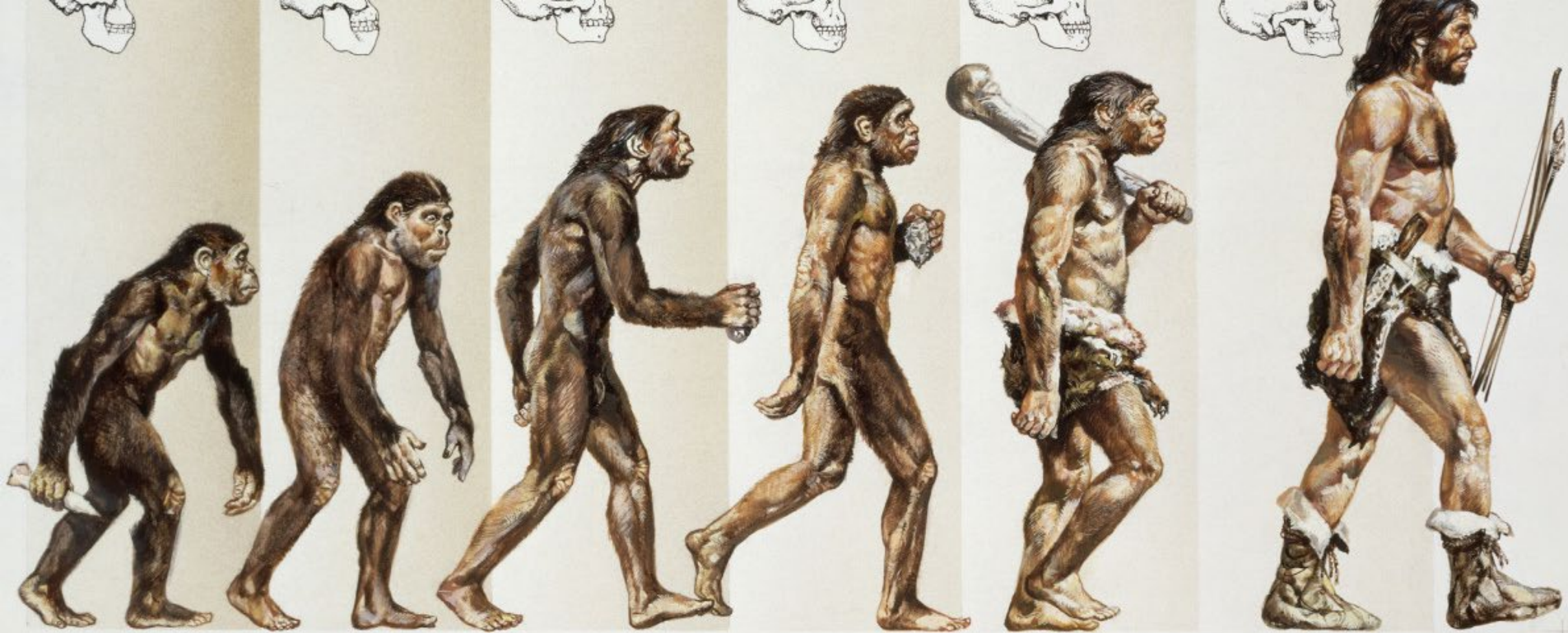


Simulated Annealing Search

Guarantee: If the temperature is decreased **slowly enough**, then simulated annealing search will find a global optimum with a probability approaching one.

However:

- This usually takes impractically long.
- We need to experiment with the cooling schedule to find one that typically avoids local optima.



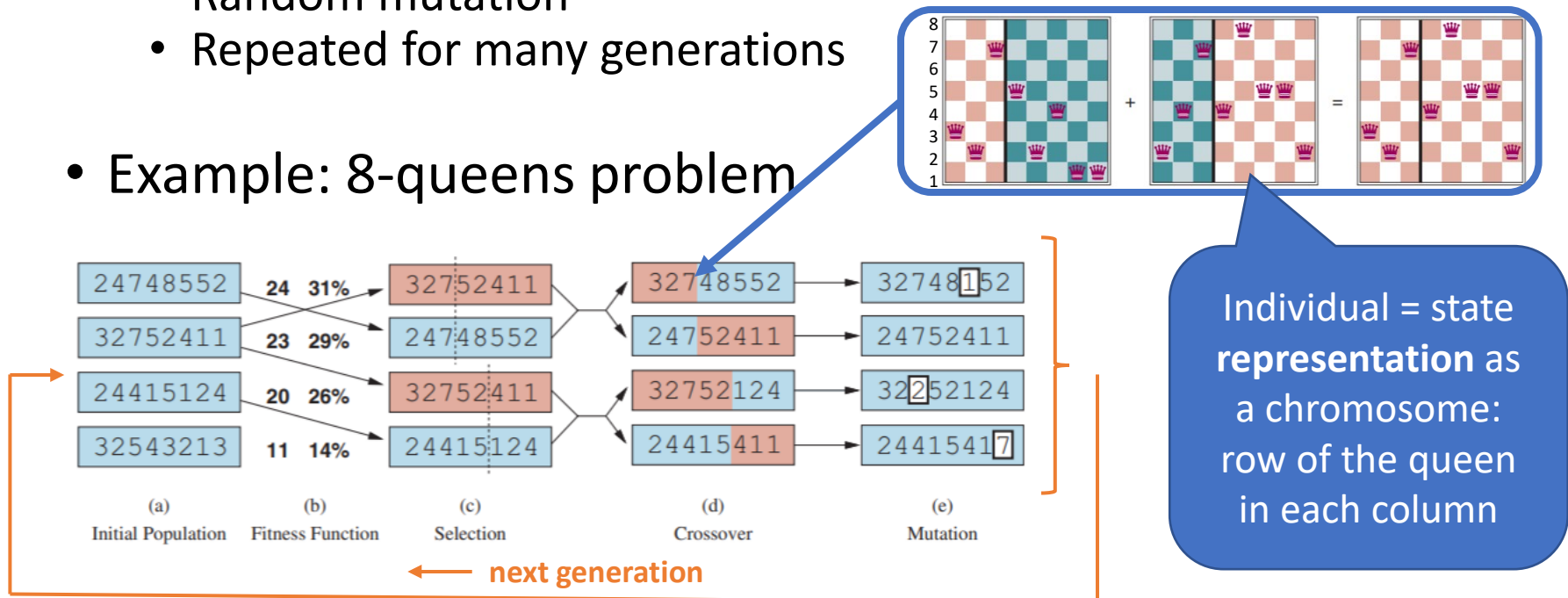
Evolutionary Algorithms

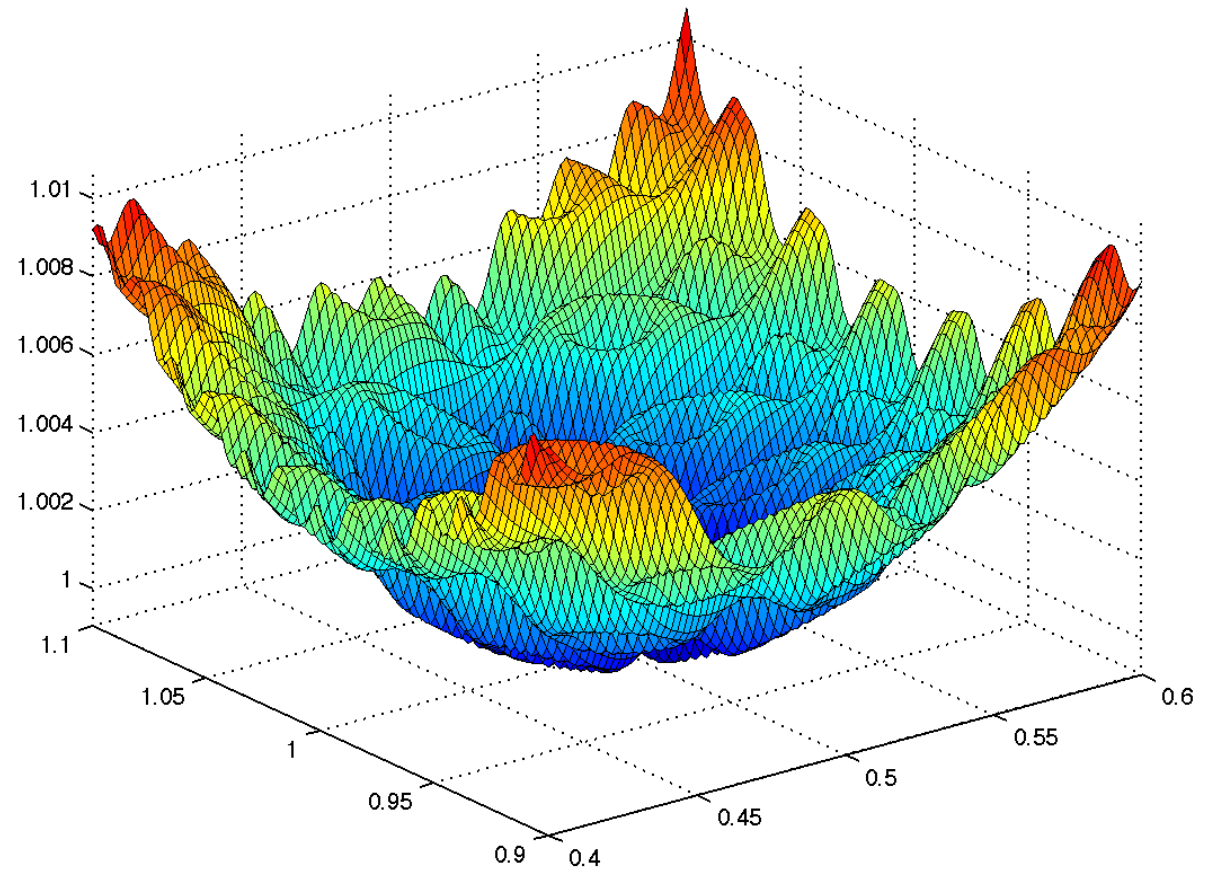
A Population-based Metaheuristics

Evolutionary Algorithms / Genetic Algorithms

- A metaheuristic for **population**-based optimization.
- Uses mechanisms inspired by biological evolution (genetics):
 - Reproduction: Random selection with probability based on a **fitness** function.
 - Random recombination (crossover)
 - Random mutation
 - Repeated for many generations

- Example: 8-queens problem

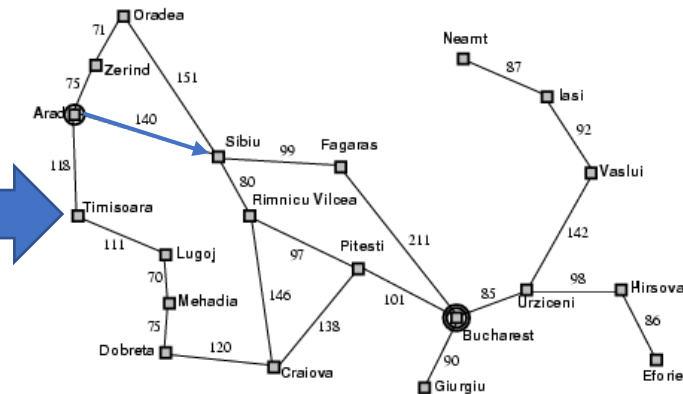




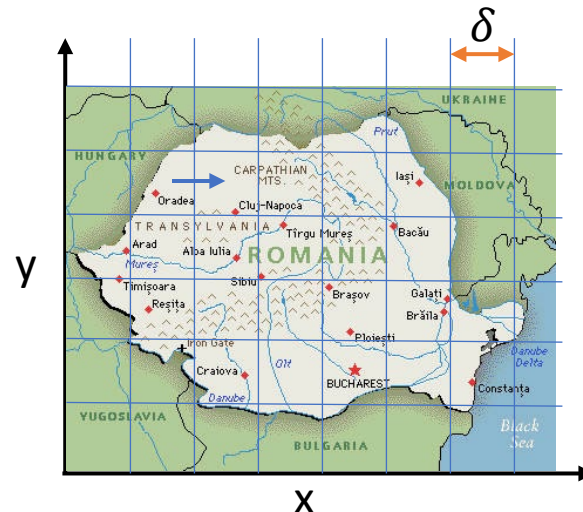
Search in Continuous Spaces

Discretization of Continuous Space

- Use atomic states and create a graph as the transition function.



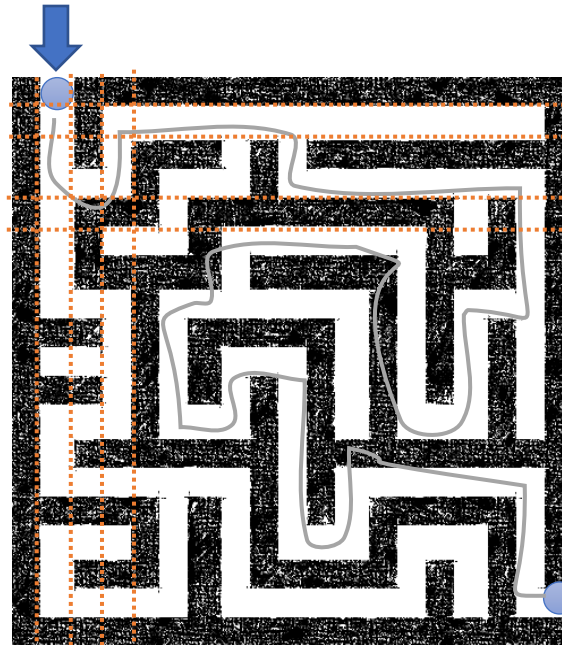
- Use a grid with spacing of size δ
Note: You probably need a way finer grid!



Discretization of Continuous Space

How did we discretize this space?

Initial state



..... Discretization grid

Search in Continuous Spaces: Gradient Descent

State representation: $\mathbf{x} = (x_1, x_2, \dots, x_k)$

State space size: infinite

Objective function: $\min f(\mathbf{x}) = f(x_1, x_2, \dots, x_k)$

Local neighborhood: small changes in x_1, x_2, \dots, x_k

Gradient at point \mathbf{x} : $\nabla f(\mathbf{x}) = \left(\frac{\partial f(\mathbf{x})}{\partial x_1}, \frac{\partial f(\mathbf{x})}{\partial x_2}, \dots, \frac{\partial f(\mathbf{x})}{\partial x_k} \right)$
(=evaluation of the Jacobian matrix at \mathbf{x})

Find optimum by solving: $\nabla f(\mathbf{x}) = 0$

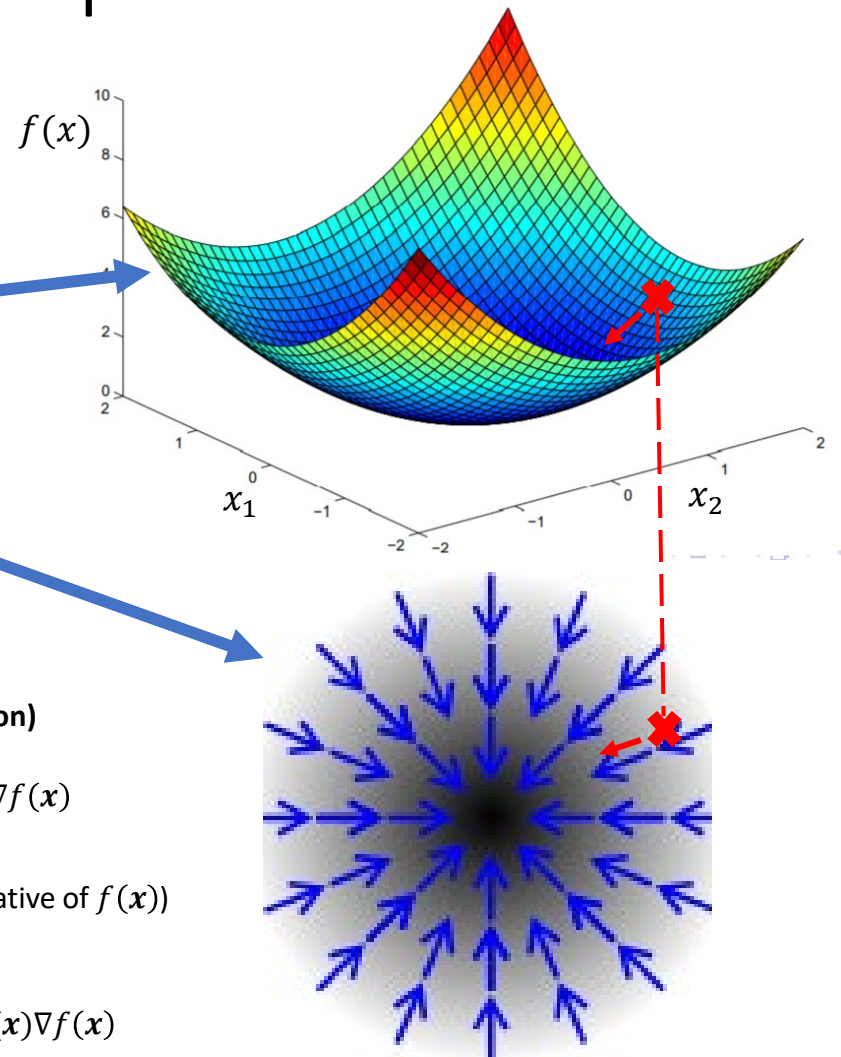
- **Gradient descent (= Steepest-ascent hill climbing for minimization)**
with step size α (typically reduced over time)

$$\text{Repeat: } \mathbf{x} \leftarrow \mathbf{x} - \alpha \nabla f(\mathbf{x})$$

- **Newton-Raphson method**
uses the inverse of the Hessian matrix (second-order partial derivative of $f(\mathbf{x})$)

$$H_{ij} = \frac{\partial^2 f}{\partial x_i \partial x_j} \text{ as the optimal step size}$$

$$\text{Repeat: } \mathbf{x} \leftarrow \mathbf{x} - \mathbf{H}_f^{-1}(\mathbf{x}) \nabla f(\mathbf{x})$$



Note: May get stuck in a local optimum if the search space is non-convex! Use simulated annealing, momentum or other methods to escape local optima.

Search in Continuous Spaces: Stochastic Gradient Descent

- What if the mathematical formulation of the objective function is not known?
 - We may have objective values at fixed points, called the **training data**.
 - In this case, we can perform gradient descent with an approximation of the gradient using the data points. This is called **stochastic gradient descent (SGD)**.
- We will talk more about search in continuous spaces with loss functions using gradient descent when we talk about **parameter learning for learning from examples (machine learning)**.



Conclusion

- Local search provides a fast method to find good solutions to many difficult optimization problems.
- Local optima are a big issue that can be addressed with random restarts and simulated annealing.