## Virtual Memory A Project for CS854

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University of Waterloo

February 25, 2016

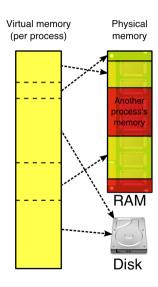
#### **Abstract**

In short:

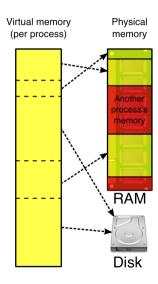
#### **Abstract**

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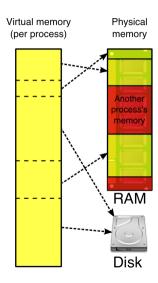
• We propose to study virtual memory!



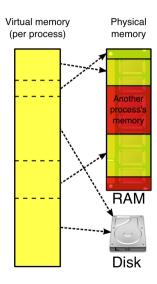
• (on x86) Instructions operate on virtual addresses



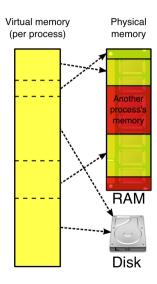
- (on x86) Instructions operate on virtual addresses
- Data may be stored:



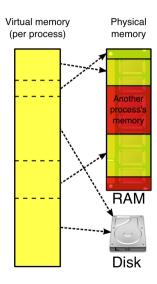
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- Data may be stored:
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- Each process has a page table
  - $\bullet \ \ \text{Maps virtual} \to \text{physical addresses}$

Our proposal has 3 parts:

1 Literature Review

- 1 Literature Review
- 2 Experimental Design

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- 2 Experimental Design
- 3 Implementation

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- 3 Implementation

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5/20

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For each OS, we wish to answer the following questions:

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- How is memory freed?

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- Are there data structures for physical pages, separate from the page tables?
- How are contiguous regions of memory managed?
- How is memory freed?
  - What happens when the kernel runs out of memory?
- Do they do anything special on Non-Uniform Memory Access (NUMA) architectures?

- 1 Literature Review
- 2 Experimental Design
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- Example:
  - Implement data structures from different VMs
  - Test performance

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# Proposal: Implementation

Optional

- Optional
- Implement a memory management system for KOS

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  - Experiments

## High-level design

Now we'll summarize the VM design of:

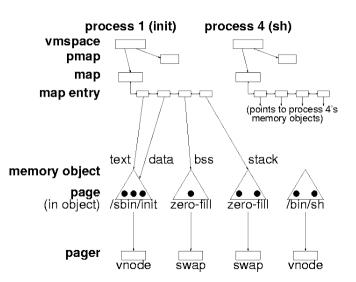
- Linux
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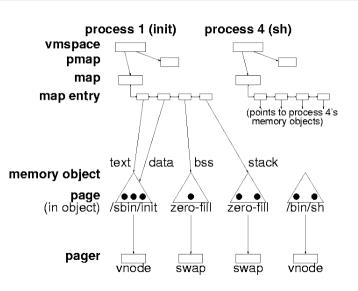
## High-level: Linux

vm\_area\_struct

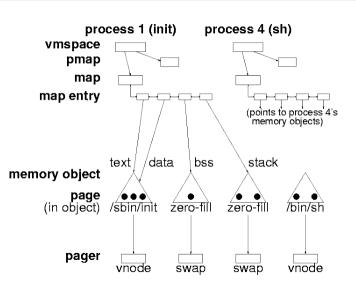
```
44 struct vm area struct {
45
       struct mm struct * vm mm:
                                                                         task struct
46
       unsigned long vm start;
47
       unsigned long vm end:
                                                                                                mm struct
                                                                                                                        vm area struct
49
                                                                          mm
                                                                                                                           vm end
                                                                                               count
50
       /* linked list of VM areas per task, sorted by address */
                                                                                                                           vm_start
51
       struct vm area struct *vm next:
                                                                                               pgd
                                                                                                                                                    Data
52
                                                                                                                          ym flags
53
       pgprot t vm page prot;
                                                                                                                          vm inode
       unsigned long vm flags;
                                                                                                                          vm ons
55
                                                                                                                                                                     0×8059BB8
56
       rb node t vm rb;
                                                                                               mman
57
                                                                                                                           vm_next
                                                                                               mmap avl
63
       struct vm area struct *vm next share;
       struct vm area struct **vm pprev share;
                                                                                               mmap sem
65
       /* Function pointers to deal with this struct. */
67
       struct vm operations struct * vm ops;
                                                                                                                                                    Code
                                                                                                                        vm area struct
69
       /* Information about our backing store: */
                                                                                                                          vm end
70
       unsigned long vm pgoff:
                                                                                                                          vm start
72
       struct file * vm file;
                                                                                                                                                                     0×8048000
                                                                                                                          ym flaes
73
       unsigned long vm raend:
74
       void * vm private data:
                                                                                                                          vm inode
75 1:
                                                                                                                          vm ons
                                                                                                                          vm_next
                                                                                                                                                                     0×0000000
```

Processes Virtual Memory

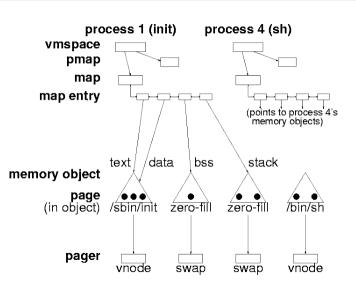




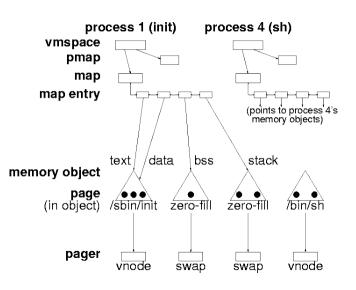
VM based on:



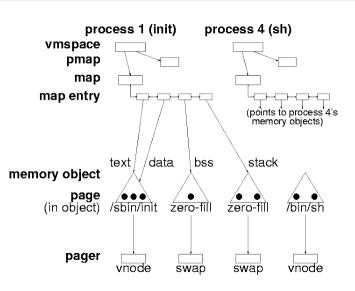
- VM based on:
  - 386BSD



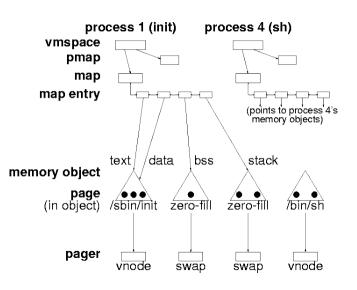
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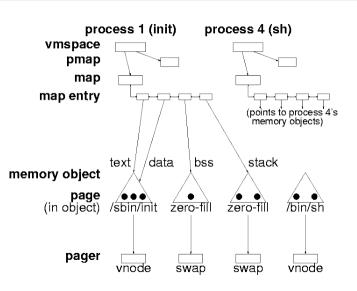
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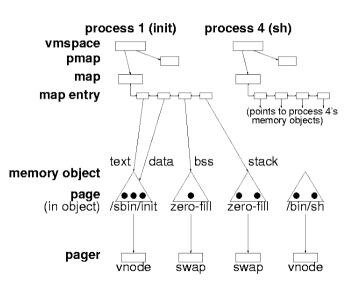
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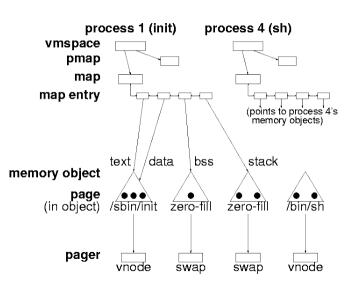
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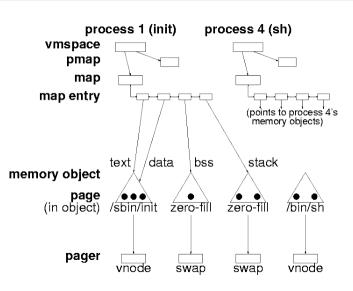
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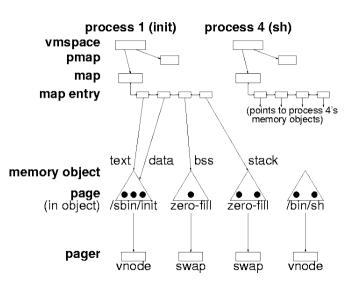
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- Minor modifications since then

# History: OpenIndiana

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Open source fork of OpenSolaris after Oracle take over

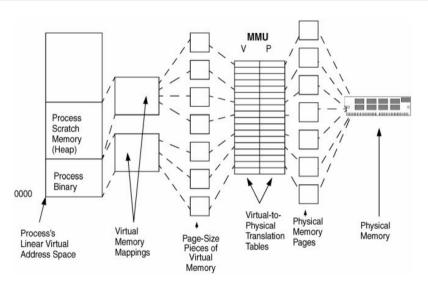
## History: OpenIndiana

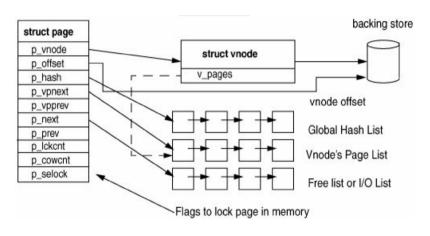
- Open source fork of OpenSolaris after Oracle take over
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- 3 Memory management to manage pages is basically swapping and demand paging





 Page table structure different from x86 hardware page table structure

#### Differences

We have found some significant differences so far:

#### **Differences**

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 What happens when the kernel runs out of memory?

#### Differences

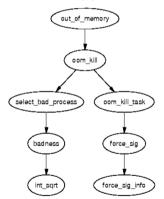
We have found some significant differences so far:

- What happens when the kernel runs out of memory?
- What are the copy-on-write mechanisms?

# What happens when the kernel runs out of memory?

#### Linux:

Start killing processes



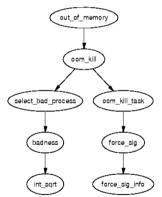
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#### Linux:

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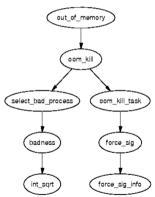
Panic!



# What happens when the kernel runs out of memory?

#### Linux:

Start killing processes



#### NetBSD:

Panic!

#### OpenIndiana:

- Periodically checks kernel space, and "snaps" data to user space if kernel space is low
- If kernel runs out of memory, crashes as far as I can tell

# What are the copy-on-write mechanisms?

#### Linux:

Page-based copy

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#### Linux:

Page-based copy

#### OpenIndiana:

Anonymous maps

# What are the copy-on-write mechanisms?

Linux:

Page-based copy

OpenIndiana:

Anonymous maps

NetBSD:

Copied SunOS/Solaris

## Summary

- 1 Literature Review
  - High-level design
  - Differences
- 2 Experimental Design
- 3 Implementation

#### References

UVM dissertation:

http://vorpal.math.drexel.edu/course/opsys2/uvm-project/uvm.pdf

• UVM paper:

https://www.usenix.org/legacy/event/usenix99/full\_papers/cranor/cranor.pdf

• UBC paper:

https://www.usenix.org/legacy/publications/library/proceedings/ usenix2000/freenix/silvers.html

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- McDougall, Richard, and Jim Mauro. Solaris internals: Solaris 10 and OpenSolaris kernel architecture. Pearson Education, 2006.

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- NetBSD data structure diagram from:
  - http://usenix.org/legacy/publications/library/proceedings/usenix99/full\_papers/cranor/cranor\_html/index.html
  - Linux vm\_area\_struct source from:
  - Linux data structures diagram from:
    - Linux OOM diagram from:
  - Linux OOM diagram from: ???
  - Solaris VM diagram: McDougall, Richard, and Jim Mauro. Solaris internals: Solaris 10 and OpenSolaris kernel architecture. Pearson Education, 2006.

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