Virtual Memory A Project for CS854

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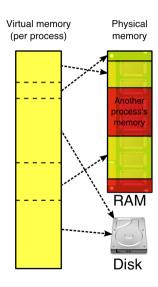
Abstract

In short:

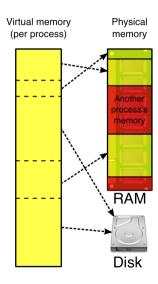
Abstract

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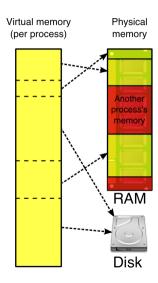
• We propose to study virtual memory!



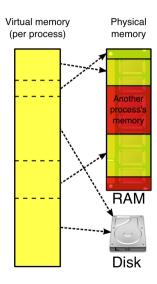
• (on x86) Instructions operate on virtual addresses



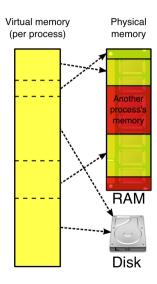
- (on x86) Instructions operate on virtual addresses
- Data may be stored:



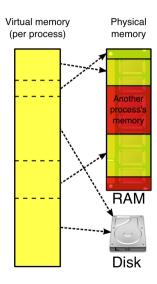
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 - $\bullet \ \ \text{Maps virtual} \to \text{physical addresses}$

Our proposal has 3 parts:

1 Literature Review

- 1 Literature Review
- 2 Experimental Design

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- 2 Experimental Design
- 3 Implementation

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5/20

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1 Linux

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- 2 NetBSD

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For each OS, we wish to answer the following questions:

How is physical memory managed?

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- How is memory freed?

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 - What happens when the kernel runs out of memory?

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- How is physical memory managed?
- Are there data structures for physical pages, separate from the page tables?
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- How is memory freed?
 - What happens when the kernel runs out of memory?
- Do they do anything special on Non-Uniform Memory Access (NUMA) architectures?

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- Design simple experiments to test this hypothesis
- Example:
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 - Test performance

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Proposal: Implementation

Optional

- Optional
- Implement a memory management system for KOS

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 - Lit review
 - Experiments

High-level design

Now we'll summarize the VM design of:

- Linux
- NetBSD
- OpenIndiana

High-level: Linux

vm_area_struct

```
44 struct vm area struct {
45
       struct mm struct * vm mm:
                                                                         task struct
46
       unsigned long vm start;
47
       unsigned long vm end:
                                                                                                mm struct
                                                                                                                        vm area struct
49
                                                                          mm
                                                                                                                           vm end
                                                                                               count
50
       /* linked list of VM areas per task, sorted by address */
                                                                                                                           vm_start
51
       struct vm area struct *vm next:
                                                                                               pgd
                                                                                                                                                    Data
52
                                                                                                                          ym flags
53
       pgprot t vm page prot;
                                                                                                                          vm inode
       unsigned long vm flags;
                                                                                                                          vm ons
55
                                                                                                                                                                     0×8059BB8
56
       rb node t vm rb;
                                                                                               mman
57
                                                                                                                           vm_next
                                                                                               mmap avl
63
       struct vm area struct *vm next share;
       struct vm area struct **vm pprev share;
                                                                                               mmap sem
65
       /* Function pointers to deal with this struct. */
67
       struct vm operations struct * vm ops;
                                                                                                                                                    Code
                                                                                                                        vm area struct
69
       /* Information about our backing store: */
                                                                                                                          vm end
70
       unsigned long vm pgoff:
                                                                                                                          vm start
72
       struct file * vm file;
                                                                                                                                                                     0×8048000
                                                                                                                          ym flaes
73
       unsigned long vm raend:
74
       void * vm private data:
                                                                                                                          vm inode
75 1:
                                                                                                                          vm ons
                                                                                                                          vm_next
                                                                                                                                                                     0×0000000
```

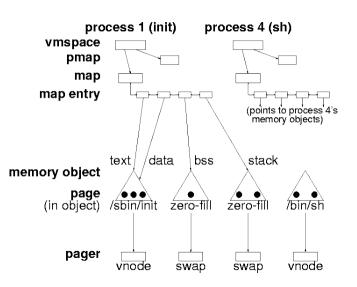
Processes Virtual Memory

 Based on 386BSD, 4.4BSD-Lite

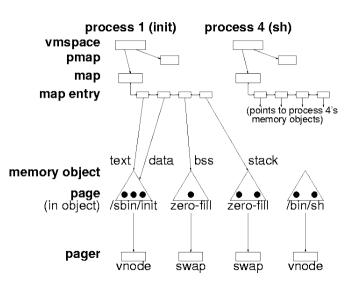
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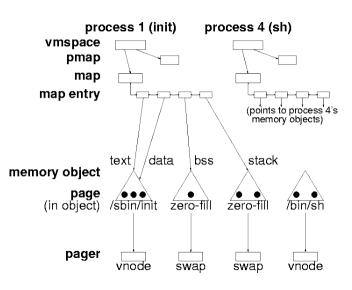
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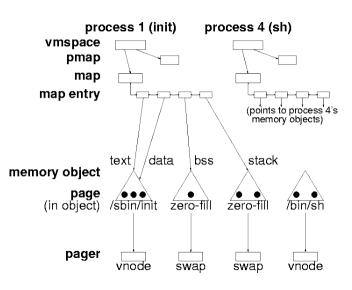
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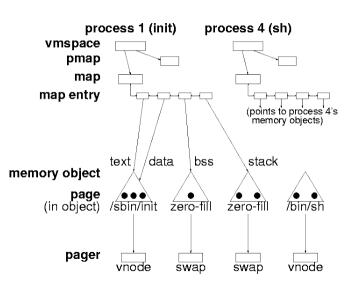
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- Minor modifications since then

History: OpenIndiana

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Open source fork of OpenSolaris after Oracle take over

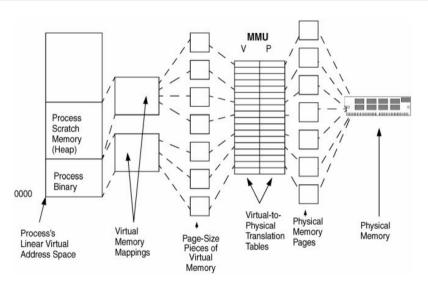
History: OpenIndiana

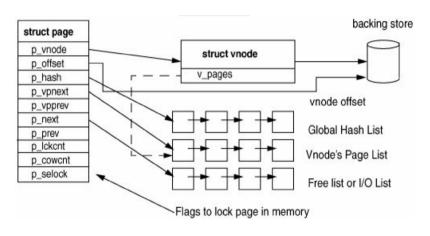
- Open source fork of OpenSolaris after Oracle take over
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- 2 Hardware MMU maps pages to physical memory using platform-specific translation tables
- 3 Memory management to manage pages is basically swapping and demand paging





 Page table structure different from x86 hardware page table structure

Differences

We have found some significant differences so far:

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 What happens when the kernel runs out of memory?

Differences

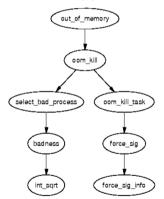
We have found some significant differences so far:

- What happens when the kernel runs out of memory?
- What are the copy-on-write mechanisms?

What happens when the kernel runs out of memory?

Linux:

Start killing processes



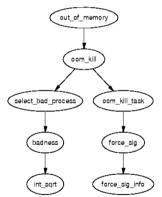
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NetBSD:

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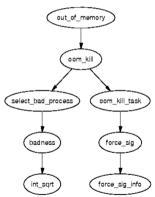
Panic!



What happens when the kernel runs out of memory?

Linux:

Start killing processes



NetBSD:

Panic!

OpenIndiana:

- Periodically checks kernel space, and "snaps" data to user space if kernel space is low
- If kernel runs out of memory, crashes as far as I can tell

What are the copy-on-write mechanisms?

Linux:

Page-based copy

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Linux:

Page-based copy

OpenIndiana:

Anonymous maps

What are the copy-on-write mechanisms?

Linux:

Page-based copy

OpenIndiana:

Anonymous maps

NetBSD:

Copied SunOS/Solaris

Summary

- 1 Literature Review
 - High-level design
 - Differences
- 2 Experimental Design
- 3 Implementation

References

UVM dissertation:

http://vorpal.math.drexel.edu/course/opsys2/uvm-project/uvm.pdf

• UVM paper:

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• UBC paper:

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- McDougall, Richard, and Jim Mauro. Solaris internals: Solaris 10 and OpenSolaris kernel architecture. Pearson Education, 2006.

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- NetBSD data structure diagram from:
 - http://usenix.org/legacy/publications/library/proceedings/usenix99/full_papers/cranor/cranor_html/index.html
 - Linux vm_area_struct source from:
 - Linux data structures diagram from:
 - Linux OOM diagram from:
 - Linux OOM diagram from: ???
 - Solaris VM diagram: McDougall, Richard, and Jim Mauro. Solaris internals: Solaris 10 and OpenSolaris kernel architecture. Pearson Education, 2006.

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