Virtual Memory A Project for CS854

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Our proposal has 3 parts:

1 Literature Review

- 1 Literature Review
- 2 Experimental Design

- 1 Literature Review
- 2 Experimental Design
- 3 Implementation

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- 2 Experimental Design
- 3 Implementation

We wish to investigate the following operating systems:

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1 Linux

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For each OS, we wish to answer the following questions:

How is physical memory managed?

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- How is physical memory managed?
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- How are contiguous regions of memory managed?

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- How is physical memory managed?
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- How are contiguous regions of memory managed?
- How is memory freed?

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- How is physical memory managed?
- Are there data structures for physical pages, separate from the page tables?
- How are contiguous regions of memory managed?
- How is memory freed?
 - What happens when the kernel runs out of memory?

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- How is physical memory managed?
- Are there data structures for physical pages, separate from the page tables?
- How are contiguous regions of memory managed?
- How is memory freed?
 - What happens when the kernel runs out of memory?
- Do they do anything special on Non-Uniform Memory Access (NUMA) architectures?

- 1 Literature Review
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Proposal: Experimental Design

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Make a testable hypothesis based on lit. review

Proposal: Experimental Design

- Make a testable hypothesis based on lit. review
- Design simple experiments to test this hypothesis

- 1 Literature Review
- 2 Experimental Design
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- 1 Literature Review
- 2 Experimental Design
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Proposal: Implementation

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Proposal: Implementation

 Implement a memory management system for KOS

Progress: high-level

We have made some progress:

- Linux
- OpenBSD
- OpenIndiana

Progress: Linux

vm_area_struct

```
44 struct vm area struct {
45
       struct mm struct * vm mm:
                                                                         task struct
46
       unsigned long vm start;
47
       unsigned long vm end:
                                                                                                mm struct
                                                                                                                        vm area struct
49
                                                                          mm
                                                                                                                          vm end
                                                                                               count
50
       /* linked list of VM areas per task, sorted by address */
                                                                                                                          vm_start
51
       struct vm area struct *vm next:
                                                                                               pgd
                                                                                                                                                    Data
52
                                                                                                                          vm flags
53
       pgprot t vm page prot;
                                                                                                                          vm inode
       unsigned long vm flags;
                                                                                                                          vm ons
55
                                                                                                                                                                     0×8059BB8
56
       rb node t vm rb;
                                                                                               mman
57
                                                                                                                          vm_next
                                                                                               mmap avl
63
       struct vm area struct *vm next share;
       struct vm area struct **vm pprev share;
                                                                                               mmap sem
65
       /* Function pointers to deal with this struct. */
67
       struct vm operations struct * vm ops;
                                                                                                                                                    Code
                                                                                                                        vm area struct
69
       /* Information about our backing store: */
                                                                                                                          vm end
70
       unsigned long vm pgoff:
                                                                                                                          vm start
72
       struct file * vm file;
                                                                                                                                                                     0×8048000
                                                                                                                          ym flaes
73
       unsigned long vm raend:
74
       void * vm private data:
                                                                                                                          vm inode
75 1:
                                                                                                                          vm ons
                                                                                                                          vm_next
                                                                                                                                                                     0×0000000
```

Processes Virtual Memory

386BSD↓NetBSD↓OpenBSD

386BSD↓NetBSD↓OpenBSD

VM based on NetBSD

```
386BSD↓NetBSD↓OpenBSD
```

- VM based on NetBSD
 - Rewritten in 1998

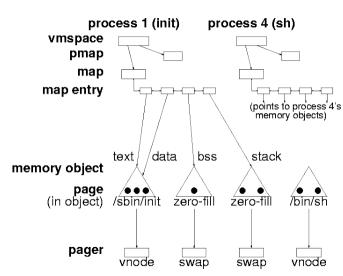
```
    386BSD

 NetBSD
 OpenBSD
```

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Progress: OpenIndiana

- Open source fork of OpenSolaris after Oracle take over
- Stewarded by the Illumos Foundation
- VM uses the ast package by AT&T, written by Kiem-Phong Vo
- Based on paper "Vmalloc: A General and Efficient Memory Allocator"

Progress: OpenIndiana

- Legacy malloc function is old, has shortcomings
- 2 malloc not designed for modern environments
- 3 Vmalloc a memory allocation library that is flexible and allows a wide range of memory operations
 - Regions to organize memory
 - Obtain memory by application definable disciplines
 - 3 Customize memory management

Differences

We have found some significant differences so far:

- What happens when the kernel runs out of memory
- How the kernel accesses user memory
- Copy-on-write mechanisms

What happens when the kernel runs out of memory

Linux:

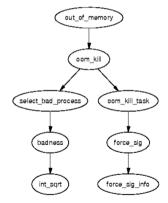
BSD:

Start killing processes

Panic!

OpenIndiana:

• ???



How the kernel accesses user memory

Linux:

 Map all of physical memory into kernel's virtual address space

BSD:

• ???

OpenIndiana:

• ???

Copy-on-write mechanisms

Linux:

• ???

BSD:

Anonymous maps

OpenIndiana:

• ???

Summary

- 1 Literature Review
 - Progress: High-level
 - Progress: Differences
- 2 Experimental Design
- 3 Implementation

References

UVM dissertation:

http://vorpal.math.drexel.edu/course/opsys2/uvm-project/uvm.pdf

Understanding the Linux Virtual Memory Manager
 https://www.kernel.org/doc/gorman/html/understand/index.html

Vmalloc: A General and Efficient Memory Allocator:

http://onlinelibrary.wiley.com/doi/10.1002/%28SICI%291097-024X% 28199603%2926:3%3C357::AID-SPE15%3E3.0.CO;2-%23/abstract

Attribution

• OpenBSD data structure diagram from: http://usenix.org/legacy/publications/library/proceedings/usenix99/full_papers/cranor/cranor_html/index.html

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