

## **Preview Fall 2019 CS120 Project 2. Multiple Access**

**Handout Date: Oct. 9, 2019 Due Date: Nov. 10, 2019**

(16 points + 6 points)

**Please read the following instructions carefully:**

- This project is to be completed by each group **individually**.
- Submit your code through Blackboard. The submission is performed by one of the group members.
- Each group needs to submit the code **once and only once**. Immediately after TAs' checking.

### **Overview**

This project will augment the physical data communication link from Project 1 to support mutual communication among multiple nodes. The design goal of this project is very similar to Ethernet. We call it *Athernet*. As Figure 1 shows, a core component of *Athernet* is the mechanism to manage the transmission of each transceiver so that the communication medium can be shared efficiently, i.e., the Medium Access Control (MAC) protocol.

In order to finish this project and remaining projects, the following necessary and sufficient accessories are provided. Each group is eligible to borrow one set from TAs, which includes:

1. Tee\*6 <https://item.jd.com/5005384.html>
2. USB Sound Card\*3 <https://item.jd.com/1804882.html>
3. Line\*6 <https://item.jd.com/1192674.html>

Contact TAs to borrow the accessory set before Oct. 25.

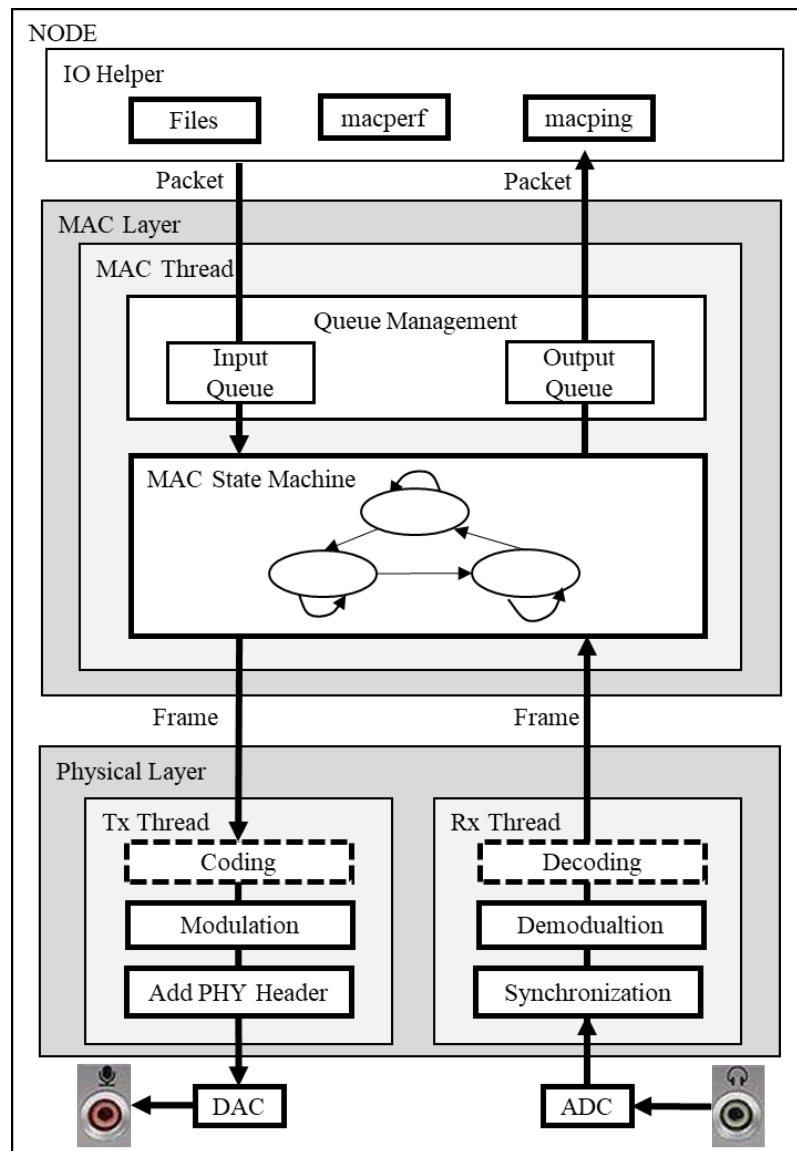


Figure 1 Project2 Overview

(Hierarchy of Tasks. Tasks are graded according to their hierarchy. A full score of one part automatically guarantees the full score of its subparts. In this project, Part 1 is a subpart of Part 2 and Part 3. Part 2 is a subpart of Part 3. The hierarchy is denoted as Part1<Part2<Part3.)

## Part 1. (4 points) From Wireless to Wired

The main purpose of moving from wireless link to wired link is to reduce the heterogeneity and complexity in acoustic hardware, e.g., ringing effect from mechanical vibration, distortions from nonlinear amplifiers, uneven response, echos, etc. The method to establish a wired link between two nodes is connecting their 3.5mm analog audio interfaces, e.g., connecting the MIC port to the Speaker port (See the red line in Figure 2). Note that the signal flowing between the two ports is no longer an acoustic signal. It is an electric wave generated by the DAC.

The electric signal output from the Speaker port is normally used to stimulate the mechanical hardware, e.g., the speaker, to generate sound. On the other hand, the ADC in the MIC port will transform any the input electric signal into digital samples, no matter the signal is generated by a DAC from other device (red) or a microphone (black). Generally speaking, through the wired connection between DAC and ADC, the received signal contains much fewer distortions than the signal in Project1. Your code of Project1 should work seamlessly and very likely much better in the wired situation.

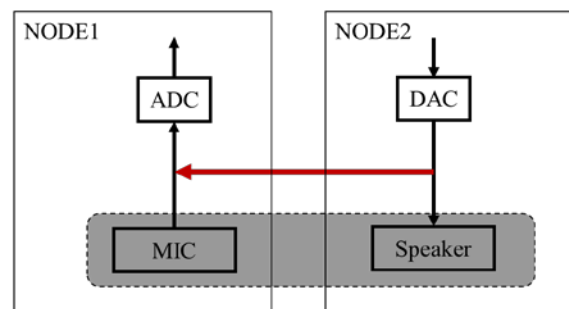


Figure 2 Wire Connection

In this part, your task is to set up the wired link between two nodes. This task is almost the same as Project1.Part3, but we expect higher performance (>5kbps).

### Checkpoints:

The group provides two devices: NODE1 and NODE2, connected by audio cables.

CK(4 points). TAs provide a binary file "INPUT.bin" which contains 6250 Bytes. NODE1 sends bits according to this file. NODE2 stores the received bytes into a binary file "OUTPUT.bin".

The transmission must be finished within 10 seconds.

Transmission Time	Points
<10s	100%
>10s	0%

TAs compare the difference between INPUT.bin and OUTPUT.bin:

Accuracy	Points
<80%	0%
80% to 99%	80%
>99%	100%

Tips:

- You are recommended to unplug the chargers/power adapters of your devices when doing wired experiments. The reason is that different AC-to-DC adapters (especially the ones without ground plugin) may have different voltages of their “ground”, so do the “ground” of the charging devices. When interconnecting two devices with wires, they may have different “ground”. In some cases, the problem just results in biased DC levels in the received signal from connected devices. Sometimes it may bring harmful damages. If you have to connect two devices with different “ground” (e.g., a charging laptop and a desktop), you should be very careful.
- The throughput can be increased by modifying your physical layer (PHY) of project 1, e.g., decrease the time for each symbol. You can also develop a new PHY suitable for the wired situation.
- You can design shorter preamble to reduce the header overhead.
- You can use the audio card of your own devices. Sometime they have better performance (in terms of the quality of the generated/received signal) than the offered ones.
- In some laptops and smartphones, the MIC in and speaker ports are combined into a single port, you may need an adapter to split them.

## Part 2. (4 points) A Simple Reliable Link

To build a CSMA MAC from ground up, a good warm up is a simple send-and-pray protocol. A state machine is helpful in describing a protocol. The state machine of the send-and-pray protocol is shown in Figure 3. The node can only be in one of the states at any moment. The transition of the state is driven by some events. The arrows connecting states identify the event and the transition direction.

From the state machine in Figure, it is clear that the node switches between Tx and Rx states, meaning that the node cannot transmit and receive at the same time (i.e., half duplex). Supporting full duplex may require more complex designs. Half duplex is the assumption of *Athenet* nodes.

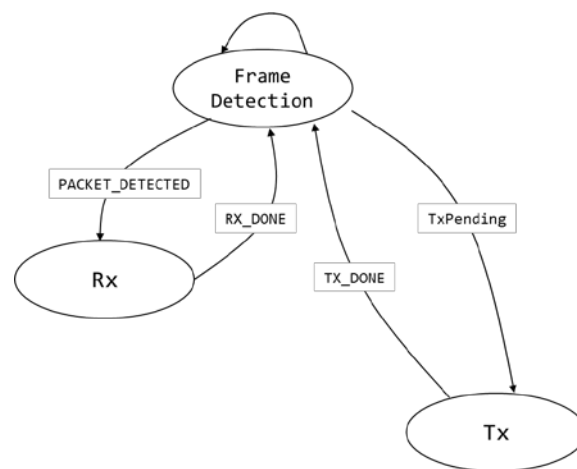


Figure 3 Send-and-pray State Machine

The goal of this task is to help you polish your physical layer code. Making a clear abstraction of the functionalities of the physical layer from project 1 helps a lot in implementing MAC protocols. For example, the following primitives are necessary ones that the physical layer interface must support for implementing the above simple protocol:

- Function: `PHYSend(MACFrame)`, is provided by `PHY.Tx` thread. MAC thread can call this function to send one MAC frame. The finishing of the transmission is notified by `TX_DONE` event.
- Event: `TX_DONE`, is issued by `PHY.Tx` thread, MAC thread can capture this event and process some necessary jobs (e.g., dequeuer of the FIFO). Then the state proceeds from Tx to FrameDetection.
- Event: `PACKET_DETECTED`. If a packet header is detected in `PHY.Rx` thread, the thread will issue this event. Once the MAC thread receives this event, it will switch to the Rx state.
- Buffer: `PHYRxFrame`, is shared between `PHY.Rx` thread and MAC thread. Once a frame is correctly received by the physical layer, `PHY.Rx` thread will put the frame into this buffer.
- Event: `RX_DONE`, is issued by `PHY.Rx` thread, it notifies MAC thread to take the frame in `PHYRxFrame` and proceed to FrameDetection state.

It is easy to extend the send-and-pray protocol to a more effective one: stop-and-wait ACK protocol. “Upgrading” the physical link in Project1 with ACK ability will provide reliability in data communication. A reference stop-and-wait protocol is shown in Figure 4.

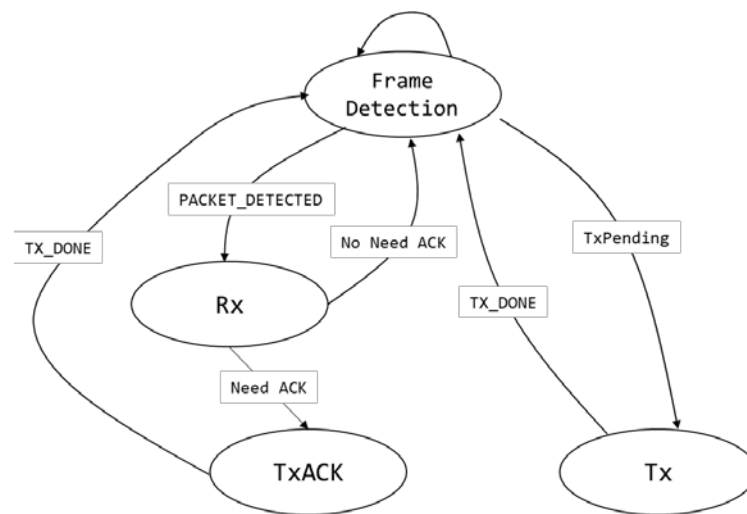


Figure 4 Stop-and-wait State Machine

The state machine contains more states and events than send-and-pray protocol. One obvious difference is that, after correctly receiving a frame, the node needs to reply an ACK immediately. Note that the state machine does not contain full details of this protocol. For example, when handling TX\_DONE event, MAC thread should set timeout counter for waiting ACK, and if the ACK is not received, the packet should be retransmitted.

#### Checkpoints:

The group provides two devices: NODE1 and NODE2, connected by two cables.

CK1(2 points). Similar to Part 1. TAs provide a binary file “INPUT.bin” which contains 6250 bytes. NODE1 sends bits according to this file. NODE2 stores the received bytes into a binary file “OUTPUT.bin”.

The transmission must be finished within 10 seconds.

<10s	-0%
>10s	-100%

TAs compare the difference between INPUT.bin and OUTPUT.bin:

<100%	-100%
100%	-0%

CK2(2 points). Redo CK2. TAs can unplug one of the wires, and the transmitter should be able to identify the event and display “link error” (according to retransmission times).

#### Tips:

- You may want to use thread safe data structures in delivering data between threads.
- The inter-frame time and time-out time should be carefully designed (refer to Lec5 and Lec6).

### Part 3. (4 points) CSMA

In order to efficiently support multiple nodes in the shared communication medium, their transmissions should be carefully coordinated. There are many different MAC protocols. The one we choose in *Athern* is CSMA, which is widely used in Ethernet and Wi-Fi.

**CSMA Protocol.** The intuition of CSMA protocol can be simply put as “listen before talk”. The lecture and the textbook have covered it quite well. The state machine of CSMA protocol is shown in Figure 5. Check more information from references [1-5].

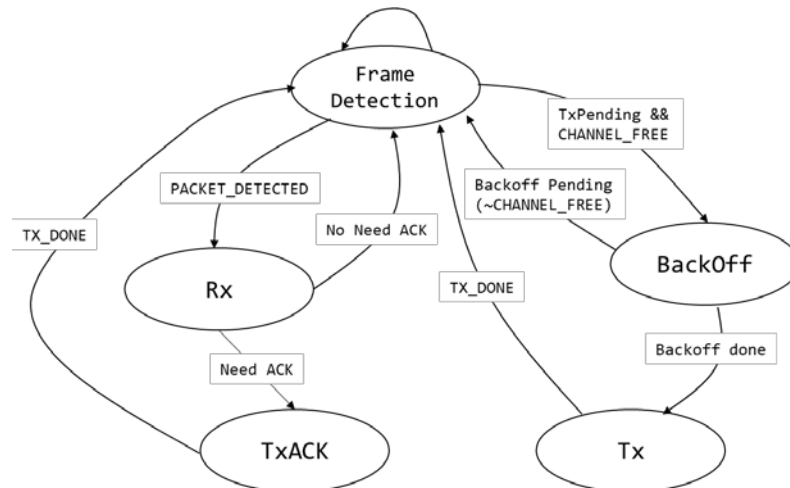


Figure 5 CSMA+ACK State Machine [1]

**Frame Format.** In order to support multiple nodes, each node must have a unique destination identifier, which contributes the destination and source address in MAC frames. Using 8 bits for each address field is suggested. A type field is very useful for differentiating different kinds of frames, e.g., ACKs and normal data frames. Using 4 bits for type field is suggested. Finally, the maximum MAC payload must be limited, otherwise, the network may have large delay (see Part5).

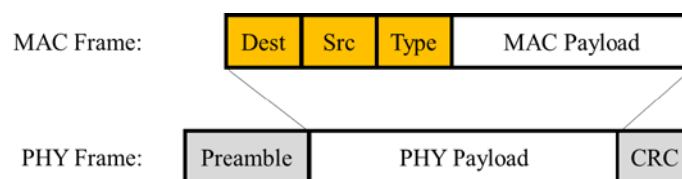


Figure 6 Frame Format

**Network with Jamming.** Using your accessories to connect 2 nodes like Part1, so the node is able to hear the transmissions from the other node. In addition to that, as shown in Figure 7, each node should be able to hear the output from a Jamming source, which intermittently plays high-power noise signals. As CSMA is enabled in the two nodes, the communication should be able to work under the silent intervals of the Jamming source. **XXXX how to connect jamming src, needs four sound card and two tee. Be aware of the voltage, change jamming source to 100+rand(100)**

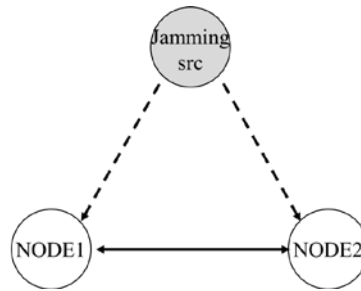


Figure 7 Network Topology

**Checkpoints:**

The group provides two devices: NODE1, NODE2 connected by wires.

TAs provide one Jamming source.

The group provides the interface to connect the jamming source.

The topology is shown in Figure 7.

CK1(2 points). “INPUT1to2.bin” of 6250 Bytes and “INPUT2to1.bin” of 5000 Bytes are provided.

NODE1: transmits INPUT1to2.bin to NODE2, store the received file to OUTPUT2to1.bin

NODE2: transmits INPUT2to1.bin to NODE1, store the received file to OUTPUT1to2.bin

Jamming Src: play “Jamming.wav”

Jamming Src starts first. The transmission of NODE2 must be started immediately after the start of NODE1. The transmission of NODE2 must be finished before NODE1. The received files must be error free. The rank of Part7 is determined by the transmission time of NODE1.

The transmission of NODE1 must be finished within 60 seconds.

<60s	100%
>60s	0%

TAs compare the difference between INPUT\*.bin and OUTPUT\*.bin:

100%	100%
<100%	0%

CK2(2 points). Redo CK1. TAs can unplug one of the wires, and transmitter should be able to identify the event and display “link error”.



**Part 4. (2 point) macperf Utility**

macperf utility is used to measure the throughput between two nodes in *Athenet*. The utility is similar to the *iperf* utility, but it is specialized for *Athenet* and works in a different layer.

The working flow of macperf utility:

- macperf is invoked by executing `macperf dest_address`
- By default, macperf frame is generated with random MAC payload (choose your own payload size) and the type field of the MAC frame is set to type: DATA.
- The sender tries its best to send out macperf packets
- The sender counts and prints the throughput on the screen every one second.

Checkpoints:

The group provides two devices: NODE1, NODE2. The topology is shown in Figure 7 without jamming source.

CK(2 points). TAs will check the functionality of macping utility in the following case:

NODE1: `macperf NODE2`

NODE2: `macperf NODE1`

TAs record the throughput of NODE1 and NODE3.

If NODE1 TH >1.5kbps && NODE2 TH>1.5kbps	100%
otherwise	0%

**Part 5. (2 points) macping Utility**

macping utility is used to measure the round trip delay between two nodes in *Athern*. The utility is similar to ping utility in current operating systems, but it is specialized for *Athern* and works in a different layer.

The working flow of macping utility:

- macping is invoked by executing `macping dest_address`
- A macping frame is generated with zero MAC payload and the type field of the MAC frame is set to type: `MACPING_REQ`.
- The macping frame is timestamped when sending into the PHY layer.
- The node with `dest_address` is responsible for automatically replying a frame targeting the sender and having its type field set to type: `MACPING_REPLY`.
- If the sender receives `MACPING_REPLY`, it calculates and prints the round trip delay on the screen.
- If the sender does not receive `MACPING_REPLY` after 2 seconds, the sender prints `TIMEOUT` on screen.

Checkpoints:

The group provides two devices: `NODE1`, `NODE2`.

The topology is shown in Figure 7.

CK(2 points). TAs will check the functionality of macping utility in the following case:

`NODE1: macping NODE2`

`NODE2: macperf NODE1`

TAs record the RTT of `NODE1` to `NODE2` and the throughput of `NODE3`

<code>NODE1 RTT &lt; 0.15s &amp;&amp; NODE2 TH &gt; 2.7kbps</code>	100%
<code>otherwise</code>	0%

(Tasks with “Optional” tag are optional tasks. The instructor is responsible for checking and grading the optional tasks. Contact the instructor to check if you have finished one or more of them)

### **Part 6. (Optional + 2 points) Collision Detection**

Network performance is determined by many factors. High throughput physical layer does not necessarily indicate high network performance. To further increase the efficiency of CSMA protocol, PHY.Tx thread can be aborted if PHY.Rx thread indicates high power, i.e., CSMA/CD.

Checkpoints:

Check settings are the same as Part3.

The throughput of NODE1 should be higher when collision detection is enabled.

**Part 7. (Optional + 4 points) Performance Rank**

Network systems always demand high performance, but high performance is hard to achieve. This task is to reward groups who will have achieved high throughput performance. In the first round, top 5 groups are selected according to their throughput performance in Part3. The second round will be held on Nov. 21 after the lecture, the final rank is determined by the performance in the second round.

**Checkpoints:**

Check settings are almost the same as Part3.

TAs provide one Jamming source.

The group provides the interface to connect the jamming source.

The topology is shown in Figure 7.

CK(4 points). "INPUT1to2.bin" of 2MB and "INPUT2to1.bin" of 1MB are provided.

NODE1: transmits INPUT1to2.bin to NODE2, store the received file to OUTPUT2to1.bin

NODE2: transmits INPUT2to1.bin to NODE1, store the received file to OUTPUT1to2.bin

Jamming Src: play "Jamming.wav"

The transmission of NODE2 must be started immediately after the start of NODE1. The transmission of NODE2 must be finished before NODE1. The received files must be error free. The rank is determined by the transmission time of NODE1.

Rank	Points
1	100%
2	50%
3	25%
4	12.5%
5	0%

## Reference and Useful Links

- [1] Sora: High Performance Software Radio Using General Purpose Multi-core Processors  
[https://www.usenix.org/legacy/events/nsdi09/tech/full\\_papers/tan/tan.pdf](https://www.usenix.org/legacy/events/nsdi09/tech/full_papers/tan/tan.pdf)
- [2] WARP CSMAMAC <https://warpproject.org/trac/wiki/CSMAMAC>
- [3] WARP CSMAMAC Src  
<https://warpproject.org/trac/browser/ResearchApps/MAC/CSMAMAC/csmaMac.c>
- [4] NS3 MAC Low [https://www.nsnam.org/doxygen/mac-low\\_8cc\\_source.html](https://www.nsnam.org/doxygen/mac-low_8cc_source.html)
- [5] GNURadio MAC Report <http://gnumac.wikispaces.com/file/view/final-project.pdf>