

# **Standard Code Library**

Shanghai Jiao Tong University

October, 2015

# Contents

<b>1</b>	<b>数论算法</b>	<b>3</b>
1.1	快速数论变换	3
1.2	多项式求逆	3
1.3	中国剩余定理	4
1.4	Miller Rabin	4
1.5	Pollard Rho	5
1.6	坚固的逆元	5
1.7	直线下整点个数	5
<b>2</b>	<b>数值算法</b>	<b>7</b>
2.1	快速傅立叶变换	7
2.2	单纯形法求解线性规划	7
2.3	自适应辛普森	9
<b>3</b>	<b>数据结构</b>	<b>10</b>
3.1	Splay 普通操作版	10
3.2	Splay 区间操作版	11
3.3	坚固的 Treap	15
3.4	k-d 树	18
3.5	树链剖分	19
3.5.1	点操作版本	19
3.5.2	链操作版本	22
3.6	Link-Cut-Tree	22
<b>4</b>	<b>图论</b>	<b>25</b>
4.1	强连通分量	25
4.2	点双连通分量	25
4.2.1	坚固的点双连通分量	25
4.3	2-SAT 问题	27
4.4	二分图最大匹配	28
4.4.1	Hungary 算法	28
4.4.2	Hopcroft Karp 算法	29
4.5	二分图最大权匹配	30
4.6	最大流	30
4.6.1	Dinic	30
4.6.2	ISAP	31
4.6.3	SAP	32
4.7	上下界网络流	33
4.7.1	无源汇的上下界可行流	33
4.7.2	有源汇的上下界可行流	33
4.7.3	有源汇的上下界最大流	33
4.7.4	有源汇的上下界最小流	33
4.8	最小费用最大流	33
4.8.1	稀疏图	33
4.8.2	稠密图	34
4.9	一般图最大匹配	35
4.10	无向图全局最小割	37
4.11	最小树形图	38
4.12	有根树的同构	39
4.13	度限制生成树	39
4.14	弦图相关	41
4.14.1	弦图的判定	41

4.14.2	弦图的团数	42
4.15	哈密尔顿回路 (ORE 性质的图)	43
<b>5</b>	<b>字符串</b>	<b>46</b>
5.1	模式串匹配	46
5.2	坚固的模式串匹配	46
5.3	AC 自动机	47
5.4	后缀数组	48
5.5	广义后缀自动机	48
5.6	Manacher 算法	49
5.7	回文树	49
5.8	循环串最小表示	50
<b>6</b>	<b>计算几何</b>	<b>51</b>
6.1	二维基础	51
6.1.1	点类	51
6.1.2	凸包	51
6.1.3	半平面交	52
6.1.4	最近点对	52
6.2	三维基础	53
6.2.1	点类	53
6.2.2	凸包	54
6.2.3	绕轴旋转	56
6.3	多边形	57
6.3.1	判断点在多边形内部	57
6.3.2	多边形内整点计数	57
6.4	圆	58
6.4.1	最小覆盖圆	58
6.4.2	最小覆盖球	58
6.4.3	多边形与圆的交面积	61
<b>7</b>	<b>其它</b>	<b>62</b>
7.1	STL 使用方法	62
7.1.1	nth_element	62
7.1.2	next_permutation	62
7.2	博弈论相关	62
7.2.1	巴什博弈	62
7.2.2	威佐夫博弈	62
7.2.3	阶梯博弈	62
7.2.4	图上删边游戏	62
7.3	Java Reference	63
<b>8</b>	<b>数学公式</b>	<b>65</b>
8.1	常用数学公式	65
8.1.1	求和公式	65
8.1.2	斐波那契数列	65
8.1.3	错排公式	65
8.1.4	莫比乌斯函数	65
8.1.5	Burnside 引理	66
8.1.6	五边形数定理	66
8.1.7	树的计数	66
8.1.8	欧拉公式	66
8.1.9	皮克定理	66
8.1.10	牛顿恒等式	67
8.2	平面几何公式	67
8.2.1	三角形	67
8.2.2	四边形	67
8.2.3	正 $n$ 边形	68
8.2.4	圆	68
8.2.5	棱柱	68
8.2.6	棱锥	68
8.2.7	棱台	69
8.2.8	圆柱	69
8.2.9	圆锥	69

8.2.10	圆台 . . . . .	69
8.2.11	球 . . . . .	69
8.2.12	球台 . . . . .	70
8.2.13	球扇形 . . . . .	70
8.3	立体几何公式 . . . . .	70
8.3.1	球面三角公式 . . . . .	70
8.3.2	四面体体积公式 . . . . .	70

# Chapter 1

## 数论算法

### 1.1 快速数论变换

使用条件及注意事项： $mod$  必须要是一个形如  $a2^b + 1$  的数， $pri$  表示  $mod$  的原根。

```
1  const int mod = 998244353;
2  const int pri = 3;
3  int prepare(int n) {
4      int len = 1;
5      for (; len <= 2 * n; len <<= 1);
6      for (int i = 0; i <= len; i++) {
7          e[0][i] = fpm(pri, (mod - 1) / len * i, mod);
8          e[1][i] = fpm(pri, (mod - 1) / len * (len - i), mod);
9      }
10     return len;
11 }
12 void DFT(int *a, int n, int f) {
13     for (int i = 0, j = 0; i < n; i++) {
14         if (i > j) std::swap(a[i], a[j]);
15         for (int t = n >> 1; (j ^= t) < t; t >>= 1);
16     }
17     for (int i = 2; i <= n; i <<= 1)
18         for (int j = 0; j < n; j += i)
19             for (int k = 0; k < (i >> 1); k++) {
20                 int A = a[j + k];
21                 int B = (long long)a[j + k + (i >> 1)] * e[f][n / i * k] % mod;
22                 a[j + k] = (A + B) % mod;
23                 a[j + k + (i >> 1)] = (A - B + mod) % mod;
24             }
25     if (f == 1) {
26         long long rev = fpm(n, mod - 2, mod);
27         for (int i = 0; i < n; i++) {
28             a[i] = (long long)a[i] * rev % mod;
29         }
30     }
31 }
```

### 1.2 多项式求逆

使用条件及注意事项：求一个多项式在模意义下的逆元。

```
1  void getInv(int *a, int *b, int n) {
2      static int tmp[100000];
3      std::fill(b, b + n, 0);
4      b[0] = fpm(a[0], mod - 2, mod);
5      for (int c = 1; c <= n; c <<= 1) {
6          for (int i = 0; i < c; i++) tmp[i] = a[i];
7          std::fill(b + c, b + (c << 1), 0);
```

```

8         std::fill(tmp + c, tmp + (c << 1), 0);
9         DFT(tmp, c << 1, 0);
10        DFT(b, c << 1, 0);
11        for (int i = 0; i < (c << 1); i++) {
12            b[i] = (long long)(2 - (long long)tmp[i] * b[i] % mod + mod) * b[i] % mod;
13        }
14        DFT(b, c << 1, 1);
15        std::fill(b + c, b + (c << 1), 0);
16    }
17 }

```

### 1.3 中国剩余定理

使用条件及注意事项：模数可以不互质。

```

1 bool solve(int n, std::pair<long long, long long> input[],
2           std::pair<long long, long long> &output) {
3     output = std::make_pair(1, 1);
4     for (int i = 0; i < n; ++i) {
5         long long number, useless;
6         euclid(output.second, input[i].second, number, useless);
7         long long divisor = std::__gcd(output.second, input[i].second);
8         if ((input[i].first - output.first) % divisor) {
9             return false;
10        }
11        number = (input[i].first - output.first) / divisor;
12        fix(number, input[i].second);
13        output.first += output.second * number;
14        output.second *= input[i].second / divisor;
15        fix(output.first, output.second);
16    }
17    return true;
18 }

```

### 1.4 Miller Rabin

```

1 const int BASE[12] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37};
2
3 bool check(const long long &prime, const long long &base) {
4     long long number = prime - 1;
5     for (; ~number & 1; number >= 1);
6     long long result = power_mod(base, number, prime);
7     for (; number != prime - 1 && result != 1 && result != prime - 1; number <= 1) {
8         result = multiply_mod(result, result, prime);
9     }
10    return result == prime - 1 || (number & 1) == 1;
11 }
12
13 bool miller_rabin(const long long &number) {
14     if (number < 2) {
15         return false;
16     }
17     if (number < 4) {
18         return true;
19     }
20     if (~number & 1) {
21         return false;
22     }
23     for (int i = 0; i < 12 && BASE[i] < number; ++i) {
24         if (!check(number, BASE[i])) {

```

```

25         return false;
26     }
27 }
28 return true;
29 }

```

## 1.5 Pollard Rho

```

1  long long pollard_rho(const long long &number, const long long &seed) {
2      long long x = rand() % (number - 1) + 1, y = x;
3      for (int head = 1, tail = 2; ; ) {
4          x = multiply_mod(x, x, number);
5          x = add_mod(x, seed, number);
6          if (x == y) {
7              return number;
8          }
9          long long answer = std::__gcd(abs(x - y), number);
10         if (answer > 1 && answer < number) {
11             return answer;
12         }
13         if (++head == tail) {
14             y = x;
15             tail <= 1;
16         }
17     }
18 }
19
20 void factorize(const long long &number, std::vector<long long> &divisor) {
21     if (number > 1) {
22         if (miller_rabin(number)) {
23             divisor.push_back(number);
24         } else {
25             long long factor = number;
26             for (; factor >= number;
27                 factor = pollard_rho(number, rand() % (number - 1) + 1));
28             factorize(number / factor, divisor);
29             factorize(factor, divisor);
30         }
31     }
32 }

```

## 1.6 坚固的逆元

```

1  long long inverse(const long long &x, const long long &mod) {
2      if (x == 1) {
3          return 1;
4      } else {
5          return (mod - mod / x) * inverse(mod % x, mod) % mod;
6      }
7  }

```

## 1.7 直线下整点个数

```

1  long long solve(const long long &n, const long long &a,
2                 const long long &b, const long long &m) {
3      if (b == 0) {
4          return n * (a / m);
5      }

```

```
6   if (a >= m) {  
7       return n * (a / m) + solve(n, a % m, b, m);  
8   }  
9   if (b >= m) {  
10      return (n - 1) * n / 2 * (b / m) + solve(n, a, b % m, m);  
11  }  
12  return solve((a + b * n) / m, (a + b * n) % m, m, b);  
13 }
```



## Chapter 2

# 数值算法

### 2.1 快速傅立叶变换

```
1  int prepare(int n) {
2      int len = 1;
3      for (; len <= 2 * n; len <= 1);
4      for (int i = 0; i < len; i++) {
5          e[0][i] = Complex(cos(2 * pi * i / len), sin(2 * pi * i / len));
6          e[1][i] = Complex(cos(2 * pi * i / len), -sin(2 * pi * i / len));
7      }
8      return len;
9  }
10
11 void DFT(Complex *a, int n, int f) {
12     for (int i = 0, j = 0; i < n; i++) {
13         if (i > j) std::swap(a[i], a[j]);
14         for (int t = n >> 1; (j ^= t) < t; t >>= 1);
15     }
16     for (int i = 2; i <= n; i <= 1)
17         for (int j = 0; j < n; j += i)
18             for (int k = 0; k < (i >> 1); k++) {
19                 Complex A = a[j + k];
20                 Complex B = e[f][n / i * k] * a[j + k + (i >> 1)];
21                 a[j + k] = A + B;
22                 a[j + k + (i >> 1)] = A - B;
23             }
24     if (f == 1) {
25         for (int i = 0; i < n; i++)
26             a[i].a /= n;
27     }
28 }
```

### 2.2 单纯形法求解线性规划

使用条件及注意事项：返回结果为  $\max\{c_{1 \times m} \cdot x_{m \times 1} \mid x_{m \times 1} \geq 0_{m \times 1}, a_{n \times m} \cdot x_{m \times 1} \leq b_{n \times 1}\}$

```
1  std::vector<double> solve(const std::vector<std::vector<double>> &a,
2                          const std::vector<double> &b, const std::vector<double> &c) {
3      int n = (int)a.size(), m = (int)a[0].size() + 1;
4      std::vector<std::vector<double>> value(n + 2, std::vector<double>(m + 1));
5      std::vector<int> index(n + m);
6      int r = n, s = m - 1;
7      for (int i = 0; i < n + m; ++i) {
8          index[i] = i;
9      }
10     for (int i = 0; i < n; ++i) {
11         for (int j = 0; j < m - 1; ++j) {
```

```

12         value[i][j] = -a[i][j];
13     }
14     value[i][m - 1] = 1;
15     value[i][m] = b[i];
16     if (value[r][m] > value[i][m]) {
17         r = i;
18     }
19 }
20 for (int j = 0; j < m - 1; ++j) {
21     value[n][j] = c[j];
22 }
23 value[n + 1][m - 1] = -1;
24 for (double number; ; ) {
25     if (r < n) {
26         std::swap(index[s], index[r + m]);
27         value[r][s] = 1 / value[r][s];
28         for (int j = 0; j <= m; ++j) {
29             if (j != s) {
30                 value[r][j] *= -value[r][s];
31             }
32         }
33         for (int i = 0; i <= n + 1; ++i) {
34             if (i != r) {
35                 for (int j = 0; j <= m; ++j) {
36                     if (j != s) {
37                         value[i][j] += value[r][j] * value[i][s];
38                     }
39                 }
40                 value[i][s] *= value[r][s];
41             }
42         }
43     }
44     r = s = -1;
45     for (int j = 0; j < m; ++j) {
46         if (s < 0 || index[s] > index[j]) {
47             if (value[n + 1][j] > eps || value[n + 1][j] > -eps && value[n][j] > eps) {
48                 s = j;
49             }
50         }
51     }
52     if (s < 0) {
53         break;
54     }
55     for (int i = 0; i < n; ++i) {
56         if (value[i][s] < -eps) {
57             if (r < 0
58                 || (number = value[r][m] / value[r][s] - value[i][m] / value[i][s]) < -eps
59                 || number < eps && index[r + m] > index[i + m]) {
60                 r = i;
61             }
62         }
63     }
64     if (r < 0) {
65         // Solution is unbounded.
66         return std::vector<double>();
67     }
68 }
69 if (value[n + 1][m] < -eps) {
70     // No solution.
71     return std::vector<double>();
72 }
73 std::vector<double> answer(m - 1);
74 for (int i = m; i < n + m; ++i) {

```

```
75         if (index[i] < m - 1) {
76             answer[index[i]] = value[i - m][m];
77         }
78     }
79     return answer;
80 }
```

## 2.3 自适应辛普森

```
1  double area(const double &left, const double &right) {
2      double mid = (left + right) / 2;
3      return (right - left) * (calc(left) + 4 * calc(mid) + calc(right)) / 6;
4  }
5
6  double simpson(const double &left, const double &right,
7                const double &eps, const double &area_sum) {
8      double mid = (left + right) / 2;
9      double area_left = area(left, mid);
10     double area_right = area(mid, right);
11     double area_total = area_left + area_right;
12     if (std::abs(area_total - area_sum) < 15 * eps) {
13         return area_total + (area_total - area_sum) / 15;
14     }
15     return simpson(left, mid, eps / 2, area_left)
16         + simpson(mid, right, eps / 2, area_right);
17 }
18
19 double simpson(const double &left, const double &right, const double &eps) {
20     return simpson(left, right, eps, area(left, right));
21 }
```

# Chapter 3

## 数据结构

### 3.1 Splay 普通操作版

使用条件及注意事项：

1. 插入  $x$  数
2. 删除  $x$  数 (若有多个相同的数，因只删除一个)
3. 查询  $x$  数的排名 (若有多个相同的数，因输出最小的排名)
4. 查询排名为  $x$  的数
5. 求  $x$  的前驱 (前驱定义为小于  $x$ ，且最大的数)
6. 求  $x$  的后继 (后继定义为大于  $x$ ，且最小的数)

```
1  int pred(int x) {
2      splay(x, -1);
3      for (x = c[x][0]; c[x][1]; x = c[x][1]);
4      return x;
5  }
6  int succ(int x) {
7      splay(x, -1);
8      for (x = c[x][1]; c[x][0]; x = c[x][0]);
9      return x;
10 }
11 void remove(int x) {
12     if (b[x] > 1) {b[x]--; splay(x, -1); return;}
13     splay(x, -1);
14     if (!c[x][0] && !c[x][1]) rt = 0;
15     else if (c[x][0] && !c[x][1]) f[rt = c[x][0]] = -1;
16     else if (!c[x][0] && c[x][1]) f[rt = c[x][1]] = -1;
17     else{
18         int t = pred(x); f[rt = c[x][0]] = -1;
19         c[t][1] = c[x][1]; f[c[x][1]] = t;
20         splay(c[x][1], -1);
21     }
22     c[x][0] = c[x][1] = f[x] = d[x] = s[x] = b[x] = 0;
23 }
24 int find(int z) {
25     int x=rt;
26     while (d[x]!=z)
27         if (c[x][d[x]<z]) x=c[x][d[x]<z];
28         else break;
29     return x;
30 }
31 void insert(int z) {
32     if (!rt) {
33         f[rt = ++size] = -1;
34         d[size] = z; b[size] = 1;
```

```

35     splay(size, -1);
36     return;
37 }
38 int x = find(z);
39 if (d[x] == z) {
40     b[x]++;
41     splay(x, -1);
42     return;
43 }
44 c[x][d[x]<z] = ++size; f[size] = x;
45 d[size] = z; b[size] = s[size] = 1;
46 splay(size, -1);
47 }
48 int select(int z) {
49     int x = rt;
50     while (z < s[c[x][0]] + 1 || z > s[c[x][0]] + b[x])
51         if (z > s[c[x][0]] + b[x]) {
52             z -= s[c[x][0]] + b[x];
53             x = c[x][1];
54         }
55         else x = c[x][0];
56     return x;
57 }
58 int main() {
59     scanf("%d",&n);
60     for (int i = 1; i <= n; i++) {
61         int opt, x;
62         scanf("%d%d", &opt, &x);
63         if (opt == 1) insert(x);
64         else if (opt == 2) remove(find(x)); //删除 x 数 (若有多个相同的数, 因只删除一个)
65         else if (opt == 3) { // 查询 x 数的排名 (若有多个相同的数, 因输出最小的排名)
66             insert(x);
67             printf("%d\n", s[c[find(x)][0]] + 1);
68             remove(find(x));
69         }
70         else if (opt == 4) printf("%d\n", d[select(x)]);
71         else if (opt == 5) {
72             insert(x);
73             printf("%d\n", d[pred(find(x))]);
74             remove(find(x));
75         }
76         else if (opt == 6) {
77             insert(x);
78             printf("%d\n", d[succ(find(x))]);
79             remove(find(x));
80         }
81     }
82     return 0;
83 }

```

## 3.2 Splay 区间操作版

使用条件及注意事项：

这是为 NOI2005 维修数列的代码，仅供区间操作用的 splay 参考。

```

1  const int INF = 100000000;
2  const int Maxspace = 500000;
3  struct SplayNode{
4      int ls, rs, zs, ms;
5      SplayNode() {
6          ms = 0;
7          ls = rs = zs = -INF;

```

```

8     }
9     SplayNode(int d) {
10         ms = zs = ls = rs = d;
11     }
12     SplayNode operator +(const SplayNode &p)const {
13         SplayNode ret;
14         ret.ls = max(ls, ms + p.ls);
15         ret.rs = max(rs + p.ms, p.rs);
16         ret.zs = max(rs + p.ls, max(zs, p.zs));
17         ret.ms = ms + p.ms;
18         return ret;
19     }
20 }t[MAXN], d[MAXN];
21 int n, m, rt, top, a[MAXN], f[MAXN], c[MAXN][2], g[MAXN], h[MAXN], z[MAXN];
22 bool r[MAXN], b[MAXN];
23 void makesame(int x, int s) {
24     if (!x) return;
25     b[x] = true;
26     d[x] = SplayNode(g[x] = s);
27     t[x].zs = t[x].ms = g[x] * h[x];
28     t[x].ls = t[x].rs = max(g[x], g[x] * h[x]);
29 }
30 void makerev(int x) {
31     if (!x) return;
32     r[x] ^= 1;
33     swap(c[x][0], c[x][1]);
34     swap(t[x].ls, t[x].rs);
35 }
36 void pushdown(int x) {
37     if (!x) return;
38     if (r[x]) {
39         makerev(c[x][0]);
40         makerev(c[x][1]);
41         r[x]=0;
42     }
43     if (b[x]) {
44         makesame(c[x][0], g[x]);
45         makesame(c[x][1], g[x]);
46         b[x]=g[x]=0;
47     }
48 }
49 void updata(int x) {
50     if (!x) return;
51     h[x]=h[c[x][0]]+h[c[x][1]]+1;
52     t[x]=t[c[x][0]]+d[x]+t[c[x][1]];
53 }
54 void rotate(int x,int k) {
55     pushdown(x);pushdown(c[x][k]);
56     int y = c[x][k]; c[x][k] = c[y][k^1]; c[y][k^1] = x;
57     if (f[x] != -1) c[f[x]][c[f[x]][1] == x] = y; else rt = y;
58     f[y] = f[x]; f[x] = y; f[c[x][k]] = x;
59     updata(x); updata(y);
60 }
61 void splay(int x, int s) {
62     while (f[x] != s) {
63         if (f[f[x]]!=s) {
64             pushdown(f[f[x]]);
65             rotate(f[f[x]], (c[f[f[x]]][1] == f[x]) ^ r[f[f[x]]]);
66         }
67         pushdown(f[x]);
68         rotate(f[x], (c[f[x]][1]==x) ^ r[f[x]]);
69     }
70 }

```

```

71 void build(int &x,int l,int r) {
72     if (l > r) {x = 0; return;}
73     x = z[top--];
74     if (l < r) {
75         build(c[x][0],l,(l+r>>1)-1);
76         build(c[x][1],(l+r>>1)+1,r);
77     }
78     f[c[x][0]] = f[c[x][1]] = x;
79     d[x] = SplayNode(a[l+r>>1]);
80     updata(x);
81 }
82 void init() {
83     d[0] = SplayNode();
84     f[rt=2] = -1;
85     f[1] = 2; c[2][0] = 1;
86     int x;
87     build(x,1,n);
88     c[1][1] = x; f[x] = 1;
89     splay(x, -1);
90 }
91 int find(int z) {
92     int x = rt; pushdown(x);
93     while (z != h[c[x][0]] + 1) {
94         if (z > h[c[x][0]] + 1) {
95             z -= h[c[x][0]] + 1;
96             x = c[x][1];
97         }
98         else x = c[x][0];
99         pushdown(x);
100     }
101     return x;
102 }
103 void getrange(int &x,int &y) {
104     y = x + y - 1;
105     x = find(x);
106     y = find(y + 2);
107     splay(y, -1);
108     splay(x, y);
109 }
110 void recycle(int x) {
111     if (!x) return;
112     recycle(c[x][0]);
113     recycle(c[x][1]);
114     z[++top]=x;
115     t[x] = d[x] = SplayNode();
116     r[x] = b[x] = g[x] = f[x] = h[x] = 0;
117     c[x][0] = c[x][1]=0;
118 }
119 int main() {
120     scanf("%d%d",&n,&m);
121     for (int i = 1; i <= n; i++) scanf("%d",a+i);
122     for (int i = Maxspace; i>=3; i--) z[++top] = i;
123     init();
124     for (int i = 1; i <= m; i++) {
125         char op[10];
126         int x, y, tmp;
127         scanf("%s", op);
128         if (!strcmp(op, "INSERT")) {
129             scanf("%d%d", &x, &y);
130             n += y;
131             if (!y) continue;
132             for (int i = 1; i <= y; i++) scanf("%d",a+i);
133             build(tmp, 1, y);

```

```

134     x = find(x + 1); pushdown(x);
135     if (!c[x][1]) {c[x][1] = tmp; f[tmp] = x;}
136     else{
137         x = c[x][1]; pushdown(x);
138         while (c[x][0]) {
139             x = c[x][0];
140             pushdown(x);
141         }
142         c[x][0] = tmp; f[tmp] = x;
143     }
144     splay(tmp, -1);
145 }
146 else if (!strcmp(op, "DELETE")) {
147     scanf("%d%d", &x, &y); n -= y;
148     if (!y) continue;
149     getrange(x, y);
150     int k = (c[y][0] == x);
151     recycle(c[x][k]);
152     f[c[x][k]] = 0;
153     c[x][k] = 0;
154     splay(x, -1);
155 }
156 else if (!strcmp(op, "REVERSE")) {
157     scanf("%d%d", &x, &y);
158     if (!y) continue;
159     getrange(x, y);
160     int k = (c[y][0] == x);
161     makerev(c[x][k]);
162     splay(c[x][k], -1);
163 }
164 else if (!strcmp(op, "GET-SUM")) {
165     scanf("%d%d", &x, &y);
166     if (!y) {
167         printf("0\n");
168         continue;
169     }
170     getrange(x, y);
171     int k = (c[y][0] == x);
172     printf("%d\n", t[c[x][k]].ms);
173     splay(c[x][k], -1);
174 }
175 else if (!strcmp(op, "MAX-SUM")) {
176     x = 1; y = n;
177     getrange(x, y);
178     int k = (c[y][0] == x);
179     printf("%d\n", t[c[x][k]].zs);
180     splay(c[x][k], -1);
181 }
182 else if (!strcmp(op, "MAKE-SAME")) {
183     scanf("%d%d%d", &x, &y, &tmp);
184     if (!y) continue;
185     getrange(x, y);
186     int k = (c[y][0] == x);
187     makesame(c[x][k], tmp);
188     splay(c[x][k], -1);
189 }
190 }
191 return 0;
192 }

```



### 3.3 坚固的 Treap

使用条件及注意事项：题目来源 UVA 12358

```

1  const int INF = 100000000;
2  const int Maxspace = 500000;
3  struct SplayNode{
4      int ls, rs, zs, ms;
5      SplayNode() {
6          ms = 0;
7          ls = rs = zs = -INF;
8      }
9      SplayNode(int d) {
10         ms = zs = ls = rs = d;
11     }
12     SplayNode operator +(const SplayNode &p)const {
13         SplayNode ret;
14         ret.ls = max(ls, ms + p.ls);
15         ret.rs = max(rs + p.ms, p.rs);
16         ret.zs = max(rs + p.ls, max(zs, p.zs));
17         ret.ms = ms + p.ms;
18         return ret;
19     }
20 }t[MAXN], d[MAXN];
21 int n, m, rt, top, a[MAXN], f[MAXN], c[MAXN][2], g[MAXN], h[MAXN], z[MAXN];
22 bool r[MAXN], b[MAXN];
23 void makesame(int x, int s) {
24     if (!x) return;
25     b[x] = true;
26     d[x] = SplayNode(g[x] = s);
27     t[x].zs = t[x].ms = g[x] * h[x];
28     t[x].ls = t[x].rs = max(g[x], g[x] * h[x]);
29 }
30 void makerev(int x) {
31     if (!x) return;
32     r[x] ^= 1;
33     swap(c[x][0], c[x][1]);
34     swap(t[x].ls, t[x].rs);
35 }
36 void pushdown(int x) {
37     if (!x) return;
38     if (r[x]) {
39         makerev(c[x][0]);
40         makerev(c[x][1]);
41         r[x]=0;
42     }
43     if (b[x]) {
44         makesame(c[x][0],g[x]);
45         makesame(c[x][1],g[x]);
46         b[x]=g[x]=0;
47     }
48 }
49 void updata(int x) {
50     if (!x) return;
51     h[x]=h[c[x][0]]+h[c[x][1]]+1;
52     t[x]=t[c[x][0]]+d[x]+t[c[x][1]];
53 }
54 void rotate(int x,int k) {
55     pushdown(x);pushdown(c[x][k]);
56     int y = c[x][k]; c[x][k] = c[y][k^1]; c[y][k^1] = x;
57     if (f[x] != -1) c[f[x]][c[f[x]][1] == x] = y; else rt = y;
58     f[y] = f[x]; f[x] = y; f[c[x][k]] = x;
59     updata(x); updata(y);

```

```

60 }
61 void splay(int x, int s) {
62     while (f[x] != s) {
63         if (f[f[x]] != s) {
64             pushdown(f[f[x]]);
65             rotate(f[f[x]], (c[f[f[x]]][1] == f[x]) ^ r[f[f[x]]]);
66         }
67         pushdown(f[x]);
68         rotate(f[x], (c[f[x]][1] == x) ^ r[f[x]]);
69     }
70 }
71 void build(int &x, int l, int r) {
72     if (l > r) {x = 0; return;}
73     x = z[top--];
74     if (l < r) {
75         build(c[x][0], l, (l+r>>1)-1);
76         build(c[x][1], (l+r>>1)+1, r);
77     }
78     f[c[x][0]] = f[c[x][1]] = x;
79     d[x] = SplayNode(a[l+r>>1]);
80     updata(x);
81 }
82 void init() {
83     d[0] = SplayNode();
84     f[rt=2] = -1;
85     f[1] = 2; c[2][0] = 1;
86     int x;
87     build(x, 1, n);
88     c[1][1] = x; f[x] = 1;
89     splay(x, -1);
90 }
91 int find(int z) {
92     int x = rt; pushdown(x);
93     while (z != h[c[x][0]] + 1) {
94         if (z > h[c[x][0]] + 1) {
95             z -= h[c[x][0]] + 1;
96             x = c[x][1];
97         }
98         else x = c[x][0];
99         pushdown(x);
100     }
101     return x;
102 }
103 void getrange(int &x, int &y) {
104     y = x + y - 1;
105     x = find(x);
106     y = find(y + 2);
107     splay(y, -1);
108     splay(x, y);
109 }
110 void recycle(int x) {
111     if (!x) return;
112     recycle(c[x][0]);
113     recycle(c[x][1]);
114     z[++top] = x;
115     t[x] = d[x] = SplayNode();
116     r[x] = b[x] = g[x] = f[x] = h[x] = 0;
117     c[x][0] = c[x][1] = 0;
118 }
119 int main() {
120     scanf("%d%d", &n, &m);
121     for (int i = 1; i <= n; i++) scanf("%d", a+i);
122     for (int i = Maxspace; i >= 3; i--) z[++top] = i;

```

```

123     init();
124     for (int i = 1; i <= m; i++) {
125         char op[10];
126         int x, y, tmp;
127         scanf("%s", op);
128         if (!strcmp(op, "INSERT")) {
129             scanf("%d%d", &x, &y);
130             n += y;
131             if (!y) continue;
132             for (int i = 1; i <= y; i++) scanf("%d", &a[i]);
133             build(tmp, 1, y);
134             x = find(x + 1); pushdown(x);
135             if (!c[x][1]) {c[x][1] = tmp; f[tmp] = x;}
136             else{
137                 x = c[x][1]; pushdown(x);
138                 while (c[x][0]) {
139                     x = c[x][0];
140                     pushdown(x);
141                 }
142                 c[x][0] = tmp; f[tmp] = x;
143             }
144             splay(tmp, -1);
145         }
146         else if (!strcmp(op, "DELETE")) {
147             scanf("%d%d", &x, &y); n -= y;
148             if (!y) continue;
149             getrange(x, y);
150             int k = (c[y][0] == x);
151             recycle(c[x][k]);
152             f[c[x][k]] = 0;
153             c[x][k] = 0;
154             splay(x, -1);
155         }
156         else if (!strcmp(op, "REVERSE")) {
157             scanf("%d%d", &x, &y);
158             if (!y) continue;
159             getrange(x, y);
160             int k = (c[y][0] == x);
161             makerev(c[x][k]);
162             splay(c[x][k], -1);
163         }
164         else if (!strcmp(op, "GET-SUM")) {
165             scanf("%d%d", &x, &y);
166             if (!y) {
167                 printf("0\n");
168                 continue;
169             }
170             getrange(x, y);
171             int k = (c[y][0] == x);
172             printf("%d\n", t[c[x][k]].ms);
173             splay(c[x][k], -1);
174         }
175         else if (!strcmp(op, "MAX-SUM")) {
176             x = 1; y = n;
177             getrange(x, y);
178             int k = (c[y][0] == x);
179             printf("%d\n", t[c[x][k]].zs);
180             splay(c[x][k], -1);
181         }
182         else if (!strcmp(op, "MAKE-SAME")) {
183             scanf("%d%d%d", &x, &y, &tmp);
184             if (!y) continue;
185             getrange(x, y);

```

```

186         int k = (c[y][0] == x);
187         makesame(c[x][k], tmp);
188         splay(c[x][k], -1);
189     }
190 }
191 return 0;
192 }

```

### 3.4 k-d 树

使用条件及注意事项：这是求  $k$  远点的代码，要求  $k$  近点的话把堆的比较函数改一改，把朝左儿子或者是右儿子的方向改一改。

```

1  struct Heapnode{
2      long long d;
3      int pos;
4      bool operator <(const Heapnode &p)const {
5          return d > p.d || (d == p.d && pos < p.pos);
6      }
7  };
8
9  struct MsgNode{
10     int xmin, xmax, ymin, ymax;
11     MsgNode() {}
12     MsgNode(const Point &a) : xmin(a.x), xmax(a.x), ymin(a.y), ymax(a.y) {}
13     long long dist(const Point &a) {
14         int dx = std::max(std::abs(a.x - xmin), std::abs(a.x - xmax));
15         int dy = std::max(std::abs(a.y - ymin), std::abs(a.y - ymax));
16         return (long long)dx * dx + (long long)dy * dy;
17     }
18     MsgNode operator +(const MsgNode &rhs)const {
19         MsgNode ret;
20         ret.xmin = std::min(xmin, rhs.xmin);
21         ret.xmax = std::max(xmax, rhs.xmax);
22         ret.ymin = std::min(ymin, rhs.ymin);
23         ret.ymax = std::max(ymax, rhs.ymax);
24         return ret;
25     }
26 };
27
28 struct TNode{
29     int l, r;
30     Point p;
31     MsgNode d;
32 }tree[MAXN];
33
34 void buildtree(int &rt, int l, int r, int pivot) {
35     if (l > r) return;
36     rt = ++size;
37     int mid = l + r >> 1;
38     if (pivot == 1) std::nth_element(p + l, p + mid, p + r + 1, cmpx);
39     if (pivot == 0) std::nth_element(p + l, p + mid, p + r + 1, cmpy);
40     tree[rt].d = MsgNode(tree[rt].p = p[mid]);
41     buildtree(tree[rt].l, l, mid - 1, pivot ^ 1);
42     buildtree(tree[rt].r, mid + 1, r, pivot ^ 1);
43     if (tree[rt].l) tree[rt].d = tree[rt].d + tree[tree[rt].l].d;
44     if (tree[rt].r) tree[rt].d = tree[rt].d + tree[tree[rt].r].d;
45 }
46
47 void query(int rt, const Point &a, int k, int pivot) {
48     Heapnode now = (Heapnode){dist(a, tree[rt].p), tree[rt].p.pos};
49     if (heap.size() < k) heap.push(now);

```

```

50     else if (now < heap.top()) {heap.pop(); heap.push(now);}
51     int lson = tree[rt].l, rson = tree[rt].r;
52     if (pivot == 1 && cmpx(a, tree[rt].p)) std::swap(lson, rson);
53     if (pivot == 0 && cmpy(a, tree[rt].p)) std::swap(lson, rson);
54     if (lson && (heap.size() < k || tree[lson].d.dist(a) >= heap.top().d)) query(lson, a, k,
↪ pivot ^ 1);
55     if (rson && (heap.size() < k || tree[rson].d.dist(a) >= heap.top().d)) query(rson, a, k,
↪ pivot ^ 1);
56 }
57
58 int main() {
59     for (int i = 1; i <= q; i++) {
60         int k;
61         Point now;
62         now.read();
63         scanf("%d", &k);
64         while (!heap.empty()) heap.pop();
65         query(rt, now, k, 1);
66         printf("%d\n", heap.top().pos);
67     }
68     return 0;
69 }

```

## 3.5 树链剖分

### 3.5.1 点操作版本

使用条件及注意事项：树上最大（非空）子段和，注意一条路径询问的时候信息统计的顺序。

```

1  struct Node{
2      int asum, lsum, rsum, zsum;
3      Node() {
4          asum = 0;
5          lsum = -INF;
6          rsum = -INF;
7          zsum = -INF;
8      }
9      Node(int d) : asum(d), lsum(d), rsum(d), zsum(d) {}
10     Node operator +(const Node &rhs) const {
11         Node ret;
12         ret.asum = asum + rhs.asum;
13         ret.lsum = std::max(lsum, asum + rhs.lsum);
14         ret.rsum = std::max(rsum + rhs.asum, rhs.rsum);
15         ret.zsum = std::max(zsum, rhs.zsum);
16         ret.zsum = std::max(ret.zsum, rsum + rhs.lsum);
17         return ret;
18     }
19 }tree[MAXN * 6];
20
21 int n, q, cnt, tot, h[MAXN], d[MAXN], t[MAXN], f[MAXN], s[MAXN], z[MAXN], w[MAXN], o[MAXN],
↪ a[MAXN];
22 std::pair<bool, int> flag[MAXN * 6];
23
24 void addedge(int x, int y) {
25     cnt++; e[cnt] = (Edge){y, h[x]}; h[x] = cnt;
26     cnt++; e[cnt] = (Edge){x, h[y]}; h[y] = cnt;
27 }
28
29 void makesame(int n, int l, int r, int d) {
30     flag[n] = std::make_pair(true, d);
31     tree[n].asum = d * (r - l + 1);
32     if (d > 0) {

```

```

33     tree[n].lsum = d * (r - l + 1);
34     tree[n].rsum = d * (r - l + 1);
35     tree[n].zsum = d * (r - l + 1);
36 }
37 else{
38     tree[n].lsum = d;
39     tree[n].rsum = d;
40     tree[n].zsum = d;
41 }
42 }
43
44 void pushdown(int n, int l, int r) {
45     if (flag[n].first) {
46         makesame(n << 1, l, l + r >> 1, flag[n].second);
47         makesame(n << 1 ^ 1, (l + r >> 1) + 1, r, flag[n].second);
48         flag[n] = std::make_pair(false, 0);
49     }
50 }
51
52 void modify(int n, int l, int r, int x, int y, int d) {
53     if (x <= l && r <= y) {
54         makesame(n, l, r, d);
55         return;
56     }
57     pushdown(n, l, r);
58     if ((l + r >> 1) < x) modify(n << 1 ^ 1, (l + r >> 1) + 1, r, x, y, d);
59     else if ((l + r >> 1) + 1 > y) modify(n << 1, l, l + r >> 1, x, y, d);
60     else{
61         modify(n << 1, l, l + r >> 1, x, y, d);
62         modify(n << 1 ^ 1, (l + r >> 1) + 1, r, x, y, d);
63     }
64     tree[n] = tree[n << 1] + tree[n << 1 ^ 1];
65 }
66
67 Node query(int n, int l, int r, int x, int y) {
68     if (x <= l && r <= y) return tree[n];
69     pushdown(n, l, r);
70     if ((l + r >> 1) < x) return query(n << 1 ^ 1, (l + r >> 1) + 1, r, x, y);
71     else if ((l + r >> 1) + 1 > y) return query(n << 1, l, l + r >> 1, x, y);
72     else{
73         Node left = query(n << 1, l, l + r >> 1, x, y);
74         Node right = query(n << 1 ^ 1, (l + r >> 1) + 1, r, x, y);
75         return left + right;
76     }
77 }
78
79 void modify(int x, int y, int val) {
80     int fx = t[x], fy = t[y];
81     while (fx != fy) {
82         if (d[fx] > d[fy]) {
83             modify(1, 1, n, w[fx], w[x], val);
84             x = f[fx]; fx = t[x];
85         }
86         else{
87             modify(1, 1, n, w[fy], w[y], val);
88             y = f[fy]; fy = t[y];
89         }
90     }
91     if (d[x] < d[y]) modify(1, 1, n, w[x], w[y], val);
92     else modify(1, 1, n, w[y], w[x], val);
93 }
94
95 Node query(int x, int y) {

```

```

96     int fx = t[x], fy = t[y];
97     Node left = Node(), right = Node();
98     while (fx != fy) {
99         if (d[fx] > d[fy]) {
100             left = query(1, 1, n, w[fx], w[x]) + left;
101             x = f[fx]; fx = t[x];
102         }
103         else{
104             right = query(1, 1, n, w[fy], w[y]) + right;
105             y = f[fy]; fy = t[y];
106         }
107     }
108     if (d[x] < d[y]) {
109         right = query(1, 1, n, w[x], w[y]) + right;
110     }
111     else{
112         left = query(1, 1, n, w[y], w[x]) + left;
113     }
114     std::swap(left.lsum, left.rsum);
115     return left + right;
116 }
117
118 void predfs(int x) {
119     s[x] = 1; z[x] = 0;
120     for (int i = h[x]; i; i = e[i].next) {
121         if (e[i].node == f[x]) continue;
122         f[e[i].node] = x;
123         d[e[i].node] = d[x] + 1;
124         predfs(e[i].node);
125         s[x] += s[e[i].node];
126         if (s[z[x]] < s[e[i].node]) z[x] = e[i].node;
127     }
128 }
129
130 void getanc(int x, int anc) {
131     t[x] = anc; w[x] = ++tot; o[tot] = x;
132     if (z[x]) getanc(z[x], anc);
133     for (int i = h[x]; i; i = e[i].next) {
134         if (e[i].node == f[x] || e[i].node == z[x]) continue;
135         getanc(e[i].node, e[i].node);
136     }
137 }
138
139 void buildtree(int n, int l, int r) {
140     if (l == r) {
141         tree[n] = Node(a[o[l]]);
142         return;
143     }
144     buildtree(n << 1, l, l + r >> 1);
145     buildtree(n << 1 ^ 1, (l + r >> 1) + 1, r);
146     tree[n] = tree[n << 1] + tree[n << 1 ^ 1];
147 }
148
149 int main() {
150     scanf("%d", &n);
151     for (int i = 1; i <= n; i++) scanf("%d", a + i);
152     for (int i = 1; i < n; i++) {
153         int x, y; scanf("%d%d", &x, &y);
154         addedge(x, y);
155     }
156     predfs(1);
157     getanc(1, 1);
158     buildtree(1, 1, n);

```

```

159     scanf("%d", &q);
160     for (int i = 1; i <= q; i++) {
161         int op, x, y, c;
162         scanf("%d", &op);
163         if (op == 1) {
164             scanf("%d%d", &x, &y);
165             Node ret = query(x, y);
166             printf("%d\n", std::max(0, ret.zsum));
167         }
168         else{
169             scanf("%d%d%d", &x, &y, &c);
170             modify(x, y, c);
171         }
172     }
173     return 0;
174 }

```

### 3.5.2 链操作版本

```

1  void modify(int x, int y) {
2      int fx = t[x], fy = t[y];
3      while (fx != fy) {
4          if (d[fx] > d[fy]) {
5              modify(1, 1, n, w[fx], w[x]);
6              x = f[fx]; fx = t[x];
7          }
8          else{
9              modify(1, 1, n, w[fy], w[y]);
10             y = f[fy]; fy = t[y];
11         }
12     }
13     if (x != y) {
14         if (d[x] < d[y]) modify(1, 1, n, w[z[x]], w[y]);
15         else modify(1, 1, n, w[z[y]], w[x]);
16     }
17 }

```

## 3.6 Link-Cut-Tree

```

1  struct MsgNode{
2      int leftColor, rightColor, answer;
3      MsgNode() {
4          leftColor = -1;
5          rightColor = -1;
6          answer = 0;
7      }
8      MsgNode(int c) {
9          leftColor = rightColor = c;
10         answer = 1;
11     }
12     MsgNode operator +(const MsgNode &p) const {
13         if (answer == 0) return p;
14         if (p.answer == 0) return *this;
15         MsgNode ret;
16         ret.leftColor = leftColor;
17         ret.rightColor = p.rightColor;
18         ret.answer = answer + p.answer - (rightColor == p.leftColor);
19         return ret;
20     }
21 }d[MAXN], g[MAXN];

```



```

22  int n, m, c[MAXN][2], f[MAXN], p[MAXN], s[MAXN], flag[MAXN];
23  bool r[MAXN];
24  void init(int x, int value) {
25      d[x] = g[x] = MsgNode(value);
26      c[x][0] = c[x][1] = 0;
27      f[x] = p[x] = flag[x] = -1;
28      s[x] = 1;
29  }
30  void update(int x) {
31      s[x] = s[c[x][0]] + s[c[x][1]] + 1;
32      g[x] = MsgNode();
33      if (c[x][0] ^ r[x]) g[x] = g[x] + g[c[x][0] ^ r[x]];
34      g[x] = g[x] + d[x];
35      if (c[x][1] ^ r[x]) g[x] = g[x] + g[c[x][1] ^ r[x]];
36  }
37  void makesame(int x, int c) {
38      flag[x] = c;
39      d[x] = MsgNode(c);
40      g[x] = MsgNode(c);
41  }
42  void pushdown(int x) {
43      if (r[x]) {
44          std::swap(c[x][0], c[x][1]);
45          r[c[x][0]] ^= 1;
46          r[c[x][1]] ^= 1;
47          std::swap(g[c[x][0]].leftColor, g[c[x][0]].rightColor);
48          std::swap(g[c[x][1]].leftColor, g[c[x][1]].rightColor);
49          r[x] = false;
50      }
51      if (flag[x] != -1) {
52          if (c[x][0]) makesame(c[x][0], flag[x]);
53          if (c[x][1]) makesame(c[x][1], flag[x]);
54          flag[x] = -1;
55      }
56  }
57  void rotate(int x, int k) {
58      pushdown(x); pushdown(c[x][k]);
59      int y = c[x][k]; c[x][k] = c[y][k ^ 1]; c[y][k ^ 1] = x;
60      if (f[x] != -1) c[f[x]][c[f[x]][1] == x] = y;
61      f[y] = f[x]; f[x] = y; f[c[x][k]] = x; std::swap(p[x], p[y]);
62      update(x); update(y);
63  }
64  void splay(int x, int s = -1) {
65      pushdown(x);
66      while (f[x] != s) {
67          if (f[f[x]] != s) rotate(f[f[x]], (c[f[f[x]]][1] == f[x]) ^ r[f[f[x]]]);
68          rotate(f[x], (c[f[x]][1] == x) ^ r[f[x]]);
69      }
70      update(x);
71  }
72  void access(int x) {
73      int y = 0;
74      while (x != -1) {
75          splay(x); pushdown(x);
76          f[c[x][1]] = -1; p[c[x][1]] = x;
77          c[x][1] = y; f[y] = x; p[y] = -1;
78          update(x); x = p[y = x];
79      }
80  }
81  void setroot(int x) {
82      access(x);
83      splay(x);
84      r[x] ^= 1;

```

```
85     std::swap(g[x].leftColor, g[x].rightColor);
86 }
87 void link(int x, int y) {
88     setroot(x);
89     p[x] = y;
90 }
```

# Chapter 4

## 图论

### 4.1 强连通分量

```
1  int stamp, comps, top;
2  int dfn[N], low[N], comp[N], stack[N];
3
4  void tarjan(int x) {
5      dfn[x] = low[x] = ++stamp;
6      stack[top++] = x;
7      for (int i = 0; i < (int)edge[x].size(); ++i) {
8          int y = edge[x][i];
9          if (!dfn[y]) {
10             tarjan(y);
11             low[x] = std::min(low[x], low[y]);
12         } else if (!comp[y]) {
13             low[x] = std::min(low[x], dfn[y]);
14         }
15     }
16     if (low[x] == dfn[x]) {
17         comps++;
18         do {
19             int y = stack[--top];
20             comp[y] = comps;
21         } while (stack[top] != x);
22     }
23 }
24
25 void solve() {
26     stamp = comps = top = 0;
27     std::fill(dfn, dfn + n, 0);
28     std::fill(comp, comp + n, 0);
29     for (int i = 0; i < n; ++i) {
30         if (!dfn[i]) {
31             tarjan(i);
32         }
33     }
34 }
```

### 4.2 点双连通分量

#### 4.2.1 坚固的点双连通分量

```
1  int n, m, x, y, ans1, ans2, tot1, tot2, flag, size, ind2, dfn[N], low[N], block[M], vis[N];
2  vector<int> a[N];
3  pair<int, int> stack[M];
4  void tarjan(int x, int p) {
5      dfn[x] = low[x] = ++ind2;
```

```

6   for (int i = 0; i < a[x].size(); ++i)
7       if (dfn[x] > dfn[a[x][i]] && a[x][i] != p){
8           stack[++size] = make_pair(x, a[x][i]);
9           if (i == a[x].size() - 1 || a[x][i] != a[x][i + 1])
10              if (!dfn[a[x][i]]){
11                  tarjan(a[x][i], x);
12                  low[x] = min(low[x], low[a[x][i]]);
13                  if (low[a[x][i]] >= dfn[x]){
14                      tot1 = tot2 = 0;
15                      ++flag;
16                      for (; ; ){
17                          if (block[stack[size].first] != flag) {
18                              ++tot1;
19                              block[stack[size].first] = flag;
20                          }
21                          if (block[stack[size].second] != flag) {
22                              ++tot1;
23                              block[stack[size].second] = flag;
24                          }
25                          if (stack[size].first == x && stack[size].second == a[x][i])
26                              break;
27                          ++tot2;
28                          --size;
29                      }
30                      for (; stack[size].first == x && stack[size].second == a[x][i]; --size)
31                          ++tot2;
32                      if (tot2 < tot1)
33                          ans1 += tot2;
34                      if (tot2 > tot1)
35                          ans2 += tot2;
36                  }
37              }
38           else
39               low[x] = min(low[x], dfn[a[x][i]]);
40       }
41 }
42 int main(){
43     for (; ; ){
44         scanf("%d%d", &n, &m);
45         if (n == 0 && m == 0) return 0;
46         for (int i = 1; i <= n; ++i) {
47             a[i].clear();
48             dfn[i] = 0;
49         }
50         for (int i = 1; i <= m; ++i){
51             scanf("%d%d", &x, &y);
52             ++x, ++y;
53             a[x].push_back(y);
54             a[y].push_back(x);
55         }
56         for (int i = 1; i <= n; ++i)
57             sort(a[i].begin(), a[i].end());
58         ans1 = ans2 = ind2 = 0;
59         for (int i = 1; i <= n; ++i)
60             if (!dfn[i]) {
61                 size = 0;
62                 tarjan(i, 0);
63             }
64         printf("%d %d\n", ans1, ans2);
65     }
66     return 0;
67 }

```

## 4.2.2 朴素的点双连通分量

```

1  void tarjan(int x){
2      dfn[x] = low[x] = ++ind2;
3      v[x] = 1;
4      for (int i = nt[x]; pt[i]; i = nt[i])
5          if (!dfn[pt[i]]){
6              tarjan(pt[i]);
7              low[x] = min(low[x], low[pt[i]]);
8              if (dfn[x] <= low[pt[i]])
9                  ++v[x];
10         }
11         else
12             low[x] = min(low[x], dfn[pt[i]]);
13     }
14     int main(){
15         for (; ; ){
16             scanf("%d%d", &n, &m);
17             if (n == 0 && m == 0)
18                 return 0;
19             for (int i = 1; i <= ind; ++i)
20                 nt[i] = pt[i] = 0;
21             ind = n;
22             for (int i = 1; i <= ind; ++i)
23                 last[i] = i;
24             for (int i = 1; i <= m; ++i){
25                 scanf("%d%d", &x, &y);
26                 ++x, ++y;
27                 edge(x, y), edge(y, x);
28             }
29             memset(dfn, 0, sizeof(dfn));
30             memset(v, 0, sizeof(v));
31             ans = num = ind2 = 0;
32             for (int i = 1; i <= n; ++i)
33                 if (!dfn[i]){
34                     root = i;
35                     size = 0;
36                     ++num;
37                     tarjan(i);
38                     --v[root];
39                 }
40             for (int i = 1; i <= n; ++i)
41                 if (v[i] + num - 1 > ans)
42                     ans = v[i] + num - 1;
43             printf("%d\n", ans);
44         }
45         return 0;
46     }

```

## 4.3 2-SAT 问题

```

1  int stamp, comps, top;
2  int dfn[N], low[N], comp[N], stack[N];
3
4  void add(int x, int a, int y, int b) {
5      edge[x << 1 | a].push_back(y << 1 | b);
6  }
7
8  void tarjan(int x) {
9      dfn[x] = low[x] = ++stamp;
10     stack[top++] = x;

```

```

11     for (int i = 0; i < (int)edge[x].size(); ++i) {
12         int y = edge[x][i];
13         if (!dfn[y]) {
14             tarjan(y);
15             low[x] = std::min(low[x], low[y]);
16         } else if (!comp[y]) {
17             low[x] = std::min(low[x], dfn[y]);
18         }
19     }
20     if (low[x] == dfn[x]) {
21         comps++;
22         do {
23             int y = stack[--top];
24             comp[y] = comps;
25         } while (stack[top] != x);
26     }
27 }
28
29 bool solve() {
30     int counter = n + n + 1;
31     stamp = top = comps = 0;
32     std::fill(dfn, dfn + counter, 0);
33     std::fill(comp, comp + counter, 0);
34     for (int i = 0; i < counter; ++i) {
35         if (!dfn[i]) {
36             tarjan(i);
37         }
38     }
39     for (int i = 0; i < n; ++i) {
40         if (comp[i << 1] == comp[i << 1 | 1]) {
41             return false;
42         }
43         answer[i] = (comp[i << 1 | 1] < comp[i << 1]);
44     }
45     return true;
46 }

```

## 4.4 二分图最大匹配

### 4.4.1 Hungary 算法

时间复杂度:  $\mathcal{O}(V \cdot E)$

```

1  int n, m, stamp;
2  int match[N], visit[N];
3
4  bool dfs(int x) {
5      for (int i = 0; i < (int)edge[x].size(); ++i) {
6          int y = edge[x][i];
7          if (visit[y] != stamp) {
8              visit[y] = stamp;
9              if (match[y] == -1 || dfs(match[y])) {
10                 match[y] = x;
11                 return true;
12             }
13         }
14     }
15     return false;
16 }
17
18 int solve() {
19     std::fill(match, match + m, -1);

```

```

20     int answer = 0;
21     for (int i = 0; i < n; ++i) {
22         stamp++;
23         answer += dfs(i);
24     }
25     return answer;
26 }

```

#### 4.4.2 Hopcroft Karp 算法

时间复杂度:  $\mathcal{O}(\sqrt{V} \cdot E)$

```

1  int matchx[N], matchy[N], level[N];
2
3  bool dfs(int x) {
4      for (int i = 0; i < (int)edge[x].size(); ++i) {
5          int y = edge[x][i];
6          int w = matchy[y];
7          if (w == -1 || level[x] + 1 == level[w] && dfs(w)) {
8              matchx[x] = y;
9              matchy[y] = x;
10             return true;
11         }
12     }
13     level[x] = -1;
14     return false;
15 }
16
17 int solve() {
18     std::fill(matchx, matchx + n, -1);
19     std::fill(matchy, matchy + m, -1);
20     for (int answer = 0; ; ) {
21         std::vector<int> queue;
22         for (int i = 0; i < n; ++i) {
23             if (matchx[i] == -1) {
24                 level[i] = 0;
25                 queue.push_back(i);
26             } else {
27                 level[i] = -1;
28             }
29         }
30         for (int head = 0; head < (int)queue.size(); ++head) {
31             int x = queue[head];
32             for (int i = 0; i < (int)edge[x].size(); ++i) {
33                 int y = edge[x][i];
34                 int w = matchy[y];
35                 if (w != -1 && level[w] < 0) {
36                     level[w] = level[x] + 1;
37                     queue.push_back(w);
38                 }
39             }
40         }
41         int delta = 0;
42         for (int i = 0; i < n; ++i) {
43             if (matchx[i] == -1 && dfs(i)) {
44                 delta++;
45             }
46         }
47         if (delta == 0) {
48             return answer;
49         } else {
50             answer += delta;
51         }
52     }
53 }

```

```

52     }
53 }

```

## 4.5 二分图最大权匹配

时间复杂度:  $\mathcal{O}(V^4)$

```

1  int DFS(int x){
2      visx[x] = 1;
3      for (int y = 1; y <= ny; y++){
4          if (visy[y]) continue;
5          int t = lx[x] + ly[y] - w[x][y];
6          if (t == 0) {
7              visy[y] = 1;
8              if (link[y] == -1 || DFS(link[y])){
9                  link[y] = x;
10                 return 1;
11             }
12         }
13         else slack[y] = min(slack[y], t);
14     }
15     return 0;
16 }
17 int KM(){
18     int i, j;
19     memset(link, -1, sizeof(link));
20     memset(ly, 0, sizeof(ly));
21     for (i = 1; i <= nx; i++){
22         for (j = 1, lx[i] = -inf; j <= ny; j++){
23             lx[i] = max(lx[i], w[i][j]);
24         }
25         for (int x = 1; x <= nx; x++){
26             for (i = 1; i <= ny; i++) slack[i] = inf;
27             while (true) {
28                 memset(visx, 0, sizeof(visx));
29                 memset(visy, 0, sizeof(visy));
30                 if (DFS(x)) break;
31                 int d = inf;
32                 for (i = 1; i <= ny; i++){
33                     if (!visy[i] && d > slack[i]) d = slack[i];
34                 }
35                 for (i = 1; i <= nx; i++){
36                     if (visx[i]) lx[i] -= d;
37                 }
38                 for (i = 1; i <= ny; i++){
39                     if (visy[i]) ly[i] += d;
40                     else slack[i] -= d;
41                 }
42             }
43         }
44     }
45     int res = 0;
46     for (i = 1; i <= ny; i++){
47         if (link[i] > -1) res += w[link[i]][i];
48     }
49     return res;
50 }

```

## 4.6 最大流

### 4.6.1 Dinic

使用方法以及注意事项:  $n$  个点,  $m$  条边,  $inf$  为一个很大的值, 源点  $s$ , 汇点  $t$ , 图中最大点的编号为  $t$ 。

邻接表:  $p$  数组记录节点,  $next$  数组指向下一个位置,  $c$  数组记录可增广量,  $h$  数组记录表头 (初始全为-1)。

时间复杂度:  $\mathcal{O}(V^2 \cdot E)$



```

1  int bfs(){
2      for (int i = 1; i <= t; i++) d[i] = -1;
3      int l, r;
4      q[l = r = 0] = s, d[s] = 0;
5      for (; l <= r; l++)
6          for (int k = h[q[l]]; k > -1; k = nxt[k])
7              if (d[p[k]] == -1 && c[k] > 0) d[p[k]] = d[q[l]] + 1, q[++ r] = p[k];
8      return d[t] > -1 ? 1 : 0;
9  }
10 int dfs(int u, int ext){
11     if (u == t) return ext;
12     int k = w[u], ret = 0;
13     for (; k > -1; k = nxt[k], w[u] = k){          //w 数组为当前弧
14         if (ext == 0) break;
15         if (d[p[k]] == d[u] + 1 && c[k] > 0){
16             int flow = dfs(p[k], min(c[k], ext));
17             if (flow > 0){
18                 c[k] -= flow, c[k ^ 1] += flow;
19                 ret += flow, ext -= flow;          //ret 累计增广量, ext 记录还可增广的量
20             }
21         }
22     }
23     if (k == -1) d[u] = -1;
24     return ret;
25 }
26 void dinic() {
27     while (bfs()) {
28         for (int i = 1; i <= t; i++) w[i] = h[i];
29         dfs(s, inf);
30     }
31 }

```

#### 4.6.2 ISAP

时间复杂度:  $\mathcal{O}(V^2 \cdot E)$

```

1  int Maxflow_Isap(int s, int t, int n) {
2      std::fill(pre + 1, pre + n + 1, 0);
3      std::fill(d + 1, d + n + 1, 0);
4      std::fill(gap + 1, gap + n + 1, 0);
5      for (int i = 1; i <= n; i++) cur[i] = h[i];
6      gap[0] = n;
7      int u = pre[s] = s, v, maxflow = 0;
8      while (d[s] < n) {
9          v = n + 1;
10         for (int i = cur[u]; i; i = e[i].next)
11             if (e[i].flow && d[u] == d[e[i].node] + 1) {
12                 v = e[i].node; cur[u] = i; break;
13             }
14         if (v <= n) {
15             pre[v] = u; u = v;
16             if (v == t) {
17                 int dflow = INF, p = t; u = s;
18                 while (p != s) {
19                     p = pre[p];
20                     dflow = std::min(dflow, e[cur[p]].flow);
21                 }
22                 maxflow += dflow; p = t;
23                 while (p != s) {
24                     p = pre[p];
25                     e[cur[p]].flow -= dflow;

```

```

26         e[e[cur[p]].opp].flow += dflow;
27     }
28 }
29 }
30 else{
31     int mindist = n + 1;
32     for (int i = h[u]; i; i = e[i].next)
33         if (e[i].flow && mindist > d[e[i].node]) {
34             mindist = d[e[i].node]; cur[u] = i;
35         }
36     if (!--gap[d[u]]) return maxflow;
37     gap[d[u] = mindist + 1]++; u = pre[u];
38 }
39 }
40 return maxflow;
41 }

```

### 4.6.3 SAP

时间复杂度:  $\mathcal{O}(V^2 \cdot E)$

```

1  int Maxflow_Isap(int s, int t, int n) {
2      std::fill(pre + 1, pre + n + 1, 0);
3      std::fill(d + 1, d + n + 1, 0);
4      std::fill(gap + 1, gap + n + 1, 0);
5      for (int i = 1; i <= n; i++) cur[i] = h[i];
6      gap[0] = n;
7      int u = pre[s] = s, v, maxflow = 0;
8      while (d[s] < n) {
9          v = n + 1;
10         for (int i = cur[u]; i; i = e[i].next)
11             if (e[i].flow && d[u] == d[e[i].node] + 1) {
12                 v = e[i].node; cur[u] = i; break;
13             }
14         if (v <= n) {
15             pre[v] = u; u = v;
16             if (v == t) {
17                 int dflow = INF, p = t; u = s;
18                 while (p != s) {
19                     p = pre[p];
20                     dflow = std::min(dflow, e[cur[p]].flow);
21                 }
22                 maxflow += dflow; p = t;
23                 while (p != s) {
24                     p = pre[p];
25                     e[cur[p]].flow -= dflow;
26                     e[e[cur[p]].opp].flow += dflow;
27                 }
28             }
29         }
30         else{
31             int mindist = n + 1;
32             for (int i = h[u]; i; i = e[i].next)
33                 if (e[i].flow && mindist > d[e[i].node]) {
34                     mindist = d[e[i].node]; cur[u] = i;
35                 }
36             if (!--gap[d[u]]) return maxflow;
37             gap[d[u] = mindist + 1]++; u = pre[u];
38         }
39     }
40     return maxflow;
41 }

```

## 4.7 上下界网络流

$B(u, v)$  表示边  $(u, v)$  流量的下界,  $C(u, v)$  表示边  $(u, v)$  流量的上界,  $F(u, v)$  表示边  $(u, v)$  的流量。设  $G(u, v) = F(u, v) - B(u, v)$ , 显然有

$$0 \leq G(u, v) \leq C(u, v) - B(u, v)$$

### 4.7.1 无源汇的上下界可行流

建立超级源点  $S^*$  和超级汇点  $T^*$ , 对于原图每条边  $(u, v)$  在新网络中连如下三条边:  $S^* \rightarrow v$ , 容量为  $B(u, v)$ ;  $u \rightarrow T^*$ , 容量为  $B(u, v)$ ;  $u \rightarrow v$ , 容量为  $C(u, v) - B(u, v)$ 。最后求新网络的最大流, 判断从超级源点  $S^*$  出发的边是否都满流即可, 边  $(u, v)$  的最终解中的实际流量为  $G(u, v) + B(u, v)$ 。

### 4.7.2 有源汇的上下界可行流

从汇点  $T$  到源点  $S$  连一条上界为  $\infty$ , 下界为 0 的边。按照无源汇的上下界可行流一样做即可, 流量即为  $T \rightarrow S$  边上的流量。

### 4.7.3 有源汇的上下界最大流

1. 在有源汇的上下界可行流中, 从汇点  $T$  到源点  $S$  的边改为连一条上界为  $\infty$ , 下届为  $x$  的边。  $x$  满足二分性质, 找到最大的  $x$  使得新网络存在无源汇的上下界可行流即为原图的最大流。
2. 从汇点  $T$  到源点  $S$  连一条上界为  $\infty$ , 下界为 0 的边, 变成无源汇的网络。按照无源汇的上下界可行流的方法, 建立超级源点  $S^*$  和超级汇点  $T^*$ , 求一遍  $S^* \rightarrow T^*$  的最大流, 再将从汇点  $T$  到源点  $S$  的这条边拆掉, 求一次  $S \rightarrow T$  的最大流即可。

### 4.7.4 有源汇的上下界最小流

1. 在有源汇的上下界可行流中, 从汇点  $T$  到源点  $S$  的边改为连一条上界为  $x$ , 下界为 0 的边。  $x$  满足二分性质, 找到最小的  $x$  使得新网络存在无源汇的上下界可行流即为原图的最小流。
2. 按照无源汇的上下界可行流的方法, 建立超级源点  $S^*$  与超级汇点  $T^*$ , 求一遍  $S^* \rightarrow T^*$  的最大流, 但是注意这一次不加上汇点  $T$  到源点  $S$  的这条边, 即不使之改为无源汇的网络去求解。求完后, 再加上那条汇点  $T$  到源点  $S$  上界  $\infty$  的边。因为这条边下界为 0, 所以  $S^*, T^*$  无影响, 再直接求一次  $S^* \rightarrow T^*$  的最大流。若超级源点  $S^*$  出发的边全部满流, 则  $T \rightarrow S$  边上的流量即为原图的最小流, 否则无解。

## 4.8 最小费用最大流

### 4.8.1 稀疏图

时间复杂度:  $\mathcal{O}(V \cdot E^2)$

```

1  struct EdgeList {
2      int size;
3      int last[N];
4      int succ[M], other[M], flow[M], cost[M];
5      void clear(int n) {
6          size = 0;
7          std::fill(last, last + n, -1);
8      }
9      void add(int x, int y, int c, int w) {
10         succ[size] = last[x];
11         last[x] = size;
12         other[size] = y;
13         flow[size] = c;
14         cost[size++] = w;
15     }
16 } e;
17
18 int n, source, target;
19 int prev[N];
20
21 void add(int x, int y, int c, int w) {
22     e.add(x, y, c, w);

```

```

23     e.add(y, x, 0, -w);
24 }
25
26 bool augment() {
27     static int dist[N], occur[N];
28     std::vector<int> queue;
29     std::fill(dist, dist + n, INT_MAX);
30     std::fill(occur, occur + n, 0);
31     dist[source] = 0;
32     occur[source] = true;
33     queue.push_back(source);
34     for (int head = 0; head < (int)queue.size(); ++head) {
35         int x = queue[head];
36         for (int i = e.last[x]; ~i; i = e.succ[i]) {
37             int y = e.other[i];
38             if (e.flow[i] && dist[y] > dist[x] + e.cost[i]) {
39                 dist[y] = dist[x] + e.cost[i];
40                 prev[y] = i;
41                 if (!occur[y]) {
42                     occur[y] = true;
43                     queue.push_back(y);
44                 }
45             }
46         }
47         occur[x] = false;
48     }
49     return dist[target] < INT_MAX;
50 }
51
52 std::pair<int, int> solve() {
53     std::pair<int, int> answer = std::make_pair(0, 0);
54     while (augment()) {
55         int number = INT_MAX;
56         for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
57             number = std::min(number, e.flow[prev[i]]);
58         }
59         answer.first += number;
60         for (int i = target; i != source; i = e.other[prev[i] ^ 1]) {
61             e.flow[prev[i]] -= number;
62             e.flow[prev[i] ^ 1] += number;
63             answer.second += number * e.cost[prev[i]];
64         }
65     }
66     return answer;
67 }

```

#### 4.8.2 稠密图

使用条件：费用非负

时间复杂度： $\mathcal{O}(V \cdot E^2)$

```

1  int aug(int no, int res) {
2      if (no == t) return cost += pi1 * res, res;
3      v[no] = true;
4      int flow = 0;
5      for (int i = h[no]; ~i; i = nxt[i])
6          if (cap[i] && !expense[i] && !v[p[i]]) {
7              int d = aug(p[i], min(res, cap[i]));
8              cap[i] -= d, cap[i ^ 1] += d, flow += d, res -= d;
9              if (!res) return flow;
10         }
11     return flow;
12 }

```

```

13 bool modlabel() {
14     int d = maxint;
15     for(int i = 1; i <= t; ++ i)
16         if(v[i]) {
17             for(int j = h[i]; ~ j; j = nxt[j])
18                 if(cap[j] && !v[p[j]] && expense[j] < d) d = expense[j];
19         }
20     if(d == maxint) return false;
21     for(int i = 1; i <= t; ++ i)
22         if(v[i]) {
23             for(int j = h[i]; ~ j; j = nxt[j])
24                 expense[j] -= d, expense[j ^ 1] += d;
25         }
26     pi1 += d;
27     return true;
28 }
29 void minimum_cost_flow_zkw() {
30     cost = 0;
31     do{
32         do{
33             memset(v, false, sizeof v);
34         }while (aug(s, maxint));
35     }while (modlabel());
36 }

```

## 4.9 一般图最大匹配

时间复杂度:  $\mathcal{O}(V^3)$

```

1  int match[N], belong[N], next[N], mark[N], visit[N];
2  std::vector<int> queue;
3
4  int find(int x) {
5      if (belong[x] != x) {
6          belong[x] = find(belong[x]);
7      }
8      return belong[x];
9  }
10
11 void merge(int x, int y) {
12     x = find(x);
13     y = find(y);
14     if (x != y) {
15         belong[x] = y;
16     }
17 }
18
19 int lca(int x, int y) {
20     static int stamp = 0;
21     stamp++;
22     while (true) {
23         if (x != -1) {
24             x = find(x);
25             if (visit[x] == stamp) {
26                 return x;
27             }
28             visit[x] = stamp;
29             if (match[x] != -1) {
30                 x = next[match[x]];
31             } else {
32                 x = -1;
33             }

```

```

34     }
35     std::swap(x, y);
36 }
37 }
38
39 void group(int a, int p) {
40     while (a != p) {
41         int b = match[a], c = next[b];
42         if (find(c) != p) {
43             next[c] = b;
44         }
45         if (mark[b] == 2) {
46             mark[b] = 1;
47             queue.push_back(b);
48         }
49         if (mark[c] == 2) {
50             mark[c] = 1;
51             queue.push_back(c);
52         }
53         merge(a, b);
54         merge(b, c);
55         a = c;
56     }
57 }
58
59 void augment(int source) {
60     queue.clear();
61     for (int i = 0; i < n; ++i) {
62         next[i] = visit[i] = -1;
63         belong[i] = i;
64         mark[i] = 0;
65     }
66     mark[source] = 1;
67     queue.push_back(source);
68     for (int head = 0; head < (int)queue.size() && match[source] == -1; ++head) {
69         int x = queue[head];
70         for (int i = 0; i < (int)edge[x].size(); ++i) {
71             int y = edge[x][i];
72             if (match[x] == y || find(x) == find(y) || mark[y] == 2) {
73                 continue;
74             }
75             if (mark[y] == 1) {
76                 int r = lca(x, y);
77                 if (find(x) != r) {
78                     next[x] = y;
79                 }
80                 if (find(y) != r) {
81                     next[y] = x;
82                 }
83                 group(x, r);
84                 group(y, r);
85             } else if (match[y] == -1) {
86                 next[y] = x;
87                 for (int u = y; u != -1; ) {
88                     int v = next[u];
89                     int mv = match[v];
90                     match[v] = u;
91                     match[u] = v;
92                     u = mv;
93                 }
94                 break;
95             } else {
96                 next[y] = x;

```

```

97         mark[y] = 2;
98         mark[match[y]] = 1;
99         queue.push_back(match[y]);
100     }
101 }
102 }
103 }
104
105 int solve() {
106     std::fill(match, match + n, -1);
107     for (int i = 0; i < n; ++i) {
108         if (match[i] == -1) {
109             augment(i);
110         }
111     }
112     int answer = 0;
113     for (int i = 0; i < n; ++i) {
114         answer += (match[i] != -1);
115     }
116     return answer;
117 }

```

## 4.10 无向图全局最小割

时间复杂度:  $\mathcal{O}(V^3)$

注意事项: 处理重边时, 应该对边权累加

```

1  int node[N], dist[N];
2  bool visit[N];
3
4  int solve(int n) {
5      int answer = INT_MAX;
6      for (int i = 0; i < n; ++i) {
7          node[i] = i;
8      }
9      while (n > 1) {
10         int max = 1;
11         for (int i = 0; i < n; ++i) {
12             dist[node[i]] = graph[node[0]][node[i]];
13             if (dist[node[i]] > dist[node[max]]) {
14                 max = i;
15             }
16         }
17         int prev = 0;
18         memset(visit, 0, sizeof(visit));
19         visit[node[0]] = true;
20         for (int i = 1; i < n; ++i) {
21             if (i == n - 1) {
22                 answer = std::min(answer, dist[node[max]]);
23                 for (int k = 0; k < n; ++k) {
24                     graph[node[k]][node[prev]] =
25                         (graph[node[prev]][node[k]] += graph[node[k]][node[max]]);
26                 }
27                 node[max] = node[--n];
28             }
29             visit[node[max]] = true;
30             prev = max;
31             max = -1;
32             for (int j = 1; j < n; ++j) {
33                 if (!visit[node[j]]) {
34                     dist[node[j]] += graph[node[prev]][node[j]];
35                     if (max == -1 || dist[node[max]] < dist[node[j]]) {

```

```

36         max = j;
37     }
38 }
39 }
40 }
41 }
42 return answer;
43 }

```

## 4.11 最小树形图

```

1  int n, m, used[N], pass[N], eg[N], more, queue[N];
2  double g[N][N];
3
4  void combine(int id, double &sum) {
5      int tot = 0, from, i, j, k;
6      for (; id != 0 && !pass[id]; id = eg[id]) {
7          queue[tot++] = id;
8          pass[id] = 1;
9      }
10
11     for (from = 0; from < tot && queue[from] != id; from++);
12     if (from == tot) return;
13     more = 1;
14     for (i = from; i < tot; i++) {
15         sum += g[eg[queue[i]]][queue[i]];
16         if (i != from) {
17             used[queue[i]] = 1;
18             for (j = 1; j <= n; j++) if (!used[j]) {
19                 if (g[queue[i]][j] < g[id][j]) g[id][j] = g[queue[i]][j];
20             }
21         }
22     }
23
24     for (i = 1; i <= n; i++) if (!used[i] && i != id) {
25         for (j = from; j < tot; j++) {
26             k = queue[j];
27             if (g[i][id] > g[i][k] - g[eg[k]][k]) g[i][id] = g[i][k] - g[eg[k]][k];
28         }
29     }
30 }
31
32 double mdst(int root) {
33     int i, j, k;
34     double sum = 0;
35     memset(used, 0, sizeof(used));
36     for (more = 1; more; ) {
37         more = 0;
38         memset(eg, 0, sizeof(eg));
39         for (i = 1; i <= n; i++) if (!used[i] && i != root) {
40             for (j = 1, k = 0; j <= n; j++) if (!used[j] && i != j)
41                 if (k == 0 || g[j][i] < g[k][i]) k = j;
42             eg[i] = k;
43         }
44
45         memset(pass, 0, sizeof(pass));
46         for (i = 1; i <= n; i++) if (!used[i] && !pass[i] && i != root) combine(i, sum);
47     }
48
49     for (i = 1; i <= n; i++) if (!used[i] && i != root) sum += g[eg[i]][i];
50     return sum;

```



51 } }

## 4.12 有根树的同构

时间复杂度:  $\mathcal{O}(V \log V)$

```

1  const unsigned long long MAGIC = 4423;
2
3  unsigned long long magic[N];
4  std::pair<unsigned long long, int> hash[N];
5
6  void solve(int root) {
7      magic[0] = 1;
8      for (int i = 1; i <= n; ++i) {
9          magic[i] = magic[i - 1] * MAGIC;
10     }
11     std::vector<int> queue;
12     queue.push_back(root);
13     for (int head = 0; head < (int)queue.size(); ++head) {
14         int x = queue[head];
15         for (int i = 0; i < (int)son[x].size(); ++i) {
16             int y = son[x][i];
17             queue.push_back(y);
18         }
19     }
20     for (int index = n - 1; index >= 0; --index) {
21         int x = queue[index];
22         hash[x] = std::make_pair(0, 0);
23
24         std::vector<std::pair<unsigned long long, int> > value;
25         for (int i = 0; i < (int)son[x].size(); ++i) {
26             int y = son[x][i];
27             value.push_back(hash[y]);
28         }
29         std::sort(value.begin(), value.end());
30
31         hash[x].first = hash[x].first * magic[1] + 37;
32         hash[x].second++;
33         for (int i = 0; i < (int)value.size(); ++i) {
34             hash[x].first = hash[x].first * magic[value[i].second] + value[i].first;
35             hash[x].second += value[i].second;
36         }
37         hash[x].first = hash[x].first * magic[1] + 41;
38         hash[x].second++;
39     }
40 }
```

## 4.13 度限制生成树

```

1  int n, m, S, K, ans, cnt, Best[N], fa[N], FE[N];
2  int f[N], p[M], t[M], c[M], o, Cost[N];
3  bool u[M], d[M];
4  pair<int, int> MinCost[N];
5  struct Edge {
6      int a, b, c;
7      bool operator < (const Edge & E) const { return c < E.c; }
8  }E[M];
9  vector<int> SE;
10 inline int F(int x) {
11     return fa[x] == x ? x : fa[x] = F(fa[x]);
```

```

12 }
13 inline void AddEdge(int a, int b, int C) {
14     p[++o] = b; c[o] = C;
15     t[o] = f[a]; f[a] = o;
16 }
17 void dfs(int i, int father) {
18     fa[i] = father;
19     if (father == S) Best[i] = -1;
20     else {
21         Best[i] = i;
22         if (~Best[father] && Cost[Best[father]] > Cost[i]) Best[i] = Best[father];
23     }
24     for (int j = f[i]; j; j = t[j])
25         if (!d[j] && p[j] != father) {
26             Cost[p[j]] = c[j];
27             FE[p[j]] = j;
28             dfs(p[j], i);
29         }
30 }
31 inline bool Kruskal() {
32     cnt = n - 1, ans = 0; o = 1;
33     for (int i = 1; i <= n; i++) fa[i] = i, f[i] = 0;
34     sort(E + 1, E + m + 1);
35     for (int i = 1; i <= m; i++) {
36         if (E[i].b == S) swap(E[i].a, E[i].b);
37         if (E[i].a != S && F(E[i].a) != F(E[i].b)) {
38             fa[F(E[i].a)] = F(E[i].b);
39             ans += E[i].c;
40             cnt--;
41             u[i] = true;
42             AddEdge(E[i].a, E[i].b, E[i].c);
43             AddEdge(E[i].b, E[i].a, E[i].c);
44         }
45     }
46     for (int i = 1; i <= n; i++) MinCost[i] = make_pair(INF, INF);
47     for (int i = 1; i <= m; i++)
48         if (E[i].a == S) {
49             SE.push_back(i);
50             MinCost[F(E[i].b)] = min(MinCost[F(E[i].b)], make_pair(E[i].c, i));
51         }
52     int dif = 0;
53     for (int i = 1; i <= n; i++)
54         if (i != S && fa[i] == i) {
55             if (MinCost[i].second == INF) return false;
56             if (++dif > K) return false;
57             dfs(E[MinCost[i].second].b, S);
58             u[MinCost[i].second] = true;
59             ans += MinCost[i].first;
60         }
61     return true;
62 }
63 bool Solve() {
64     memset(d, false, sizeof d);
65     memset(u, false, sizeof u);
66     if (!Kruskal()) return false;
67     for (int i = cnt + 1; i <= K && i <= n; i++) {
68         int MinD = INF, MinID = -1;
69         for (int j = (int) SE.size() - 1; j >= 0; j--)
70             if (u[SE[j]])
71                 SE.erase(SE.begin() + j);
72         for (int j = 0; j < (int) SE.size(); j++) {
73             int tmp = E[SE[j]].c - Cost[Best[E[SE[j]].b]];
74             if (tmp < MinD) {

```

```

75         MinD = tmp;
76         MinID= SE[j];
77     }
78 }
79 if (MinID == -1) return true;
80 if (MinD >= 0) break;
81 ans += MinD;
82 u[MinID] = true;
83 d[FE[Best[E[MinID].b]]] = d[FE[Best[E[MinID].b]] ^ 1] = true;
84 dfs(E[MinID].b, S);
85 }
86 return true;
87 }
88 int main(){
89     Solve();
90     return 0;
91 }

```

## 4.14 弦图相关

### 4.14.1 弦图的判定

```

1  int n, m, first[1001], l, next[2000001], where[2000001], f[1001], a[1001], c[1001], L[1001],
    ↪ R[1001],
2  v[1001], idx[1001], pos[1001];
3  bool b[1001][1001];
4
5  inline void makelist(int x, int y){
6      where[++l] = y;
7      next[l] = first[x];
8      first[x] = l;
9  }
10
11 bool cmp(const int &x, const int &y){
12     return(idx[x] < idx[y]);
13 }
14
15 int main(){
16     for (;;)
17     {
18         n = read(); m = read();
19         if (!n && !m) return 0;
20         memset(first, 0, sizeof(first)); l = 0;
21         memset(b, false, sizeof(b));
22         for (int i = 1; i <= m; i++)
23         {
24             int x = read(), y = read();
25             if (x != y && !b[x][y])
26             {
27                 b[x][y] = true; b[y][x] = true;
28                 makelist(x, y); makelist(y, x);
29             }
30         }
31         memset(f, 0, sizeof(f));
32         memset(L, 0, sizeof(L));
33         memset(R, 255, sizeof(R));
34         L[0] = 1; R[0] = n;
35         for (int i = 1; i <= n; i++) c[i] = i, pos[i] = i;
36         memset(idx, 0, sizeof(idx));
37         memset(v, 0, sizeof(v));
38         for (int i = n; i; --i)
39         {

```

```

40     int now = c[i];
41     R[f[now]]--;
42     if (R[f[now]] < L[f[now]]) R[f[now]] = -1;
43     idx[now] = i; v[i] = now;
44     for (int x = first[now]; x; x = next[x])
45         if (!idx[where[x]])
46             {
47                 swap(c[pos[where[x]]], c[R[f[where[x]]]]);
48                 pos[c[pos[where[x]]]] = pos[where[x]];
49                 pos[where[x]] = R[f[where[x]]];
50                 L[f[where[x]] + 1] = R[f[where[x]]]--;
51                 if (R[f[where[x]]] < L[f[where[x]]]) R[f[where[x]]] = -1;
52                 if (R[f[where[x]] + 1] == -1)
53                     R[f[where[x]] + 1] = L[f[where[x]] + 1];
54                 ++f[where[x]];
55             }
56     }
57     bool ok = true;
58     //v 是完美消除序列.
59     for (int i = 1; i <= n && ok; i++)
60     {
61         int cnt = 0;
62         for (int x = first[v[i]]; x; x = next[x])
63             if (idx[where[x]] > i) c[++cnt] = where[x];
64         sort(c + 1, c + cnt + 1, cmp);
65         bool can = true;
66         for (int j = 2; j <= cnt; j++)
67             if (!b[c[1]][c[j]])
68                 {
69                     ok = false;
70                     break;
71                 }
72     }
73     if (ok) printf("Perfect\n");
74     else printf("Imperfect\n");
75     printf("\n");
76 }
77 }

```

#### 4.14.2 弦图的团数

```

1  int n, m, first[100001], next[2000001], where[2000001], l, L[100001], R[100001], c[100001],
   ↪ f[100001],
2  pos[100001], idx[100001], v[100001], ans;
3
4  inline void makelist(int x, int y){
5      where[++l] = y;
6      next[l] = first[x];
7      first[x] = l;
8  }
9
10 int read(){
11     char ch;
12     for (ch = getchar(); ch < '0' || ch > '9'; ch = getchar());
13     int cnt = 0;
14     for (; ch >= '0' && ch <= '9'; ch = getchar()) cnt = cnt * 10 + ch - '0';
15     return(cnt);
16 }
17
18 int main(){
19     //freopen("1006.in", "r", stdin);
20     //freopen("1006.out", "w", stdout);

```

```

21     memset(first, 0, sizeof(first)); l = 0;
22     n = read(); m = read();
23     for (int i = 1; i <= m; i++)
24     {
25         int x, y;
26         x = read(); y = read();
27         makelist(x, y); makelist(y, x);
28     }
29     memset(L, 0, sizeof(L));
30     memset(R, 255, sizeof(R));
31     memset(f, 0, sizeof(f));
32     memset(idx, 0, sizeof(idx));
33     for (int i = 1; i <= n; i++) c[i] = i, pos[i] = i;
34     L[0] = 1; R[0] = n; ans = 0;
35     for (int i = n; i; --i)
36     {
37         int now = c[i], cnt = 1;
38         idx[now] = i; v[i] = now;
39         if (--R[f[now]] < L[f[now]]) R[f[now]] = -1;
40         for (int x = first[now]; x; x = next[x])
41             if (!idx[where[x]])
42             {
43                 swap(c[pos[where[x]]], c[R[f[where[x]]]]);
44                 pos[c[pos[where[x]]]] = pos[where[x]];
45                 pos[where[x]] = R[f[where[x]]];
46                 L[f[where[x]] + 1] = R[f[where[x]]]--;
47                 if (R[f[where[x]]] < L[f[where[x]]]) R[f[where[x]]] = -1;
48                 if (R[f[where[x]] + 1] == -1) R[f[where[x]] + 1] = L[f[where[x]] + 1];
49                 ++f[where[x]];
50             }
51         else ++cnt;
52         ans = max(ans, cnt);
53     }
54     printf("%d\n", ans);
55 }

```

## 4.15 哈密尔顿回路 (ORE 性质的图)

ORE 性质:

$$\forall x, y \in V \wedge (x, y) \notin E \quad s.t. \quad deg_x + deg_y \geq n$$

返回结果: 从顶点 1 出发的一个哈密尔顿回路

使用条件:  $n \geq 3$

```

1     int left[N], right[N], next[N], last[N];
2
3     void cover(int x) {
4         left[right[x]] = left[x];
5         right[left[x]] = right[x];
6     }
7
8     int adjacent(int x) {
9         for (int i = right[0]; i <= n; i = right[i]) {
10             if (graph[x][i]) {
11                 return i;
12             }
13         }
14         return 0;
15     }
16
17     std::vector<int> solve() {
18         for (int i = 1; i <= n; ++i) {

```

```

19     left[i] = i - 1;
20     right[i] = i + 1;
21 }
22 int head, tail;
23 for (int i = 2; i <= n; ++i) {
24     if (graph[1][i]) {
25         head = 1;
26         tail = i;
27         cover(head);
28         cover(tail);
29         next[head] = tail;
30         break;
31     }
32 }
33 while (true) {
34     int x;
35     while (x = adjacent(head)) {
36         next[x] = head;
37         head = x;
38         cover(head);
39     }
40     while (x = adjacent(tail)) {
41         next[tail] = x;
42         tail = x;
43         cover(tail);
44     }
45     if (!graph[head][tail]) {
46         for (int i = head, j; i != tail; i = next[i]) {
47             if (graph[head][next[i]] && graph[tail][i]) {
48                 for (j = head; j != i; j = next[j]) {
49                     last[next[j]] = j;
50                 }
51                 j = next[head];
52                 next[head] = next[i];
53                 next[tail] = i;
54                 tail = j;
55                 for (j = i; j != head; j = last[j]) {
56                     next[j] = last[j];
57                 }
58                 break;
59             }
60         }
61     }
62     next[tail] = head;
63     if (right[0] > n) {
64         break;
65     }
66     for (int i = head; i != tail; i = next[i]) {
67         if (adjacent(i)) {
68             head = next[i];
69             tail = i;
70             next[tail] = 0;
71             break;
72         }
73     }
74 }
75 std::vector<int> answer;
76 for (int i = head; ; i = next[i]) {
77     if (i == 1) {
78         answer.push_back(i);
79         for (int j = next[i]; j != i; j = next[j]) {
80             answer.push_back(j);
81         }

```

```
82         answer.push_back(i);
83         break;
84     }
85     if (i == tail) {
86         break;
87     }
88 }
89 return answer;
90 }
```

## Chapter 5

# 字符串

### 5.1 模式串匹配

```
1 void build(char *pattern) {
2     int length = (int)strlen(pattern + 1);
3     fail[0] = -1;
4     for (int i = 1, j; i <= length; ++i) {
5         for (j = fail[i - 1]; j != -1 && pattern[i] != pattern[j + 1]; j = fail[j]);
6         fail[i] = j + 1;
7     }
8 }
9
10 void solve(char *text, char *pattern) {
11     int length = (int)strlen(text + 1);
12     for (int i = 1, j; i <= length; ++i) {
13         for (j = match[i - 1]; j != -1 && text[i] != pattern[j + 1]; j = fail[j]);
14         match[i] = j + 1;
15     }
16 }
```

### 5.2 坚固的模式串匹配

```
1 lenA = strlen(A); lenB = strlen(B);
2 nxt[0] = lenB, nxt[1] = lenB - 1;
3 for (int i = 0; i <= lenB; i++)
4     if (B[i] != B[i + 1]) {nxt[1] = i; break;}
5 int j, k = 1, p, L;
6 for (int i = 2; i < lenB; i++) {
7     p = k + nxt[k] - 1; L = nxt[i - k];
8     if (i + L <= p) nxt[i] = L;
9     else {
10         j = p - i + 1;
11         if (j < 0) j = 0;
12         while (i + j < lenB && B[i + j] == B[j]) j++;
13         nxt[i] = j; k = i;
14     }
15 }
16 int minlen = lenA <= lenB ? lenA : lenB; ex[0] = minlen;
17 for (int i = 0; i < minlen; i++)
18     if (A[i] != B[i]) {ex[0] = i; break;}
19 k = 0;
20 for (int i = 1; i < lenA; i++){
21     p = k + ex[k] - 1; L = next[i - k];
22     if (i + L <= p) ex[i] = L;
23     else {
24         j = p - i + 1;
25         if (j < 0) j = 0;
```



```

26         while (i + j < lenA && j < lenB && A[i + j] == B[j]) j++;
27         ex[i] = j; k = i;
28     }
29 }

```

## 5.3 AC 自动机

```

1  int size, c[MAXT][26], f[MAXT], fail[MAXT], d[MAXT];
2
3  int alloc() {
4      size++;
5      std::fill(c[size], c[size] + 26, 0);
6      f[size] = fail[size] = d[size] = 0;
7      return size;
8  }
9
10 void insert(char *s) {
11     int len = strlen(s + 1), p = 1;
12     for (int i = 1; i <= len; i++) {
13         if (c[p][s[i] - 'a']) p = c[p][s[i] - 'a'];
14         else{
15             int newnode = alloc();
16             c[p][s[i] - 'a'] = newnode;
17             d[newnode] = s[i] - 'a';
18             f[newnode] = p;
19             p = newnode;
20         }
21     }
22 }
23
24 void buildfail() {
25     static int q[MAXT];
26     int left = 0, right = 0;
27     fail[1] = 0;
28     for (int i = 0; i < 26; i++) {
29         c[0][i] = 1;
30         if (c[1][i]) q[++right] = c[1][i];
31     }
32     while (left < right) {
33         left++;
34         int p = fail[f[q[left]]];
35         while (!c[p][d[q[left]]]) p = fail[p];
36         fail[q[left]] = c[p][d[q[left]]];
37         for (int i = 0; i < 26; i++) {
38             if (c[q[left]][i]) {
39                 q[++right] = c[q[left]][i];
40             }
41         }
42     }
43     for (int i = 1; i <= size; i++)
44         for (int j = 0; j < 26; j++) {
45             int p = i;
46             while (!c[p][j]) p = fail[p];
47             c[i][j] = c[p][j];
48         }
49 }

```

## 5.4 后缀数组

```

1  namespace suffix_array{
2      int wa[MAXN], wb[MAXN], ws[MAXN], wv[MAXN];
3      bool cmp(int *r, int a, int b, int l) {
4          return r[a] == r[b] && r[a + l] == r[b + l];
5      }
6      void DA(int *r, int *sa, int n, int m) {
7          int *x = wa, *y = wb, *t;
8          for (int i = 0; i < m; i++) ws[i] = 0;
9          for (int i = 0; i < n; i++) ws[x[i]]++;
10         for (int i = 1; i < m; i++) ws[i] += ws[i - 1];
11         for (int i = n - 1; i >= 0; i--) sa[--ws[x[i]]] = i;
12         for (int i, j = 1, p = 1; p < n; j <= 1, m = p) {
13             for (p = 0, i = n - j; i < n; i++) y[p++] = i;
14             for (i = 0; i < n; i++) if (sa[i] >= j) y[p++] = sa[i] - j;
15             for (i = 0; i < n; i++) wv[i] = x[y[i]];
16             for (i = 0; i < m; i++) ws[i] = 0;
17             for (i = 0; i < n; i++) ws[wv[i]]++;
18             for (i = 1; i < m; i++) ws[i] += ws[i - 1];
19             for (i = n - 1; i >= 0; i--) sa[--ws[wv[i]]] = y[i];
20             for (t = x, x = y, y = t, p = 1, x[sa[0]] = 0, i = 1; i < n; i++)
21                 x[sa[i]] = cmp(y, sa[i - 1], sa[i], j) ? p - 1 : p++;
22         }
23     }
24     void getheight(int *r, int *sa, int *rk, int *h, int n) {
25         for (int i = 1; i <= n; i++) rk[sa[i]] = i;
26         for (int i = 0, j, k = 0; i < n; h[rk[i++]] = k)
27             for (k ? k-- : 0, j = sa[rk[i] - 1]; r[i + k] == r[j + k]; k++);
28     }
29 };

```

## 5.5 广义后缀自动机

```

1  // Generalized Suffix Automaton
2  void add(int x, int &last) {
3      int lastnode = last;
4      if (c[lastnode][x]) {
5          int nownode = c[lastnode][x];
6          if (l[nownode] == l[lastnode] + 1) last = nownode;
7          else{
8              int auxnode = ++size; l[auxnode] = l[lastnode] + 1;
9              for (int i = 0; i < 26; i++) c[auxnode][i] = c[nownode][i];
10             f[auxnode] = f[nownode]; f[nownode] = auxnode;
11             for (; lastnode && c[lastnode][x] == nownode; lastnode = f[lastnode]) {
12                 c[lastnode][x] = auxnode;
13             }
14             last = auxnode;
15         }
16     }
17     else{
18         int newnode = ++size; l[newnode] = l[lastnode] + 1;
19         for (; lastnode && !c[lastnode][x]; lastnode = f[lastnode]) c[lastnode][x] = newnode;
20         if (!lastnode) f[newnode] = 1;
21         else{
22             int nownode = c[lastnode][x];
23             if (l[lastnode] + 1 == l[nownode]) f[newnode] = nownode;
24             else{
25                 int auxnode = ++size; l[auxnode] = l[lastnode] + 1;
26                 for (int i = 0; i < 26; i++) c[auxnode][i] = c[nownode][i];
27                 f[auxnode] = f[nownode]; f[nownode] = f[newnode] = auxnode;

```

```

28         for (; lastnode && c[lastnode][x] == nownode; lastnode = f[lastnode]) {
29             c[lastnode][x] = auxnode;
30         }
31     }
32 }
33 last = newnode;
34 }
35 }

```

## 5.6 Manacher 算法

```

1 void manacher(char *text, int length) {
2     palindrome[0] = 1;
3     for (int i = 1, j = 0; i < length; ++i) {
4         if (j + palindrome[j] <= i) {
5             palindrome[i] = 0;
6         } else {
7             palindrome[i] = std::min(palindrome[(j << 1) - i], j + palindrome[j] - i);
8         }
9         while (i - palindrome[i] >= 0 && i + palindrome[i] < length
10             && text[i - palindrome[i]] == text[i + palindrome[i]]) {
11             palindrome[i]++;
12         }
13         if (i + palindrome[i] > j + palindrome[j]) {
14             j = i;
15         }
16     }
17 }

```

## 5.7 回文树

```

1 struct Palindromic_Tree{
2     int nTree, nStr, last, c[MAXT][26], fail[MAXT], r[MAXN], l[MAXN], s[MAXN];
3     int allocate(int len) {
4         l[nTree] = len;
5         r[nTree] = 0;
6         fail[nTree] = 0;
7         memset(c[nTree], 0, sizeof(c[nTree]));
8         return nTree++;
9     }
10    void init() {
11        nTree = nStr = 0;
12        int newEven = allocate(0);
13        int newOdd = allocate(-1);
14        last = newEven;
15        fail[newEven] = newOdd;
16        fail[newOdd] = newEven;
17        s[0] = -1;
18    }
19    void add(int x) {
20        s[++nStr] = x;
21        int nownode = last;
22        while (s[nStr - l[nownode] - 1] != s[nStr]) nownode = fail[nownode];
23        if (!c[nownode][x]) {
24            int newnode = allocate(l[nownode] + 2), &newfail = fail[newnode];
25            newfail = fail[nownode];
26            while (s[nStr - l[newfail] - 1] != s[nStr]) newfail = fail[newfail];
27            newfail = c[newfail][x];
28            c[nownode][x] = newnode;

```

```
29     }
30     last = c[nownode][x];
31     r[last]++;
32 }
33 void count() {
34     for (int i = nTree - 1; i >= 0; i--) {
35         r[fail[i]] += r[i];
36     }
37 }
38 }
```

## 5.8 循环串最小表示

```
1  int solve(char *text, int length) {
2      int i = 0, j = 1, delta = 0;
3      while (i < length && j < length && delta < length) {
4          char tokeni = text[(i + delta) % length];
5          char tokenj = text[(j + delta) % length];
6          if (tokeni == tokenj) {
7              delta++;
8          } else {
9              if (tokeni > tokenj) {
10                 i += delta + 1;
11             } else {
12                 j += delta + 1;
13             }
14             if (i == j) {
15                 j++;
16             }
17             delta = 0;
18         }
19     }
20     return std::min(i, j);
21 }
```

# Chapter 6

## 计算几何

### 6.1 二维基础

#### 6.1.1 点类

```
1 struct Point{
2     double x, y;
3     Point() {}
4     Point(double x, double y):x(x), y(y) {}
5     Point operator +(const Point &p)const {
6         return Point(x + p.x, y + p.y);
7     }
8     Point operator -(const Point &p)const {
9         return Point(x - p.x, y - p.y);
10    }
11    Point operator *(const double &p)const {
12        return Point(x * p, y * p);
13    }
14    Point operator /(const double &p)const {
15        return Point(x / p, y / p);
16    }
17    int read() {
18        return scanf("%lf%lf", &x, &y);
19    }
20 };
21
22 struct Line{
23     Point a, b;
24     Line() {}
25     Line(Point a, Point b):a(a), b(b) {}
26 };
```

#### 6.1.2 凸包

```
1 bool Pair_Comp(const Point &a, const Point &b) {
2     if (dcmp(a.x - b.x) < 0) return true;
3     if (dcmp(a.x - b.x) > 0) return false;
4     return dcmp(a.y - b.y) < 0;
5 }
6
7 int Convex_Hull(int n, Point *P, Point *C) {
8     sort(P, P + n, Pair_Comp);
9     int top = 0;
10    for (int i = 0; i < n; i++) {
11        while (top >= 2 && dcmp(det(C[top - 1] - C[top - 2], P[i] - C[top - 2])) <= 0) top--;
12        C[top++] = P[i];
13    }
14    int lasttop = top;
```

```

15     for (int i = n - 1; i >= 0; i--) {
16         while (top > lasttop && dcmp(det(C[top - 1] - C[top - 2], P[i] - C[top - 2])) <= 0)
    ↪ top--;
17         C[top++] = P[i];
18     }
19     return top;
20 }

```

### 6.1.3 半平面交

```

1  bool isOnLeft(const Point &x, const Line &l) {
2      double d = det(x - l.a, l.b - l.a);
3      return dcmp(d) <= 0;
4  }
5  // 传入一个线段的集合 L, 传出 A, 并且返回 A 的大小
6  int getIntersectionOfHalfPlane(int n, Line *L, Line *A) {
7      Line *q = new Line[n + 1];
8      Point *p = new Point[n + 1];
9      sort(L, L + n, Polar_Angle_Comp_Line);
10     int l = 1, r = 0;
11     for (int i = 0; i < n; i++) {
12         while (l < r && !isOnLeft(p[r - 1], L[i])) r--;
13         while (l < r && !isOnLeft(p[l], L[i])) l++;
14         q[++r] = L[i];
15         if (l < r && is_Colinear(q[r], q[r - 1])) {
16             r--;
17             if (isOnLeft(L[i].a, q[r])) q[r] = L[i];
18         }
19         if (l < r) p[r - 1] = getIntersection(q[r - 1], q[r]);
20     }
21     while (l < r && !isOnLeft(p[r - 1], q[l])) r--;
22     if (r - l + 1 <= 2) return 0;
23     int tot = 0;
24     for (int i = l; i <= r; i++) A[tot++] = q[i];
25     return tot;
26 }

```

### 6.1.4 最近点对

```

1  7bool comparex(const Point &a, const Point &b) {
2      return sgn(a.x - b.x) < 0;
3  }
4
5  bool comparey(const Point &a, const Point &b) {
6      return sgn(a.y - b.y) < 0;
7  }
8
9  double solve(const std::vector<Point> &point, int left, int right) {
10     if (left == right) {
11         return INF;
12     }
13     if (left + 1 == right) {
14         return dist(point[left], point[right]);
15     }
16     int mid = left + right >> 1;
17     double result = std::min(solve(left, mid), solve(mid + 1, right));
18     std::vector<Point> candidate;
19     for (int i = left; i <= right; ++i) {
20         if (std::abs(point[i].x - point[mid].x) <= result) {
21             candidate.push_back(point[i]);
22         }

```

```

23     }
24     std::sort(candidate.begin(), candidate.end(), comparey);
25     for (int i = 0; i < (int)candidate.size(); ++i) {
26         for (int j = i + 1; j < (int)candidate.size(); ++j) {
27             if (std::abs(candidate[i].y - candidate[j].y) >= result) {
28                 break;
29             }
30             result = std::min(result, dist(candidate[i], candidate[j]));
31         }
32     }
33     return result;
34 }
35
36 double solve(std::vector<Point> point) {
37     std::sort(point.begin(), point.end(), comparex);
38     return solve(point, 0, (int)point.size() - 1);
39 }

```

## 6.2 三维基础

### 6.2.1 点类

```

1  int dcmp(const double &x) {
2      return fabs(x) < eps ? 0 : (x > 0 ? 1 : -1);
3  }
4
5  struct TPoint{
6      double x, y, z;
7      TPoint() {}
8      TPoint(double x, double y, double z) : x(x), y(y), z(z) {}
9      TPoint operator +(const TPoint &p) const {
10         return TPoint(x + p.x, y + p.y, z + p.z);
11     }
12     TPoint operator -(const TPoint &p) const {
13         return TPoint(x - p.x, y - p.y, z - p.z);
14     }
15     TPoint operator *(const double &p) const {
16         return TPoint(x * p, y * p, z * p);
17     }
18     TPoint operator /(const double &p) const {
19         return TPoint(x / p, y / p, z / p);
20     }
21     bool operator <(const TPoint &p) const {
22         int dX = dcmp(x - p.x), dY = dcmp(y - p.y), dZ = dcmp(z - p.z);
23         return dX < 0 || (dX == 0 && (dY < 0 || (dY == 0 && dZ < 0)));
24     }
25     bool read() {
26         return scanf("%lf%lf%lf", &x, &y, &z) == 3;
27     }
28 };
29
30 double sqrdist(const TPoint &a) {
31     double ret = 0;
32     ret += a.x * a.x;
33     ret += a.y * a.y;
34     ret += a.z * a.z;
35     return ret;
36 }
37 double sqrdist(const TPoint &a, const TPoint &b) {
38     double ret = 0;
39     ret += (a.x - b.x) * (a.x - b.x);
40     ret += (a.y - b.y) * (a.y - b.y);

```





```

34         if (!check) return false;
35         check = false;
36         for (int i = 2; i < (int)p.size(); i++) {
37             if (dcmp(sqrdist(det(p[i] - p[0], p[1] - p[0])))) {
38                 std::swap(p[2], p[i]);
39                 check = true;
40                 break;
41             }
42         }
43         if (!check) return false;
44         check = false;
45         for (int i = 3; i < (int)p.size(); i++) {
46             if (dcmp(detdot(p[0], p[1], p[2], p[i]))) {
47                 std::swap(p[3], p[i]);
48                 check = true;
49                 break;
50             }
51         }
52         if (!check) return false;
53         for (int i = 0; i < (int)p.size(); i++)
54             for (int j = 0; j < (int)p.size(); j++) {
55                 whe[i][j] = -1;
56             }
57         return true;
58     }
59
60     void pushface(const int &a, const int &b, const int &c) {
61         nFace++;
62         tmp[nFace] = Face(a, b, c);
63         tmp[nFace].isOnConvex = true;
64         whe[a][b] = nFace;
65         whe[b][c] = nFace;
66         whe[c][a] = nFace;
67     }
68
69     bool deal(const std::vector<TPoint> &p, const std::pair<int, int> &now, const TPoint &base)
70     {
71         int id = whe[now.second][now.first];
72         if (!tmp[id].isOnConvex) return true;
73         if (isVisible(p, tmp[id], base)) {
74             queue[++right] = tmp[id];
75             tmp[id].isOnConvex = false;
76             return true;
77         }
78         return false;
79     }
80
81     std::vector<Triangle> getConvex(std::vector<TPoint> &p) {
82         static std::vector<Triangle> ret;
83         ret.clear();
84         if (!init(p)) return ret;
85         if (!isVisible(p, Face(0, 1, 2), p[3])) pushface(0, 1, 2); else pushface(0, 2, 1);
86         if (!isVisible(p, Face(0, 1, 3), p[2])) pushface(0, 1, 3); else pushface(0, 3, 1);
87         if (!isVisible(p, Face(0, 2, 3), p[1])) pushface(0, 2, 3); else pushface(0, 3, 2);
88         if (!isVisible(p, Face(1, 2, 3), p[0])) pushface(1, 2, 3); else pushface(1, 3, 2);
89         for (int a = 4; a < (int)p.size(); a++) {
90             TPoint base = p[a];
91             for (int i = 1; i <= nFace; i++) {
92                 if (tmp[i].isOnConvex && isVisible(p, tmp[i], base)) {
93                     left = 0, right = 0;
94                     queue[++right] = tmp[i];
95                     tmp[i].isOnConvex = false;
96                     while (left < right) {

```

```

96         Face now = queue[++left];
97         if (!deal(p, std::make_pair(now.a, now.b), base)) pushface(now.a,
↪ now.b, a);
98         if (!deal(p, std::make_pair(now.b, now.c), base)) pushface(now.b,
↪ now.c, a);
99         if (!deal(p, std::make_pair(now.c, now.a), base)) pushface(now.c,
↪ now.a, a);
100     }
101     break;
102 }
103 }
104 }
105 for (int i = 1; i <= nFace; i++) {
106     Face now = tmp[i];
107     if (now.isOnConvex) {
108         ret.push_back(Triangle(p[now.a], p[now.b], p[now.c]));
109     }
110 }
111 return ret;
112 }
113 };
114
115 int n;
116 std::vector<TPoint> p;
117 std::vector<Triangle> answer;
118
119 int main() {
120     scanf("%d", &n);
121     for (int i = 1; i <= n; i++) {
122         TPoint a;
123         a.read();
124         p.push_back(a);
125     }
126     answer = Convex_Hull::getConvex(p);
127     double areaCounter = 0.0;
128     for (int i = 0; i < (int)answer.size(); i++) {
129         areaCounter += answer[i].getArea();
130     }
131     printf("%.3f\n", areaCounter);
132     return 0;
133 }

```

### 6.2.3 绕轴旋转

使用方法及注意事项：逆时针绕轴  $AB$  旋转  $\theta$  角

```

1  Matrix getTrans(const double &a, const double &b, const double &c) {
2      Matrix ret;
3      ret.a[0][0] = 1; ret.a[0][1] = 0; ret.a[0][2] = 0; ret.a[0][3] = 0;
4      ret.a[1][0] = 0; ret.a[1][1] = 1; ret.a[1][2] = 0; ret.a[1][3] = 0;
5      ret.a[2][0] = 0; ret.a[2][1] = 0; ret.a[2][2] = 1; ret.a[2][3] = 0;
6      ret.a[3][0] = a; ret.a[3][1] = b; ret.a[3][2] = c; ret.a[3][3] = 1;
7      return ret;
8  }
9  Matrix getRotate(const double &a, const double &b, const double &c, const double &theta) {
10     Matrix ret;
11     ret.a[0][0] = a * a * (1 - cos(theta)) + cos(theta);
12     ret.a[0][1] = a * b * (1 - cos(theta)) + c * sin(theta);
13     ret.a[0][2] = a * c * (1 - cos(theta)) - b * sin(theta);
14     ret.a[0][3] = 0;
15
16     ret.a[1][0] = b * a * (1 - cos(theta)) - c * sin(theta);
17     ret.a[1][1] = b * b * (1 - cos(theta)) + cos(theta);

```

```

18     ret.a[1][2] = b * c * (1 - cos(theta)) + a * sin(theta);
19     ret.a[1][3] = 0;
20
21     ret.a[2][0] = c * a * (1 - cos(theta)) + b * sin(theta);
22     ret.a[2][1] = c * b * (1 - cos(theta)) - a * sin(theta);
23     ret.a[2][2] = c * c * (1 - cos(theta)) + cos(theta);
24     ret.a[2][3] = 0;
25
26     ret.a[3][0] = 0;
27     ret.a[3][1] = 0;
28     ret.a[3][2] = 0;
29     ret.a[3][3] = 1;
30     return ret;
31 }
32 Matrix getRotate(const double &ax, const double &ay, const double &az, const double &bx, const
↪ double &by, const double &bz, const double &theta) {
33     double l = dist(Point(0, 0, 0), Point(bx, by, bz));
34     Matrix ret = getTrans(-ax, -ay, -az);
35     ret = ret * getRotate(bx / l, by / l, bz / l, theta);
36     ret = ret * getTrans(ax, ay, az);
37     return ret;
38 }

```

## 6.3 多边形

### 6.3.1 判断点在多边形内部

```

1 bool point_on_line(const Point &p, const Point &a, const Point &b) {
2     return sgn(det(p, a, b)) == 0 && sgn(dot(p, a, b)) <= 0;
3 }
4
5 bool point_in_polygon(const Point &p, const std::vector<Point> &polygon) {
6     int counter = 0;
7     for (int i = 0; i < (int)polygon.size(); ++i) {
8         Point a = polygon[i], b = polygon[(i + 1) % (int)polygon.size()];
9         if (point_on_line(p, a, b)) {
10             // Point on the boundary are excluded.
11             return false;
12         }
13         int x = sgn(det(a, p, b));
14         int y = sgn(a.y - p.y);
15         int z = sgn(b.y - p.y);
16         counter += (x > 0 && y <= 0 && z > 0);
17         counter -= (x < 0 && z <= 0 && y > 0);
18     }
19     return counter;
20 }

```

### 6.3.2 多边形内整点计数

```

1 int getInside(int n, Point *P) { // 求多边形 P 内有多少个整数点
2     int OnEdge = n;
3     double area = getArea(n, P);
4     for (int i = 0; i < n - 1; i++) {
5         Point now = P[i + 1] - P[i];
6         int y = (int)now.y, x = (int)now.x;
7         OnEdge += abs(gcd(x, y)) - 1;
8     }
9     Point now = P[0] - P[n - 1];
10    int y = (int)now.y, x = (int)now.x;
11    OnEdge += abs(gcd(x, y)) - 1;

```

```

12     double ret = area - (double)OnEdge / 2 + 1;
13     return (int)ret;
14 }

```

## 6.4 圆

### 6.4.1 最小覆盖圆

```

1  Point getmid(Point a,Point b) {
2      return Point((a.x + b.x) / 2, (a.y + b.y) / 2);
3  }
4  Point getcross(Point a, Point vA, Point b, Point vB) {
5      Point u = a - b;
6      double t = det(vB, u) / det(vA, vB);
7      return a + vA * t;
8  }
9  Point getcir(Point a,Point b,Point c) {
10     Point midA = getmid(a,b), vA = Point(-(b - a).y, (b - a).x);
11     Point midB = getmid(b,c), vB = Point(-(c - b).y, (c - b).x);
12     return getcross(midA, vA, midB, vB);
13 }
14 double mincir(Point *p,int n) {
15     std::random_shuffle(p + 1, p + n + 1);
16     Point O = p[1];
17     double r = 0;
18     for (int i = 2; i <= n; i++) {
19         if (dist(O, p[i]) <= r) continue;
20         O = p[i]; r = 0;
21         for (int j = 1; j < i; j++) {
22             if (dist(O, p[j]) <= r) continue;
23             O = getmid(p[i], p[j]); r = dist(O,p[i]);
24             for (int k = 1; k < j; k++) {
25                 if (dist(O,p[k]) <= r) continue;
26                 O = getcir(p[i], p[j], p[k]);
27                 r = dist(O,p[i]);
28             }
29         }
30     }
31     return r;
32 }

```

### 6.4.2 最小覆盖球

```

1  double eps(1e-8);
2  int sign(const double & x) {
3      return (x > eps) - (x + eps < 0);
4  }
5  bool equal(const double & x, const double & y) {
6      return x + eps > y and y + eps > x;
7  }
8  struct Point {
9      double x, y, z;
10     Point() {}
11     Point(const double & x, const double & y, const double & z) : x(x), y(y), z(z){
12     }
13     void scan() {
14         scanf("%lf%lf%lf", &x, &y, &z);
15     }
16     double sqrlen() const {
17         return x * x + y * y + z * z;
18     }

```

```

19     }
20     double len() const {
21         return sqrt(sqrlen());
22     }
23     void print() const {
24         printf("(%lf %lf %lf)\n", x, y, z);
25     }
26 } a[33];
27 Point operator + (const Point & a, const Point & b) {
28     return Point(a.x + b.x, a.y + b.y, a.z + b.z);
29 }
30 Point operator - (const Point & a, const Point & b) {
31     return Point(a.x - b.x, a.y - b.y, a.z - b.z);
32 }
33 Point operator * (const double & x, const Point & a) {
34     return Point(x * a.x, x * a.y, x * a.z);
35 }
36 double operator % (const Point & a, const Point & b) {
37     return a.x * b.x + a.y * b.y + a.z * b.z;
38 }
39 Point operator * (const Point & a, const Point & b) {
40     return Point(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z, a.x * b.y - a.y * b.x);
41 }
42 struct Circle {
43     double r;
44     Point o;
45     Circle() {
46         o.x = o.y = o.z = r = 0;
47     }
48     Circle(const Point & o, const double & r) : o(o), r(r) {
49     }
50     void scan() {
51         o.scan();
52         scanf("%lf", &r);
53     }
54     void print() const {
55         o.print();
56         printf("%lf\n", r);
57     }
58 };
59 struct Plane {
60     Point nor;
61     double m;
62     Plane(const Point & nor, const Point & a) : nor(nor){
63         m = nor % a;
64     }
65 };
66 Point intersect(const Plane & a, const Plane & b, const Plane & c) {
67     Point c1(a.nor.x, b.nor.x, c.nor.x), c2(a.nor.y, b.nor.y, c.nor.y), c3(a.nor.z, b.nor.z,
68     ↪ c.nor.z), c4(a.m, b.m, c.m);
69     return 1 / ((c1 * c2) % c3) * Point((c4 * c2) % c3, (c1 * c4) % c3, (c1 * c2) % c4);
70 }
71 bool in(const Point & a, const Circle & b) {
72     return sign((a - b.o).len() - b.r) <= 0;
73 }
74 bool operator < (const Point & a, const Point & b) {
75     if(!equal(a.x, b.x)) {
76         return a.x < b.x;
77     }
78     if(!equal(a.y, b.y)) {
79         return a.y < b.y;
80     }
81     if(!equal(a.z, b.z)) {

```

```

81         return a.z < b.z;
82     }
83     return false;
84 }
85 bool operator == (const Point & a, const Point & b) {
86     return equal(a.x, b.x) and equal(a.y, b.y) and equal(a.z, b.z);
87 }
88 vector<Point> vec;
89 Circle calc() {
90     if(vec.empty()) {
91         return Circle(Point(0, 0, 0), 0);
92     }else if(1 == (int)vec.size()) {
93         return Circle(vec[0], 0);
94     }else if(2 == (int)vec.size()) {
95         return Circle(0.5 * (vec[0] + vec[1]), 0.5 * (vec[0] - vec[1]).len());
96     }else if(3 == (int)vec.size()) {
97         double r((vec[0] - vec[1]).len() * (vec[1] - vec[2]).len() * (vec[2] - vec[0]).len() /
→ 2 / fabs(((vec[0] - vec[2]) * (vec[1] - vec[2])).len()));
98         return Circle(intersect(Plane(vec[1] - vec[0], 0.5 * (vec[1] + vec[0])),
99                               Plane(vec[2] - vec[1], 0.5 * (vec[2] + vec[1])),
100                               Plane((vec[1] - vec[0]) * (vec[2] - vec[0]), vec[0])), r);
101     }else {
102         Point o(intersect(Plane(vec[1] - vec[0], 0.5 * (vec[1] + vec[0])),
103                           Plane(vec[2] - vec[0], 0.5 * (vec[2] + vec[0])),
104                           Plane(vec[3] - vec[0], 0.5 * (vec[3] + vec[0]))));
105         return Circle(o, (o - vec[0]).len());
106     }
107 }
108 Circle miniBall(int n) {
109     Circle res(calc());
110     for(int i(0); i < n; i++) {
111         if(!in(a[i], res)) {
112             vec.push_back(a[i]);
113             res = miniBall(i);
114             vec.pop_back();
115             if(i) {
116                 Point tmp(a[i]);
117                 memmove(a + 1, a, sizeof(Point) * i);
118                 a[0] = tmp;
119             }
120         }
121     }
122     return res;
123 }
124 int main() {
125     int n;
126     for(;;) {
127         scanf("%d", &n);
128         if(!n) {
129             break;
130         }
131         for(int i(0); i < n; i++) {
132             a[i].scan();
133         }
134         sort(a, a + n);
135         n = unique(a, a + n) - a;
136         vec.clear();
137         printf("%.10f\n", miniBall(n).r);
138     }
139 }

```

## 6.4.3 多边形与圆的交面积

```

1 // 求扇形面积
2 double getSectorArea(const Point &a, const Point &b, const double &r) {
3     double c = (2.0 * r * r - sqrdist(a, b)) / (2.0 * r * r);
4     double alpha = acos(c);
5     return r * r * alpha / 2.0;
6 }
7 // 求二次方程  $ax^2 + bx + c = 0$  的解
8 std::pair<double, double> getSolution(const double &a, const double &b, const double &c) {
9     double delta = b * b - 4.0 * a * c;
10    if (dcmp(delta) < 0) return std::make_pair(0, 0);
11    else return std::make_pair((-b - sqrt(delta)) / (2.0 * a), (-b + sqrt(delta)) / (2.0 * a));
12 }
13 // 直线与圆的交点
14 std::pair<Point, Point> getIntersection(const Point &a, const Point &b, const double &r) {
15     Point d = b - a;
16     double A = dot(d, d);
17     double B = 2.0 * dot(d, a);
18     double C = dot(a, a) - r * r;
19     std::pair<double, double> s = getSolution(A, B, C);
20     return std::make_pair(a + d * s.first, a + d * s.second);
21 }
22 // 原点到线段 AB 的距离
23 double getPointDist(const Point &a, const Point &b) {
24     Point d = b - a;
25     int sA = dcmp(dot(a, d)), sB = dcmp(dot(b, d));
26     if (sA * sB <= 0) return det(a, b) / dist(a, b);
27     else return std::min(dist(a), dist(b));
28 }
29 // a 和 b 和原点组成的三角形与半径为 r 的圆的交的面积
30 double getArea(const Point &a, const Point &b, const double &r) {
31     double dA = dot(a, a), dB = dot(b, b), dC = getPointDist(a, b), ans = 0.0;
32     if (dcmp(dA - r * r) <= 0 && dcmp(dB - r * r) <= 0) return det(a, b) / 2.0;
33     Point tA = a / dist(a) * r;
34     Point tB = b / dist(b) * r;
35     if (dcmp(dC - r) > 0) return getSectorArea(tA, tB, r);
36     std::pair<Point, Point> ret = getIntersection(a, b, r);
37     if (dcmp(dA - r * r) > 0 && dcmp(dB - r * r) > 0) {
38         ans += getSectorArea(tA, ret.first, r);
39         ans += det(ret.first, ret.second) / 2.0;
40         ans += getSectorArea(ret.second, tB, r);
41         return ans;
42     }
43     if (dcmp(dA - r * r) > 0) return det(ret.first, b) / 2.0 + getSectorArea(tA, ret.first, r);
44     else return det(a, ret.second) / 2.0 + getSectorArea(ret.second, tB, r);
45 }
46 // 求圆与多边形的交的主过程
47 double getArea(int n, Point *p, const Point &c, const double r) {
48     double ret = 0.0;
49     for (int i = 0; i < n; i++) {
50         int sgn = dcmp(det(p[i] - c, p[(i + 1) % n] - c));
51         if (sgn > 0) ret += getArea(p[i] - c, p[(i + 1) % n] - c, r);
52         else ret -= getArea(p[(i + 1) % n] - c, p[i] - c, r);
53     }
54     return fabs(ret);
55 }

```

# Chapter 7

## 其它

### 7.1 STL 使用方法

#### 7.1.1 nth\_element

用法: `nth_element(a + 1, a + id, a + n + 1);`

作用: 将排名为  $id$  的元素放在第  $id$  个位置。

#### 7.1.2 next\_permutation

用法: `next_permutation(a + 1, a + n + 1);`

作用: 以  $a$  中从小到大排序后为第一个排列, 求得当期数组  $a$  中的下一个排列, 返回值为当期排列是否为最后一个排列。

### 7.2 博弈论相关

#### 7.2.1 巴什博弈

1. 只有一堆  $n$  个物品, 两个人轮流从这堆物品中取物, 规定每次至少取一个, 最多取  $m$  个。最后取光者得胜。
2. 显然, 如果  $n = m + 1$ , 那么由于一次最多只能取  $m$  个, 所以, 无论先取者拿走多少个, 后取者都能够一次拿走剩余的物品, 后者取胜。因此我们发现了如何取胜的法则: 如果  $n = m + 1 \cdot r + s$ , ( $r$  为任意自然数,  $s \leq m$ ), 那么先取者要拿走  $s$  个物品, 如果后取者拿走  $k$  ( $k \leq m$ ) 个, 那么先取者再拿走  $m + 1 - k$  个, 结果剩下  $(m + 1)(r - 1)$  个, 以后保持这样的取法, 那么先取者肯定获胜。总之, 要保持给对手留下  $(m + 1)$  的倍数, 就能最后获胜。

#### 7.2.2 威佐夫博弈

1. 有两堆各若干个物品, 两个人轮流从某一堆或同时从两堆中取同样多的物品, 规定每次至少取一个, 多者不限, 最后取光者得胜。
2. 判断一个局势  $(a, b)$  为奇异局势 (必败态) 的方法:

$$a_k = [k(1 + \sqrt{5})/2] \quad b_k = a_k + k$$

#### 7.2.3 阶梯博弈

1. 博弈在一列阶梯上进行, 每个阶梯上放着自然数个点, 两个人进行阶梯博弈, 每一步则是将一个阶梯上的若干个 (至少一个) 移到前面去, 最后没有点可以移动的人输。
2. 解决方法: 把所有奇数阶梯看成  $N$  堆石子, 做 NIM。(把石子从奇数堆移动到偶数堆可以理解为拿走石子, 就相当于几个奇数堆的石子在做 Nim)

#### 7.2.4 图上删边游戏

##### 链的删边游戏

1. 游戏规则: 对于一条链, 其中一个端点是根, 两人轮流删边, 脱离根的部分也算被删去, 最后没边可删的人输。
2. 做法:  $sg[i] = n - dist(i) - 1$  (其中  $n$  表示总点数,  $dist(i)$  表示离根的距离)



### 树的删边游戏

1. 游戏规则：对于一棵有根树，两人轮流删边，脱离根的部分也算被删去，没边可删的人输。
2. 做法：叶子结点的  $sg = 0$ ，其他节点的  $sg$  等于儿子结点的  $sg + 1$  的异或和。

### 局部连通图的删边游戏

1. 游戏规则：在一个局部连通图上，两人轮流删边，脱离根的部分也算被删去，没边可删的人输。局部连通图的构图规则是，在一棵基础树上加边得到，所有形成的环保证不共用边，且只与基础树有一个公共点。
2. 做法：去掉所有的偶环，将所有的奇环变为长度为 1 的链，然后做树的删边游戏。

## 7.3 Java Reference

```

1  import java.io.*;
2  import java.util.*;
3  import java.math.*;
4
5  public class Main {
6      static int get(char c) {
7          if (c <= '9')
8              return c - '0';
9          else if (c <= 'Z')
10             return c - 'A' + 10;
11         else
12             return c - 'a' + 36;
13     }
14     static char get(int x) {
15         if (x <= 9)
16             return (char)(x + '0');
17         else if (x <= 35)
18             return (char)(x - 10 + 'A');
19         else
20             return (char)(x - 36 + 'a');
21     }
22     static BigInteger get(String s, BigInteger x) {
23         BigInteger ans = BigInteger.valueOf(0), now = BigInteger.valueOf(1);
24         for (int i = s.length() - 1; i >= 0; i--) {
25             ans = ans.add(now.multiply(BigInteger.valueOf(get(s.charAt(i)))));
26             now = now.multiply(x);
27         }
28         return ans;
29     }
30     public static void main(String [] args) {
31         Scanner cin = new Scanner(new BufferedInputStream(System.in));
32         for (; ; ) {
33             BigInteger x = cin.nextBigInteger();
34             if (x.compareTo(BigInteger.valueOf(0)) == 0)
35                 break;
36             String s = cin.next(), t = cin.next(), r = "";
37             BigInteger ans = get(s, x).mod(get(t, x));
38             if (ans.compareTo(BigInteger.valueOf(0)) == 0)
39                 r = "0";
40             for (; ans.compareTo(BigInteger.valueOf(0)) > 0; ) {
41                 r = get(ans.mod(x).intValue()) + r;
42                 ans = ans.divide(x);
43             }
44             System.out.println(r);
45         }
46     }
47 }
48
49 // Arrays
50 int a[];
51 .fill(a[, int fromIndex, int toIndex],val); | .sort(a[, int fromIndex, int toIndex])
52 // String
53 String s;
```

```

54 .charAt(int i); | compareTo(String) | compareToIgnoreCase () | contains(String) |
55 length () | substring(int l, int len)
56 // BigInteger
57 .abs() | .add() | bitLength () | subtract () | divide () | remainder () | divideAndRemainder
    () | modPow(b, c) |
58 pow(int) | multiply () | compareTo () |
59 gcd() | intValue () | longValue () | isProbablePrime(int c) (1 - 1/2^c) |
60 nextProbablePrime () | shiftLeft(int) | valueOf ()
61 // BigDecimal
62 .ROUND_CEILING | ROUND_DOWN_FLOOR | ROUND_HALF_DOWN | ROUND_HALF_EVEN | ROUND_HALF_UP |
    ROUND_UP
63 .divide(BigDecimal b, int scale , int round_mode) | doubleValue () | movePointLeft(int) | pow(
    int) |
64 setScale(int scale , int round_mode) | stripTrailingZeros ()
65 // StringBuilder
66 StringBuilder sb = new StringBuilder ();
67 sb.append(elem) | out.println(sb)

```

# Chapter 8

## 数学公式

### 8.1 常用数学公式

#### 8.1.1 求和公式

1.  $\sum_{k=1}^n (2k-1)^2 = \frac{n(4n^2-1)}{3}$
2.  $\sum_{k=1}^n k^3 = [\frac{n(n+1)}{2}]^2$
3.  $\sum_{k=1}^n (2k-1)^3 = n^2(2n^2-1)$
4.  $\sum_{k=1}^n k^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$
5.  $\sum_{k=1}^n k^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12}$
6.  $\sum_{k=1}^n k(k+1) = \frac{n(n+1)(n+2)}{3}$
7.  $\sum_{k=1}^n k(k+1)(k+2) = \frac{n(n+1)(n+2)(n+3)}{4}$
8.  $\sum_{k=1}^n k(k+1)(k+2)(k+3) = \frac{n(n+1)(n+2)(n+3)(n+4)}{5}$

#### 8.1.2 斐波那契数列

1.  $fib_0 = 0, fib_1 = 1, fib_n = fib_{n-1} + fib_{n-2}$
2.  $fib_{n+2} \cdot fib_n - fib_{n+1}^2 = (-1)^{n+1}$
3.  $fib_{-n} = (-1)^{n-1} fib_n$
4.  $fib_{n+k} = fib_k \cdot fib_{n+1} + fib_{k-1} \cdot fib_n$
5.  $gcd(fib_m, fib_n) = fib_{gcd(m,n)}$
6.  $fib_m | fib_n^2 \Leftrightarrow n fib_n | m$

#### 8.1.3 错排公式

1.  $D_n = (n-1)(D_{n-2} + D_{n-1})$
2.  $D_n = n! \cdot (1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + \frac{(-1)^n}{n!})$

#### 8.1.4 莫比乌斯函数

$$\mu(n) = \begin{cases} 1 & \text{若 } n = 1 \\ (-1)^k & \text{若 } n \text{ 无平方数因子, 且 } n = p_1 p_2 \dots p_k \\ 0 & \text{若 } n \text{ 有大于1的平方数因数} \end{cases}$$

$$\sum_{d|n} \mu(d) = \begin{cases} 1 & \text{若 } n = 1 \\ 0 & \text{其他情况} \end{cases}$$

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d) g(\frac{n}{d})$$

$$g(x) = \sum_{n=1}^{[x]} f\left(\frac{x}{n}\right) \Leftrightarrow f(x) = \sum_{n=1}^{[x]} \mu(n)g\left(\frac{x}{n}\right)$$

### 8.1.5 Burnside 引理

设  $G$  是一个有限群, 作用在集合  $X$  上. 对每个  $g$  属于  $G$ , 令  $X^g$  表示  $X$  中在  $g$  作用下的不动元素, 轨道数 (记作  $|X/G|$ ) 由如下公式给出:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

### 8.1.6 五边形数定理

设  $p(n)$  是  $n$  的拆分数, 有

$$p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k-1} p\left(n - \frac{k(3k-1)}{2}\right)$$

### 8.1.7 树的计数

1. 有根树计数:  $n+1$  个结点的有根树的个数为

$$a_{n+1} = \frac{\sum_{j=1}^n j \cdot a_j \cdot S_{n,j}}{n}$$

其中,

$$S_{n,j} = \sum_{i=1}^{n/j} a_{n+1-ij} = S_{n-j,j} + a_{n+1-j}$$

2. 无根树计数: 当  $n$  为奇数时,  $n$  个结点的无根树的个数为

$$a_n - \sum_{i=1}^{n/2} a_i a_{n-i}$$

当  $n$  为偶数时,  $n$  个结点的无根树的个数为

$$a_n - \sum_{i=1}^{n/2} a_i a_{n-i} + \frac{1}{2} a_{\frac{n}{2}} (a_{\frac{n}{2}} + 1)$$

3.  $n$  个结点的完全图的生成树个数为

$$n^{n-2}$$

4. 矩阵 - 树定理: 图  $G$  由  $n$  个结点构成, 设  $\mathbf{A}[G]$  为图  $G$  的邻接矩阵,  $\mathbf{D}[G]$  为图  $G$  的度数矩阵, 则图  $G$  的不同生成树的个数为  $\mathbf{C}[G] = \mathbf{D}[G] - \mathbf{A}[G]$  的任意一个  $n-1$  阶主子式的行列式值。

### 8.1.8 欧拉公式

平面图的顶点个数、边数和面的个数有如下关系:

$$V - E + F = C + 1$$

其中,  $V$  是顶点的数目,  $E$  是边的数目,  $F$  是面的数目,  $C$  是组成图形的连通部分的数目. 当图是单连通图的时候, 公式简化为:

$$V - E + F = 2$$

### 8.1.9 皮克定理

给定顶点坐标均是整点 (或正方形格点) 的简单多边形, 其面积  $A$  和内部格点数目  $i$ 、边上格点数目  $b$  的关系:

$$A = i + \frac{b}{2} - 1$$

### 8.1.10 牛顿恒等式

设

$$\prod_{i=1}^n (x - x_i) = a_n + a_{n-1}x + \cdots + a_1x^{n-1} + a_0x^n$$

$$p_k = \sum_{i=1}^n x_i^k$$

则

$$a_0p_k + a_1p_{k-1} + \cdots + a_{k-1}p_1 + ka_k = 0$$

特别地, 对于

$$|\mathbf{A} - \lambda \mathbf{E}| = (-1)^n (a_n + a_{n-1}\lambda + \cdots + a_1\lambda^{n-1} + a_0\lambda^n)$$

有

$$p_k = \text{Tr}(\mathbf{A}^k)$$

## 8.2 平面几何公式

### 8.2.1 三角形

1. 半周长

$$p = \frac{a+b+c}{2}$$

2. 面积

$$S = \frac{a \cdot H_a}{2} = \frac{ab \cdot \sin C}{2} = \sqrt{p(p-a)(p-b)(p-c)}$$

3. 中线

$$M_a = \frac{\sqrt{2(b^2 + c^2) - a^2}}{2} = \frac{\sqrt{b^2 + c^2 + 2bc \cdot \cos A}}{2}$$

4. 角平分线

$$T_a = \frac{\sqrt{bc \cdot [(b+c)^2 - a^2]}}{b+c} = \frac{2bc \cos \frac{A}{2}}{b+c}$$

5. 高线

$$H_a = b \sin C = c \sin B = \sqrt{b^2 - \left(\frac{a^2 + b^2 - c^2}{2a}\right)^2}$$

6. 内切圆半径

$$\begin{aligned} r &= \frac{S}{p} = \frac{\arcsin \frac{B}{2} \cdot \sin \frac{C}{2}}{\sin \frac{B+C}{2}} = 4R \cdot \sin \frac{A}{2} \sin \frac{B}{2} \sin \frac{C}{2} \\ &= \sqrt{\frac{(p-a)(p-b)(p-c)}{p}} = p \cdot \tan \frac{A}{2} \tan \frac{B}{2} \tan \frac{C}{2} \end{aligned}$$

7. 外接圆半径

$$R = \frac{abc}{4S} = \frac{a}{2\sin A} = \frac{b}{2\sin B} = \frac{c}{2\sin C}$$

### 8.2.2 四边形

$D_1, D_2$  为对角线,  $M$  为对角线中点连线,  $A$  为对角线夹角,  $p$  为半周长

$$1. a^2 + b^2 + c^2 + d^2 = D_1^2 + D_2^2 + 4M^2$$

$$2. S = \frac{1}{2} D_1 D_2 \sin A$$

3. 对于圆内接四边形

$$ac + bd = D_1 D_2$$

4. 对于圆内接四边形

$$S = \sqrt{(p-a)(p-b)(p-c)(p-d)}$$

### 8.2.3 正 $n$ 边形

$R$  为外接圆半径,  $r$  为内切圆半径

1. 中心角

$$A = \frac{2\pi}{n}$$

2. 内角

$$C = \frac{n-2}{n}\pi$$

3. 边长

$$a = 2\sqrt{R^2 - r^2} = 2R \cdot \sin \frac{A}{2} = 2r \cdot \tan \frac{A}{2}$$

4. 面积

$$S = \frac{nar}{2} = nr^2 \cdot \tan \frac{A}{2} = \frac{nR^2}{2} \cdot \sin A = \frac{na^2}{4 \cdot \tan \frac{A}{2}}$$

### 8.2.4 圆

1. 弧长

$$l = rA$$

2. 弦长

$$a = 2\sqrt{2hr - h^2} = 2r \cdot \sin \frac{A}{2}$$

3. 弓形高

$$h = r - \sqrt{r^2 - \frac{a^2}{4}} = r(1 - \cos \frac{A}{2}) = \frac{1}{2} \cdot \arctan \frac{A}{4}$$

4. 扇形面积

$$S_1 = \frac{rl}{2} = \frac{r^2 A}{2}$$

5. 弓形面积

$$S_2 = \frac{rl - a(r - h)}{2} = \frac{r^2}{2}(A - \sin A)$$

### 8.2.5 棱柱

1. 体积

$$V = Ah$$

$A$  为底面积,  $h$  为高

2. 侧面积

$$S = lp$$

$l$  为棱长,  $p$  为直截面周长

3. 全面积

$$T = S + 2A$$

### 8.2.6 棱锥

1. 体积

$$V = Ah$$

$A$  为底面积,  $h$  为高

2. 正棱锥侧面积

$$S = lp$$

$l$  为棱长,  $p$  为直截面周长

3. 正棱锥全面积

$$T = S + 2A$$

### 8.2.7 棱台

1. 体积

$$V = (A_1 + A_2 + \sqrt{A_1 A_2}) \cdot \frac{h}{3}$$

$A_1, A_2$  为上下底面积,  $h$  为高

2. 正棱台侧面积

$$S = \frac{p_1 + p_2}{2} l$$

$p_1, p_2$  为上下底面周长,  $l$  为斜高

3. 正棱台全面积

$$T = S + A_1 + A_2$$

### 8.2.8 圆柱

1. 侧面积

$$S = 2\pi r h$$

2. 全面积

$$T = 2\pi r(h + r)$$

3. 体积

$$V = \pi r^2 h$$

### 8.2.9 圆锥

1. 母线

$$l = \sqrt{h^2 + r^2}$$

2. 侧面积

$$S = \pi r l$$

3. 全面积

$$T = \pi r(l + r)$$

4. 体积

$$V = \frac{\pi}{3} r^2 h$$

### 8.2.10 圆台

1. 母线

$$l = \sqrt{h^2 + (r_1 - r_2)^2}$$

2. 侧面积

$$S = \pi(r_1 + r_2)l$$

3. 全面积

$$T = \pi r_1(l + r_1) + \pi r_2(l + r_2)$$

4. 体积

$$V = \frac{\pi}{3}(r_1^2 + r_2^2 + r_1 r_2)h$$

### 8.2.11 球

1. 全面积

$$T = 4\pi r^2$$

2. 体积

$$V = \frac{4}{3}\pi r^3$$

### 8.2.12 球台

1. 侧面积

$$S = 2\pi rh$$

2. 全面积

$$T = \pi(2rh + r_1^2 + r_2^2)$$

3. 体积

$$V = \frac{\pi h[3(r_1^2 + r_2^2) + h^2]}{6}$$

### 8.2.13 球扇形

1. 全面积

$$T = \pi r(2h + r_0)$$

$h$  为球冠高,  $r_0$  为球冠底面半径

2. 体积

$$V = \frac{2}{3}\pi r^2 h$$

## 8.3 立体几何公式

### 8.3.1 球面三角公式

设  $a, b, c$  是边长,  $A, B, C$  是所对的二面角, 有余弦定理

$$\cos a = \cos b \cdot \cos c + \sin b \cdot \sin c \cdot \cos A$$

正弦定理

$$\frac{\sin A}{\sin a} = \frac{\sin B}{\sin b} = \frac{\sin C}{\sin c}$$

三角形面积是  $A + B + C - \pi$

### 8.3.2 四面体体积公式

$U, V, W, u, v, w$  是四面体的 6 条棱,  $U, V, W$  构成三角形,  $(U, u), (V, v), (W, w)$  互为对棱, 则

$$V = \frac{\sqrt{(s-2a)(s-2b)(s-2c)(s-2d)}}{192uvw}$$

其中

$$\begin{cases} a &= \sqrt{xYZ}, \\ b &= \sqrt{yZX}, \\ c &= \sqrt{zXY}, \\ d &= \sqrt{xyz}, \\ s &= a + b + c + d, \\ X &= (w - U + v)(U + v + w), \\ x &= (U - v + w)(v - w + U), \\ Y &= (u - V + w)(V + w + u), \\ y &= (V - w + u)(w - u + V), \\ Z &= (v - W + u)(W + u + v), \\ z &= (W - u + v)(u - v + W) \end{cases}$$