

计算机科学中的数学基础—exercise2

陈昱衡 521021910939

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Warmup:Q6

Some of the regions defined by n lines in the plane are infinite, while others are bounded. What's the maximum possible number of bounded regions?

Answer

This question is a little different from the example from the textbook, but we can also try to solve the problem by visualizing the recurrent process.

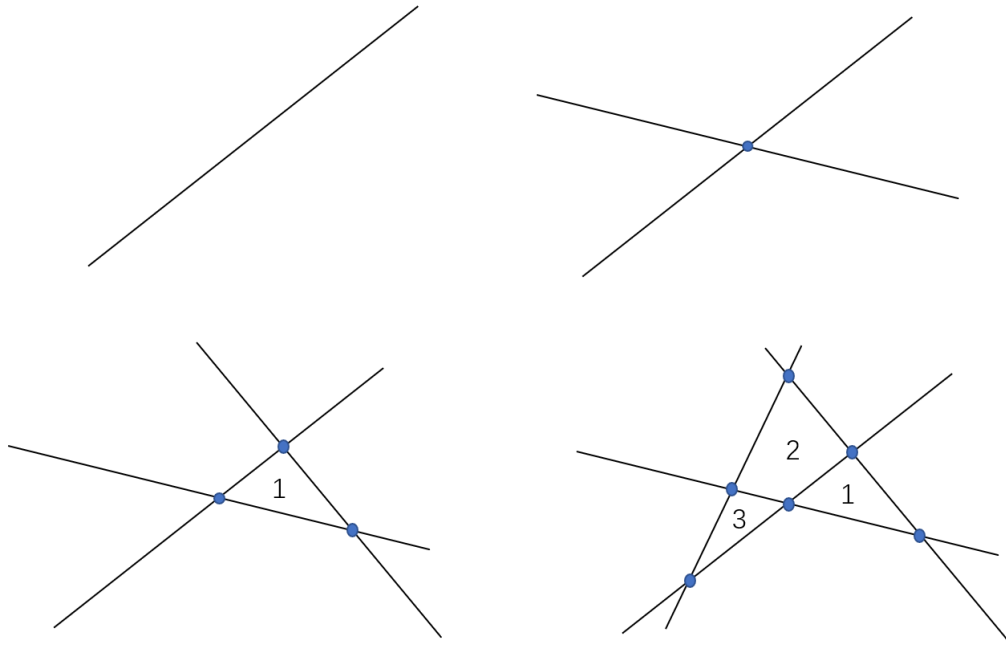


图 1: Q6-demo

So, we can kind of easily draw the conclusion that: the number of added bounded region equals the number of the lines which the new line intersect with minus one (in fact initially we can just 1). So, we get the following recursive formula:

$$T_n = 0 \quad (n = 1) \quad (1)$$

$$T_n = T_{n-1} + n - 2 \quad (n \geq 2) \quad (2)$$

We can get the answer by accumulating all the items together:

$$T_n = 0 \quad (n = 0) \quad (3)$$

$$T_n = \frac{(n-1)(n-2)}{2} \quad (n \geq 2) \quad (4)$$

Warmup:Q7

Let $H(n) = J(n+1) - J(n)$. Equation (1.8) tells us that $H(2n) = 2$, and $H(2n+1) = J(2n+2) - J(2n+1) = (2J(n+1) - 1) - (2J(n) + 1) = 2H(n) - 2$, for all $n \geq 1$. Therefore it seems possible to prove that $H(n) = 2$ for all n , by induction on n . What's wrong here?

Answer

We can easily notice that $H(0)=1 \neq 2$, so the foundation of the inductive method collapses.

Homework:Q8

Solve the recurrence

$$\begin{aligned} Q_0 &= \alpha; & Q_1 &= \beta; \\ Q_n &= (1 + Q_{n-1})/Q_{n-2}, & \text{for } n > 1. \end{aligned}$$

Assume that $Q_n \neq 0$ for all $n \geq 0$. *Hint:* $Q_4 = (1 + \alpha)/\beta$.

Answer

Well, nothing had flashed across my mind when I was trying to solve the problem initially. What could I do was keeping calculating. Then, I got the answer. The series is periodic!

$$Q_0 = \alpha \tag{5}$$

$$Q_1 = \beta \tag{6}$$

$$Q_2 = \frac{1 + \beta}{\alpha} \tag{7}$$

$$Q_3 = \frac{1 + \alpha + \beta}{\alpha\beta} \tag{8}$$

$$Q_4 = \frac{1 + \alpha}{\beta} \tag{9}$$

$$Q_5 = \alpha \tag{10}$$

Homework:Q9

Sometimes it's possible to use induction backwards, proving things from n to $n - 1$ instead of vice versa! For example, consider the statement

$$P(n) : \quad x_1 \dots x_n \leq \left(\frac{x_1 + \dots + x_n}{n} \right)^n, \quad \text{if } x_1, \dots, x_n \geq 0.$$

This is true when $n = 2$, since $(x_1 + x_2)^2 - 4x_1x_2 = (x_1 - x_2)^2 \geq 0$.

- a** By setting $x_n = (x_1 + \dots + x_{n-1})/(n - 1)$, prove that $P(n)$ implies $P(n - 1)$ whenever $n > 1$.
- b** Show that $P(n)$ and $P(2)$ imply $P(2n)$.
- c** Explain why this implies the truth of $P(n)$ for all n .

Answer

(a) If we want to prove $P(n-1)$ using $P(n)$, we can manage this by proving the following equation :

$$x_1 \cdots x_{n-1} \left(\frac{x_1 + \cdots + x_{n-1}}{n-1} \right) \leq \left(\frac{x_1 + \cdots + x_{n-1}}{n-1} \right)^n \quad (11)$$

It's quite easy, as we can use the inequality:

$$\left(\frac{x_1 + \cdots + x_{n-1}}{n-1} \right)^{n-1} \geq \left(\frac{x_1 + \cdots + x_n}{n} \right)^n \quad (12)$$

(b) From $P[n]$, we can get:

$$x_1 \cdots x_n x_{n+1} \cdots x_{2n} \leq \left(\left(\frac{x_1 + \cdots + x_n}{n} \right) \left(\frac{x_{n+1} + \cdots + x_{2n}}{n} \right) \right)^n \quad (13)$$

then using $P(2)$, we can further get :

$$\left(\left(\frac{x_1 + \cdots + x_n}{n} \right) \left(\frac{x_{n+1} + \cdots + x_{2n}}{n} \right) \right)^n \leq \left(\frac{x_1 + \cdots + x_{2n}}{2n} \right)^{2n} \quad (14)$$

then the question got proved.

(c) Using the conclusion we got from (b), since we got $P(2)$ now, then we can get $P(4)$, using (a), we can get $P(3)$. So, we can keep getting $P(2^k)$ using (b), and get $P(2^{k-1} + 1)$ to $P(2^k - 1)$.

Homework:Q10

Let Q_n be the minimum number of moves needed to transfer a tower of n disks from A to B if all moves must be *clockwise* — that is, from A to B, or from B to the other peg, or from the other peg to A. Also let R_n be the minimum number of moves needed to go from B back to A under this restriction. Prove that

$$Q_n = \begin{cases} 0, & \text{if } n = 0; \\ 2R_{n-1} + 1, & \text{if } n > 0; \end{cases} \quad R_n = \begin{cases} 0, & \text{if } n = 0; \\ Q_n + Q_{n-1} + 1, & \text{if } n > 0. \end{cases}$$

(You need not solve these recurrences; we'll see how to do that in Chapter 7.)

Answer

To solve the problem, we are supposed to be aware that the process of moving the tower from A to B equals to from B to C and C to A as well. Also, the process of moving the tower from A to C, from B to A, from C to B are the same, which is R_n under the context.

So, the first set of equations got proved, because we should move the upper $n-1$ plates from A to C which uses R_n , then move the biggest one from A to B which uses one step, finally move the n -tower from C to B which is also at the cost of R_n .

As to the second set of formula, we should move the upper $n-1$ plates from A to C, use R_{n-1} steps, then the biggest one from A to B, one step, then the $n-1$ tower from C to A, Q_{n-1} steps, then move the biggest one to C, at last move the $n-1$ tower from A to C within R_{n-1} steps. We got:

$$R_n = R_{n-1} + 1 + Q_{n-1} + R_{n-1} + 1 \quad (15)$$

With $Q_n = 2R_{n-1} + 1$ replacing Q_n , we got the right answer!