dont know	know a bit	ОК	good!	mastei	COMS 311 TOPICS				
1	2	3	4	5	BIG-OH	week4	week3	week2	week1
					Basics				х
					Definitions of big-oh, omega, theta c>0, n>=0 (7 things)				х
					big-omega				х
					big-theta				х
					Intuition/understanding (graph)				х
					tighter and weaker bounds			х	х
					how to prove O/Omega/Theta techniques				х
					for polynomials choose c >sum of coeff or coeff for omega				х
					for same type compare exponents				х
					take log				х
					Application to Algorithms				х
					ram model (vs actual)				Х
					instances and runtime graphs			х	Х
					WCET, BCET, ACET			х	х
					Big-oh of code segments			х	х
					problem complexity and algorithmic complexity				Х
					code examples of different Os				Х
					Big-oh in real world			х	
					real code times (matrix mult)			х	
					effect of cache/pipelining etc			х	
					choosing algo in real-world vs big-Oh			x	
					constants might matter more in real world than O			x	
					easier implementation might make the diff			x	
					casier imprementation imgre make the airi			_^_	
					Dominance Relationships				х
					logs beat constants				X
					poly beats all logs				X
					exp beats all poly				
					fact beats all exp				X
					n^n beats fact				X
					II II beats fact				_^
1	2	3	4	5	DATA STRUCTURES				
		,			Basic			Х	х
					arrays (sorted/unsorted)	1		X	X
					linked lists (singly/doubly; sorted/unsorted)		<del>                                     </del>	X	X
					comparison of arrays and linked lists		<del>                                     </del>	X	X
	$\vdash \vdash$				ostripulison of arrays and infined iists		-	<del>  ^</del>	
					 Basic Abstract Data Types		-	х	х
	$\vdash \vdash$				Stack, Queue (implementations using array/linkedlists)		-	X	X
					comparison of operations of diff impl of stack/queue			X	X
					reasons for differences	+	<del>                                     </del>	X	X
					Dictionaries		X	X	_ <u> </u>
					Hash Tables		X		
	$\vdash \vdash \vdash$				וומטוו ומטוכט		<del>  ^</del>	Х	
					 Other ADTs and their Java Implementations		x	х	
	ш		<u> </u>	<u> </u>	other Ap 13 and then Java Implementations			_ ^	

				1	D: C 1.T	I	1	I	ı
$\vdash$					Binary Search Trees		Х	Х	
					Priority Search Queues		Х	Х	
					Heap impl		Х	Х	
					Fast Heap Impl and Analysis		Х		
					Storing points, graphs, sets etc		Х	X	
$\longrightarrow$					graphs	X	Х	Х	
					sets	X	Х	Х	
					his at af an anti-man an data at materials				
					big-oh of operations on data structures		Х	Х	
					algorithms on data structures (BST, HEAP etc)		Х	Х	
1	2	3	4		CRADUS (vervoirebted)				
1	2	3	4	5	GRAPHS (unweighted)	.,,			
$\longrightarrow$					BFS Separated components	X	X		
					connected components	X	X		
					two-color problem	Х	X		
			-		DFS on undirected graphs (tree and back edges)	X	Х		
					articulation vertices (parent, root, bridge cutnodes)  DFS on directed graphs(tree,back,cross,forward)	X			
					topological sorting/DAG	X			
-					· · ·	X			
					strongly connected components following algorithm by marking graph	X			
+					big-oh of graph algorithms	X	, , , , , , , , , , , , , , , , , , ,		
-					big-on or graph algorithms	X	Х		
1	2	3	4	5	P-NP				
		<u> </u>	4	J	Intro Concepts				Х
					the diagram and four classes of problems	X			X
					informal (solvable, probab intract, provably intract, prov unsolvable)	X			X
					examples of problems in four classes	X			X
					Halting Problem	X			X
					Hamiltonian Cycle Enumeration problem	X			X
					Hamiltonian Cycle Search problem	X			X
					Independent Set problem	X			X
					Search/Sort problems	X			X
					Jos. S. y Sort prositions				
					Classes of problems				
					Undecidable (prove Halting problem is undecidable)	X	<del>                                     </del>		х
I			1	1			-		X
					IP	X			. ^
					P NP	X			х
					NP	X			Х
					NP prove P is a subset of NP				
					NP prove P is a subset of NP NP-Complete (probably intractable class)				X
					NP prove P is a subset of NP NP-Complete (probably intractable class) why if a NP-C problem is in P, then P=NP				х
					NP prove P is a subset of NP NP-Complete (probably intractable class) why if a NP-C problem is in P, then P=NP provably intractable classes				
					NP prove P is a subset of NP NP-Complete (probably intractable class) why if a NP-C problem is in P, then P=NP				х
					NP prove P is a subset of NP NP-Complete (probably intractable class) why if a NP-C problem is in P, then P=NP provably intractable classes NP-Hard problems				х
					NP prove P is a subset of NP NP-Complete (probably intractable class) why if a NP-C problem is in P, then P=NP provably intractable classes NP-Hard problems  Reduction				х
					NP prove P is a subset of NP NP-Complete (probably intractable class) why if a NP-C problem is in P, then P=NP provably intractable classes NP-Hard problems  Reduction optimization, search, decision and reductions	X			х
					NP prove P is a subset of NP NP-Complete (probably intractable class) why if a NP-C problem is in P, then P=NP provably intractable classes NP-Hard problems  Reduction	x			х

				December 1997		1		1
_				how to prove X <=p Y (the three steps)	X			
_				· · · · · · · · · · · · · · · · · · ·	X			
					X			
					X			
				bipartite matching <=p IS	X			
_								
			ا	·				
				· · · · · · · · · · · · · · · · · · ·				
_				·				
_				·				
				· ·				
				prove VC is NP-C				
_								
2	3	4						
_								Х
_				<u>'</u>				X
_				·				Х
				why is proving so important?				Х
4								-
_			- (					PreRec
_								PreRec
_								PreRec
_				form of deduction proofs				PreRec
_								
				-				PreRec
_				·		X		PreRec
+								PreRec
+				, ,				PreRec
+							<del> </del>	PreRec
+				contrapositive	X		X	PreRec
+				EVANADI EC INI CI ACC				PreRec
+						, , , , , , , , , , , , , , , , , , ,		PreRec
+						×		
+	-+		$\vdash$	, ,				PreRec
+				1 ,				PreRec PreRec
+	-+		$\vdash$	·	-			PreRec
+	-+		$\vdash$	Contrapositive	X	-		rieked
+			<del>    </del> ,	 Proofs in class		1		
-								PreRec
+						, x	<del>                                     </del>	Frened
+					<b>^</b>			+
+	-			Select Jobs satisfies greedy choice + opt substructuring				
+				 Proving Code correct				1
+	-+		<del>                                     </del>		+		-	
+				·				+
+	$\dashv$			proof of recursive codes		<del>                                     </del>		+
2	3	4	5	ALGORITHMIC TECHNIQUES				
			,					
				2 3 4 5	interval sched <=p IS  vector cover <=p IS  IS <=p vector cover  bipartite matching <=p IS  NP-Complete  Definition and Intuition  The first NP-C problem (circuitSAT) & Cook-Levin's theorem how to prove that a problem is NP-C prove 3-SAT is NP-C prove IS is NP-C prove IS is NP-C prove VC is NP-C  Why is proving important?  Iob selection problem solutions to JS problem why is proving so important?  CALCULUS propositional logic rules predicate logic rules form of deduction proofs  DIFFERENT TECHNIQUES  contradiction (and correct form) induction (form) direct proof (see form of deduction proofs) trivial/vacuous contrapositive  EXAMPLES IN CLASS contradiction (and correct form) induction (form) direct proof (see form of deduction proofs) trivial/vacuous contrapositive  Proofs in class Halting problem is undecidable VC <=p IS and IS <=p VC select jobs satisfies greedy choice + opt substructuring Proving Code correct Loop invariants proof of recursive codes	interval sched <=p IS	interval sched <=p IS	Interval sched <=p IS

search space for different problems			х	х
recursion tree for brute force approach			х	Х
back track algo from text book			х	Х
iterative way to generate all subsets				
recursive way to generate all subsets			х	х
recursive way to generate all perms			х	Х
recursive way to gen size k subsets			х	х
DIVIDE AND CONQUER				
recurrence formula	X	х		
divide and conquer approach	х	х		
mergesort + analysis	X	х		
quicksort + analysis	х	х		
max	х	х		
max sum of sequence	х			
counting inversions	х			
finding sink in graph				
Recurrence Formulae	X	х		-
general form of recurrence formula & masters theorem	X	х		
how to derive recurrence tree, term for each level, sum	х	х		
2T(n/2) + c	X	х		
T(n/2) + cn	X	х		
2T(n/2) + cn	X	х		
3T(n/2) + cn	X	х		
2T(n/2) + cn^2	х	х		
5T(n/2) + cn^2	X	Х		
GREEDY TECHNIQUE				
greedy approach				
interval scheduling by greedy approach				
greedy choice property				
optimal substructuring property				
proving gc property				
proving os property				