

Keyword operation($\mathcal{H} : S \mapsto (A, T)$)

Single keyword operation

91. $\mathcal{H}[\text{from}(\langle tname \rangle)] \triangleq (A, A(\langle tname \rangle))$
92. $\mathcal{H}[\text{from}(\langle tname_1 \rangle), \langle tname_2 \rangle] \triangleq (A, \mathcal{C}[[A(\langle tname_1 \rangle), A(\langle tname_2 \rangle)], \text{cross join}])$
93. $\mathcal{H}[\text{select } *] \triangleq (\{T\}, \pi_{T, \bar{\beta}}(T))$
94. $\mathcal{H}[\text{select } \langle cname \rangle [, \langle cname \rangle \dots]] \triangleq (\{T\}, \pi_{\langle cname \rangle [, \langle cname \rangle \dots]}(T))$
95. $\mathcal{H}[\text{on } \langle bvep \rangle] \triangleq (\{T\}, \sigma_{\mathcal{B}(\langle bvep \rangle)}(T))$
96. $\mathcal{H}[\text{where } \langle bvep \rangle] \triangleq (\{T\}, \sigma_{\mathcal{B}(\langle bvep \rangle)}(T))$
97. $\mathcal{H}[\text{group by } \langle cname \rangle] \triangleq (\{T\}, (\widehat{\alpha_1}, \dots, \widehat{\alpha_k}) : \forall \widehat{\alpha_p} \in (\widehat{\alpha_1}, \dots, \widehat{\alpha_k}), \forall v_{ij} \in \pi_{\langle cname \rangle}(\widehat{\alpha_p}), (v_{ij} = v_{i1}))$
98. $\mathcal{H}[\text{having } \langle bvep \rangle] \triangleq (\{T\}, \sigma_{\mathcal{B}(\langle bvep \rangle)}(T))$
99. $\mathcal{H}[\text{order by } \langle cname \rangle \text{ asc}] \triangleq (\{T\}, T_1) \text{ where } (\forall \alpha \in T_1, \xi_\alpha(T_1) = \xi_\alpha(T)) \wedge (\forall \alpha \in T, \xi_\alpha(T) = \xi_\alpha(T_1)) \wedge (\forall v_i, v_j \in \sigma_{T_1, \langle cname \rangle}(T_1), i > j \text{ iff } v_i > v_j)$
100. $\mathcal{H}[\text{order by } \langle cname \rangle \text{ desc}] \triangleq (\{T\}, T_1) \text{ where } (\forall \alpha \in T_1, \xi_\alpha(T_1) = \xi_\alpha(T)) \wedge (\forall \alpha \in T, \xi_\alpha(T) = \xi_\alpha(T_1)) \wedge (\forall v_i, v_j \in \sigma_{T_1, \langle cname \rangle}(T_1), i < j \text{ iff } v_i < v_j)$
101. $\mathcal{H}[\langle subquery \rangle] \triangleq \mathcal{H}[\langle query \text{ expression} \rangle]$

Composite keyword operation

102. $\diamond : (1) \frac{\mathcal{H}[\text{expression}_1] \triangleq (A, T), \mathcal{H}[\text{expression}_2] \triangleq (\{T\}, T')}{\mathcal{H}[\text{expression}_1] \diamond \mathcal{H}[\text{expression}_2] \triangleq (A, T')}$
 $(2) \frac{\mathcal{H}[\text{expression}] \triangleq (A, T), \mathcal{C}[\{\{T\}, OP\}] \triangleq T'}{\mathcal{H}[\text{expression}] \diamond \mathcal{C}[\{\{T\}, OP\}] \triangleq (A, T')}$
 $(3) \frac{\mathcal{C}[\{L, OP\}] \triangleq T, \mathcal{H}[\text{expression}] \triangleq (\{T\}, T')}{\mathcal{C}[\{L, OP\}] \diamond \mathcal{H}[\text{expression}] \triangleq (L, T')}$
 $(4) \frac{\mathcal{H}[\text{expression}_1] \triangleq (A, T_1), \mathcal{H}[\text{expression}_2] \triangleq (A, T_2), \mathcal{C}[\{\{T_1, T_2\}, OP\}] \triangleq T_3}{(\mathcal{H}[\text{expression}_1], \mathcal{H}[\text{expression}_2]) \diamond \mathcal{C}[\{\{T_1, T_2\}, OP\}] \triangleq (A, T_3)}$
103. $\mathcal{H}[\langle queryexp \rangle] =$
 $\mathcal{H}[\text{select } [(\text{sop}) | (\text{af})] \langle cname_1 \rangle [, \langle cname_2 \rangle \dots] \text{from}(\langle tname_1 \rangle [, \langle tname_2 \rangle \dots]) \text{from } \langle tname_1 \rangle \text{ natural/cross join } \langle tname_2 \rangle$
 $|\text{from } \langle tname_1 \rangle \text{ left/right/full/inner join } \langle tname_2 \rangle \text{ on}(\langle bvep \rangle) [\text{where}(\langle bvep \rangle) [\text{group by } \langle cname \rangle] [\text{having } \langle bvep \rangle]$
 $[\text{order by } \langle cname \rangle [\text{asc|desc}]]]$
 \triangleq
 $\mathcal{H}[\text{from}(\langle tname_1 \rangle)]$
 $|(\mathcal{H}[\text{from } \langle tname_1 \rangle], \mathcal{H}[\text{from } \langle tname_2 \rangle]) \diamond \mathcal{C}[L, \text{natural/cross join}]$
 $|(\mathcal{H}[\text{from } \langle tname_1 \rangle], \mathcal{H}[\text{from } \langle tname_2 \rangle]) \diamond \mathcal{C}[L, \text{left/right/inner/full join}] \diamond \mathcal{H}[\text{on } \langle bvep \rangle]$
 $[\diamond \mathcal{H}[\text{where } \langle bvep \rangle]]$
 $[\diamond \mathcal{H}[\text{group by } \langle cname \rangle]]$
 $[\diamond \mathcal{H}[\text{having } \langle bvep \rangle]]$
 $\diamond \mathcal{H}[\text{select } \langle cname_1 \rangle [, \langle cname_2 \rangle \dots]] \diamond \mathcal{C}[\{T_1, \langle sop/af \rangle\}]$
 $[\diamond \mathcal{H}[\text{order by } \langle cname \rangle [\text{asc|desc}]]]$
104. $\mathcal{H}[\langle queryexp_1 \rangle \langle cop \rangle \langle queryexp_2 \rangle] \triangleq (\mathcal{H}[\langle queryexp_1 \rangle], \mathcal{H}[\langle queryexp_2 \rangle]) \diamond \mathcal{C}[\{[T_1, T_2], \langle cop \rangle\}]$

Fig. 7: The full list of semantic definitions for SQL keywords from, select, on, group by, having, order by, subquery operations, and composite semantics on SQL queries