CWI

Erasable PUFs ormal Treatment and Generic Desig

Chenglu Jin, Wayne Burleson, Marten van Dijk, and Ulrich Rührmair



UMassAmherst CONNIMU







Hardware security primitive taking challenges and generating responses



- Hardware security primitive taking challenges and generating responses
- Unique fingerprint on individual IC even if designed and fabricated in the same way



- Hardware security primitive taking challenges and generating responses
- Unique fingerprint on individual IC even if designed and fabricated in the same way
- Leveraging manufacturing process variations



- Hardware security primitive taking challenges and generating responses
- Unique fingerprint on individual IC even if designed and fabricated in the same way
- Leveraging manufacturing process variations
- Ideally, PUFs are Physical Random Functions



- Hardware security primitive taking challenges and generating responses
- Unique fingerprint on individual IC even if designed and fabricated in the same way
- Leveraging manufacturing process variations
- Ideally, PUFs are Physical Random Functions



- Hardware security primitive taking challenges and generating responses
- Unique fingerprint on individual IC even if designed and fabricated in the same way
- Leveraging manufacturing process variations
- Ideally, PUFs are Physical Random Functions

Applications

Device/Chip Authentication



- Hardware security primitive taking challenges and generating responses
- Unique fingerprint on individual IC even if designed and fabricated in the same way
- Leveraging manufacturing process variations
- Ideally, PUFs are Physical Random Functions

- Device/Chip Authentication
- Key Management/Storage



- Hardware security primitive taking challenges and generating responses
- Unique fingerprint on individual IC even if designed and fabricated in the same way
- Leveraging manufacturing process variations
- Ideally, PUFs are Physical Random Functions

- Device/Chip Authentication
- Key Management/Storage
- Cryptographic Protocols (Key Exchange, Oblivious Transfer, Bit Commitment)



- Hardware security primitive taking challenges and generating responses
- Unique fingerprint on individual IC even if designed and fabricated in the same way
- Leveraging manufacturing process variations
- Ideally, PUFs are Physical Random Functions

- Device/Chip Authentication
- Key Management/Storage
- Cryptographic Protocols (Key Exchange, Oblivious Transfer, Bit Commitment)







$$C_0, C_1, \dots, C_k$$

$$\rightarrow R_0, R_1, \dots, R_k$$







(C, R)



Public, Authenticated Physical Channel





ВОВ

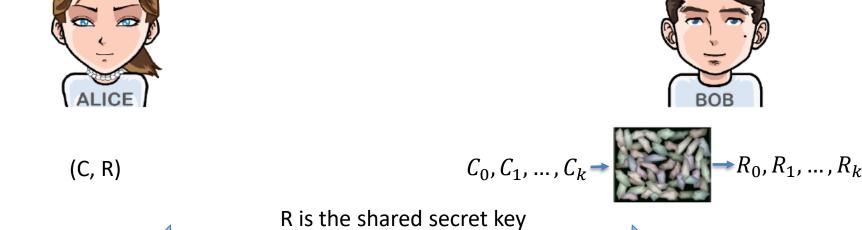
(C, R)



$$C_0, C_1, \dots, C_k$$

Public, Authenticated Communication Channel





The security of this protocol relies on the unpredictability of PUF responses given its challenges.



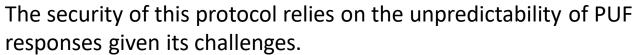




(C, R)

 C_0, C_1, \dots, C_k $\longrightarrow R_0, R_1, \dots, R_k$

R is the shared secret key

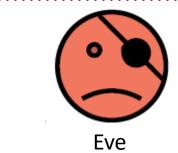


Not Complete!





(C, R)

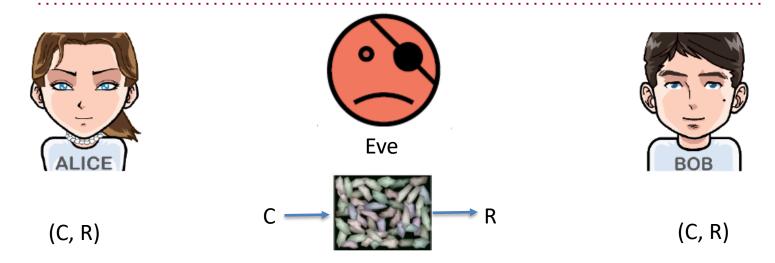






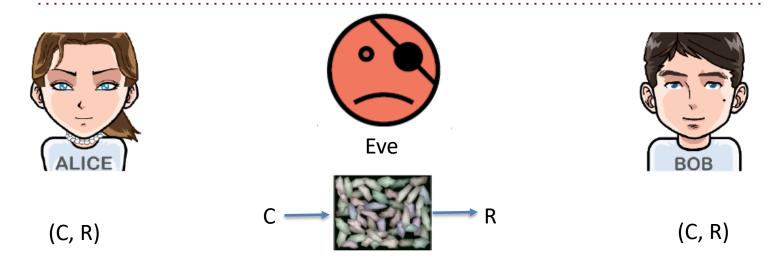
(C, R)





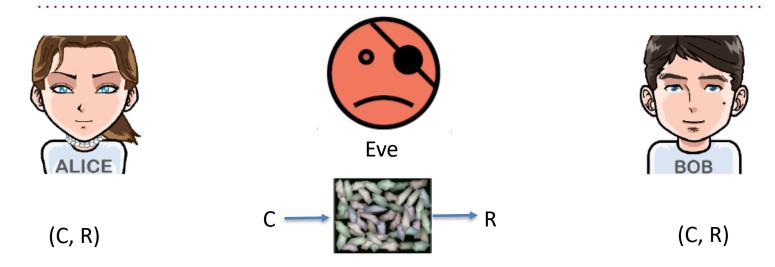
In the PUF Re-Use model, Eve can know the secret R as well.





- In the PUF Re-Use model, Eve can know the secret R as well.
- Highly realistic threat against PUF-based protocol, as destroying PUFs after one protocol execution is prohibitively uneconomic.

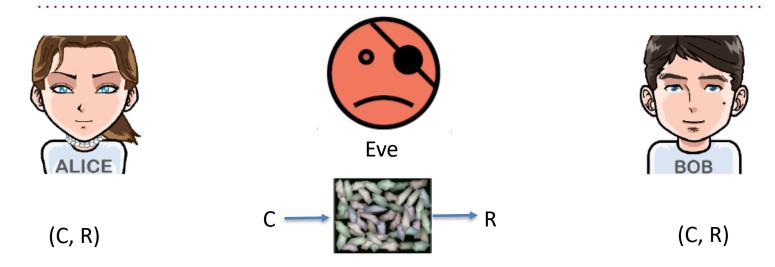




- In the PUF Re-Use model, Eve can know the secret R as well.
- Highly realistic threat against PUF-based protocol, as destroying PUFs after one protocol execution is prohibitively uneconomic.
- Actually, impossibility results of constructing PUF-based crypto protocols like KE/OT in PUF Re-Use model have been proved.

Marten van Dijk and Ulrich Rührmair. "Physical unclonable functions in cryptographic protocols: Security proofs and impossibility results." IACR ePrint (2012)





- In the PUF Re-Use model, Eve can know the secret R as well.
- Highly realistic threat against PUF-based protocol, as destroying PUFs after one protocol execution is prohibitively uneconomic.
- Actually, impossibility results of constructing PUF-based crypto protocols like KE/OT in PUF Re-Use model have been proved.
- The issue has to be solved on the hardware level.

Marten van Dijk and Ulrich Rührmair. "Physical unclonable functions in cryptographic protocols: Security proofs and impossibility results." IACR ePrint (2012)





Erase the CRP used in the protocol execution after the protocol



- Erase the CRP used in the protocol execution after the protocol
- Attackers will have no way to re-access the secret response value



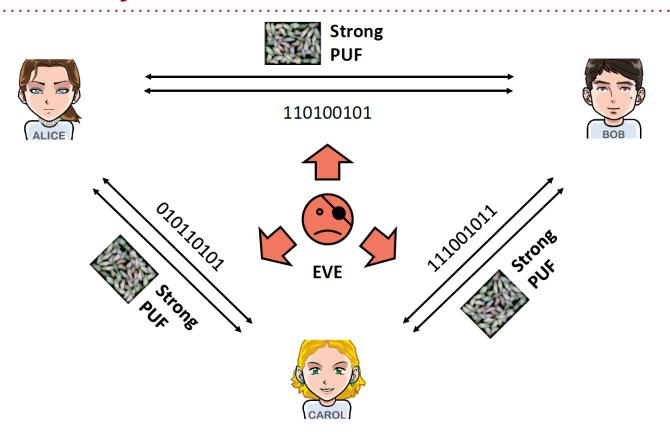
- Erase the CRP used in the protocol execution after the protocol
- Attackers will have no way to re-access the secret response value
- Can a reconfigurable PUF solve the problem?



- Erase the CRP used in the protocol execution after the protocol
- Attackers will have no way to re-access the secret response value
- Can a reconfigurable PUF solve the problem?
- A Reconfigurable PUF allows users to alter the responses of all challenges in one single operation (so-called "Reconfiguration").



Multi-Party Use Case



Using reconfigurable PUFs in crypto protocols cannot support multi-party use case.



Erasable PUFs

 Allows users to erase/alter the response of individual challenges chosen by the users



Erasable PUFs

- Allows users to erase/alter the response of individual challenges chosen by the users
- Erasable PUF-based crypto protocols can allow multiple parties to share one PUF and avoid repeated physical transfer of the PUF



Erasable PUFs

- Allows users to erase/alter the response of individual challenges chosen by the users
- Erasable PUF-based crypto protocols can allow multiple parties to share one PUF and avoid repeated physical transfer of the PUF
- Users can only erase the used CRPs after protocol execution, without affecting the other CRPs





 Add an interface around a PUF to enforce access control to the PUF



- Add an interface around a PUF to enforce access control to the PUF
- Create a list of erased challenges



- Add an interface around a PUF to enforce access control to the PUF
- Create a list of erased challenges
- If a queried challenge is in the list of erased challenges, then the interface should deny the access to the PUF



- Add an interface around a PUF to enforce access control to the PUF
- Create a list of erased challenges
- If a queried challenge is in the list of erased challenges, then the interface should deny the access to the PUF
- Otherwise, the interface will allow the PUF to be queried, and the response will be generated and outputted.



- Add an interface around a PUF to enforce access control to the PUF
- Create a list of erased challenges
- If a queried challenge is in the list of erased challenges, then the interface should deny the access to the PUF
- Otherwise, the interface will allow the PUF to be queried, and the response will be generated and outputted.
- Add new challenges into the list to erase them logically



Basic Idea to Realize Erasable PUFs

- Add an interface around a PUF to enforce access control to the PUF
- Create a list of erased challenges
- If a queried challenge is in the list of erased challenges, then the interface should deny the access to the PUF
- Otherwise, the interface will allow the PUF to be queried, and the response will be generated and outputted.
- Add new challenges into the list to erase them logically
- Drawback: The list should not be tampered with by adversaries, but the size of the list is growing when more and more challenges are erased. This implies that a large trusted memory is needed





Generic Erasable PUFs (Genie PUFs), because its just a PUF interface, and it can be integrated with any PUFs



- Generic Erasable PUFs (Genie PUFs), because its just a PUF interface, and it can be integrated with any PUFs
- Goal: Reduce the size of trusted memory in the trusted computing base (TCB)



- Generic Erasable PUFs (Genie PUFs), because its just a PUF interface, and it can be integrated with any PUFs
- Goal: Reduce the size of trusted memory in the trusted computing base (TCB)
- Key Idea: Merge Authenticated Search Tree and Red-Black Tree structure to securely outsource the list of erased challenges to untrusted memory



- Generic Erasable PUFs (Genie PUFs), because its just a PUF interface, and it can be integrated with any PUFs
- Goal: Reduce the size of trusted memory in the trusted computing base (TCB)
- Key Idea: Merge Authenticated Search Tree and Red-Black Tree structure to securely outsource the list of erased challenges to untrusted memory
- What can we achieve?



- Generic Erasable PUFs (Genie PUFs), because its just a PUF interface, and it can be integrated with any PUFs
- Goal: Reduce the size of trusted memory in the trusted computing base (TCB)
- Key Idea: Merge Authenticated Search Tree and Red-Black Tree structure to securely outsource the list of erased challenges to untrusted memory
- What can we achieve?
- Only require a constant-sized trusted memory in the TCB to store the root hash of the tree structure



- Generic Erasable PUFs (Genie PUFs), because its just a PUF interface, and it can be integrated with any PUFs
- Goal: Reduce the size of trusted memory in the trusted computing base (TCB)
- Key Idea: Merge Authenticated Search Tree and Red-Black Tree structure to securely outsource the list of erased challenges to untrusted memory
- What can we achieve?
- Only require a constant-sized trusted memory in the TCB to store the root hash of the tree structure
- Support arbitrarily large list of erased challenges



- Generic Erasable PUFs (Genie PUFs), because its just a PUF interface, and it can be integrated with any PUFs
- Goal: Reduce the size of trusted memory in the trusted computing base (TCB)
- Key Idea: Merge Authenticated Search Tree and Red-Black Tree structure to securely outsource the list of erased challenges to untrusted memory
- What can we achieve?
- Only require a constant-sized trusted memory in the TCB to store the root hash of the tree structure
- Support arbitrarily large list of erased challenges
- Using the combined tree structure, the untrusted memory can provide a O(log(N)) size proof to the TCB to prove a challenge is (not) in the list of size N



$$c_0$$
, $h_0 = H (c_0, h_1, h_2)$

$$c_1$$
, $h_1 = H (c_1, h_3, h_4)$

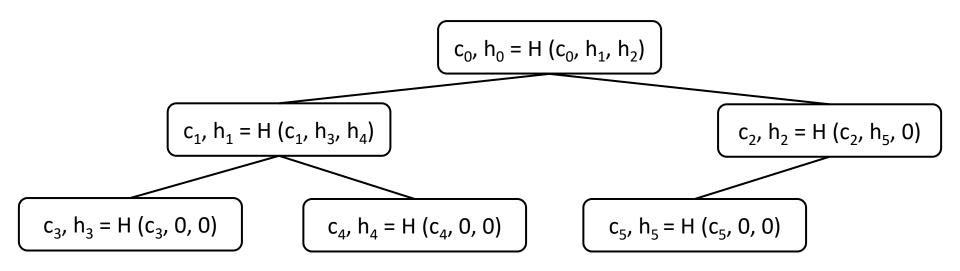
$$c_2$$
, h_2 = H (c_2 , h_5 , 0)

$$c_3$$
, $h_3 = H (c_3, 0, 0)$

$$c_4$$
, h_4 = H (c_4 , 0, 0)

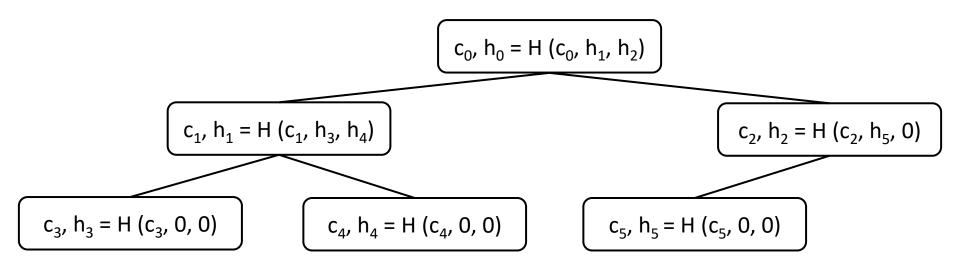
$$c_5$$
, $h_5 = H(c_5, 0, 0)$





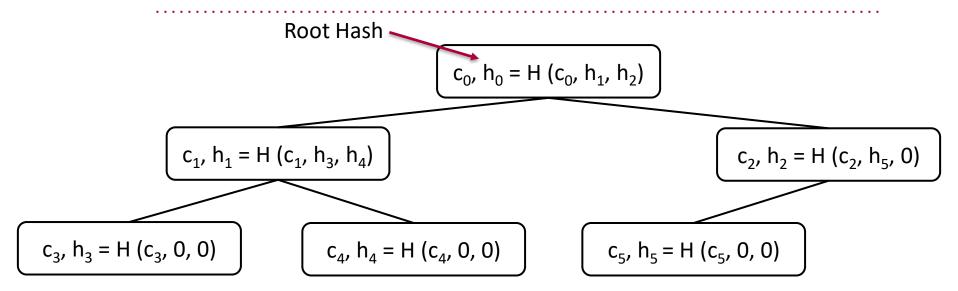
 In each node of the tree, we store one unique challenge, and the tree is sorted like a binary search tree according to the challenge value in each node





- In each node of the tree, we store one unique challenge, and the tree is sorted like a binary search tree according to the challenge value in each node
- Besides the challenge c_i , a hash value is stored in each node, where $h_i = H(c_i, hash value stored in its left child, hash value stored in its right child)$

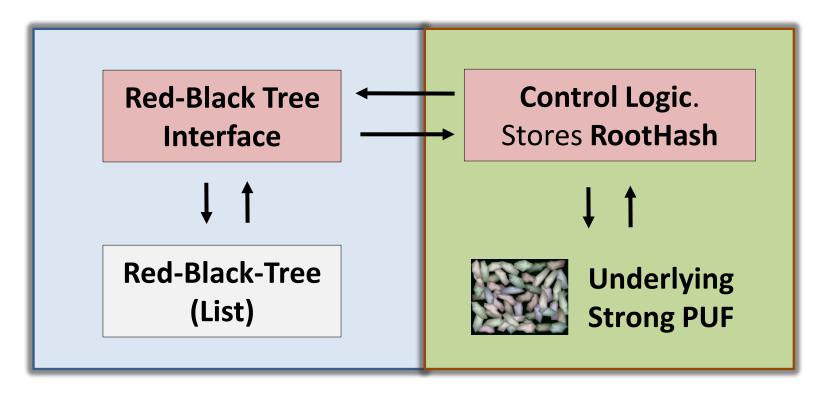




- In each node of the tree, we store one unique challenge, and the tree is sorted like a binary search tree according to the challenge value in each node
- Besides the challenge c_i , a hash value is stored in each node, where $h_i = H(c_i, hash value stored in its left child, hash value stored in its right child)$



GeniePUF Architecture

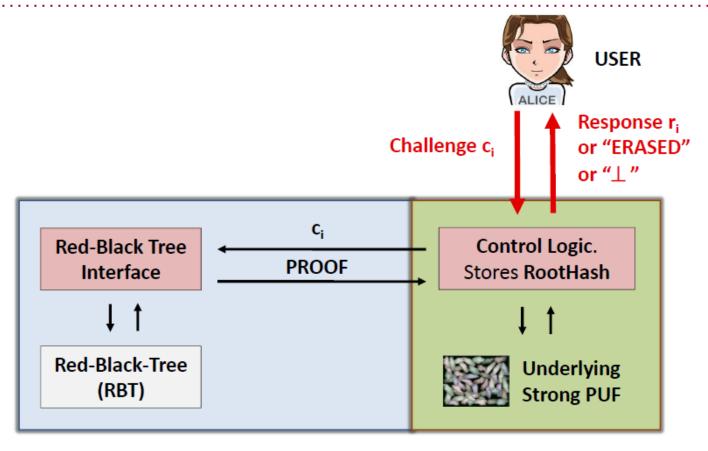


Public, Untrusted

System Part (Software)

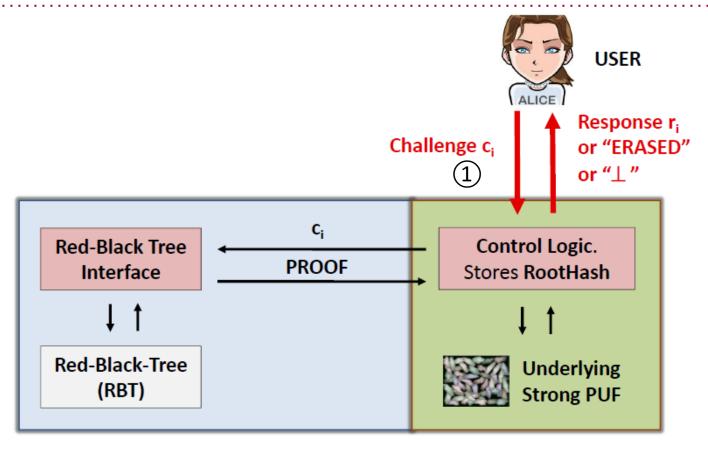
Trusted Computing Base (Hardware) of GeniePUF





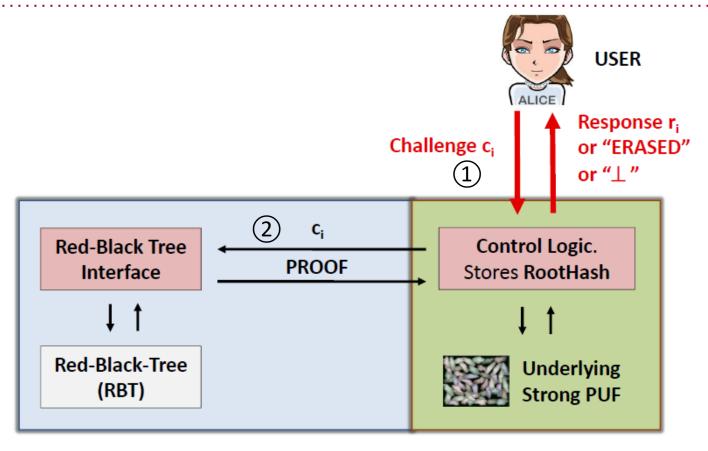
Public, Untrusted System Part





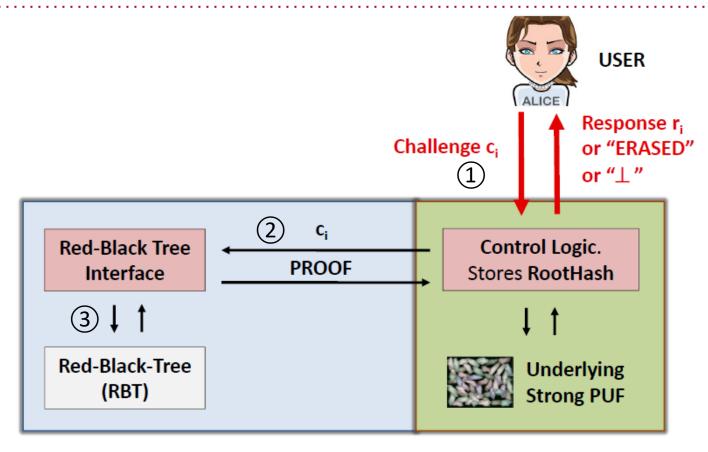
Public, Untrusted System Part





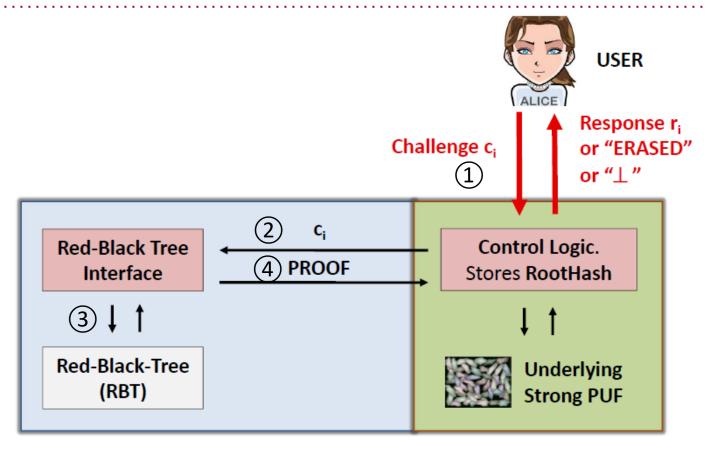
Public, Untrusted System Part





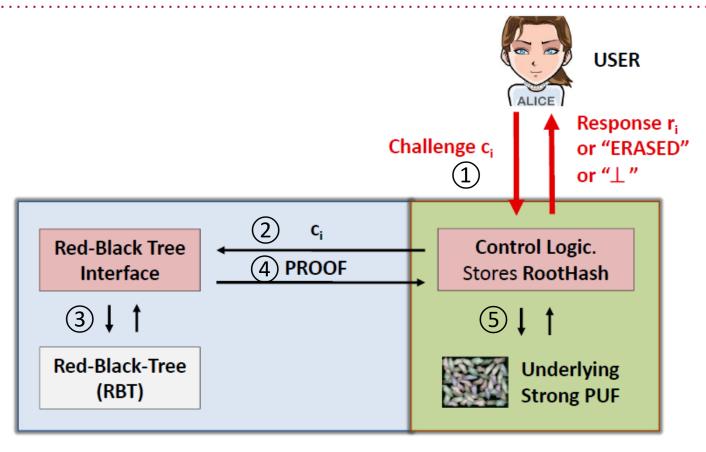
Public, Untrusted System Part





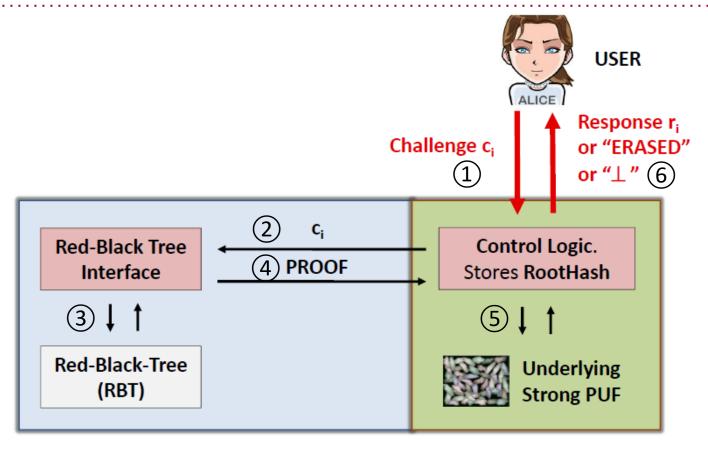
Public, Untrusted System Part





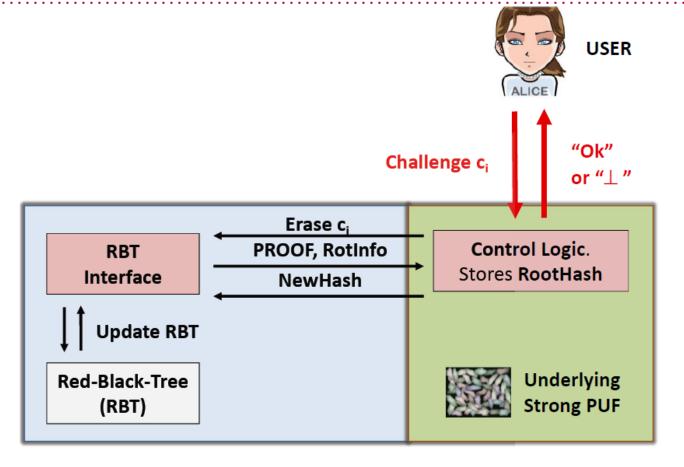
Public, Untrusted System Part





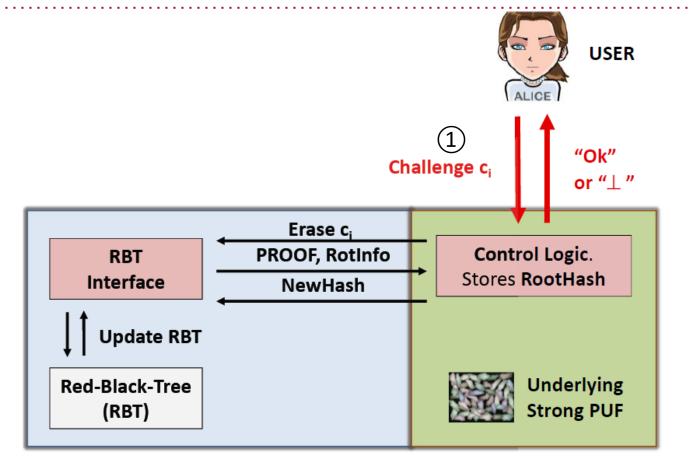
Public, Untrusted System Part





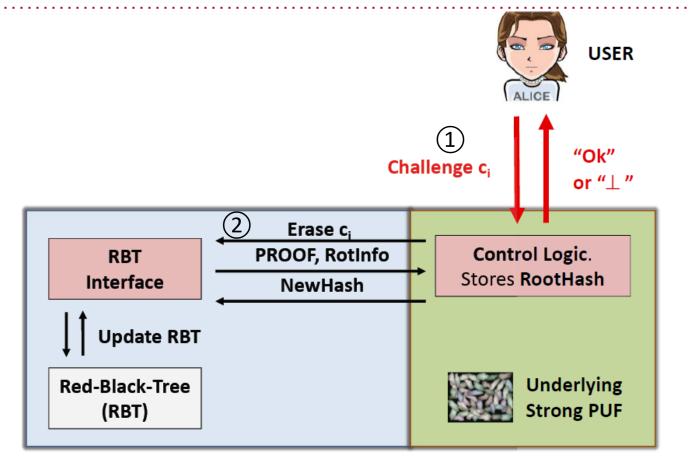
Public, Untrusted System Part





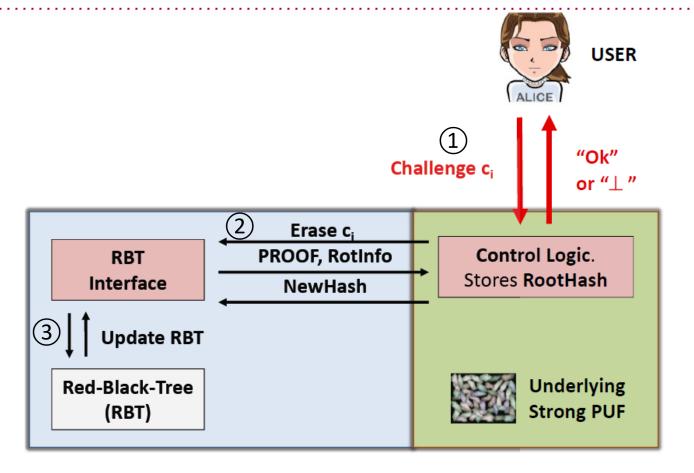
Public, Untrusted System Part





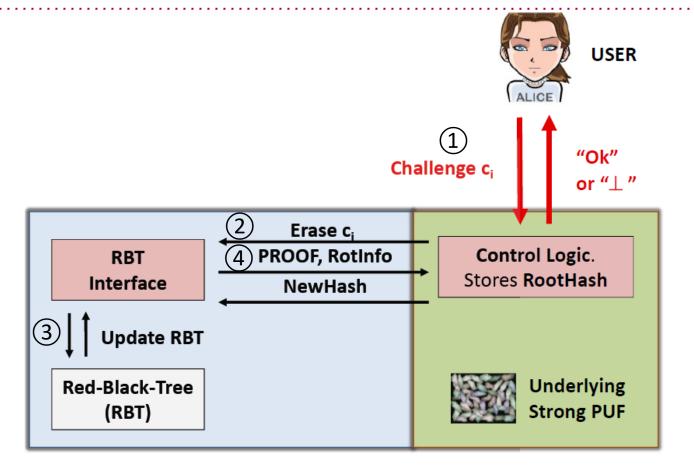
Public, Untrusted System Part





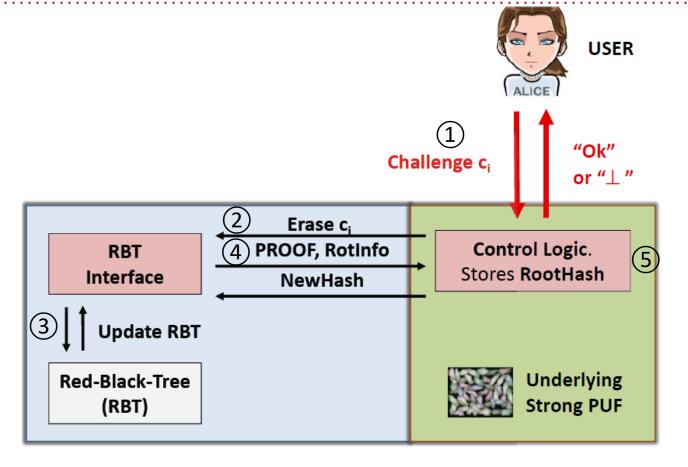
Public, Untrusted System Part





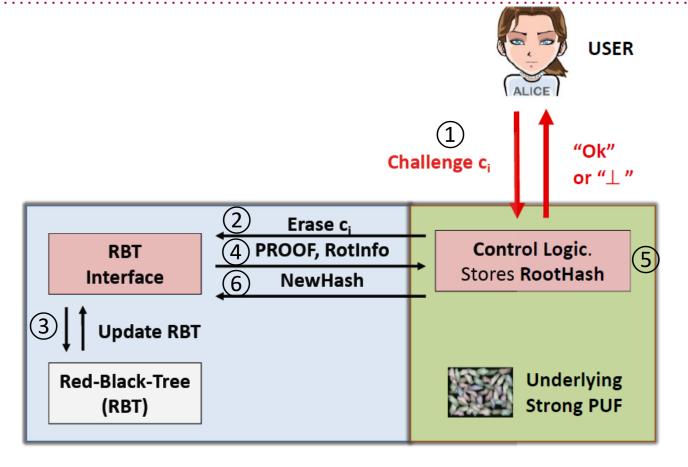
Public, Untrusted System Part





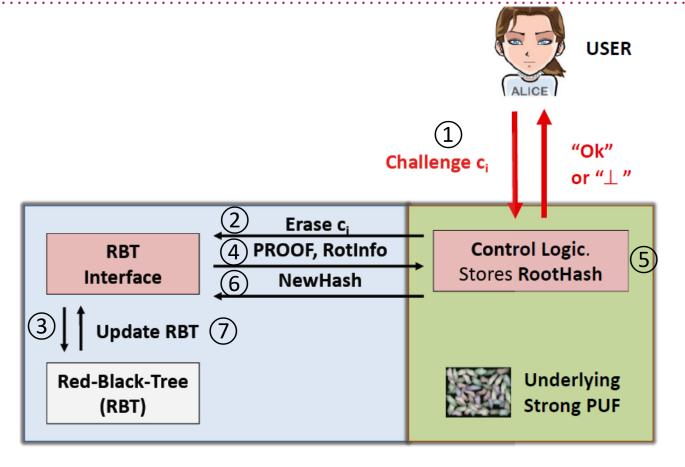
Public, Untrusted System Part





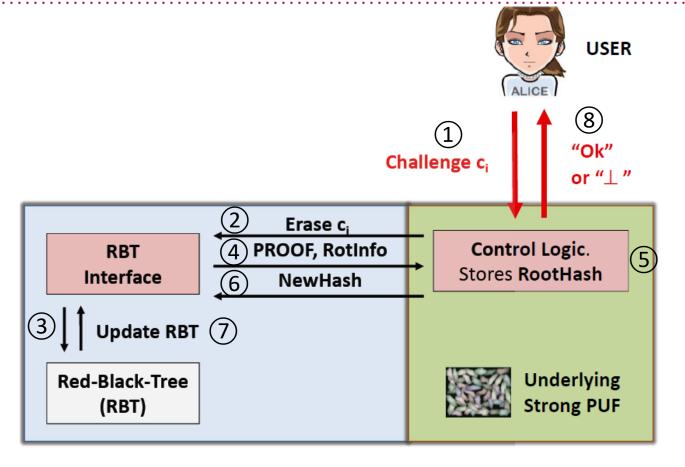
Public, Untrusted System Part





Public, Untrusted System Part

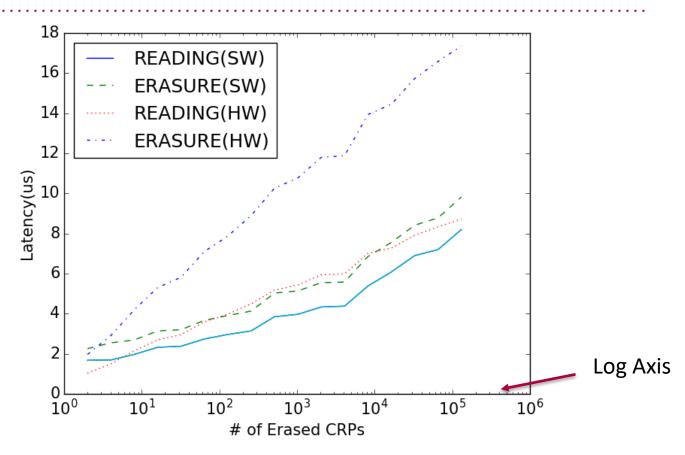




Public, Untrusted System Part



Performance Evaluation



- Implement the TCB on Zynq FPGA (HW) and the RB Tree Interface on Processor (SW)
- Latency grows logarithmically w.r.t. the number of erased challenges



Security Analysis

- Security Assumptions for Genie PUFs
- 1. Adversaries cannot circumvent the Control Logic (CL), applying their own challenges directly to the underlying Strong PUF, reading out the corresponding responses r_i .
- 2. Adversaries cannot modify the CL, for example such that it cannot correctly verify the validity of PROOF.
- 3. Adversaries may read the stored RootHash, but not modify it. It is public, but authentic.



A New Definitional Framework of PUFs

- Easily accessible, yet precise style PUF definition
- Parameterized Game-based PUF definition (ϵ, t_{att}, k)
- Intuition of Secure Erasable PUF Definition:

The security of an erasable PUF is measured by the upper bound ϵ of the accuracy of guessing one out of k randomly chosen CRPs by an attacker which takes time t_{att} for computation, physical actions, and k times game interactions with the challenger, where in each game interaction a randomly chosen CRP is erased.



Main Results of Formal Analysis

- Erasable PUFs are Strong PUFs
- Let P be a (k, t_{att}, ϵ) -secure Erasable PUF with respect to some adversary A. Then P is a (k, t_{att}, ϵ) -secure Strong PUF with respect to the same adversary A.
- The Security of Genie PUFs
- Let P be a PUF with challenge set C_P . Let A be an adversary for GeniePUF(P). Then GeniePUF(P) is $(k, t_{att}, \epsilon + \rho)$ -secure Erasable PUF with respect to A, where ρ represents the collision probability of the used hash function.



Conclusion

- Fixed the issue of PUF re-use model in PUF-based cryptographic protocols by using erasable PUFs.
- Introduced a generic erasable PUF design (Genie PUF) that can turn any strong PUFs to erasable PUFs.
- Proposed a rigorous, yet easily accessible definitional framework of PUF and proved our main theorems in the framework

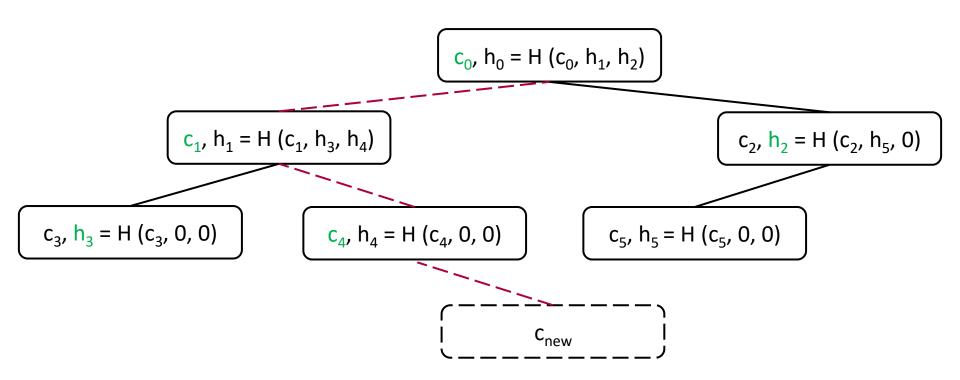


Thank you for your attention!

Questions?



Authenticated Search Tree Proof Generation



- Locate where the new challenge is supposed to be stored
- 2. Find a path from the new node for c_{new} to the root
- Fetch all the challenge values and all sibling hash values to construct a proof of (non)-existence



Red-Black Tree Background

- Self-balancing Binary Search Tree
- Guarantee O(log N) worst-case search time with a tree of size N

