

高等线性代数

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Week 15

1. 将下列对称多项式用初等对称多项式表示:

$$(x_1 + x_2 + x_3)^3 - (x_2 + x_3 - x_1)^3 - (x_1 + x_3 - x_2)^3 - (x_1 + x_2 - x_3)^3.$$

解. 观察到首项 x_1^3 的系数为零, 可能的项的系数有 $(2, 1, 0), (1, 1, 1)$, 故设

$$\text{原式} = a\sigma_1\sigma_2 + b\sigma_3$$

分别代入 $x_1 = x_2 = x_3 = 1, x_1 = x_2 = 1, x_3 = 0$ 得

$$24 = 9a + b$$

$$0 = a$$

故原式 $= 24\sigma_3$.

□

2. 用牛顿公式将 s_4 用初等对称多项式表示出来.

解. $n \geq 4$ 时,

$$s_2 = \sigma_1 s_1 - 2\sigma_2 = \sigma_1^2 - 2\sigma_2$$

$$s_3 = \sigma_1 s_2 - \sigma_2 s_1 + 3\sigma_3 = \sigma_1^3 - 3\sigma_2\sigma_1 + 3\sigma_3$$

$$s_4 = \sigma_1 s_3 - \sigma_2 s_2 + \sigma_3 s_1 - 4\sigma_4 = \sigma_1^4 - 4\sigma_2\sigma_1^2 + 4\sigma_3\sigma_1 + 2\sigma_2^2 - 4\sigma_4$$

$n = 3$ 时,

$$s_2 = \sigma_1^2 - 2\sigma_2$$

$$s_3 = \sigma_1 s_2 - \sigma_2 s_1 + 3\sigma_3 = \sigma_1^3 - 3\sigma_2\sigma_1 + 3\sigma_3$$

$$s_4 = \sigma_1 s_3 - \sigma_2 s_2 + \sigma_3 s_1 = \sigma_1^4 - 4\sigma_2\sigma_1^2 + 4\sigma_3\sigma_1 + 2\sigma_2^2$$

$n = 2$ 时,

$$s_2 = \sigma_1^2 - 2\sigma_2$$

$$s_3 = \sigma_1 s_2 - \sigma_2 s_1 = \sigma_1^3 - 3\sigma_2\sigma_1$$

$$s_4 = \sigma_1 s_3 - \sigma_2 s_2 = \sigma_1^4 - 4\sigma_1^2\sigma_2 + 2\sigma_2^2$$

$n = 1$ 时,

$$s_4 = \sigma_1^4$$

□

3. 已知 $x^3 + px^2 + qx + r = 0$ 的 3 个根为 $\alpha_1, \alpha_2, \alpha_3$, 求一个三次方程, 其根为 $\alpha_1^3, \alpha_2^3, \alpha_3^3$.

解. $f(x) = (x - \alpha_1^3)(x - \alpha_2^3)(x - \alpha_3^3)$ 即为所求. 可知

$$\sigma_1 = \alpha_1 + \alpha_2 + \alpha_3 = -p$$

$$\sigma_2 = \alpha_1\alpha_2 + \alpha_2\alpha_3 + \alpha_1\alpha_3 = q$$

$$\sigma_3 = \alpha_1\alpha_2\alpha_3 = -r$$

$$\alpha_1^3\alpha_2^3\alpha_3^3 = -r^3$$

$$\alpha_1^3\alpha_2^3 + \alpha_2^3\alpha_3^3 + \alpha_1^3\alpha_3^3 = \sigma_2^3 - 3\sigma_1\sigma_2\sigma_3 + 3\sigma_3^2 = q^3 - 3pqr + 3r^2$$

$$\sigma_1^3 + \sigma_2^3 + \sigma_3^3 = \sigma_1^3 - 3\sigma_2\sigma_1 + 3\sigma_3 = -p^3 + 3pq - 3r$$

故

$$f(x) = x^3 + (p^3 - 3pq + 3r)x^2 + (q^3 - 3pqr + 3r^2)x + r^3$$

□

4. 解下列方程组

$$\begin{cases} x_1 + x_2 + x_3 + x_4 = 4, \\ x_1^2 + x_2^2 + x_3^2 + x_4^2 = 4, \\ x_1^3 + x_2^3 + x_3^3 + x_4^3 = 4, \\ x_1^4 + x_2^4 + x_3^4 + x_4^4 = 4 \end{cases}$$

解. 即

$$\begin{cases} \sigma_1 = 4, \\ s_2 = \sigma_1^2 - 2\sigma_2 = 4, \\ s_3 = \sigma_1^3 - 3\sigma_2\sigma_1 + 3\sigma_3 = 4, \\ s_4 = \sigma_1^4 - 4\sigma_2\sigma_1^2 + 4\sigma_3\sigma_1 + 2\sigma_2^2 - 4\sigma_4 = 4 \end{cases}$$

得 $\sigma_1 = 4, \sigma_2 = 6, \sigma_3 = 4, \sigma_4 = 1$. 故

$$f(x_i) = (x_i - x_1)(x_i - x_2)(x_i - x_3)(x_i - x_4) = x_i^4 - 4x_i^3 + 6x_i^2 - 4x_i + 1 = (x - 1)^4 = 0, i = 1, 2, 3, 4$$

故唯一解为 $x_1 = x_2 = x_3 = x_4 = 1$.

□

5. 计算下列多项式的结式:

$$f(x) = x^3 + 3x^2 - x + 4, \quad g(x) = x^2 - 2x - 1$$

解.

$$\begin{vmatrix} 1 & 3 & -1 & 4 & 0 \\ 0 & 1 & 3 & -1 & 4 \\ 1 & -2 & -1 & 0 & 0 \\ 0 & 1 & -2 & -1 & 0 \\ 0 & 0 & 1 & -2 & -1 \end{vmatrix} = \begin{vmatrix} 0 & 5 & 0 & 4 & 0 \\ 0 & 1 & 7 & -9 & 0 \\ 1 & -2 & -1 & 0 & 0 \\ 0 & 1 & -2 & -1 & 0 \\ 0 & 0 & 1 & -2 & -1 \end{vmatrix} = - \begin{vmatrix} 5 & 0 & 4 \\ 1 & 7 & -9 \\ 1 & -2 & -1 \end{vmatrix} = 161$$

□

6. 设 $V = \mathbb{K}[x]$ 是数域 \mathbb{K} 上的全体多项式组成的线性空间. D 是 V 上的线性变换. 若 D 适合

$$(1) D(x) = 1;$$

$$(2) D(f(x)g(x)) = g(x)D(f(x)) + f(x)D(g(x)).$$

求证: D 就是 $\mathbb{K}[x]$ 上的求导变换.

证明. 由题意,

$$D(f+g) = D(f) + D(g), D(cf) = cD(f), c \in \mathbb{K}$$

假设 $D(x^n) = nx^{n-1}$ 成立, 根据题意, $n=1$ 时成立. 又 $n > 1$ 时,

$$D(x^{n+1}) = xD(x^n) + x^n D(x) = nx^n + x^n = (n+1)x^n$$

结合归纳假设, 又

$$D(1) = D(1 \cdot 1) = D(1) + D(1) \implies D(1) = 0$$

故对 $\mathbb{K}[x]$ 的一组基得到 D 是求导变换, 故 D 就是 $\mathbb{K}[x]$ 上的求导变换. □

7. 证明: $x^n + ax^{n-m} + b$, 其中 $n > m \geq 1, b \neq 0$, 不能有重数大于 2 的重根.

证明. 假设存在重数多于 2 的根 α , 显然 $\alpha \neq 0$. 假设 $n-m > 1$ 则分别求两次导后

$$f(x) = nx^{n-1} + a(n-m)x^{n-m-1}$$

$$g(x) = n(n-1)x^{n-2} + a(n-m)(n-m-1)x^{n-m-2}$$

则 $f(\alpha) = g(\alpha) = 0$, 整理得

$$n(n-1)\alpha^m + (n-m)(n-m-1)a = 0$$

$$n(n-1)\alpha^m + (n-m)(n-1)a = 0$$

相减得

$$am(n-m) = 0 \implies a = 0$$

代回可知重根只可能为零, 矛盾. 故不存在重数多于 2 的根.

假设 $n-m=1$, 则 $g(x)$ 的根只可能为 0, 同理矛盾, 故不存在重数多于 2 的根. \square

8. 求下列多项式的判别式 ($n > 1$):

(1) $f(x) = x^n + px + q;$

(2) $f(x) = x^{n-1} + x^{n-2} + \cdots + x + 1.$

解. (1)

(2) 设 $\omega = e^{\frac{2\pi i}{n}}$, $\Delta(f) = \prod_{1 \leq i < j \leq n-1} (\omega^i - \omega^j)^2$

$$\Delta(f) = \frac{\prod_{1 \leq i < j \leq n} (\omega^i - \omega^j)^2}{(\omega - 1)^2 (\omega^2 - 1)^2 \dots (\omega^{n-1} - 1)^2} = \frac{A}{f^2(1)} = \frac{A}{n^2}, \quad A = \prod_{1 \leq i < j \leq n} (\omega^i - \omega^j)^2$$

$$\begin{aligned} A &= \prod_{1 \leq i < j \leq n} (\omega^j - \omega^i)^2 = \prod_{1 \leq i < j \leq n} \omega^i (\omega^{j-i} - 1) \prod_{1 \leq i < j \leq n} \omega^j (1 - \omega^{n-j+i}) \\ &= 1 \cdot 1 \cdot \omega^{\frac{n(n-1)}{2}} \prod_{1 \leq k \leq n-1} (\omega^k - 1)^{n-k} \prod_{1 \leq k \leq n-1} (\omega^k - 1)^k (-1)^{\frac{n(n-1)}{2}} \\ &= f(1)^n (-1)^{\frac{n^2}{2} + \frac{n}{2} - 1} \end{aligned}$$

故 $\Delta(f) = n^{n-2} (-1)^{\frac{n^2}{2} + \frac{n}{2} - 1}$ □

9. 设 $f(x), g(x)$ 是数域 \mathbb{K} 上的互素多项式, A 是 \mathbb{K} 上的 n 阶方阵, 且满足 $f(A) = O$, 证明: $g(A)$ 是可逆阵.

证明. 由互素得, 存在 $K[x]$ 中的多项式 u, v 使得

$$fu + gv = 1$$

代入 $x = A$ 得

$$O + g(A)v(A) = I$$

故 $g(A)$ 可逆. □

注. 完全理解不能

10. 设 $f(x), g(x)$ 是次数大于 1 的多项式, 求证:

$$\Delta(f(x)g(x)) = \Delta(f(x))\Delta(g(x))R(f, g)^2$$

证明. 设 $f(x) = a_0 x^n + \dots + a_n$, $g(x) = b_0 x^m + \dots + b_m$, f, g 的所有根分别为 α_i, β_j

设 $h(x) = f(x)g(x)$, 首项为 $a_0 b_0 x^{n+m}$, 其所有根为 α_i, β_j , 重新排序为 $\gamma_i, i = 1, 2, \dots, n+m$

$$\begin{aligned} \text{右式} &= (a_0^{2n-2} \prod_{1 \leq k_1 < k_2 \leq n} (\alpha_{k_1} - \alpha_{k_2})^2) (b_0^{2m-2} \prod_{1 \leq l_1 < l_2 \leq m} (\beta_{l_1} - \beta_{l_2})^2) (a_0^{2m} b_0^{2n} \prod_{i=1}^n \prod_{j=1}^m (\alpha_i - \beta_j)^2) \\ &= (a_0 b_0)^{2n+2m-2} \prod_{1 \leq k_1 < k_2 \leq n} (\alpha_{k_1} - \alpha_{k_2})^2 \prod_{1 \leq l_1 < l_2 \leq m} (\beta_{l_1} - \beta_{l_2})^2 \left(\prod_{i=1}^n \prod_{j=1}^m (\alpha_i - \beta_j)^2 \right) \\ &= (a_0 b_0)^{2n+2m-2} \prod_{1 \leq s_1 < s_2 \leq n+m} (\gamma_{s_1} - \gamma_{s_2})^2 \\ &= \Delta(f(x)g(x)) \end{aligned}$$

□