数据库系统 2013-2014-1 (A闭)

problem1

1.foreign key: department ID

实体完整性规则:作为主键的属性或属性的值在关系表中是唯一确定的.故外键是 kepartment_ID, 主键是 ID, 从而 ID 是唯一确定的,name 和 department_ID 是与它相关的.且 department_ID 是 DEPARTMENT 表的主键,是 STUDENT 的外键.

参照完整性规则:如果表中存在外键,则外键的值必须与主表的值相同或者外键为空.故,他的值依赖于 DEPARTMENT 表中的 ID 的值来觉得.

2.(1) r1-(r1-r2)

а	b
а	1
а	3

(2)

а	b
a	1
а	3

(3)

Α	
а	

3.(1)

∵A→BDE

$$A \rightarrow D, A \rightarrow E$$

∴E→C

∴A→C

∴A→CD

- (2)A 为主键,因为其他都依赖于 A,A→BCDE
- 4.不是串行的,应为它从 read(a)和 write(a)的优化图为<T1,T2> Read(b)和 write(b)的优化图是<T2,T1> 故这不是串行的,有环会相互干扰 problem2(答案不唯一)

1 (1)

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 | Project.ID, Project.name (\sigma_{sutdent.dept_{name} = 'software'}(project \bowtie_{project.leader\_ID = S.ID} student) | Project.ID, Project.name (\sigma_{sutdent.dept_{name} = 'software'}(project \bowtie_{project.leader\_ID = S.ID} student) | Project.name (\sigma_{sutdent.dept_{name} = 'software'}(project \bowtie_{project.leader\_ID = S.ID} student) | Project.name (\sigma_{sutdent.dept_{name} = 'software'}(project \bowtie_{project.leader\_ID = S.ID} student) | Project.name (\sigma_{sutdent.dept_{name} = 'software'}(project \bowtie_{project.leader\_ID = S.ID} student) | Project.name (\sigma_{sutdent.dept_{name} = 'software'}(project \bowtie_{project.leader\_ID = S.ID} student) | Project.name (\sigma_{sutdent.dept_{name} = 'software'}(project \bowtie_{project.leader\_ID = S.ID} student) | Project.name (\sigma_{sutdent.dept_{name} = 'software'}(project.name) | Project.name (\sigma_{sutdent.dept_{name} = 'software'}(
   (2).
 student.ID, student.name (\sigma_{project.name='Sun\ light'}(project\ \bowtie_{project.leader_{ID}=S.ID} student)
2 (1) select ID, name
from project
where budget > 1000 and name like 'computer%'
(2) select S.ID, S.name
  from student s left join participate p on S.ID = P.student ID
where S.dept name='Software' and P.project ID is NULL
(3) select p.ID, p.name count(p.student-ID) as my count
from project as p, participate as q
where p.ID = Q.project-ID
group by p.ID
order by myCount desc
(4) select distinct S.name
from student s
where not exists
((select project-ID from participate , project where project-ID= ID and leader-
ID='1010128')
except
(select project-ID from student, participate where ID=student-ID)
problem 3
1.这个不知道怎么回答,知道的请说明一下,谢谢
fd1:car_VIN→all(因为车牌号是唯一的,知道车牌号就可以知道客户信息以及车的
相关信息)
fd2: cunstomer-id-)customer name
fd3:model→volume
3. 第二范式 因为不存在部分依赖,但 fd2 有传递依赖(car VIN→cunstomer-id,
cunstomer-id→customer name, car VIN→cunstomer name)
4.消除传递依赖,得到第三范式
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customer (<u>cunstomer-id</u>, customer_name)
car(car_VIN, cunstomer-id,colour,model,volume,price,sale_date)
该范式也是 BCNF 范式,因为函数依赖都是候选关键字。
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problem 4

Author(aname, address, URL)

Book(<u>ISBN</u>, title, year, price)

Publisher(pname, address, phone, URL)

Shopping_basket(basket id, email)

Warehouse(<u>code</u>, address, phone)

Customer(email, name, address, phone)

Written_by(<u>aname</u>, <u>ISBN</u>)

Published_by(pname, ISBN)

Contains(<u>ISBN</u>, <u>email</u>, <u>basket_id</u>, number)

Stocks(<u>ISBN</u>, <u>code</u>, number)

Basket_of(<u>email</u>, <u>basket_id</u>)(可省略)