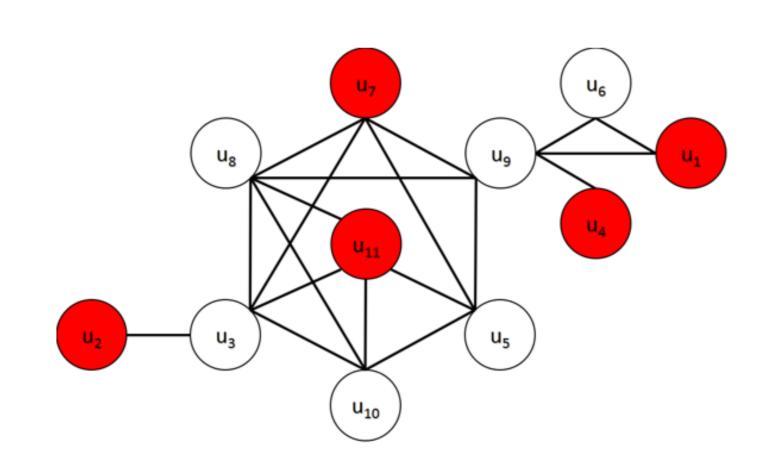
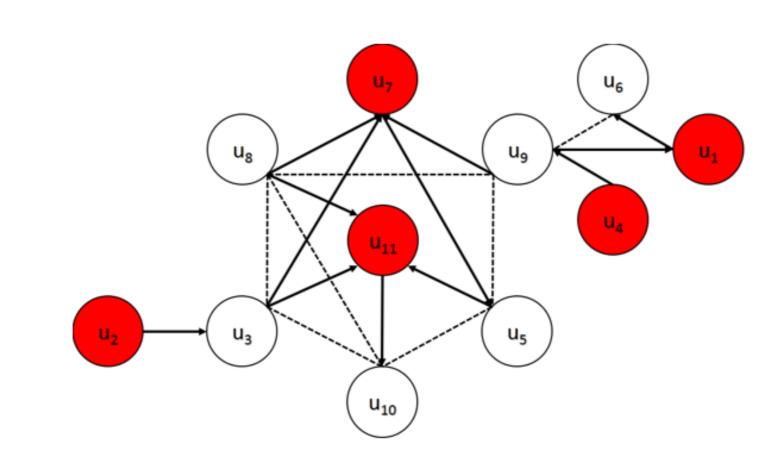
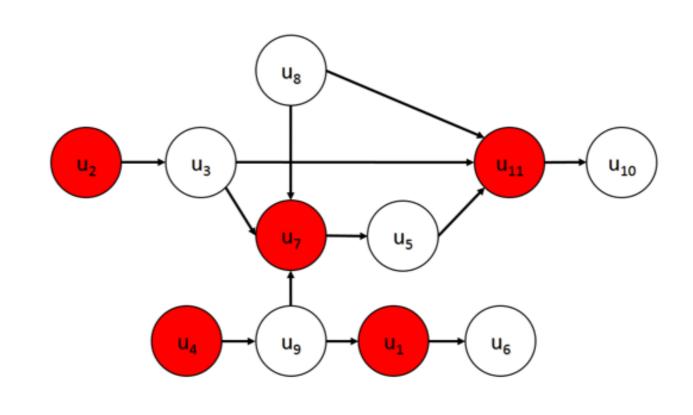
353: Computing A Near-Maximum Independent Set In Dynamic Graphs

Weiguo Zheng^{1,2} Chengzhi Piao¹ Hong Cheng¹ Jeffrey Xu Yu¹

¹The Chinese University of Hong Kong, HK SAR, China ²Fudan University, China







Background

Independent Set

Given a set of vertices in graph G, we call it an independent set of G if and only if there are no neighborhoods among them.

Maximal independent set

Can not add more vertices into it.

Maximum independent set(MIS)

A MIS has the largest size among all independent sets. Maybe not unique.

Problem Definition

Target: maintain a Near-MIS in dynamic graphs

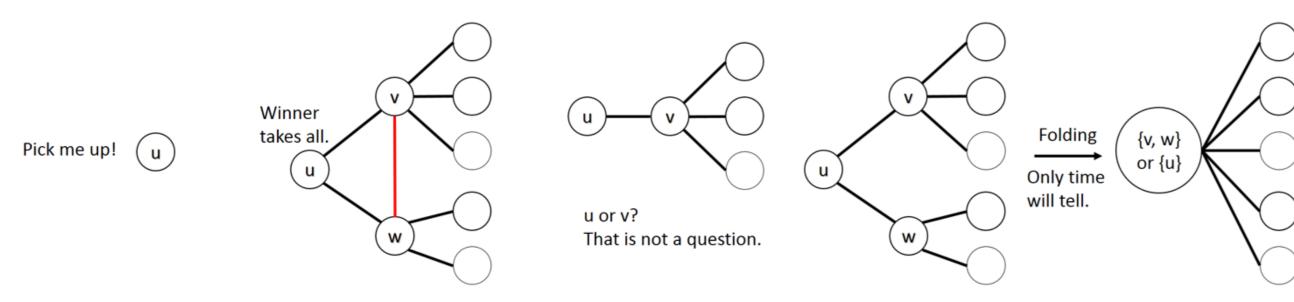
Input

An undirected graph G_0 and the initial $NearMIS_0$ A sequence of updates on the graph structure: $update_i$

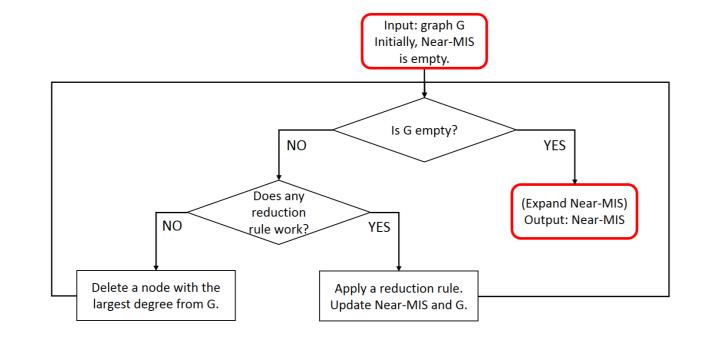
Output online maintain the Near-MIS after each update: $NearMIS_i$

Dependency Graph based Framework

Reduction Rule

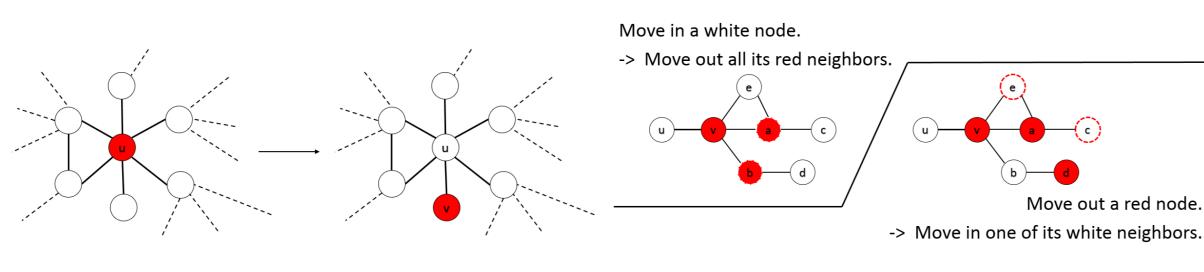


Reducing-Peeling Framework

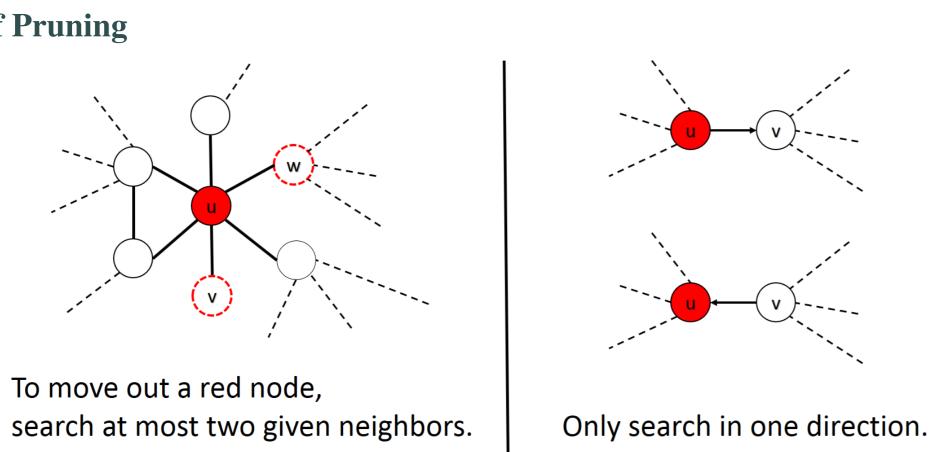


Brute Force Search: Move Out a Node from the MIS

To maintain a MIS, we must choose a neighbor to replace u.



Two Ways of Pruning



Definition of Dependency Graph

We call the nodes in the independent set reducing nodes.

Others are called dependent nodes.

- 1. Contain all edges between reducing and dependent nodes.
- 2. Each edge is directed.
- 3. Each reducing node has at most two out-neighbors.
- 4. Each dependent node has at most one in-neighbor.

DG: Construction

The reducing-peeling framework does not merely output a Near-MIS, but also implies a topological order of the graph.

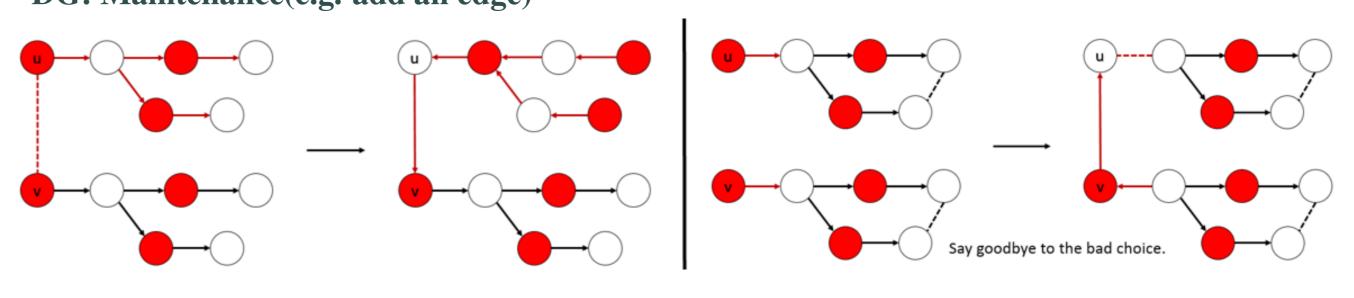
- 1. For each reducing node, its out-neighbors are decided when a reduction rule works.
- 2. Then, the others become its in-neighbors.

DG: Dealing with Updates

We mainly focus on the following three operations.

- 1. Delete a node u: If u is a reducing node, try to move out u.
- 2. Add an edge (u, v): If both u and v are reducing nodes, try to move out u or v. Simply abandon either if it fails.
- 3. Delete an edge (u, v): Try to move in the dependent node.

DG: Maintenance(e.g. add an edge)



Batch Update

- Ga: affected subgraph G_r: the remaining graph
- Crossing Edge: an edge between G_a and G_r
- For each crossing edge (u, v)
- u is in G_a and v is in G_{r.}

• Then u is called an active node.

- v is a reducing node(v is in Near-MIS).
- Delete all crossing edges which contains active nodes as few as possible Construct DGOracle_a on Ga with the help of MISa Add crossing edges back

Delete a node **not in** Near-MIS₀ with the largest degree from G. **DGOracle**

DGOracle

Use $NearMIS_0$ to guide the inexact reduction.

Results



Conclusions

(a) RGap vs. k

In this paper, we study the problem of computing the high-quality(may not be the maximum) independent sets over the graphs that are dynamically changing.

(b) Response time vs. k

Fig. 6. Effect of the batch size

Brightkite Slashdot

Based on the two widely used degree-one and degree-two reduction rules, we devise a very effective index, i.e., dependency graph, that can facilitate the computation of independent sets. The time complexity of searching on a dependency graph is merely O(d) in average case.

Extensive experiments over a wide range of graphs confirm that the proposed methods are both effective and efficient to find high-quality independent sets.