〈 题目集列表

■ 题目集概况

题目列表

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判断题(共 5 分)

1 2 3 4

A. 单选题(共 3 分) 3/3

1 2 3

② 程序填空题(共 2 分)

1

ADS-Mid-Term

※ 判断题 5

A. 单选题 3

◎ 程序填空题 1

2-1 In proving the amortized bound of a Merge operation in skew heaps, the potential of a skew heap is defined to be the total number of right heavy nodes. Then we can prove that, in an N-node skew heap, the amortized cost for a Merge operation is exactly $\underline{}$.

☆ 作者: 沈鑫

単位: 浙江大学

Hint:

5/5

Define the weight of a node, w(x), to be the number of descendants of x (including x). A non-root node is said to be *heavy* if its weight is greater than half the weight of its parent.

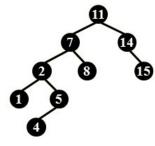
- Lemma 1: At most one child is heavy, of all children of any node.
- *Lemma 2*: On any path from node x down to a descendant y, there are at most $\lfloor log_2 \frac{w(x)}{w(y)} \rfloor$ light nodes, excluding x.

(1分)

- \bigcirc A. $|log_2N|+1$
- lacksquare B. $2\lfloor log_2 N
 floor + 1$
- \bigcirc C. $3\lfloor log_2 N \rfloor + 1$
- \bigcirc D. $4\lfloor log_2 N
 floor + 1$
- 2-2 For the result of accessing the keys 4 and 8 in order in the splay tree given in the figure, which one of the following statements is FALSE? (1分)



🞑 作者: DS课程组



- A. 8 is the root
- B. 4 and 11 are siblings
- C. 7 and 14 are siblings
- O. 4 is the parent of 7
- 2-3 To solve a problem with input size N by divide and conquer algorithm, among the following methods, $_$ is the worst. (1分)
 - 单位: 浙江大学 nquer in
 - $\ \, \square$ A. divide into 2 sub-problems of equal complexity N/3 and conquer in O(N)
 - \bigcirc B. divide into 2 sub-problems of equal complexity N/3 and conquer in O(NlogN)
 - $\ \odot$ C. divide into 3 sub-problems of equal complexity N/2 and conquer in O(N)
 - \bigcirc D. divide into 3 sub-problems of equal complexity N/3 and conquer in O(NlogN)

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