



SCHOOL OF
COMPUTING

B.S. Chenthil Hari
CH.SC.U4CSE24103

Week - 3

Date - 03/01/2026

Design and Analysis of Algorithm(23CSE211)

1. Merge Sort

Code:

```
venv>g++ -o merge merge.cpp
1 #include <stdio.h>
2 void merge(int arr[], int l, int m, int r) { int n1 = m - l + 1;
3 int n2 = r - m; int L[n1], R[n2];
4 for (int i = 0; i < n1; i++) L[i] = arr[l + i]; for (int j = 0; j < n2; j++) R[j] = arr[m + 1 + j]; int i = 0, j = 0, k = l
5 while (i < n1 && j < n2) {
6 if (L[i] <= R[j]) arr[k++] = L[i++]; else arr[k++] = R[j++];
7 }
8 while (i < n1) arr[k++] = L[i++]; while (j < n2) arr[k++] = R[j++];
9 }
10 void mergeSort(int arr[], int l, int r) { if (l < r) {
11 int m = l + (r - l) / 2; mergeSort(arr, l, m); mergeSort(arr, m + 1, r); merge(arr, l, m, r);
12 }
13 }
14 int main() {
15 int n; printf("CH.SC.U4CSE24103\n");
16 printf("Enter number of elements: "); scanf("%d", &n);
17 int arr[n];
18 printf("Enter the elements:");
19 for (int i = 0; i < n; i++) scanf("%d", &arr[i]); mergeSort(arr, 0, n - 1);
20 printf("Sorted array: ");
21 for (int i = 0; i < n; i++) printf("%d ", arr[i]); printf("\n");
22 return 0;
23 }
```

Output:

```
C:\Users\chent\OneDrive - An + ▾
e CH.SC.U4CSE24103
n Enter number of elements: 5
o Enter the elements:65
i 34
r 23
c 46
n 67
e Sorted array: 23 34 46 65 67
=
-----
Process exited after 4.233 seconds with return value 0
i Press any key to continue . . . |
```

Space Complexity: $O(n \log n)$

Justification: The algorithm recursively divides the array into halves, which takes $\log n$ levels of division. At each level, it performs a linear merge process requiring n operations, resulting in a consistent $n \times \log n$ performance.

Time Complexity: $O(n \log n)$

Justification: Merge Sort requires an auxiliary array to temporarily hold elements during the merge process. This additional memory is proportional to the size of the input array.

2. Quick Sort

Code:

```
1 #include <stdio.h>
2 void swap(int* a, int* b) { int t = *a;
3   *a = *b;
4   *b = t;
5 }
6 int partition(int arr[], int low, int high) { int pivot = arr[high];
7   int i = (low - 1);
8   for (int j = low; j <= high - 1; j++) { if (arr[j] < pivot) {
9     i++;
10    swap(&arr[i], &arr[j]);
11  }
12 }
13 swap(&arr[i + 1], &arr[high]); return (i + 1);
14 }
15 void quickSort(int arr[], int low, int high) { if (low < high) {
16   int pi = partition(arr, low, high); quickSort(arr, low, pi - 1); quickSort(arr, pi + 1, high);
17 }
18 }
19 int main() {
20   int n;
21   printf("Enter number of elements: "); scanf("%d", &n); printf("CH.SC.U4CSE24103\n");
22   int arr[n];
23   printf("Enter the elements:");
24   for (int i = 0; i < n; i++) scanf("%d", &arr[i]); quickSort(arr, 0, n - 1);
25   printf("Sorted array: ");
26   for (int i = 0; i < n; i++) printf("%d ", arr[i]); printf("\n");
27   return 0;
28 }
```

Output:

```
C:\Users\chent\OneDrive - An + ▾
Enter number of elements: 5
CH.SC.U4CSE24103
Enter the elements:43
23
67
45
2
Sorted array: 2 23 43 45 67
-----
Process exited after 5.738 seconds with return value 0
Press any key to continue . . . |
```

Space Complexity: $O(n^2)$

Justification: This occurs when the pivot consistently picks the smallest or largest element (e.g., on a sorted array). This results in highly unbalanced partitions where one side has 0 elements and the other has $n-1$, leading to n recursive levels with $O(n)$ work each.

Time Complexity: $O(n)$

Justification: In the worst-case scenario of an unbalanced partition, the recursion stack depth increases from the ideal $\log n$ to n . Each recursive call stays on the stack until it finishes, requiring memory proportional to the number of elements.