## 1 Problem

Given a directed graph G = (V, E) and one sequence of constraints  $P = (L \times L)^k$  of length k, where L is a set of labels and |L| = |V|, we want to find a sequence of sequences of adjacent mappings  $S = \langle F^1, F^2, \dots, F^k \rangle$  minimizing the total cost, such that all constraints  $p_i \in P$  are satisfied by the mapping  $F_{m_i}^i$   $(m_i = |F^i|)$ .

## 2 Adjacent Mapping

First of all, a mapping is a bijective function  $f:V\to L$  that maps the vertices of G into the labels L.

From the definition above, given two mappings f and f', we say that they are adjacent iff there is exactly one edge  $(u,v) \in E$  such that f(u) = f'(v), f(v) = f'(u) and  $\forall w \in V \setminus \{u,v\}(f(w) = f'(w))$ .

## 3 Constraints

A constraint is defined by an ordered pair (a,b) such that  $a,b \in L$ . We say that a mapping f satisfies a constraint (a,b) iff there is an edge  $(u,v) \in E$  such that f(u) = a and f(v) = b.

## 4 Output

We seek a sequence of sequences of adjacent mappings  $S = \langle F^1, F^2, \dots, F^k \rangle$ , such that its total cost is minimized.

Consider that the operation of transition from one mapping to its adjacent mapping is called swap operation. Each  $F^i \in S$  is a sequence of adjacent mappings  $\langle f_0^i, f_1^i, \ldots, f_{m_i}^i \rangle$  from an initial mapping  $(f_0^i)$  to a final mapping  $(f_{m_i}^i)$ , that should satisfy the *i*-th constraint. In other words,  $m_i - 1$  represents the number of swap operations required in order to satisfy constraint  $p_i$  from mapping  $f_0^i = f_{m_{i-1}}^{i-1}$  (note that  $f_0^1 = f^0$  – an initial mapping).

Also, given the *i*-th constraint of P  $p_i = (a, b)$ , if  $f_{m_i}^i$  does not satisfy  $p_i$ , then it should satisfy  $p'_i = (b, a)$ , using a reverse operation.

The total cost of such sequence is given by the sum of all swap and reverse operations. Given that the swap and reverse operation costs  $C_{swap}$  and  $C_{rev}$ , respectively, and the function h (Equation 1), the total cost is (Equation 2):

$$h(f,i) = \begin{cases} 0, & \text{if } f \text{ satisfies } p_i \\ C_{rev} \end{cases}$$
 (1)

$$C_{total} = \sum_{i=1}^{k} |F^i| \cdot C_{swap} + h(f_{m_i}^i, i)$$
 (2)