

Terms

Literals	$l ::= true \mid false \mid Bool \mid \mathbb{R} \mid Float \mid () \mid Unit \mid Type \mid \square$
Variables	$v ::= [a - z A - Z][a - z A - Z 1 - 9]^*$
Expressions	$e, t, k ::=$ $\mid l$ $\mid v$ $\mid \lambda v : t. e$ $\mid \forall v : t. e$ $\mid e_f \ e_a$ $\mid \text{if } e_p \text{ then } e_c \text{ else } e_a$ $\mid \text{let } v = e_v \text{ in } e_b$ $\mid \text{letrec } v : t = e_v \text{ in } e_b$ $\mid (e_l, e_r)$ $\mid e.L$ $\mid e.R$ $\mid e.R$ $\mid (t_l \wedge t_r)$ $\mid (e : t)$ $\mid (e)$

Typing

Literals

$\frac{}{\vdash true : Bool}$ BLit	$\frac{}{\vdash false : Bool}$ BLit	$\frac{}{\vdash Bool : Type}$ BTLit
$\frac{}{\vdash \mathbb{R} : Float}$ FLit	$\frac{}{\vdash Float : Type}$ FTLit	
$\frac{}{\vdash () : Unit}$ ULit	$\frac{}{\vdash Unit : Type}$ UTLit	
	$\frac{}{\vdash Type : \square}$ TLit	

Variables

$\frac{}{\Gamma, v : t \vdash v : t}$ Var

Compounds

$$\begin{array}{c}
\frac{\Gamma \vdash t : Type}{\Gamma \vdash t :^t Type} \text{IsType} \quad \frac{\Gamma \vdash t : \Box}{\Gamma \vdash t :^t \Box} \text{IsKind} \\
\\
\frac{\Gamma \vdash t_{uv} :^t k_{uv} \quad t_{uv} \Downarrow t_v \quad \Gamma, t_v : k \vdash e : t_b}{\Gamma \vdash (\lambda v : t_{uv}.e) : (\forall v : t_v.t_b)} \lambda I \\
\\
\frac{\Gamma \vdash e_f : (\forall v : t_a.t_b) \quad \Gamma \vdash e_a : t_a \quad (\lambda v : t_a.t_b)e_a \Downarrow t_o}{\Gamma \vdash e_f e_a : t_o} \lambda E \\
\\
\frac{\Gamma \vdash t_{uv} :^t k_v \quad t_{uv} \Downarrow t_v \quad \Gamma, v : t_v \vdash t_{ub} :^t k_b}{\Gamma \vdash (\forall v : t_{uv}.t_{ub}) : k_b} \forall I \\
\\
\frac{\Gamma \vdash e_p : Bool \quad \Gamma \vdash e_c : t \quad \Gamma \vdash e_a : t}{\Gamma \vdash (if\ e_p\ then\ e_c\ else\ e_a) : t} \text{If} \\
\\
\frac{\Gamma \vdash e_v : t_v \quad \Gamma, v : t_v \vdash e_b : t_b}{\Gamma \vdash (let\ v\ =\ e_v\ in\ e_b) : t_b} \text{Let} \\
\\
\frac{\Gamma \vdash t_{uv} : k \quad t_{uv} \Downarrow t_v \quad \Gamma, v : t_v \vdash e_v : t_v \quad \Gamma, v : t_v \vdash e_b : t_b}{\Gamma \vdash (letrec\ v\ :\ t_{uv}\ =\ e_v\ in\ e_b) : t_b} \text{Letrec} \\
\\
\frac{\Gamma \vdash e_l : t_l \quad \Gamma \vdash e_r : t_r}{\Gamma \vdash (e_l, e_r) : (t_l \wedge t_r)} \text{Pair} \quad \frac{\Gamma \vdash t_l :^t k \quad \Gamma \vdash t_r :^t k}{\Gamma \vdash (t_l \wedge t_r) : k} \text{PairT} \\
\\
\frac{\Gamma \vdash e : (t_l \wedge t_r)}{\Gamma \vdash e.L : t_l} \text{PairL} \quad \frac{\Gamma \vdash e : (t_l \wedge t_r)}{\Gamma \vdash e.R : t_r} \text{PairR} \\
\\
\frac{\Gamma \vdash t_u :^t k \quad t_u \Downarrow t \quad \Gamma \vdash e : t}{\Gamma \vdash (e : t_u) : t} \text{Asc}
\end{array}$$

Normalization

$$\begin{array}{c}
\frac{}{l \Downarrow l} \text{Lit} \quad \frac{}{v \Downarrow v} \text{Var} \\
\\
\frac{t \Downarrow t' \quad e \Downarrow e'}{\lambda v : t.e \Downarrow \lambda v : t'.e'} \lambda \quad \frac{t_v \Downarrow t'_v \quad t_b \Downarrow t'_b}{\forall v : t_v.t_b \Downarrow \forall v : t'_v.t'_b} \forall \\
\\
\frac{e_f \Downarrow e'_f \quad e_a \Downarrow e'_a}{e_f e_a \Downarrow e'_f e'_a} \text{NeuApp}
\end{array}$$

$$\begin{array}{c}
\frac{e_f \Downarrow \lambda v : t.e_b \quad e_b[e_a/v] \Downarrow e'_b}{e_f e_a \Downarrow e'_b} \beta \\
\\
\frac{e_p \Downarrow e'_p \quad e_c \Downarrow e'_c \quad e_a \Downarrow e'_a}{\text{if } e_p \text{ then } e_c \text{ else } e_a \Downarrow \text{if } e'_p \text{ then } e'_c \text{ else } e'_a} \text{NeuIf} \\
\\
\frac{e_p \Downarrow \text{true} \quad e_c \Downarrow e'_c}{\text{if } e_p \text{ then } e_c \text{ else } e_a \Downarrow e'_c} \text{IfT} \quad \frac{e_p \Downarrow \text{false} \quad e_a \Downarrow e'_a}{\text{if } e_p \text{ then } e_c \text{ else } e_a \Downarrow e'_a} \text{IfF} \\
\\
\frac{e_b[e_v/v] \Downarrow e'_b}{\text{let } v = e_v \text{ in } e_b \Downarrow e'_b} \text{Let} \\
\\
\frac{e_l \Downarrow e'_l \quad e_r \Downarrow e'_r}{(e_l, e_r) \Downarrow (e'_l, e'_r)} \text{Pair} \quad \frac{t_l \Downarrow t'_l \quad t_r \Downarrow t'_r}{(t_l \wedge t_r) \Downarrow (t'_l \wedge t'_r)} \text{PairT} \\
\\
\frac{e \Downarrow e'}{e.L \Downarrow e'.L} \text{NeuPairL} \quad \frac{e \Downarrow (e_l, e_r)}{e.L \Downarrow e_l} \text{PairL} \\
\\
\frac{e \Downarrow e'}{e.R \Downarrow e'.R} \text{NeuPairR} \quad \frac{e \Downarrow (e_l, e_r)}{e.R \Downarrow e_r} \text{PairR} \\
\\
\frac{e \Downarrow e'}{(e : t) \Downarrow e'} \text{Asc}
\end{array}$$