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Project Report on

"16-Bit Booth's Multiplier"

Submitted by

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# DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING PES UNIVERSITY

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# 16-Bit Booth's Multiplier

## **Problem Description**

Booth's multiplication algorithm is a multiplication algorithm that multiplies two signed binary numbers in two's complement notation.

Booth's algorithm examines adjacent pairs of bits of the 'N'-bit multiplier Y in signed two's complement representation, including an implicit bit below the least significant bit,  $y_{-1} = 0$ . For each bit  $y_i$  for i running from 0 to N - 1, the bits  $y_i$  and  $y_{i-1}$  are considered.

Where these two bits are equal, the product accumulator P is left unchanged. Where  $y_i = 0$  and  $y_{i-1} = 1$ , the multiplicand times  $2^i$  is added to P; and where  $y_i = 1$  and  $y_{i-1} = 0$ , the multiplicand times  $2^i$  is subtracted from P. The final value of P is the signed product.

### **Implementation**

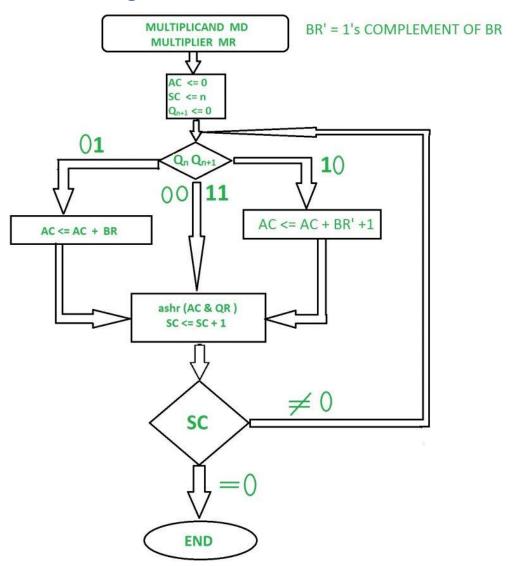
Booth's algorithm can be implemented by repeatedly adding (with ordinary unsigned binary addition) one of two predetermined values A and S to a product P, then performing a rightward arithmetic shift on P. Let m and r be the multiplicand and multiplier, respectively; and let x and y represent the number of bits in m and r.

- 1. Determine the values of A and S, and the initial value of P. All of these numbers should have a length equal to (x + y + 1).
  - a. A: Fill the most significant (leftmost) bits with the value of m. Fill the remaining (y + 1) bits with zeros.
  - b. S: Fill the most significant bits with the value of (-m) in two's complement notation. Fill the remaining (y + 1) bits with zeros.
  - c. P: Fill the most significant x bits with zeros. To the right of this, append the value of r. Fill the least significant (rightmost) bit with a zero.



- 2. Determine the two least significant (rightmost) bits of P.
  - a. If they are 01, find the value of P + A. Ignore any overflow.
  - b. If they are 10, find the value of P + S. Ignore any overflow.
  - c. If they are 00, do nothing. Use P directly in the next step.
  - d. If they are 11, do nothing. Use P directly in the next step.
- 3. Arithmetically shift the value obtained in the 2nd step by a single place to the right. Let P now equal this new value.
- 4. Repeat steps 2 and 3 until they have been done y times.
- 5. Drop the least significant (rightmost) bit from P. This is the product of m and r.

#### Flowchart Diagram



# Chetan Reddy



# Example

Let **A**: **3** and **B**: **17** 

Multiplicand -		
Decimal:	3	
Binary:	00000011	
Multiplier -		
Decimal:	17	
Binary:	00010001	
Two's Complement:	11101111	
Steps -		
Starting Out:	00000000000011	
Subtract:	1110111100000011	
Shift:	1111011110000001	
Shift:	11111011111000000	
Add:	0000110011000000	
Shift:	0000011001100000	
Shift:	0000001100110000	

Implementation of a Booth's algorithm uses various sub-modules, as described below.

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Final Product (Decimal):

#### Worst and Ideal Case

The worst case of an implementation using Booth's algorithm is when pairs of 01s or 10s occur very frequently in the multiplier.

#### Modules and Sub Modules

Implementation of a Booth's algorithm uses various sub-modules, as described below.

1. **boothmul():** This module is the main module which uses the help of other sub-modules or counterparts to solve our problem.

This module takes in two **8-bit signed** inputs, which are our multiplicand and multiplier. It has one **16-bit signed** output. Inside the module, we have **eight** 8-bit signed wires hold the value of the changed bits after shifting so that we can manipulate them later.

**2. booth\_substep()**: This sub-module does the main operation of either adding/subtracting or just shifting the bits according to the last two positions.

This module takes in an 8-bit signed **accumulator**, 8-bit signed **multiplier**, the last bit of the accumulator, 8-bit signed **multiplicand**, and the output consists of two 8-bit signed registers



containing first 8 and last 8 bits of the product, and cq0 is the changed q0 after the shift operation.

- **3.** Adder(): This sub-module adds two 8-bit register values, and gives out their sum. This uses a library module of **fa** which is nothing but a simple full adder.
- **4. Subtractor()**: This sub-module subtracts two 8-bit register values, and gives out their difference. This uses a library module of **invert** to invert each bit separately, and then uses **fa** which is nothing but a simple full adder as described above.
- 5. Lib.v:
  - a. invert(output ib,input b);
  - b. and2 (input wire i0, i1, output wire o);
  - c. or2 (input wire i0, i1, output wire o);
  - d. xor2 (input wire i0, i1, output wire o);
  - e. nand2 (input wire i0, i1, output wire o);
  - f. nor2 (input wire i0, i1, output wire o);
  - g. xnor2 (input wire i0, i1, output wire o);
  - h. and3 (input wire i0, i1, i2, output wire o);
  - i. or3 (input wire i0, i1, i2, output wire o);
  - j. nor3 (input wire i0, i1, i2, output wire o);
  - k. nand3 (input wire i0, i1, i2, output wire o);
  - I. xor3 (input wire i0, i1, i2, output wire o);
  - m. xnor3 (input wire i0, i1, i2, output wire o);
  - n. fa (input wire i0, i1, cin, output wire sum, cout);

Apart from using the above modules, we use a testbench to supply the initial values.

```
a = 8'b11110000;
b = 8'b11110000;
#10
a = 8'b10010101;
b = 8'b100000;
```

And so on and so forth.

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# Final Result on Screen

The results on the screen are printed like this:

VCD info: dumpfile tb\_boothsalgo.vcd opened for output.

-520

0	-16	Χ	-16	=	256
10	-107	X	32	=	-3424
20	7	X	0	=	0
30	1	Χ	1	=	1
40	60	X	5	=	300
50	-86	Χ	35	=	-3010
60	17	Х	28	=	476

-65

Χ