

TV CHANNEL SCHEDULER

P9

GitHub Repo Link: <https://github.com/ChhabhayaManan/CipherX>

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QUESTION DESCRIPTION

P9

We have given a text file containing M Distinct Series, and N family members, each having a favorite series and Availability slots. We have to schedule their favorite series into their availability slots such as The Scheduler needs to minimize total episode misses. There can be the same favorite.

PSEUDO CODE

Assumptions:

We have assumed that there are 168 slots each hour, and 24 slots each day. Our scheduler has already recorded all episodes of the series. We have assumed that here the number of favorite series which has been recorded cannot exceed 168 (total slots) because the recorder can record one series at a time.

Algorithm Functions:

Function storeSlotVsSeries(memTotal, slotTotal, seriesTotal):

```
slotVsSeries.resize(slotTotal)
```

```
For (i = 0 To slotTotal) then
```

```
    SlotVsSeries[i].resize(seriesTotal+1){0}
```

```
endFor
```

```
For i = 0 To slotTotal
```

```
    For j = 0 To seriesTotal
```

```
        For k = 0 To memTotal
```



```
    If (memVsSlot[k][i] == True && memVsSeries[k][j] == True ) then
        slotVsSeries[i][j]++
    Endif
    endFor
    slotVsSeries[i][seriesTotal] += slotVsSeries[i][j]
endFor
endFor
```

Function findMax(memTotal, slotTotal, seriesTotal):

Initialize mAX and cURR as empty vectors

While (mAX[2] != 0) do

Clear mAX and set mAX[0], mAX[1], mAX[2] to 0

For (l = 0 To 167) do

For (j = 0 To 167) do

If (slotVsSeries[i][j] > mAX[2]) then

mAX.clear()

mAX.push_back(i)

mAX.push_back(j)

mAX.push_back(slotVsSeries[i][j])

For (k 0 To memTotal)

If (memVsSlot[k][i] == true && memVsSeries[k][j] == true) then

mAX.push_back(k)

Endif

endfor

Endif

Else if (slotVsSeries[i][j] == mAX[2]) then

cURR.clear();

for(k = 0 To memTotal){

if(memVsSlot[k][i] == true && memVsSeries[k][j] == true) then

cURR.push_back(k);

endif

endfor

integer count1=0,count2=0



```

For ( k = 0 To slotTotal)
    For each member l in cURR:
        If (memVsSlot[k][i] == True && memVsSeries[k][j] == True ) then
            count1++
        Endif
    Endfor
    For ( l = 3 to mAX.size()) do
        If (memVsSlot[mAX[l]][k] == true && memVsSeries[mAX[l]][j] == true) then
            count2++
        Endif
    Endfor
Endfor
If count1 < count2 then
    mAX.clear()
    mAX.push_back(i)
    mAX.push_back(j)
    mAX.push_back(slotVsSeries[i][j])
    For (k from 0 to memTotal)
        If (memVsSlot[k][i] == True && memVsSeries[k][j] == True) then
            mAX.push_back(k)
        Endif
    Endfor
Endif
If (mAX[2] == 0) then
    Return
endif
pq.push(make_pair(mAX[0], mAX[1])) // pq is priority queue that stores pair slotwise.
For (l = 3 to mAX.size())
    For (k = 0 to slotTotal)
        If (memVsSlot[mAX[l]][k] == true && memVsSeries[mAX[l]][mAX[1]] == true) then
            slotVsSeries[k][mAX[1]] - -
            slotVsSeries[k][seriesTotal]- -
        Endif
    Endfor

```

```

        endfor
        memVsSeries[l][mAX[1]] = false
    endfor

    For (k = 0 to memTotal)
        memVsSlot[k][mAX[0]] = false
    endfor
    For (k = 0 to seriesTotal+1)
        slotVsSeries[mAX[0]][k] = 0
    Endfor
Endwhile

```

TIME COMPLEXITY

N = number of member

M = number of series

Time complexity of each functions:

Input() : $O(N*M*168)$

storeSlotVsSeries() : $O(N*M*168)$

findMax() : $O(168*168*M*(N+(168*N)+N))$

{As I have mentioned in assumptions, The number of series $M \leq 168$ because we can't show more than 168 series in 168 slots}

Overall Time complexity: $O(N)$ with assumption

$O(N*M)$ if M is variable (can exceed 168)

SPACE COMPLEXITY

Space complexity of arrays created into the code:

memVsSlot : $O(N*168)$

memVsSeries : $O(M*N)$

slotVsSeries: $O(168*M)$

as per our assumption

Overall Space Complexity: $O(N)$ $\{M \leq 168\}$

$O(M*N)$ if M is variable and exceeds 168

EXPLANATION OF ALGORITHM

| Series No → | 0 | 1 | 2 | 3 | 4 | Sum |
|-------------|---|---|---|----|---|-----|
| slot No ↓ | | | | | | |
| 0 | 0 | 1 | 2 | 5 | 1 | 9 |
| 1 | 2 | 3 | 5 | 16 | 6 | 17 |
| 2 | 4 | 2 | 1 | 3 | 4 | 14 |
| 3 | 2 | 1 | 2 | 3 | 1 | 9 |
| 4 | 4 | 3 | 5 | 6 | 0 | 18 |

- SlotvsSeries array contains the int number of members who are available at that slot number(Row) and also want to watch that Series(Column Number).
- The value in the last column is the sum of all values in that row.
- Number represents that if we telecast that series on that slot then that number of members is reduced from that series.
- Function **findMax()** is used to find the maximum number.
- Now we need to traverse the 2d array to find the maximum number.

Case 1-

For different numbers, we have to store the maximum number, row, column and the members.

Case 2 - (to prioritize one maximum number from two same numbers)

| Seriesno → | | K th | L th | sum |
|-------------|---|--|--|-----|
| slot no { 0 | - | 2 (M ₂ , M ₃) | 2 (M ₆ , M ₁) | - |
| 1 | - | 1 (M ₅) | 2 (M ₉ , M ₁₀) | - |
| 2 | - | 4 (M ₃ , M ₆ , M ₇ , M ₉) | 4 (M ₂ , M ₆ , M ₇ , M ₈) | X |
| 3 | - | 2 (M ₁ , M ₇) | 2 (M ₂ , M ₅) | - |
| 4 | - | 2 (M ₄ , M ₉) | 2 (M ₈ , M ₁) | - |

For the Kth series if we telecast at slot 2 then every member at that place(M3, M6, M7, M9) should be removed from that column after the telecast. (Ex. From slot 0 M3, from slot 1 M5 etc.)

Count 1 = number of members removed from Kth column.

Likewise for Lth series,

Count 2 = numbers of members removed from Lth column.

Count represents the number of ways in which the members of that place can watch that series in all slots

Reason: if we telecast series that have a higher count number then all members in that place still have more available slots so members can watch that series in other slots also. So we should prefer to telecast that series which has a lesser count(they don't have more opportunity to watch that series in other slots).

THE RATIONALE BEHIND DATA STRUCTURE

Initially, we thought of storing the input data as a 3D array (the array made using a linked list), but then we realized that it would be much better and efficient to use 2 2D-arrays to store the input data. The space complexity of a 3D array in P9 is $N \times M \times 168$. We can reduce it using 2 2D arrays of $N \times M$ and $N \times 168$.

Here we have used 1D and 2D vectors. Vector is a dynamic array. Unlike static arrays, a vector allows for efficient insertion and deletion of elements at the end or in the middle of the sequence. Vector uses array-like indexing which makes it easy to use. We can access any element of the vector in constant time (by using an index). This property makes vectors suitable for our capstone problem where fast random access to elements is required. In the findMax function initially, we don't know the number of members available for ith in slot and jth series. Hence, it is beneficial to use vectors for curr and max as it allows dynamic memory allocation.

Also, we have used a priority queue to store the pair(slot, series). The priority queue stores the pair in a slot-wise order. The time complexity for insertion in Priority Queue is $O(\log n)$ and for dequeue it is $O(1)$.

OUTPUT

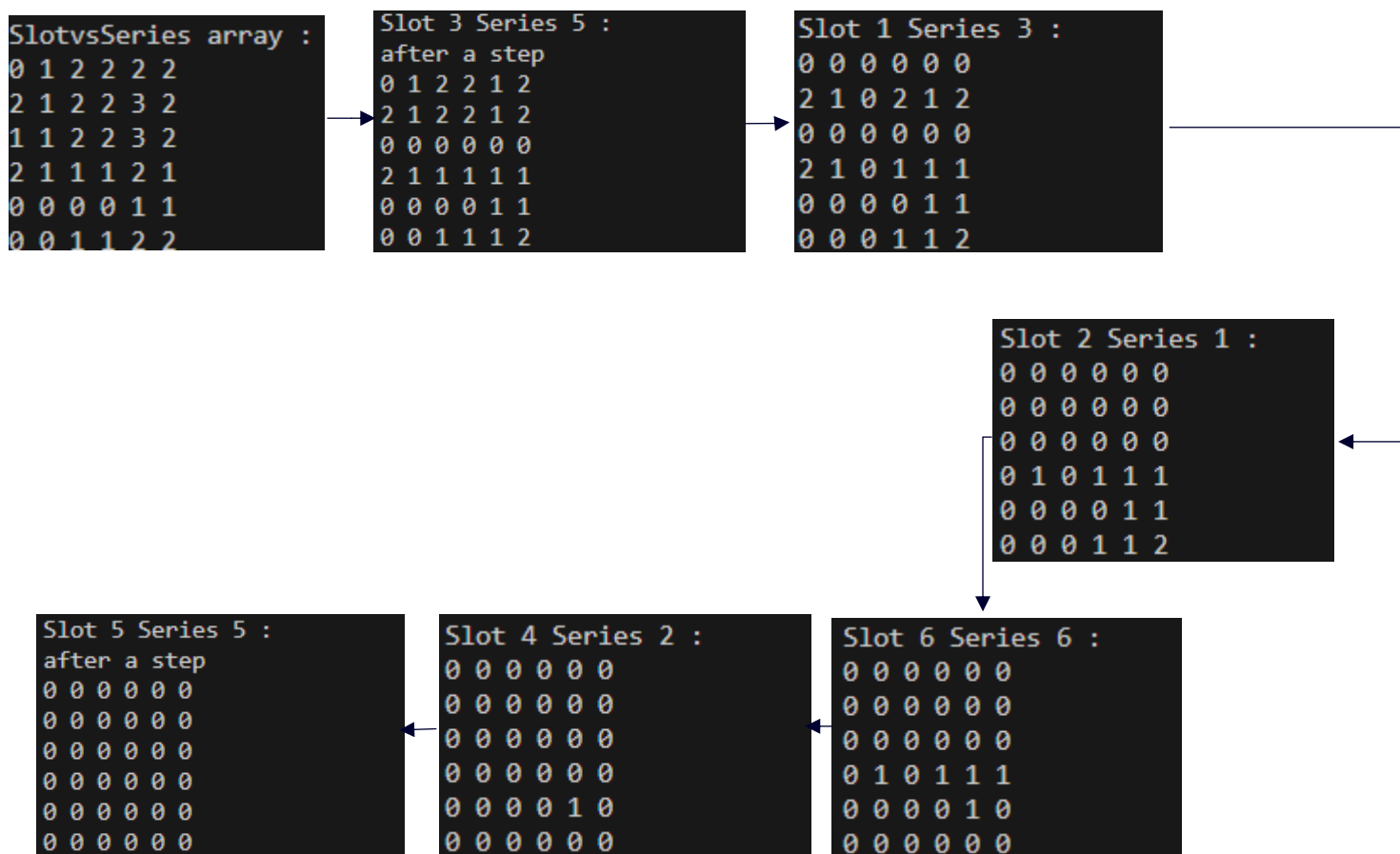
Terminal Based output:

Input file --

```
in copy.txt
1 6 6
2 TMKOC, Anupama, FRIENDS, ChotaBheem, Vir, SharkTank
3
4 Manan
5 MON 01-02, MON 02-03, MON 03-04, MON 00-01
6 Anupama, FRIENDS, ChotaBheem
7 Naman
8 MON 01-02, MON 03-04
9 TMKOC, Vir
10 Keval
11 MON 00-01, MON 01-02, MON 02-03, MON 05-06
12 FRIENDS, ChotaBheem, Vir, SharkTank
13 Parthiv
14 MON 02-03
15 Vir
16 Prince
17 MON 01-02, MON 02-03, MON 03-04
18 TMKOC, Vir, SharkTank
19 Prem
20 MON 00-01, MON 04-05, MON 05-06
21 Vir, SharkTank
```

[illegible]

Process on SlotvsSeries array :



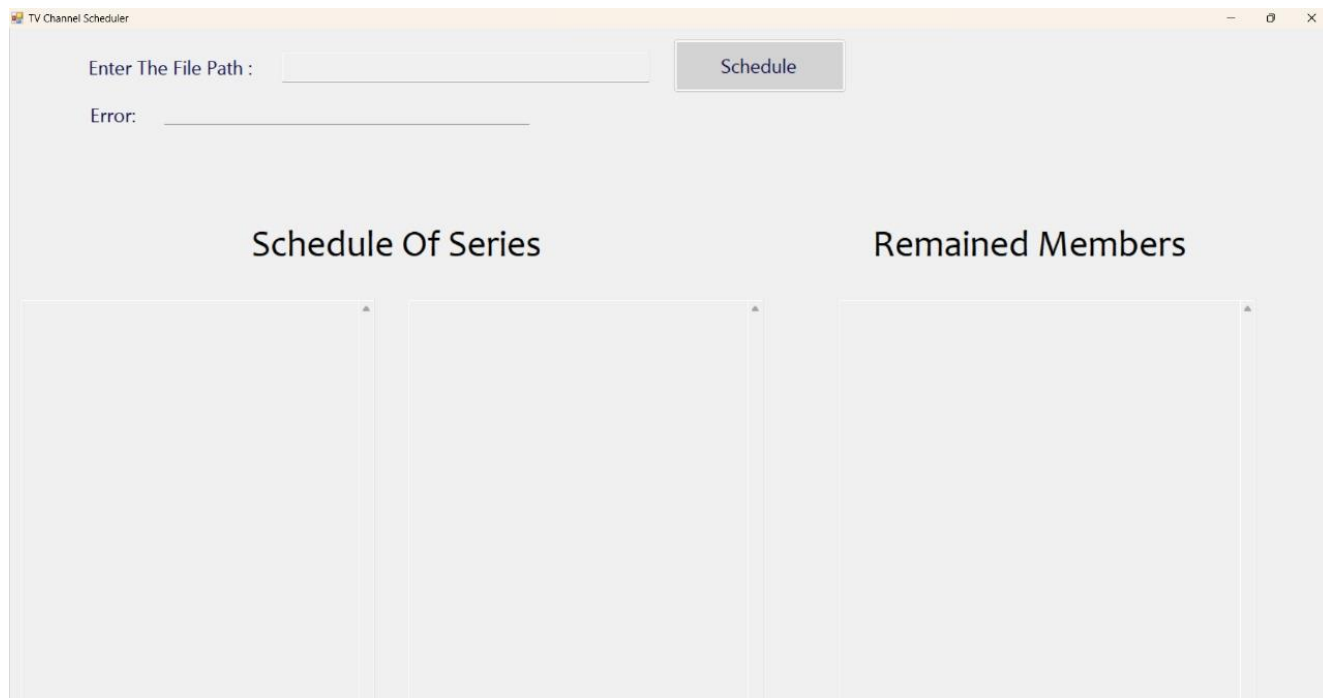
Main output :

```
=====
Remaining members:
=====
Manan      : ChotaBheem
Naman     : Vir
Keval     : ChotaBheem
Prince    : SharkTank
=====
The schedule is as follows:
=====
MON 0-1 : FRIENDS
MON 1-2 : TMKOC
MON 2-3 : Vir
MON 3-4 : Anupama
MON 4-5 : Vir
MON 5-6 : SharkTank
=====
```

It can run for any big file, but for simplicity, we used small text input.

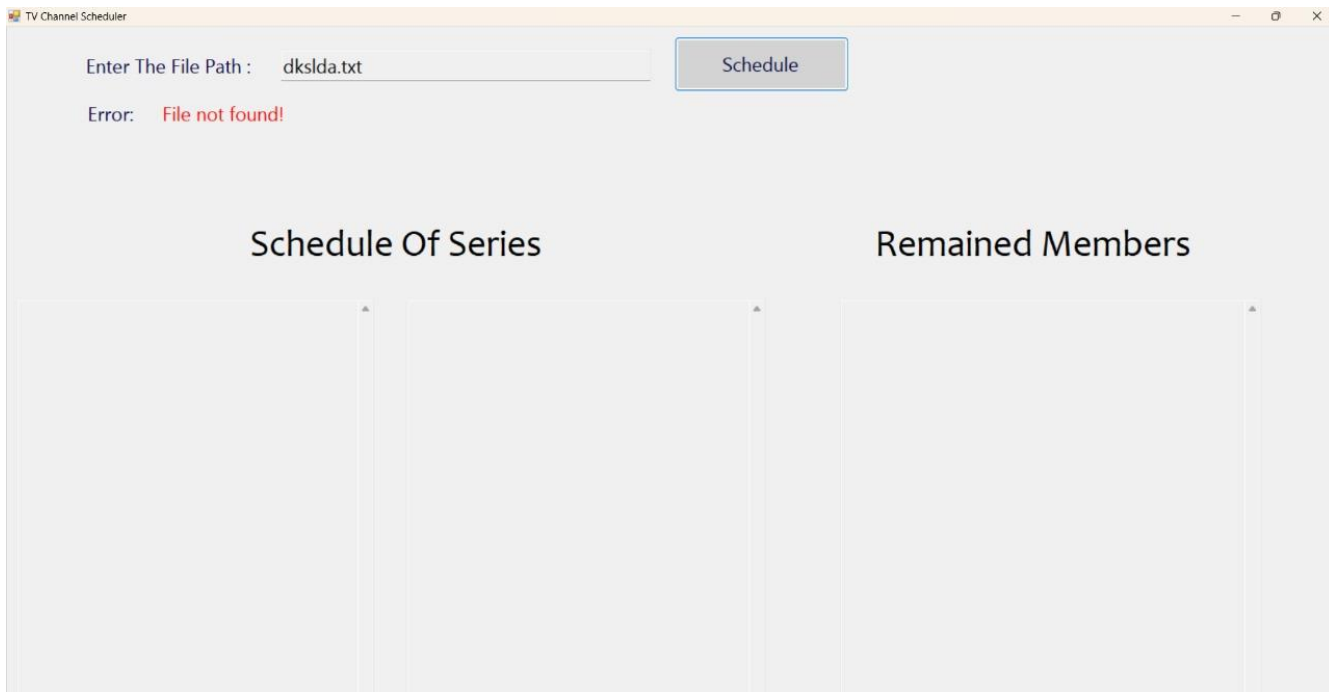
GUI based output:

Interface:



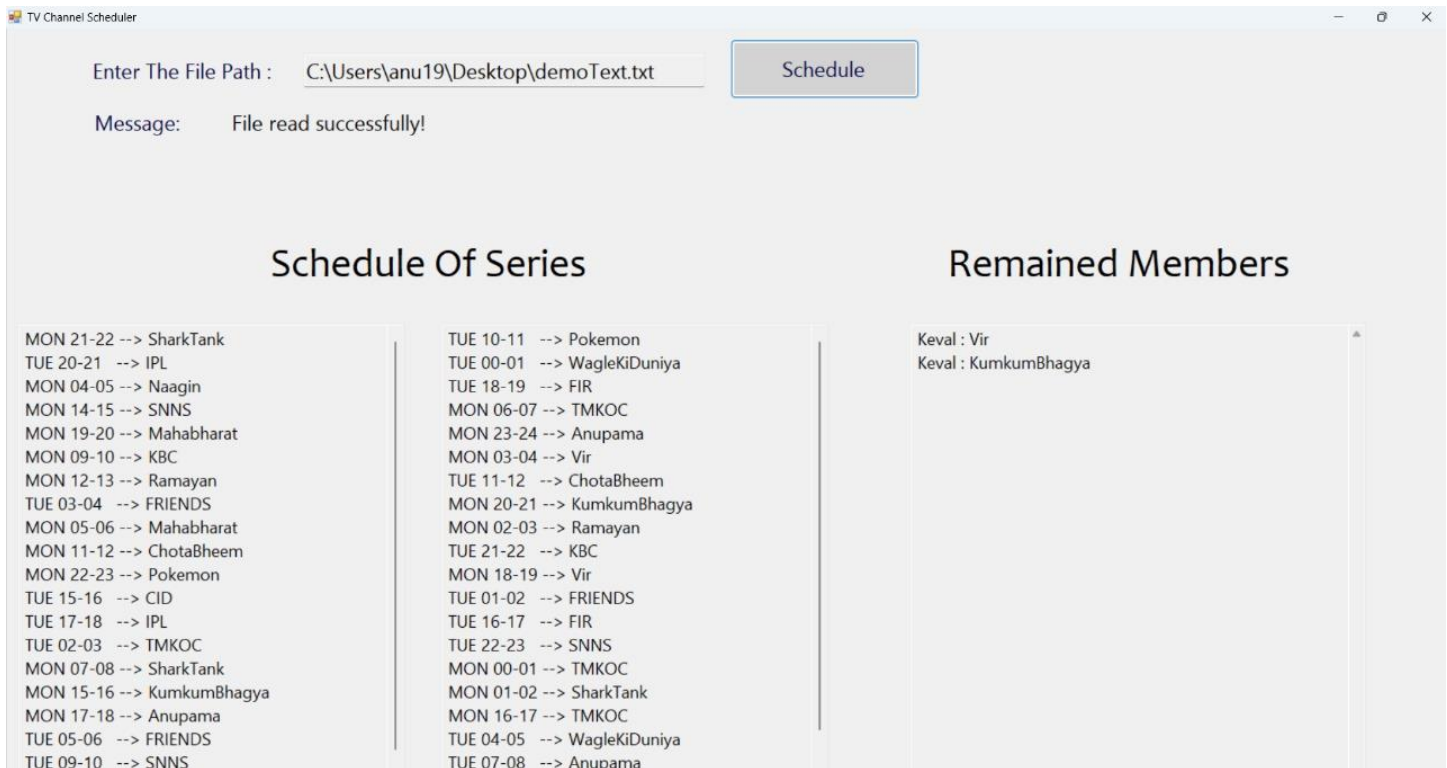
After entering a file path and push the schedule button.

If we enter wrong path it gives error –



If we enter our input file path.

DemoText.txt



Demo2.txt

TV Channel Scheduler

Enter The File Path :

Message: File read successfully!

| Schedule Of Series | | Remained Members |
|---|---|--------------------------------|
| MON 02-03 --> KBKP MON 03-04 --> GHKPM | MON 00-01 --> YRKKH MON 01-02 --> GHKPM MON 04-05 --> JAJ | kresha : ABCD kashyap : CID |

Demo3.cpp

TV Channel Scheduler

Enter The File Path :

Message: File read successfully!

| Schedule Of Series | | Remained Members |
|---|---|---|
| MON 00-01 --> FRIENDS MON 05-06 --> SharkTank MON 04-05 --> Vir | MON 02-03 --> Vir MON 01-02 --> TMKOC MON 03-04 --> Anupama | Manan : ChotaBheem Naman : Vir Keval : ChotaBheem Prince : SharkTank |

...