Operations Research

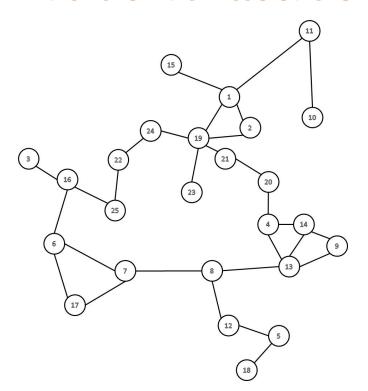
Final Presentation

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Agenda

- Research Topics
 - Traditional Allocation Problem
 - Allocation with connection : Articulation Point
- Operations Research (How we do it)
 - Data Input & Result Overview
 - Tarjan's Algorithm and Dilemmas
 - Improving the Model : Optimization Limit
 - Steps
 - Find Articulation Points
 - Iteration & Solution
- Takeaway: Future Thoughts to Generalize Model

Traditional Allocation Problem



Consider a City Network:

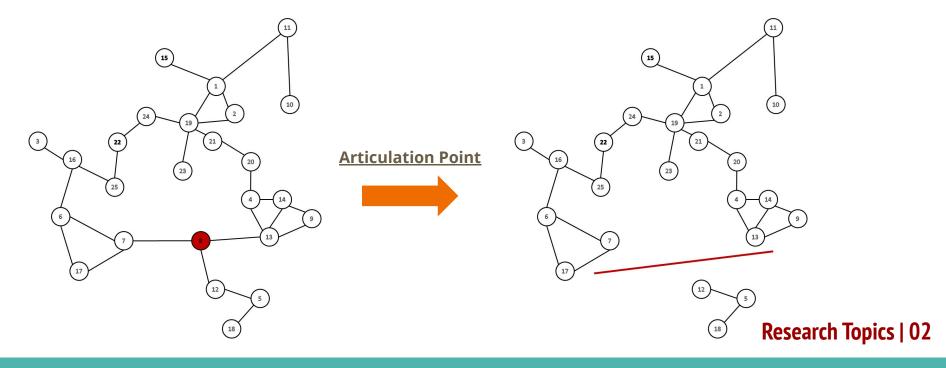
How to build a logistic center plan?

```
\max \sum_{j \in N} expected \ profit \ - costs \ under \ the \ plan s.t. Center Number \leq \delta Longest Distance \leq \beta; \forall \ every \ center : : : Other Constraints ...
```



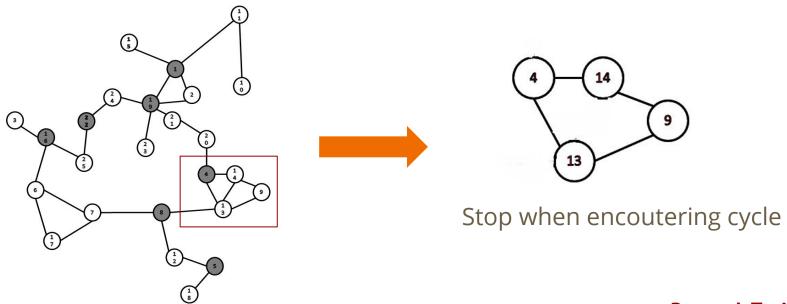
Connections are first to consider

- Circles represent cities.
- Lines represent roads connecting cities.



However we can't take all cities once because...

• <u>Capabilities</u>: We cannot build/explode in as many cities as we want



Research Goal

- Finding the articulation points, each with its own benefits defined by:
 - How many pieces it breaks the original graph?
 - How many total distance it can create?
- This is a graph theory (a branch of operations research) with iterative searching solutions (DFS or BFS) since whether the point is an articulation point will only be given at that state, we must iterate to find the optimization solution.

Related fields [edit]

Some of the fields that have considerable overlap with Operations Research and Management Science include[32]:

- Business analytics
- Data mining/Data science/Big data
- Decision analysis
- Decision intelligence
- Engineering
- Financial engineering

- Forecasting
- Game theory
- Geography/Geographic information science

Graph theory

- Industrial engineering
- Logistics

- Mathematical modeling
- · Mathematical optimization
- Probability and statistics
- Project management
- Policy analysis
- Simulation

Source: wikipedia

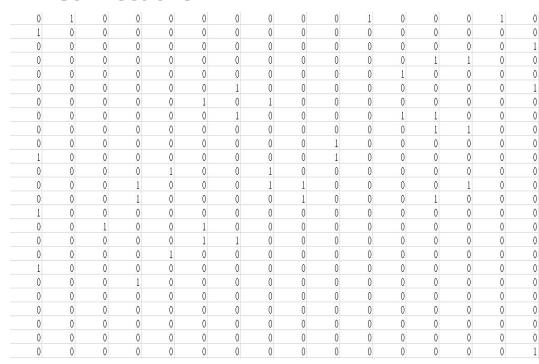
Data Overview

- Total number of cities : 25
- Adjacency Matrix : <u>25*25</u>
- Distance Matrix: 25*25
- Cost for each explode : λ
- Maximum Explosion: 10

Objective Function:

Pieces Number * Total Distance

Connections



Mathematics Formulation

Objective Function and Constraints

```
max sumt(Ptj*Dtj) t \in T, j \in length(At)

s.t. At = B(S(At-1)); \forall t = 2,...,T

A1 = B(G)

0 < t \le T

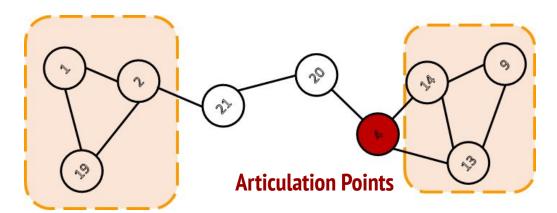
j \in At
```

- t = current level
- At = articulation set at t
- $j \in \{1,...,25\}$, index
- S(A) = a function that generates subplots from the given articulation points
- B(Fragment) = a function that finds all articulation points
- Ptj = number of pieces of articulation j at level t
- Dtj = total distance an articulation j at level t
- G = original graph
- T = maximum iterations

Research Topics | 07

Articulation Points & Biconnected Components

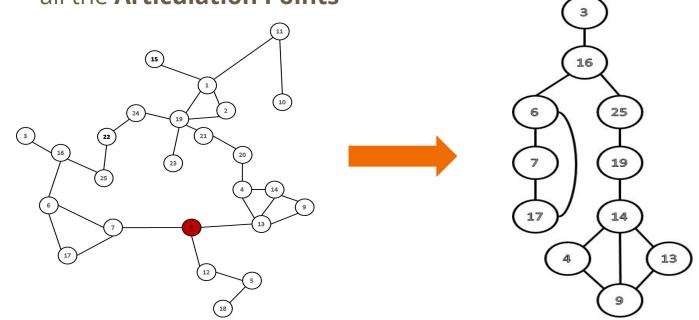
- Articulation Point disconnects an undirected graph
- Biconnected Component does not generate articulation points



Biconnected Component

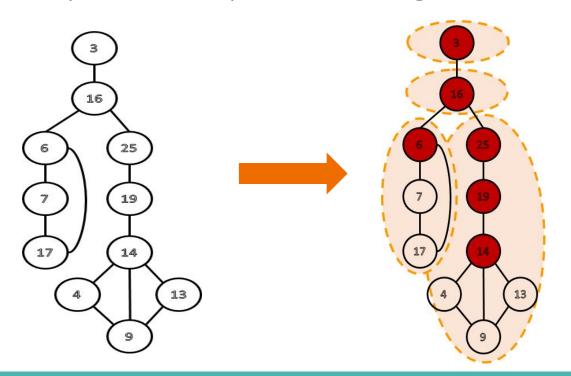
How to find articulation points: Tarjan's Algorithm

The algorithm is to find all **Biconnected Components** and thus we have
 all the **Articulation Points**



How to find articulation points: Tarjan's Algorithm

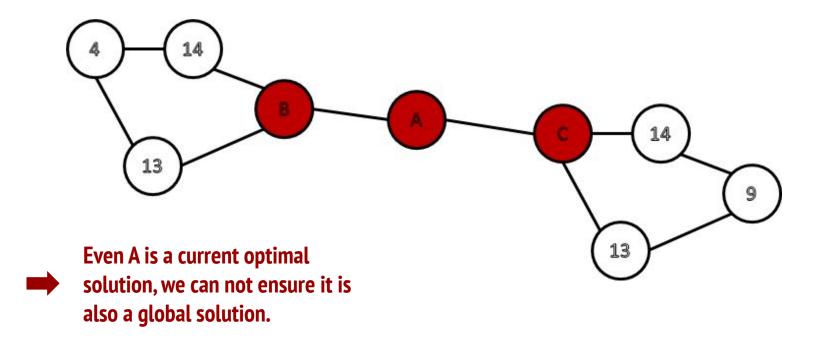
• Step 2 : DFS each point and find highest ancestor for each point



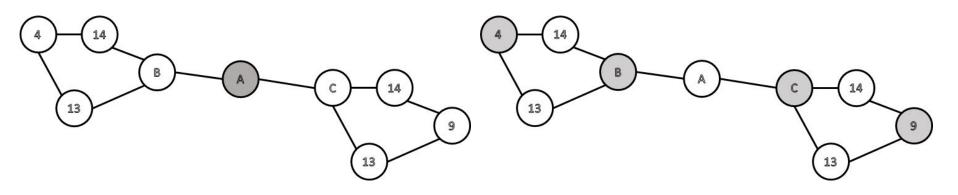
If the ancestor of a point is itself, it's a biconnected component

The overlapping part of each BCC is an articulation point

Algorithm Dilemma: Not optimization solution



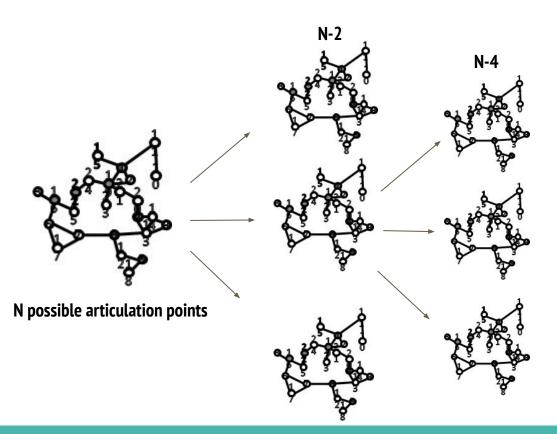
Algorithm Dilemma: Not optimization solution



No one knows which division is better considered pieces and distance



Another Dilemma: Is it solvable?



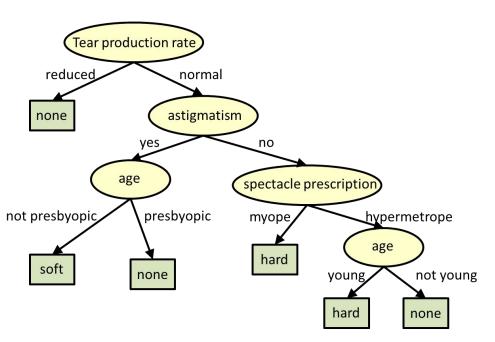
Time Complexity?

$$2N * (2n-2) * \cdots 2 = O(2^n n!)$$



... TSP only O(n!)

Improving Model: Tree-model Adjustment



- <u>Max Depth</u>: Stop if tree depth larger than this number (we can launch infinite missiles)
- Min Leaf Split: Do not accept the solution if profit less than this number (cost of each launch)
- <u>Tolerance</u>: If profit increase less than this number, stop generating new solution (cost of each launch)

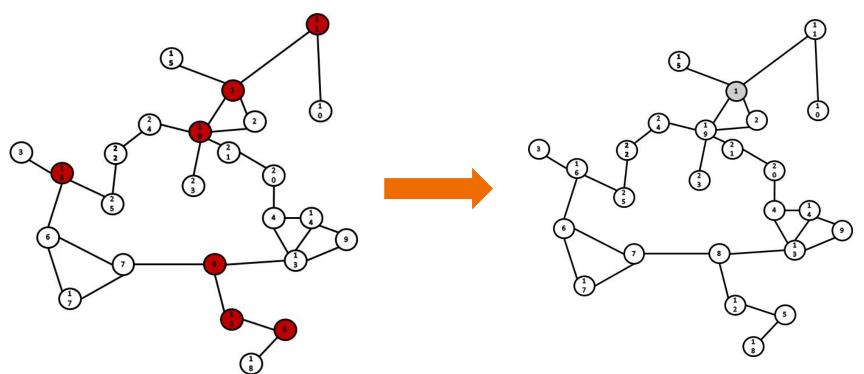
<u>Solvable when data is small but not</u> <u>satisfying</u>

Solving This is Like a Decision Tree

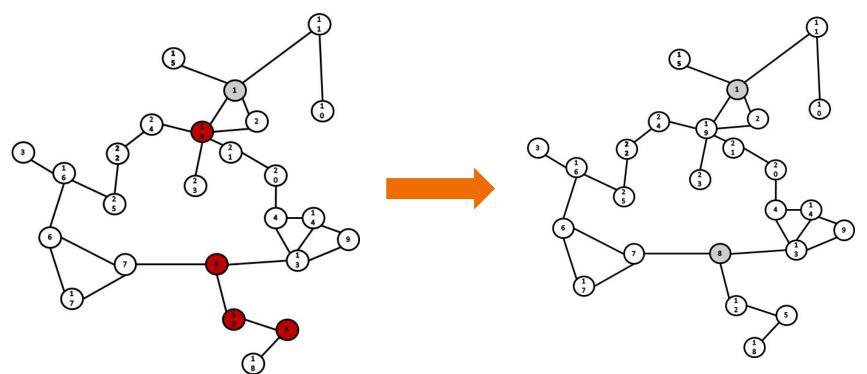
Steps: Virtual Code

```
void CityPlanning (adj matrix : graph, distance) {
     /* Articulation Set = [Articulation Point, Sub_graph, Layer] */
     Articulation_Set = Get_articulation_points(graph);
     while Articulation_Set not empty { /* BFS Find Solution */
          sub_graph = Articulation Set.top()[1];
          solution_set = Get_articulation_points(sub_graph);
          Articulation_Set.pop(); Articulation Set.push(solution set)
```

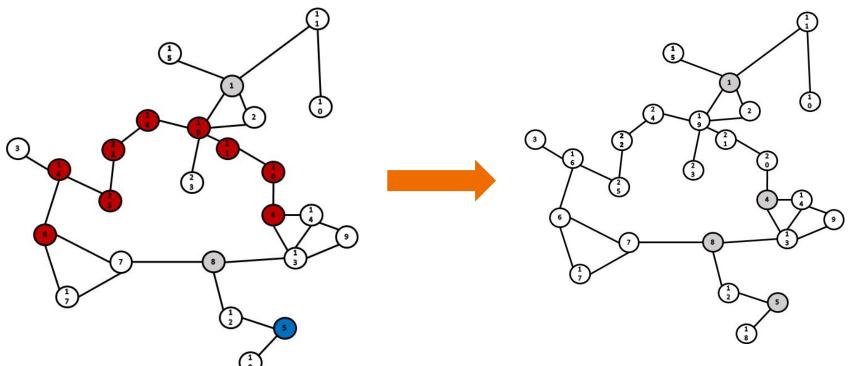
Steps: Visualization



Steps: Visualization

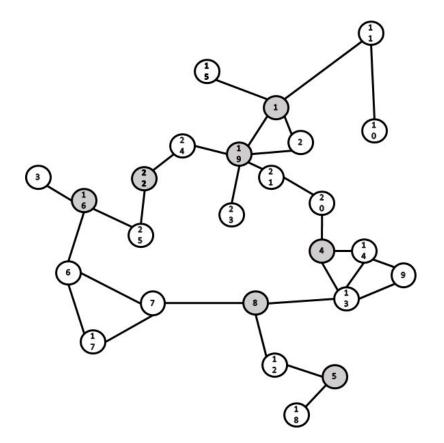


Steps: Visualization



Solution Interpretation

- Order: [1, 8, 19, 5, 4, 14, 20]
- Total Profit: 32714
- Interpretation: We follow the order to explode each city, and we do not run out all of our missiles because launching eighth missiles do not benefit in our planning.

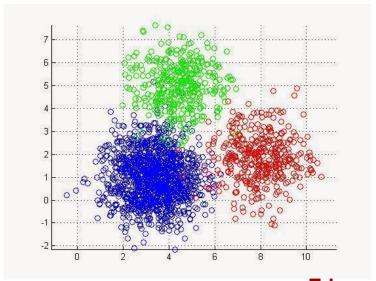


Takeaway: Future Improvement

- The model is not generalized enough, especially for big data
- Idea of the planning is like K-Means, which have :

$$l(G_1, ..., G_n) = \sum_{j=1}^k \sum_{i \in G_j} \|x_i - \mu_j\|_2^2$$
$$\bigcup_{j=1}^k G_j = \{1, 2, ..., n\}; G_j \cap G_j' = \phi$$

$$\mu_j = \frac{1}{|G_j|} \sum_{i \in G_j} x_i$$



Takeaway: Apply Community Detection

A Community is a subgraph that contains nodes more <u>densely linked</u>

$$Q = \sum_{c \in \mathcal{C}} \left(\frac{I_c}{m} - \left(\frac{D_c}{2m} \right)^2 \right)$$

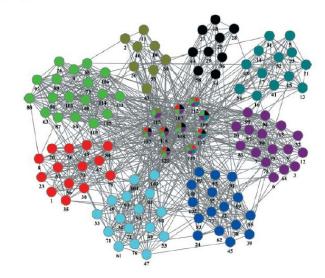
m = number of edges

 I_c = number of edges in community c

 D_c = sum of degree of all vertex in community

By BMLPA Algorithm, it could be done in O(nlogn)

overlapping communities



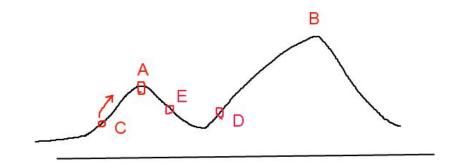
Takeaway: Simulated Annealing (Other Searching)

• SA algorithm can somewhat improve the TSP dilemma.

```
Accept new state by a specific probability (not definitely optimized solution)
```

```
if \Delta E > 0:
Accept new state else:
```

Accept by probability $e^{\frac{\Delta E}{T}}$



Thank You!

Final Presentation

Members: 周永堂、鍾孟芳、林家毅、童安弘