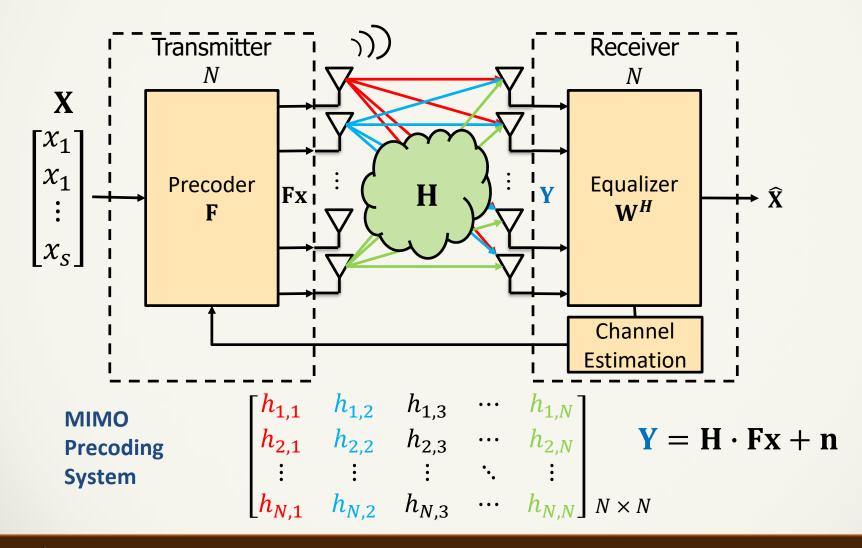
Design and Implementation of Matrix Decomposition

Professor: Pei-Yun Tsai

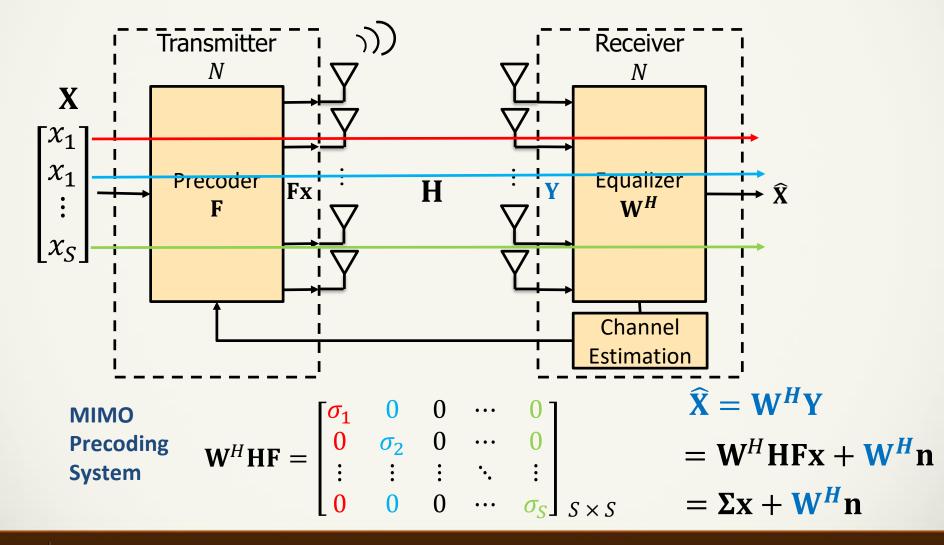
Department of Electrical Engineering,

National Central University

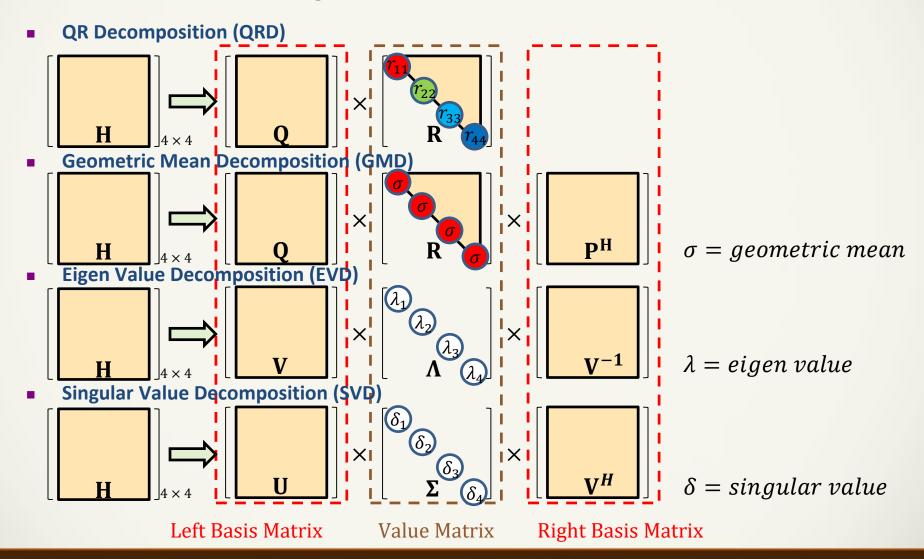
Applications in Communication Systems



Channel Decomposition



Matrix Decomposition

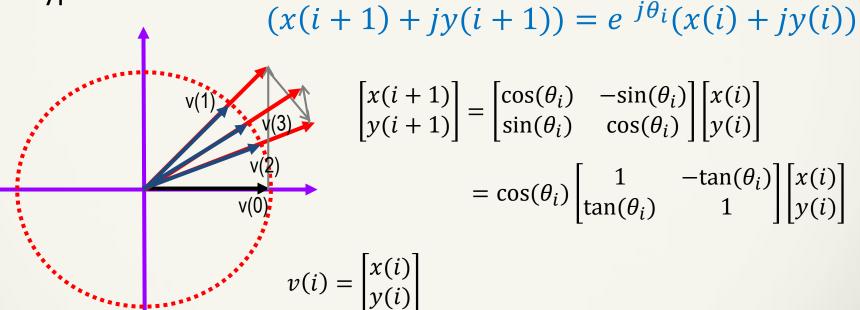


Methods

- Coordinate rotation digital computer (CORDIC)
- Matrix decomposition
 - QR decomposition
 - Eigenvalue decomposition
 - ■Singular value decomposition

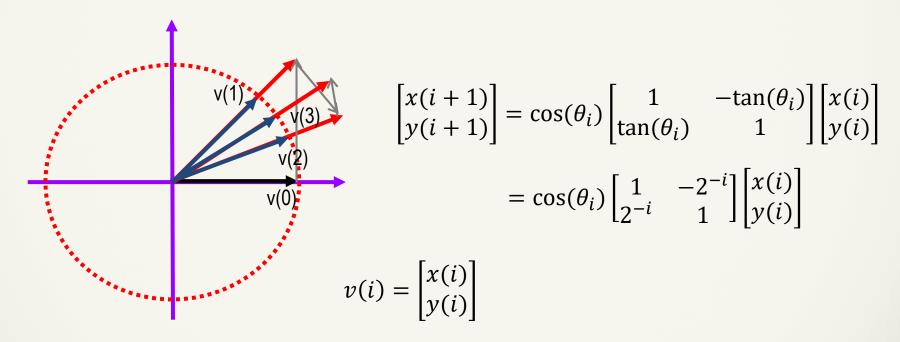
CORDIC (1/2) (Rotation)

- CORDIC (coordinate rotation digital computer)
 - Linear Mode
 - ■Circular Mode
 - Hyperbolic Mode



CORDIC (2/2) (Rotation)

- ■In circular mode,
 - ■Angle of the *i*th micro rotation is defined as $\theta_i = tan^{-1}(\frac{1}{2^i})$



CORDIC Arctangent Function (1/2) (Vectoring)

Initialization

$$x(0) = x, y(0) = y, \theta'(0) = 0$$

Direction

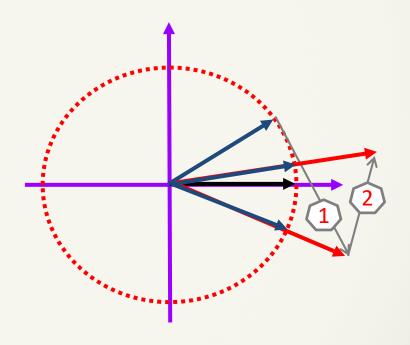
$$\mu_i = -sign(y(i))$$

Micro rotation

$$\begin{bmatrix} x(i+1) \\ y(i+1) \end{bmatrix} = \begin{bmatrix} 1 & -\mu_i 2^{-i} \\ \mu_i 2^{-i} & 1 \end{bmatrix} \begin{bmatrix} x(i) \\ y(i) \end{bmatrix}$$

Angle accumulation

$$\theta'(i+1) = \theta'(i) - \mu_i \theta_i$$



CORDIC Arctangent Function (2/2) (Vectoring)

■To restore the original vector length, we need scaling factor

$$S(\infty) = \prod_{i=0}^{\infty} \sqrt{(1+2^{-2i})} = 1.6468$$

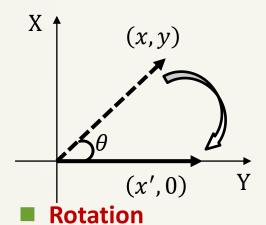
■The converge range of CORDIC algorithm is

$$\theta_{MAX} = \sum_{i=0}^{\infty} tan^{-1} (2^{-i}) = 1.7433 \dots (\sim 99^{\circ})$$

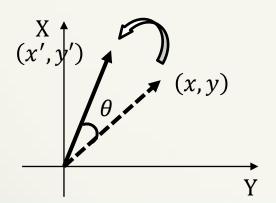
Hence, initial mapping of the vectors in the second and the third quadrant is required.

Givens Rotation for Real Numbers

- Givens rotation for two real numbers
 - Vectoring (Nullification, Real to Zero (R2Z))



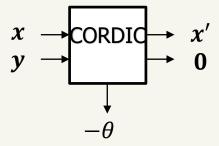
$$\begin{bmatrix} x' \\ 0 \end{bmatrix} = \begin{bmatrix} \cos\theta & \sin\theta \\ -\sin\theta & \cos\theta \end{bmatrix} \cdot \begin{bmatrix} x \\ y \end{bmatrix} \ \left(\theta = \tan^{-1}\frac{y}{x}\right)$$



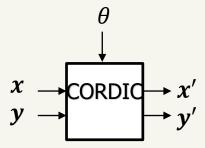
$$\begin{bmatrix} x' \\ y' \end{bmatrix} = \begin{bmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{bmatrix} \cdot \begin{bmatrix} x \\ y \end{bmatrix} \quad (\theta \text{ is given})$$

CORDIC for Real Numbers

Vectoring



Rotation



Givens Rotation for One Complex Number

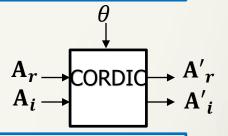
Givens rotation for one complex number

$$e^{-i\theta} \cdot (\mathbf{A}_r + j\mathbf{A}_i) = \begin{bmatrix} \cos\theta & \sin\theta \\ -\sin\theta & \cos\theta \end{bmatrix} \cdot \begin{bmatrix} \mathbf{A}_r \\ \mathbf{A}_i \end{bmatrix} \qquad \begin{matrix} \mathbf{A}_r \rightarrow \mathsf{CORDIC} \\ \mathbf{A}_i \rightarrow \mathbf{0} \end{matrix} \rightarrow \mathbf{0}$$

Vectoring (Nullification, Complex to Real (C2R))

$$\mathbf{A'}_{r} = e^{-i\theta} \cdot (\mathbf{A}_{r} + j\mathbf{A}_{i}) = \begin{bmatrix} \cos\theta & \sin\theta \\ -\sin\theta & \cos\theta \end{bmatrix} \cdot \begin{bmatrix} \mathbf{A}_{r} \\ \mathbf{A}_{i} \end{bmatrix} \quad \left(\theta = \tan^{-1}\frac{\mathbf{A}_{i}}{\mathbf{A}_{r}}\right)$$

Rotation of one complex number



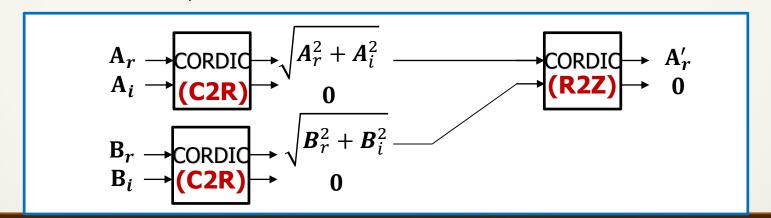
$$\mathbf{A'}_{r} + j\mathbf{A'}_{i} = e^{i\theta} \cdot (\mathbf{A}_{r} + j\mathbf{A}_{i}) = \begin{bmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{bmatrix} \cdot \begin{bmatrix} \mathbf{A}_{r} \\ \mathbf{A}_{i} \end{bmatrix} \quad (\theta \text{ is given})$$

Givens Rotation for Two Complex Numbers (1/2)

- Givens rotation for two complex numbers
 - Vectoring (Nullification, Complex to Real (C2R) and Real to Zero (R2Z))

$$\mathbf{Gu} = \begin{bmatrix} \cos\theta_z & \sin\theta_z \\ -\sin\theta_z & \cos\theta_z \end{bmatrix} \begin{bmatrix} e^{-j\theta_x} & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & e^{-j\theta_y} \end{bmatrix} \cdot \begin{bmatrix} \mathbf{A}_r + j\mathbf{A}_i \\ \mathbf{B}_r + j\mathbf{B}_i \end{bmatrix} = \begin{bmatrix} \mathbf{A}'_r \\ \mathbf{0} \end{bmatrix}$$

$$\boldsymbol{\theta}_z = \tan^{-1} \frac{\sqrt{\boldsymbol{B}_r^2 + \boldsymbol{B}_i^2}}{\sqrt{\boldsymbol{A}_r^2 + \boldsymbol{A}_i^2}} \qquad \boldsymbol{\theta}_x = \tan^{-1} \frac{\mathbf{A}_i}{\mathbf{A}_r} \qquad \boldsymbol{\theta}_y = \tan^{-1} \frac{\mathbf{B}_i}{\mathbf{B}_r}$$



Givens Rotation for Two Complex Numbers (2/2)

Rotation of two complex numbers

$$\begin{bmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{bmatrix} \cdot \begin{bmatrix} \mathbf{A_r} + j\mathbf{A_i} \\ \mathbf{B_r} + j\mathbf{B_i} \end{bmatrix} = \begin{bmatrix} (\mathbf{A_r}\cos\theta - \mathbf{B_r}\sin\theta) + j(\mathbf{A_i}\cos\theta - \mathbf{B_i}\sin\theta) \\ (\mathbf{A_r}\sin\theta + \mathbf{B_r}\cos\theta) + j(\mathbf{A_i}\sin\theta + \mathbf{B_i}\cos\theta) \end{bmatrix}$$

$$3$$

$$\begin{array}{c|c} \vdots & \begin{bmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{bmatrix} \cdot \begin{bmatrix} \mathbf{A}_r + j\mathbf{A}_i \\ \mathbf{B}_r + j\mathbf{B}_i \end{bmatrix} = \\ & \mathbf{B}_r \\ & \mathbf{B}_i \\ & \mathbf{A}_i \\ & \mathbf{B}_i \\ & \mathbf{A}_i \\ & \mathbf{A$$

QR Decomposition

QR Decomposition Algorithm

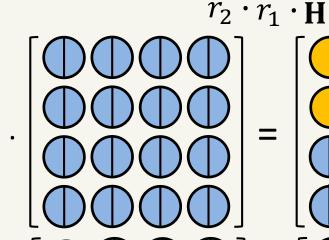
- Gram Schmidt algorithm
 - Modified Gram Schmidt algorithm
- Givens rotation
- Householder transformation
- With QR decomposition, $\mathbf{H} = \mathbf{Q}\mathbf{R}$
 - Q is a unitary matrix
 - R is an upper triangular matrix

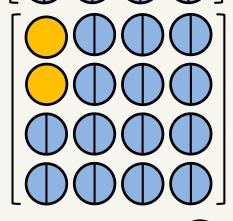
Givens Rotation for QR Decomposition (1/2)

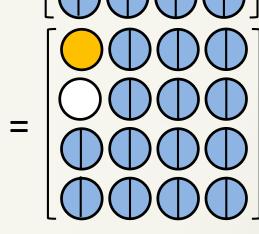
$$\left[egin{array}{ccccc} e^{j heta_1} & 0 & 0 & 0 \ 0 & e^{j heta_2} & 0 & 0 \ 0 & 0 & 1 & 0 \ 0 & 0 & 0 & 1 \end{array}
ight.$$

$$\begin{bmatrix} \cos \theta_3 & -\sin \theta_3 & 0 & 0 \\ \sin \theta_3 & \cos \theta_3 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$









real

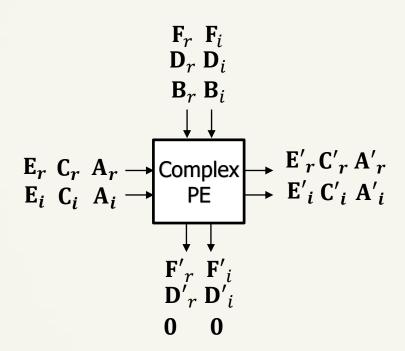


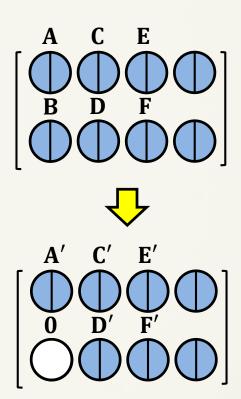
zero

Givens Rotation for QR Decomposition

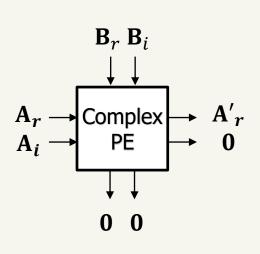
(2/2) $r_N \cdot \cdots r_5 \cdot r_4 \cdot r_3 \cdot r_2 \cdot r_1 \cdot \mathbf{H}$

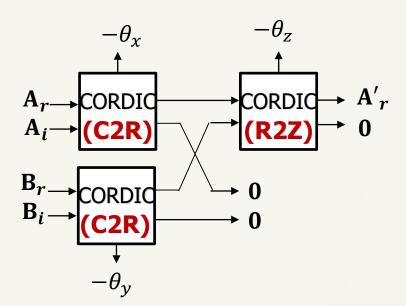
Architecture Design Complex Processing Element (Complex PE)





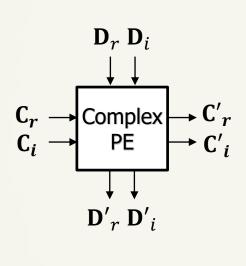
Complex PE Vectoring Mode

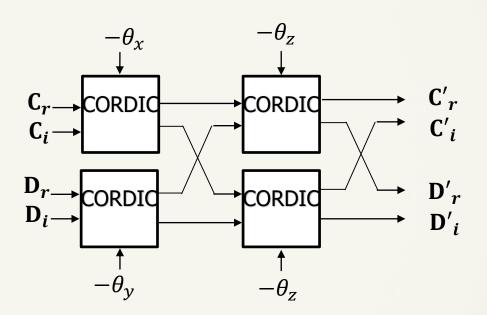




$$\begin{bmatrix} \cos\theta_z & \sin\theta_z \\ -\sin\theta_z & \cos\theta_z \end{bmatrix} \begin{bmatrix} e^{-j\theta_x} & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & e^{-j\theta_y} \end{bmatrix} \cdot \begin{bmatrix} \mathbf{A_r} + j\mathbf{A_i} \\ \mathbf{B_r} + j\mathbf{B_i} \end{bmatrix} = \begin{bmatrix} \mathbf{A'_r} \\ \mathbf{0} \end{bmatrix}$$

Complex PE Rotation Mode

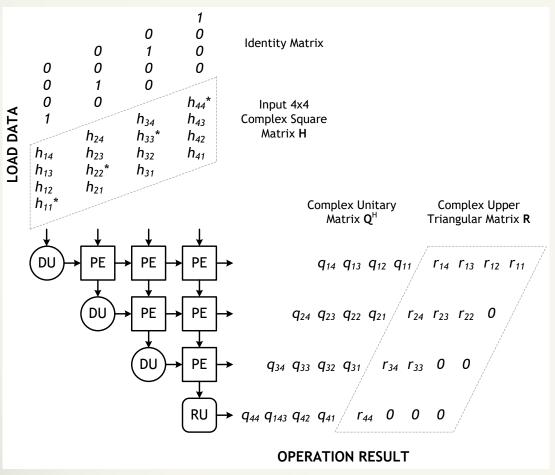




$$\begin{bmatrix} \cos\theta_z & \sin\theta_z \\ -\sin\theta_z & \cos\theta_z \end{bmatrix} \begin{bmatrix} e^{-j\theta_x} & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & e^{-j\theta_y} \end{bmatrix} \cdot \begin{bmatrix} \mathbf{C}_r + j\mathbf{C}_i \\ \mathbf{D}_r + j\mathbf{D}_i \end{bmatrix} = \begin{bmatrix} \mathbf{C}'_r + j\mathbf{C}'_i \\ \mathbf{D}'_r + j\mathbf{D}'_i \end{bmatrix}$$

[1]

Conventional Systolic Array for QR Decomposition (1/2)



- The input matrix is loaded in the skewed triangular shape.
- Input identity matrix, we can obtain the unitary matrix.
- DU (Delay Unit)



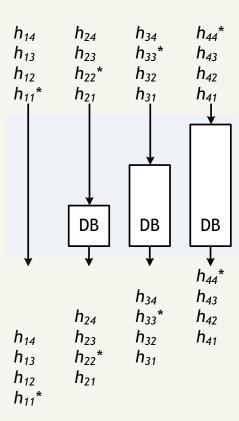
- Reads data, delays them, and then passes to the output.
- RU (Rotation Unit)



 Eliminate complex lowest diagonal element of matrix by additional rotation.

Conventional Systolic Array for QR Decomposition (2/2)

- In the conventional design, four input signal streams enter into the systolic array in skewed form. Hence, delay buffers are required.
- Delays data for period equal to PE operation time.
- If quantity of pipeline stage in CORDIC is large, a lot of delay buffers are introduced to hardware at input.



Eigen Value Decomposition

Eigen Value Decomposition

- For Hermitian symmetric matrix $\mathbf{A} = \mathbf{H}^{\mathbf{H}}\mathbf{H}$
 - $\blacksquare A = V\Lambda V^{-1}$

$$\Lambda = \begin{bmatrix} \lambda_1 & \mathbf{0} & \mathbf{0} & \mathbf{0} \\ \mathbf{0} & \lambda_2 & \mathbf{0} & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \ddots & \mathbf{0} \\ \mathbf{0} & \mathbf{0} & \mathbf{0} & \lambda_S \end{bmatrix} \qquad \mathbf{V} = \begin{bmatrix} \mathbf{v_1} & \mathbf{v_2} & \cdots & \mathbf{v_S} \end{bmatrix}$$

lacksquare λ_S is the eigenvalue and \mathbf{v}_S is the eigen vector

Eigen Value Decomposition Algorithm [3]

- Power iteration
- Inverse power iteration
- Iterative QR algorithm (with shift)

Power Method (1/2)

Algorithm: Power Method

Given Hermitian symmetric matrix $\mathbf{A} \in \mathbb{C}^{N \times N}$

1.
$$A(1) = A$$

2. For
$$s = 1:S$$

3.
$$q_s^{(0)} = [1 \ 0 \ ... \ 0]^T$$

4. For
$$i = 1: I_1$$

5.
$$\mathbf{z}_{s}^{(i)} = \mathbf{A}(s)\mathbf{q}_{s}^{(i-1)}$$

6.
$$\mathbf{q}_{s}^{(i)} = \mathbf{z}_{s}^{(i)} / \|\mathbf{z}_{s}^{(i)}\|_{2}$$

7. End

8.
$$\lambda_s = \mathbf{q}_s^{(I_1)H} \mathbf{A}(s) \mathbf{q}_s^{(I_1)}$$

9.
$$\mathbf{A}(s+1) = \left(\mathbf{I} - \mathbf{q}_s^{(I_1)} \mathbf{q}_s^{(I_1)H}\right) \mathbf{A}(s)$$

10. End

11.
$$\Lambda = [\mathbf{q}_1^{(l_1)} \dots \mathbf{q}_S^{(l_1)}],$$

Eigenvalue

Deflation

Power Method (2/2)

- Power method has linear convergence rate
 - **■**Slow
 - Especially when two eigenvalues are close

- Deflation is helpful to find the remaining eigenvectors
 - Eigenvector corresponding to the largest eigenvalue is obtained first.

Inverse Power Iteration (1/2)

Algorithm: Inverse Power Method

Given Hermitian symmetric matrix $\mathbf{A} \in \mathbb{C}^{N \times N}$

1.
$$A(1) = A$$

2. For
$$s = 1:S$$

3.
$$q_s^{(0)} = [1 \ 0 \ ... \ 0]^T$$

4. For
$$i = 1: I_1$$

5.
$$\mathbf{z}_{s}^{(i)} = (\mathbf{A}(s) - \lambda \mathbf{I})^{-1} \mathbf{q}_{s}^{(i-1)}$$

6.
$$\mathbf{q}_{S}^{(i)} = \mathbf{z}_{S}^{(i)} / \|\mathbf{z}_{S}^{(i)}\|_{2}$$

7. End

8.
$$\lambda_S = \mathbf{q}_S^{(l_1)H} \mathbf{A}(S) \mathbf{q}_S^{(l_1)}$$

9.
$$\mathbf{A}(s+1) = \left(\mathbf{I} - \mathbf{q}_s^{(I_1)} \mathbf{q}_s^{(I_1)H}\right) \mathbf{A}(s)$$

10. End

11.
$$\Lambda = [\mathbf{q}_1^{(l_1)} \quad ... \quad \mathbf{q}_S^{(l_1)}],$$

Shift λ

Eigenvalue

Deflation

Inverse Power Method (2/2)

■ Given that $(\mathbf{q}_k, \lambda_k)$ is eigen pair of matrix \mathbf{A} ,

$$\mathbf{q}_{S}^{(i)} = \sum_{k=1}^{K} a_{k}^{(i)} \mathbf{q}_{k} \qquad \mathbf{A} \mathbf{q}_{k} = \lambda_{k} \mathbf{q}_{k}$$
$$(\mathbf{A} - \lambda)^{-1} \mathbf{q}_{S}^{(i-1)} = \sum_{k=1}^{K} \frac{a_{k}^{(i-1)}}{\lambda_{k} - \lambda} \mathbf{q}_{k}$$

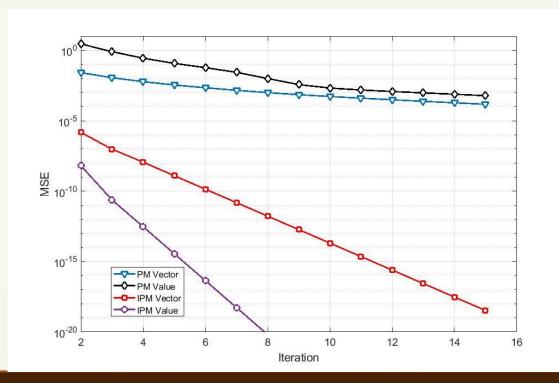
Thus, if λ is very close to an eigenvalue λ_s of A, the inverse iteration will make $\mathbf{q}_s^{(i)}$ close to the direction of \mathbf{q}_s .

Comparison

■ Use 4x4 complex channel matrix to generate Hermitian symmetric matrix A

■ Inverse power method uses true eigenvalue rounding to 10⁻¹ as

the shift



Iterative QR Algorithm

Algorithm: Iterative QR

Given Hermitian symmetric matrix $\mathbf{C} \in \mathbb{C}^{N \times N}$ // First phase

- 1. $[\mathbf{U}_{EVD}^{(0)}, \mathbf{A}^{(0)}]$ =HessenbergReduction(C) // Second phase i =0,
- 1. while (!converged)
- 2. $T^{(i)} = A^{(i)} \mu_i I$
- 3. $\left[\mathbf{Q}^{(i)}, \mathbf{R}^{(i)}\right] = QRD(\mathbf{T}^{(i)})$
- 4. $T^{(i+1)} = R^{(i)}Q^{(i)}$
- 5. $\mathbf{A}^{(i+1)} = \mathbf{T}^{(i+1)} + \mu_i \mathbf{I}$
- 6. $\mathbf{U}_{EVD}^{(i+1)} = \mathbf{U}_{EVD}^{(i)} \mathbf{Q}^{(i)}$
- 7. i = i + 1

End

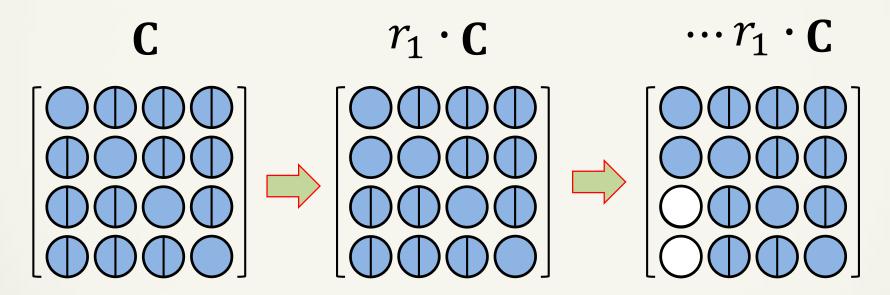
Phase 1: Triangularization

Phase 2: Iterative QR

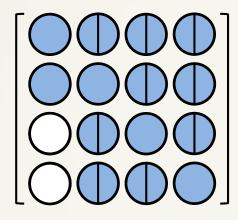
Phase 2: Shift to accelerate convergence

Hessenberg Reduction (1/3)

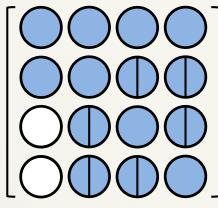
 Hessenberg reduction (tridiagonalization) helps to reduce computation complexity in the iterative procedure



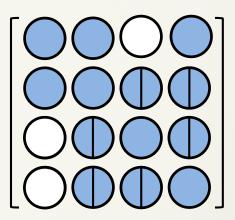
Hessenberg Reduction (2/3)



$$\cdot \left[\begin{array}{ccccc}
1 & 0 & 0 & 0 \\
0 & e^{-j\theta_1} & 0 & 0 \\
0 & 0 & e^{-j\theta_2} & 0 \\
0 & 0 & 0 & e^{-j\theta_3}
\end{array} \right]$$

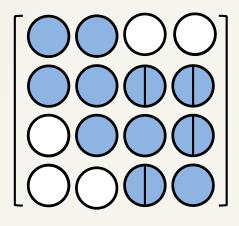


$$.\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & \cos\theta_4 & -\sin\theta_4 & 0 \\ 0 & \sin\theta_4 & \cos\theta_4 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

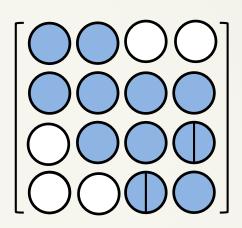


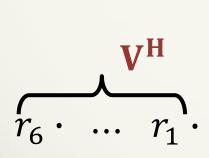
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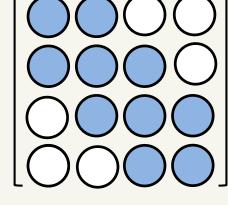
Hessenberg Reduction (3/3)



$$\begin{bmatrix} 1 & 0 & 0 & 0 \ 0 & 1 & 0 & 0 \ 0 & 0 & e^{-j heta_3} & 0 \ 0 & 0 & 0 & e^{-j heta_4} \end{bmatrix}$$







Eigen Value Decomposition Based on Iterative QR Algorithm

$$T^{(0)} = V \Lambda V^{-1}$$

$$\xrightarrow{QR} = R^{(0)} \Rightarrow R^{(0)} \cdot Q^{(0)} = T^{(1)} (= Q^{(0)H} \cdot T^{(0)} \cdot Q^{(0)})$$

$$T^{(1)} \xrightarrow{QR} = R^{(1)} \Rightarrow R^{(1)} \cdot Q^{(1)} = T^{(2)} (= Q^{(1)H} \cdot T^{(1)} \cdot Q^{(1)})$$

$$T^{(2)} \xrightarrow{QR} = R^{(2)} \Rightarrow R^{(2)} \cdot Q^{(2)} = T^{(3)} (= Q^{(2)H} \cdot T^{(2)} \cdot Q^{(2)})$$

$$T^{(3)} \xrightarrow{QR} = R^{(3)} \Rightarrow R^{(3)} \cdot Q^{(3)} = T^{(4)} (= Q^{(3)H} \cdot T^{(3)} \cdot Q^{(3)})$$

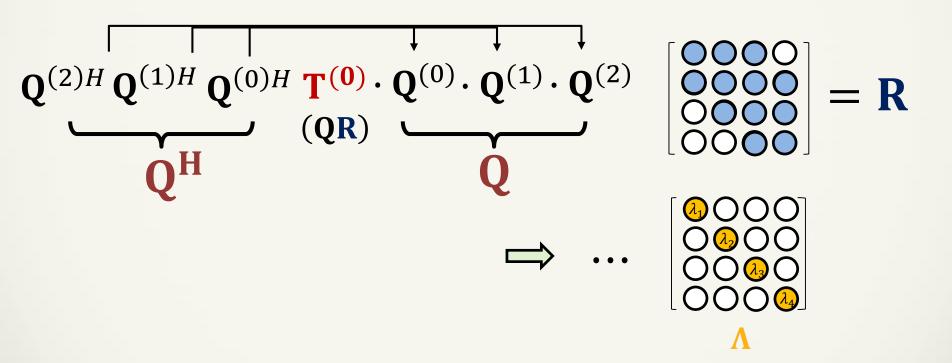
$$\vdots$$

$$T^{(\infty)} = Q^{(\infty)H} \qquad Q^{(1)H} \cdot Q^{(0)H} \cdot V \Lambda V^{-1} \qquad Q^{(0)} \qquad Q^{(\infty)}$$

$$= \Lambda$$

QR and RQ Operation

Row operation (left multiplication) and column operation (right multiplication)

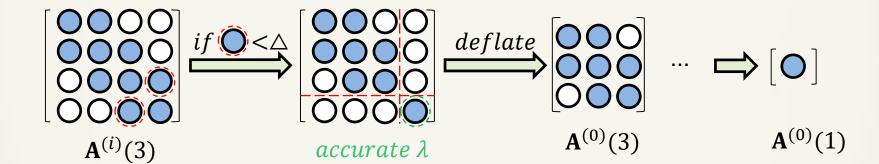


Shift

- In iterative QR algorithm, shift μ_i can be
 - Zero shift
 - $\blacksquare \mu_i = 0$
 - ■Slow convergence
 - Rayleigh quotient shift
 - $\blacksquare \mu_i = \mathbf{A}_{N,N}^{(i)}$, the lower right corner entry of $\mathbf{A}^{(i)}$
 - ■Medium convergence rate
 - Wilkinson shift
 - Eeigenvalue of the lower right 2x2 submatrix close to $\mathbf{A}_{N,N}^{(i)}$
 - ■Fast convergence

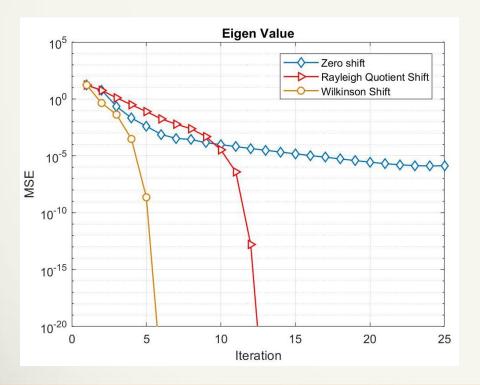
Deflation

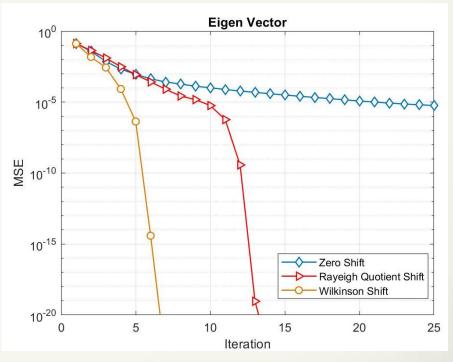
Deflate to accelerate convergence



Results

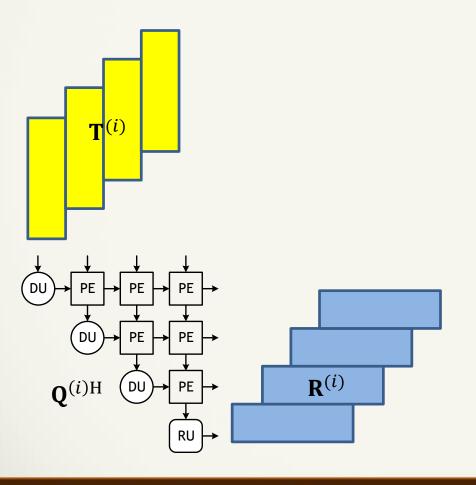
- Use 4x4 complex channel matrix to generate Hermitian
 Symmetric matrix A
- With deflation

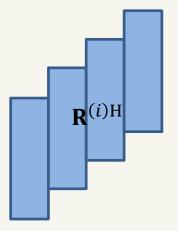


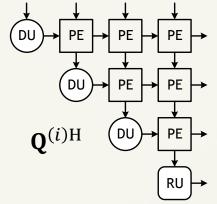


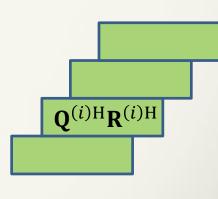
Architecture

Based on QR systolic array









Singular Value Decomposition

Singular Value Decomposition

- Jacobi algorithm
 - Real Two-sided Jacobi algorithm [2]
 - Complex Two-sided Jacobi algorithm [4]
- Two phase algorithm
 - Golub-Kahan algorithm [3]

Singular Value Decomposition

- For real/complex matrix $\mathbf{A} \in \mathbb{C}^{M \times N}$
 - $\blacksquare A = U\Sigma V^{H}$
 - $\mathbf{\Sigma} \in \mathbb{R}^{M \times N}$ is a real non-negative diagonal matrix with entries, known as singular values. Usually, they are arranged in descending order.

$$\mathbf{\Sigma} = \begin{bmatrix} \sigma_1 & 0 & & & & 0 \\ 0 & \sigma_1 & 0 & & & 0 \\ & 0 & \ddots & 0 & & 0 \\ 0 & & 0 & \sigma_M & 0 & 0 \end{bmatrix}$$

- $\blacksquare \mathbf{U} \in \mathbb{C}^{M \times M}$ is a real/complex unitary matrix. The column of \mathbf{U} is called left singular vector
- $\blacksquare V \in \mathbb{C}^{N \times N}$ is a real/complex unitary matrix. The column of V is called right singular vector

Real Two-Sided Jacobi Algorithm (1/2)

- Diagonalization
- For symmetric real matrix $\mathbf{A} \in \mathbb{R}^{2 \times 2}$

$$\mathbf{A} = \begin{bmatrix} a_{1,1} & a_{1,2} \\ a_{2,1} & a_{2,2} \end{bmatrix}$$

$$\mathbf{A} = \begin{bmatrix} a_{1,1} & a_{1,2} \\ a_{2,1} & a_{2,2} \end{bmatrix} \qquad \mathbf{J} = \begin{bmatrix} \cos\theta & \sin\theta \\ -\sin\theta & \cos\theta \end{bmatrix} = \begin{bmatrix} c & s \\ -s & c \end{bmatrix}$$

$$\mathbf{B} = \mathbf{J}^T \mathbf{A} \mathbf{J} = \begin{bmatrix} b_{1,1} & b_{1,2} \\ b_{2,1} & b_{2,2} \end{bmatrix} = \begin{bmatrix} \cos\theta & \sin\theta \\ -\sin\theta & \cos\theta \end{bmatrix}^T \begin{bmatrix} a_{1,1} & a_{1,2} \\ a_{2,1} & a_{2,2} \end{bmatrix} \begin{bmatrix} \cos\theta & \sin\theta \\ -\sin\theta & \cos\theta \end{bmatrix}$$

If
$$b_{1,2} = b_{2,1} = 0 = cs(a_{1,1} - a_{2,2}) + a_{1,2}(c^2 - s^2)$$

we have

$$\rho = \frac{a_{2,2} - a_{1,1}}{2a_{1,2}} \qquad tan\theta = \frac{sgn(\rho)}{|\rho| + \sqrt{1 + \rho^2}}$$

Real Two-Sided Jacobi Algorithm (2/2)

Algorithm: Real Two-Sided Jacobi

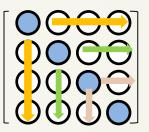
```
Given real symmetric matrix \mathbf{A} \in \mathbb{R}^{N \times N}
1. for sweep = 1: S
```

for p = 1: N - 1for q = p: N3.

2.

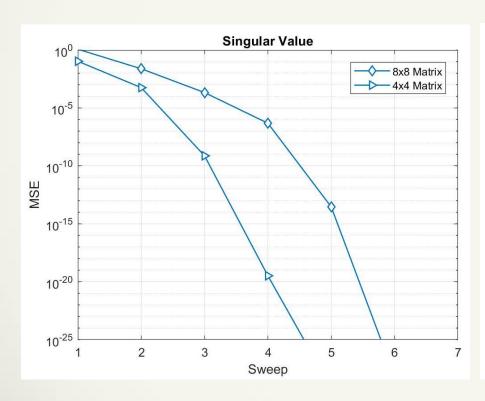
- $\rho = \frac{a_{q,q} a_{p,p}}{2a_{p,q}}$ 4.
- $\theta = \tan^{-1}\left(\frac{sgn(\rho)}{|\rho| + \sqrt{1 + \rho^2}}\right),\,$ 5.
- $\mathbf{J} = \mathbf{I}_N$, $\mathbf{J}_{p,p} = \mathbf{J}_{q,q} = \cos\theta$,
- $\mathbf{J}_{p,q} = -\mathbf{J}_{q,p} = \sin\theta$ 7.
- $\mathbf{A} = \mathbf{J}^T \mathbf{A} \mathbf{J}$ 8.
- 9. V = VI
- 10. end
- 11. end end

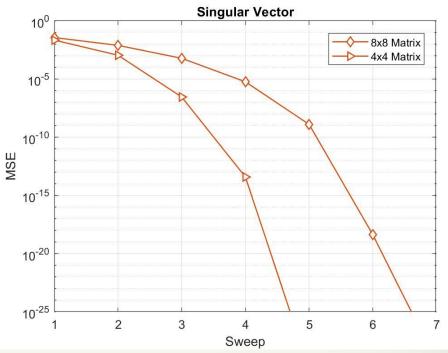
One sweep



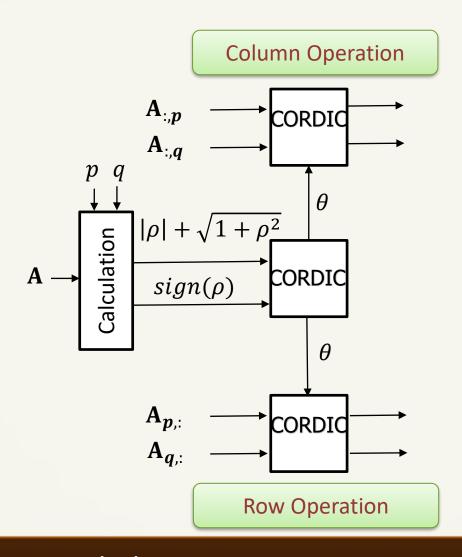
Performance

Fast convergence but large complexity for each sweep.





Operations of Real Two-Sided Jacobi Algorithm



Complex Two-Sided Jacobi Algorithm (1/3)

- Two-step two-sided rotation [4]
 - Step 1: upper-triangularization

$$\mathbf{C} = \begin{bmatrix} a_{1,1}e^{j\theta_{1,1}} & a_{1,2}e^{j\theta_{1,2}} \\ a_{2,1}e^{j\theta_{2,1}} & a_{2,2}e^{j\theta_{2,2}} \end{bmatrix}$$

$$\begin{bmatrix} a_{1,1}e^{j\theta_{1,1}} & a_{1,2}e^{j\theta_{1,2}} \\ a_{2,1}e^{j\theta_{2,1}} & a_{2,2}e^{j\theta_{2,2}} \end{bmatrix} \cdot \begin{bmatrix} e^{j\gamma_1} & 0 \\ 0 & e^{j\varepsilon_1} \end{bmatrix} \begin{bmatrix} \cos\varphi & \sin\varphi \\ -\sin\varphi & \cos\varphi \end{bmatrix} = \begin{bmatrix} \overline{a}_{1,1}e^{j\overline{\theta}_{1,1}} & \overline{a}_{1,2}e^{j\overline{\theta}_{1,2}} \\ 0 & \overline{a}_{2,2} \end{bmatrix}$$

$$\mathbf{C} \cdot \mathbf{D}(\mathbf{C}_{2,1}, \mathbf{C}_{2,2}) \mathbf{G}(\mathbf{C}_{2,1}, \mathbf{C}_{2,2}) = \overline{\mathbf{C}}$$

$$\gamma_1 = - heta_{2,1}$$
 , $\varepsilon_1 = - heta_{2,2}$ $\varphi = tan^{-1}(rac{a_{2,1}}{a_{2,2}})$

$$\begin{bmatrix} e^{j\alpha_1} & 0 \\ 0 & e^{j\beta_1} \end{bmatrix} \begin{bmatrix} \overline{a}_{1,1}e^{j\overline{\theta}_{1,1}} & \overline{a}_{1,2}e^{j\overline{\theta}_{1,2}} \\ 0 & \overline{a}_{2,2} \end{bmatrix} \begin{bmatrix} e^{j\alpha_2} & 0 \\ 0 & e^{j\beta_2} \end{bmatrix} = \begin{bmatrix} r_{1,1} & r_{1,2} \\ 0 & r_{2,2} \end{bmatrix}$$

$$\mathbf{D}_{l}(\bar{\mathbf{C}}_{1,1},\bar{\mathbf{C}}_{1,2})\bar{\mathbf{C}}\mathbf{D}_{r}(\bar{\mathbf{C}}_{1,1},\bar{\mathbf{C}}_{1,2}) = \mathbf{R}$$

$$\alpha_1 = -\frac{\bar{\theta}_{1,1} + \bar{\theta}_{1,2}}{2}$$

$$\beta_1 = \alpha_2 = -\beta_2 = \frac{\bar{\theta}_{1,2} - \bar{\theta}_{1,1}}{2}$$

Complex Two-Sided Jacobi Algorithm (2/3)

■ Step 2: diagonalization

$$\mathbf{R} = \begin{bmatrix} r_{1,1} & r_{1,2} \\ 0 & r_{2,2} \end{bmatrix}$$

$$\begin{bmatrix} \cos\theta_1 & \sin\theta_1 \\ -\sin\theta_1 & \cos\theta_1 \end{bmatrix} \cdot \begin{bmatrix} r_{1,1} & r_{1,2} \\ 0 & r_{2,2} \end{bmatrix} \cdot \begin{bmatrix} \cos\theta_2 & -\sin\theta_2 \\ \sin\theta_2 & \cos\theta_2 \end{bmatrix} = \begin{bmatrix} \boldsymbol{\delta_1} & 0 \\ 0 & \boldsymbol{\delta_2} \end{bmatrix}$$

$$G_l(R)RG_r(R) = \Sigma$$

$$\theta_1 = \frac{1}{2} \cdot \left(\tan^{-1} \left(\frac{r_{1,2}}{r_{1,1} - r_{2,2}} \right) + \tan^{-1} \left(\frac{-r_{1,2}}{r_{1,1} + r_{2,2}} \right) \right)$$

$$\theta_2 = \frac{1}{2} \cdot \left(\tan^{-1} \left(\frac{r_{1,2}}{r_{1,1} - r_{2,2}} \right) - \tan^{-1} \left(\frac{-r_{1,2}}{r_{1,1} + r_{2,2}} \right) \right)$$

Complex Two-Sided Jacobi Algorithm (3/3)

Algorithm: Complex Two-Sided Jacobi Given Complex matrix $\mathbf{C} \in \mathbb{C}^{N \times N}$ 1. for sweep = 1: S2. for p = 1: N - 1

 $\widetilde{\mathbf{D}} = \mathbf{I}_{N}$

Extend the dimension of

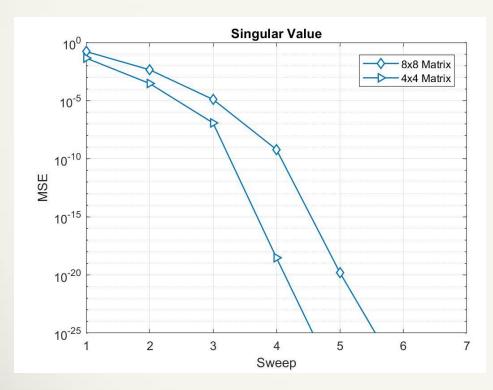
rotation matrixes

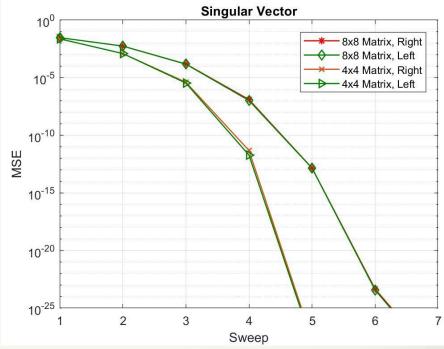
$$\widetilde{\mathbf{D}}_{[p,q],[p,q]}(\cdot,\cdot) = \mathbf{D}(\cdot,\cdot)$$

```
for q = p: N
3.
                                                        \overline{\mathbf{C}} = \mathbf{C} \cdot \widetilde{\mathbf{D}}(\mathbf{C}_{a,p}, \mathbf{C}_{a,a}) \widetilde{\mathbf{G}}(\mathbf{C}_{a,p}, \mathbf{C}_{a,a})
4.
                                                       \mathbf{V} = \mathbf{V}\widetilde{\mathbf{D}}(\mathbf{C}_{a.p}, \mathbf{C}_{a.a})\widetilde{\mathbf{G}}(\mathbf{C}_{a.p}, \mathbf{C}_{a.a})
5.
                                                        \mathbf{R} = \widetilde{\mathbf{D}}_{l}(\overline{\mathbf{C}}_{n,n}, \overline{\mathbf{C}}_{n,a})\overline{\mathbf{C}}\widetilde{\mathbf{D}}_{r}(\overline{\mathbf{C}}_{n,n}, \overline{\mathbf{C}}_{n,a})
6.
                                                       \mathbf{V} = \mathbf{V}\widetilde{\mathbf{D}}_r(\overline{\mathbf{C}}_{n.n}, \overline{\mathbf{C}}_{n.a}), \mathbf{U}^H = \widetilde{\mathbf{D}}_l(\overline{\mathbf{C}}_{p,p}, \overline{\mathbf{C}}_{p,q})\mathbf{U}^H
7.
                                                       \mathbf{C} = \tilde{\mathbf{G}}_l(\mathbf{R}_{[p,q],[p,q]})\mathbf{R}\tilde{\mathbf{G}}_r(\mathbf{R}_{[p,q],[p,q]})
8.
                                                       \mathbf{V} = \mathbf{V}\tilde{\mathbf{G}}_r(\mathbf{R}_{[n,a],[n,a]}), \mathbf{U}^H = \tilde{\mathbf{G}}_l(\mathbf{R}_{[n,a],[n,a]})\mathbf{U}^H
9.
10.
                                          end
11.
                           end
               end
```

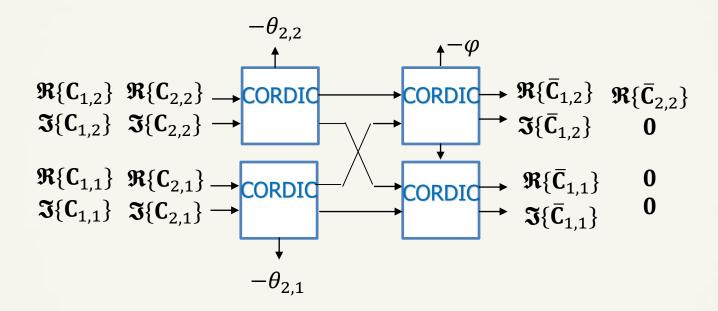
Performance

Fast convergence but large complexity for each sweep.



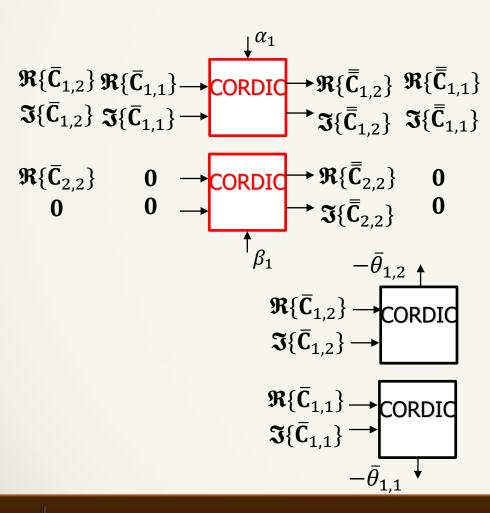


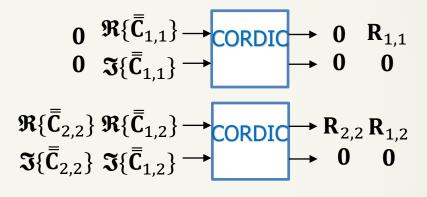
Operations of Complex Two-Sided Jacobi Algorithm (1/3)



$$\mathbf{C} \cdot \mathbf{D}(\mathbf{C}_{2,1}, \mathbf{C}_{2,2}) \mathbf{G}(\mathbf{C}_{2,1}, \mathbf{C}_{2,2}) = \overline{\mathbf{C}}$$

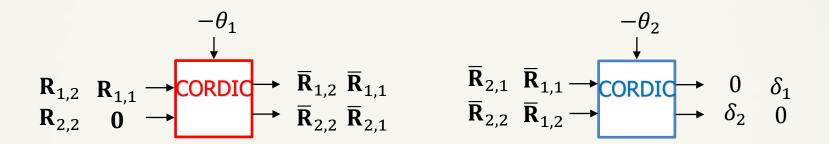
Operations of Complex Two-Sided Jacobi Algorithm (2/3)



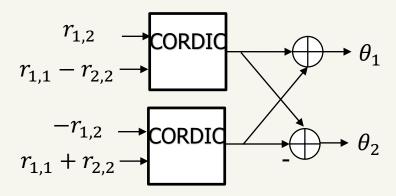


$$\mathbf{D}_{l}(\bar{\mathbf{C}}_{1,1},\bar{\mathbf{C}}_{1,2})\bar{\mathbf{C}}\mathbf{D}_{r}(\bar{\mathbf{C}}_{1,1},\bar{\mathbf{C}}_{1,2}) = \mathbf{R}$$

Operations of Complex Two-Sided Jacobi Algorithm (3/3)



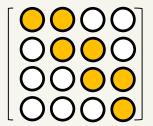
$$G_l(R)RG_r(R) = \Sigma$$



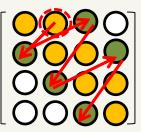
Two-Phase Algorithm

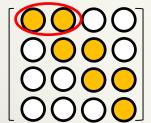
- Bidiagonalization
 - $\blacksquare H = QBP^H$
 - **B** is an upper bidiagonal matrix with real elements

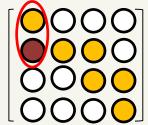
Real element

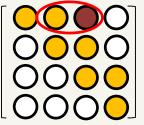


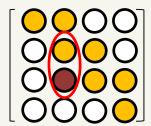
- Implicit QR algorithm (Golub-Kahan Algorithm)
 - Diagonalization
 - ■Chasing
 - ■Eliminating the bulge



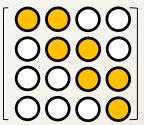








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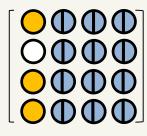
Bidiagonalization (1/2)

$$\begin{bmatrix} e^{-j\theta_1} & 0 & 0 & 0 \\ 0 & e^{-j\theta_2} & 0 & 0 \\ 0 & 0 & e^{-j\theta_3} & 0 \\ 0 & 0 & 0 & e^{-j\theta_4} \end{bmatrix}$$

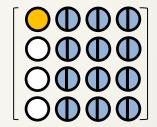
$$\begin{bmatrix} \Phi \Phi \Phi \Phi \\ \Phi \Phi \Phi \Phi \\ \Phi \Phi \Phi \Phi \end{bmatrix}$$

Complex element

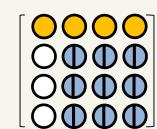
$$\left[\begin{array}{cccc} cos\theta_5 & sin\theta_5 & 0 & 0 \\ -sin\theta_5 & cos\theta_5 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{array}\right]$$





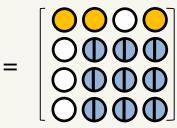


$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & e^{-j\theta_8} & 0 & 0 \\ 0 & 0 & e^{-j\theta_9} & 0 \\ 0 & 0 & 0 & e^{-j\theta_{10}} \end{bmatrix}$$



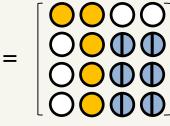
Bidiagonalization (2/2)

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & \cos\theta_{11} & -\sin\theta_{11} & 0 \\ 0 & \sin\theta_{11} & \cos\theta_{11} & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

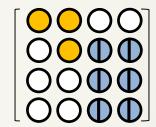




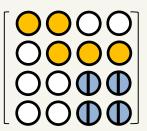
$$\left[egin{array}{ccccc} 1 & 0 & 0 & 0 & 0 \ 0 & e^{-j heta_{12}} & 0 & 0 & 0 \ 0 & 0 & e^{-j heta_{13}} & 0 & 0 \ 0 & 0 & 0 & e^{-j heta_{14}} \end{array}
ight]$$



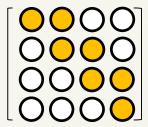












Golub-Kahan SVD Algorithm

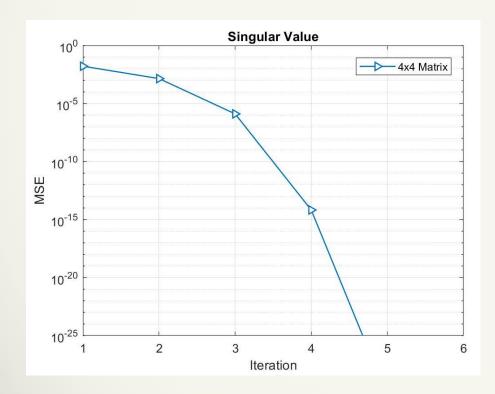
Algorithm: Golub-Kahan Algorithm

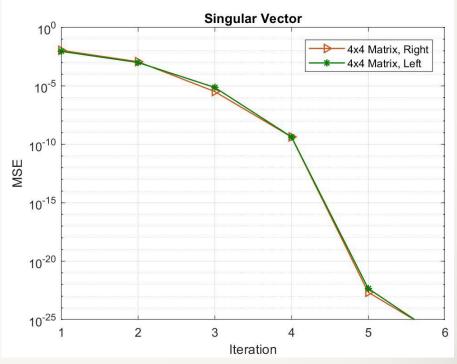
Given bidiagonal matrix $\mathbf{B} \in \mathbb{R}^{N \times N}$ with nonzero bidiagonal elements

```
1. While (N > 1)
            \mathbf{T} = \mathbf{B}^T \mathbf{B}, \mu is the eigenvalue of \mathbf{T}_{N-1:N,N-1:N} close to \mathbf{T}_{N,N}
2.
         for l = 1: N - 1
3.
                   if (l == 1) // implicit shifted QR
4.
                        \alpha = (\mathbf{B}_{11})^2 - \mu, \ \beta = \mathbf{B}_{11} \mathbf{B}_{12}
5.
6.
            else
7. \alpha = \mathbf{B}_{l-1,l}, \ \beta = \mathbf{B}_{l-1,l+1}
8.
                 end
       \theta = tan^{-1} \left( \frac{\beta}{\alpha} \right), \mathbf{B}_{:,l:l+1} = \mathbf{B}_{:,l:l+1} \begin{bmatrix} cos\theta & -sin\theta \\ sin\theta & cos\theta \end{bmatrix}
9.
        \alpha = \mathbf{B}_{l,l}, \beta = \mathbf{B}_{l+1,l}
10.
                  \theta = tan^{-1} \left(\frac{\beta}{\alpha}\right), \mathbf{B}_{l:l+1,:} = \begin{bmatrix} cos\theta & sin\theta \\ -sin\theta & cos\theta \end{bmatrix} \mathbf{B}_{l:l+1,:}
11.
12.
             end
            if |\mathbf{B}_{N-1,N}| < \delta // deflate
13.
                     \mathbf{B} = \mathbf{B}_{1:N-1,1:N-1}, N = N-1
14.
15.
              end
16. end
```

Performance of Golub-Kahan Algorithm

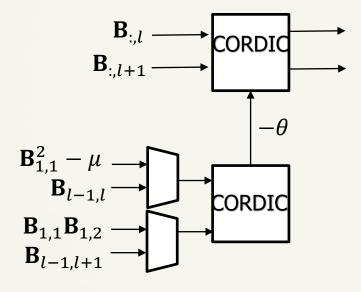
- The complexity of iteration is low because of sparse real elements. However, the complexity of bidiagonalization is high.
- Wilkinson shift can accelerate convergence.



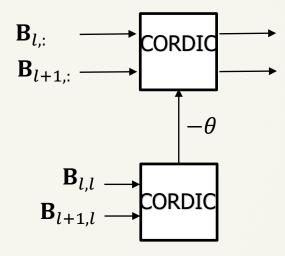


Operations of GK SVD

Column Operation



Row Operation



Summary

- Usually, iterative algorithms are required for EVD and SVD
- Shift and deflation is helpful to accelerate convergence.
- CORDIC is effective to handle Givens rotation for matrix decomposition.
- Tradeoff is required between performance and complexity.

Reference

- [1] Zheng-Yu Huang and Pei-Yun Tsai, "Efficient Implementation of QR Decomposition for Gigabit MIMO-OFDM Systems," IEEE Transactions on Circuits and Systems I: Regular paper, vol. 58, pp. 2531-2542, Oct. 2011.
- [2] R. Brent, F. Luk, and C. Van Loan, "Computation of the singular value decomposition using mesh-connected processors," Dept. Comput. Science, Cornell Univ., Ithaca, NY, USA Tech. Rep., 1983. [Online]. Available: http://hdl.handle.net/1813/6367
- [3] G. H. Golub and C. F. Van Loan, *Matrix Computation*, 4th ed. Baltimore, MD, USA: The Johns Hopkins Univ. Press, 2013.
- [4] C.-H. Yang, C.-W. Chou, C.-S. Hsu, and C.-E. Chen, "A systolic array based GTD processor with a parallel algorithm," *IEEE Trans. Circuits Syst. I, Reg. Papers*, vol. 62, no. 4, pp. 1099–1108, Apr. 2015.