CS321 Introduction to Theory of Computation Assignment No. 1, Due: Friday January 19, 2024

- 1. Prove that $\overline{S_1 \cup S_2} = \overline{S}_1 \cap \overline{S}_2$ where S_1 and S_2 are sets and \overline{S} is the complement of the set S.
- 2. A tree is a graph with no cycle. Show by induction that a tree with n nodes contains n-1 edges.
- 3. Prove by induction that the sum of the first k odd integers is equal to k^2 . For example, $1 = 1^2$, $1+3=4=2^2$, $1+3+5=9=3^2$, $1+3+5+7=16=4^2$, and so on. (Hint: The kth odd integer is 2k-1).
- 4. A rational number is of the form m/n where m and n are integers. For example, 2/3, 3/4, 2/5, 4/7, 3/8, 5/9, 11/18, 9/25 are some rational numbers. Show by contradiction that $\sqrt{2}$ is not a rational number.
- 5. Let the input symbols in a finite automata be $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$. Design a DFA that accepts all integers which are divisible by 3. (Hint: An integer is divisible by 3 if the sum of the digits is divisible by 3).
- 6. For this problem assume that the input symbols are {0, 1}. Design a DFA that accepts the binary string if it is divisible by 3.

1. Prove that $\overline{S}, \overline{US} = \overline{S}, \overline{\Lambda S}$, where S, and S_2 are sets and \overline{S} is the complement of the set S.

The complement \overline{S} contains all the elements not in S. If $X \in \overline{S}$, it means that $X \notin S$.

To proof SIUS = SINS.

if LI-IS = XE $S_1 U S_2$, it means that X is neither in S_1 nor in S_2 .

if RHS = $X \in S_1 \cap S_2$, it means X is in both \overline{S} and \overline{S}_2 , implying that X is not in S_1 and also not in S_2 .

Since shown both the subset and superset relationships. We can conclude that $\overline{S_1US_2} = \overline{S_1} \cap \overline{S_2}$, This completes the proof.

2. A tree is a graph with no cycle. Show by induction that a tree with n nodes contains n-1 edges.

Base Case: N=1

Consider a tree with a Single node. If there is no other nodes to connect to. there can't by any edges. Therefore, a tree with one node has 1-1=0 edges.

Inductive Step: Prove for n = k+1Assume a tree with k nodes has k-1 edges. It means a tree with k+1 nodes will have (k+1)-1 = k edge. Prove by in duction that the sum of the first k odd integers is equal to k^2 . For example, $l=l^2$, $l+3=4=2^2$, $l+3+5=9=3^2$, $l+3+5+7=16=4^2$, and so on. (Hint: the kth odd integer is 2k-1)

Base case: K=1

For k=1, the sum of the first k odd integer is 1. Which equals 1^2 .

This satisfies the base case.

Inductive Step:

Assume the statement is true for some positive integer K, that is, the Sum of the first K odd integer is K^2 . We need to prove that the Sum of the first K+1 odd integer $(K+1)^2$

The sum of the first k odd integers is k^2 the (k+1) th odd integer is 2(k+1)-1 Therefore, The sum of the first k+1 odd integers is: $k^2+(2(k+1)-1)$

$$=> k^2 + 2k + |$$

This completes the proof that the sum of the first k odd integers is indeed k^2 for all positive integers k.

4. A rational number is of the form m/n where m and n are integer.

For example, 3/3, 3/4, 3/5, 4/7, 3/8, 5/9, 11/18, 9/xs are some rational numbers. Show by contradiction that 12 is not a rational number.

Assume that
$$\sqrt{2}$$
 is a rational number
$$\sqrt{2} = \frac{m}{n}$$

$$= 2 = \frac{m^2}{n^2}$$

$$= 2n^2 = m^2$$

The equation (shows that m' is an even number, since it is

two times of some integer n²
if m² is even, then m is even. Therefore, m must be 2 times
of some integer k.

m = 21c

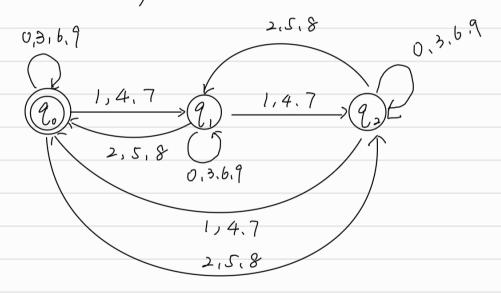
$$= 2n^2 = (2k)^2$$

$$=> 2n^2 = 4k^2$$

=>
$$n^2 = 2k^2$$

So The equation @ Shows that n' is also even. Therefore, n must also be even.

5. Let the input symbols in finite automata be 90,1,2,3,4,5,6,7,8,99. Design a DFA that accepts all integers which are divisible by 3.



6. For this problem assume that the input symbols are \$0,17, Design a DFA that accepts the binary string if it is divisible by 3.

