## CS321 Introduction to Theory of Computation Assignment No. 4, Due: Saturday February 24, 2024

- 1. Determine whether or not the following languages are regular. If the language is regular then give an NFA or regular expression for the language. Otherwise, use the pumping lemma for regular languages or closure properties to prove the language is not regular.
  - (a)  $L = \{a^n b^k : k \le n \le 2k\}$
  - (b)  $L = \{b^n a^k : n > 0, k > 0\} \cup \{a^n b^k : k > 0, n > 0\}$
  - (c)  $L = \{a^n : n = 3k \text{ for some } k > 0\}$
  - (d)  $L = \{a^n : n = k^3 \text{ for some } k \ge 0\}$
  - (e)  $L = \{w : n_a(w) > n_b(w), w \in \{a, b\}^*\}$
- 2. Give context-free grammars that generate the following language
  - (a)  $L_1 = \{ w \in \{0,1\}^* \mid \text{w contains at least three 1's } \}.$
  - (b)  $L_2 = \{a^i b^j c^k | i, j, k \ge 0, and \ i = j \text{ or } i = k\}$
- 3. Consider the following grammar  $G = (\{S, A\}, \{a, b\}, S, P\}$  where P is defined below.

$$S \to SS|AAA|\lambda$$
  
 $A \to aA|Aa|b$ 

- (a) Give a left-most derivation for the string abbaba.
- (b) Show that the grammar is ambiguous by exhibiting two distinct derivation trees for some terminal string.
- 4. Using the CYK algorithm show that the string baabba is in the context free language generated by the following production rules. Here,  $V = \{S, A, B, C, D\}$  and  $T = \{a, b\}$

$$S \to AB|BA$$

$$A \rightarrow AS|a$$

$$B \rightarrow BS|b$$

$$A \rightarrow BC$$

$$B \rightarrow AD$$

$$C \rightarrow AA$$

$$D \rightarrow BB$$

Show at least twol derivation trees.

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(R)

Assume that m=5 and w= a3b2, So IWI >5

$$\omega = \frac{aaa}{x} \frac{b}{Y} \frac{b}{z}$$

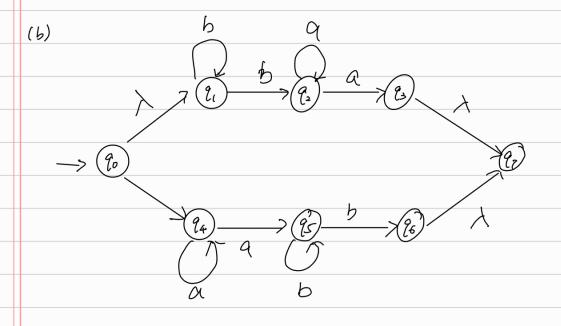
1xy155 and 1y121

$$W_1 = xy^2z = aaab^2b = a^3b^3$$

$$W_3 = XY^3Z = aaab^3b = a^3b^4$$
  $W_3$ 

So for all i , Wix L

So it's proven that it is not a regular language.



$$(\mathcal{C})$$

$$\Rightarrow (\mathcal{C})$$

$$\alpha \Rightarrow (\mathcal{C})$$

$$\alpha \Rightarrow (\mathcal{C})$$

$$W = \underbrace{aaaa}_{X} \underbrace{a}_{Y} \underbrace{aaa}_{Z}$$

1×41 < 8 141>1

Wi = XYiZ

Wz = xyz = aaaa aa aaa

= a <sup>q</sup> ∉ L , So W<sub>2</sub> G L

For all i, W; EL

So it's proven that it is not a regular language

Assume m=y, w= aaabb & L

iwl > y

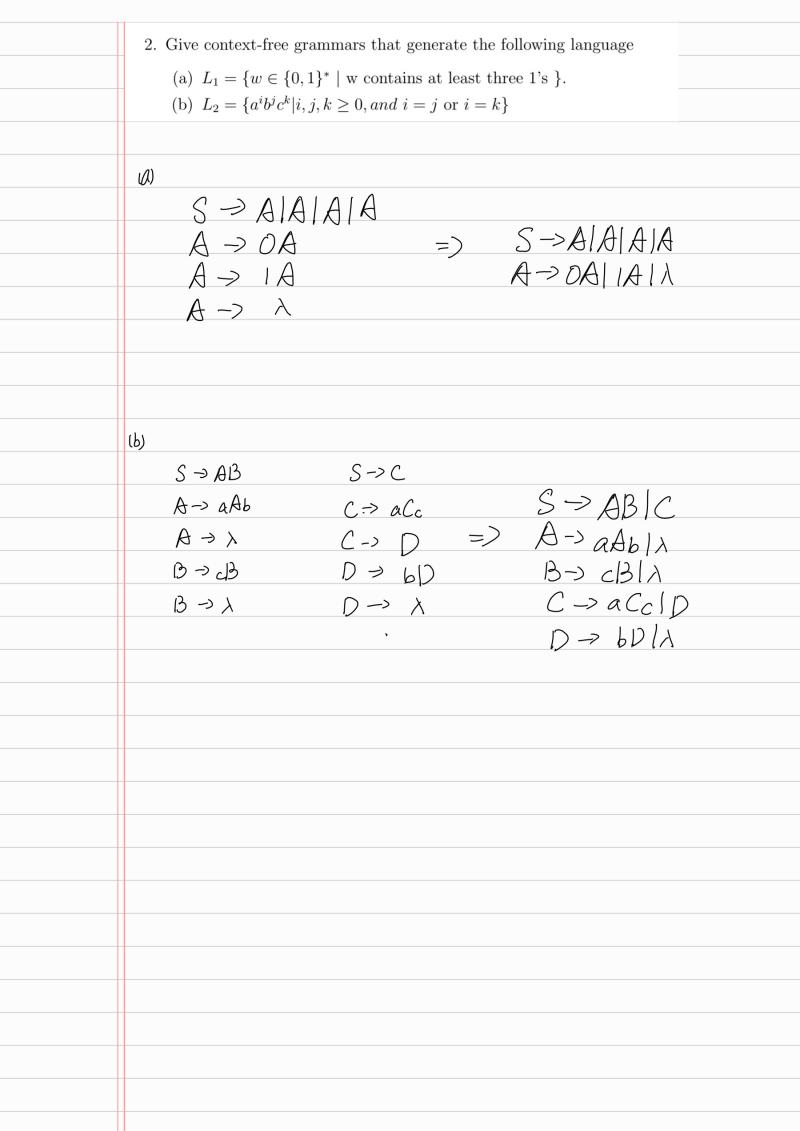
 $W = \underbrace{aaa}_{X} \underbrace{b}_{Y} \underbrace{b}_{Z}$ 

W; = XYiZ

 $U_2 = \chi y^2 Z = aaab^2b = a^3b^3 \qquad W_2 \not\in L$ 

So for all i, Wi & L

So it's proven that it is not a regular language.

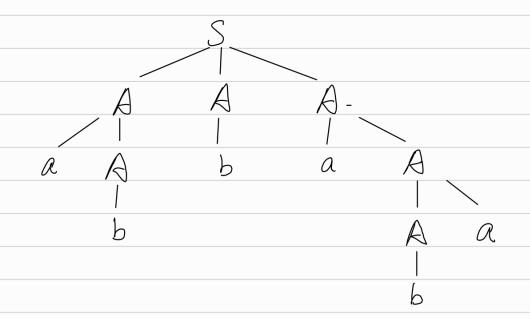


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## a. Left-most derivation for the string "abbaba"



S=> AAA => aAAA => abAA => abAaA => abbaA => abbaAa => abbaba

